

# Qualcomm<sup>®</sup> Hexagon<sup>™</sup> V73

## Programmer's Reference Manual

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# 1 Introduction

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The Qualcomm Hexagon™ processor is a general-purpose digital signal processor designed for high performance and low power across a wide variety of multimedia and modem applications. V73 is a member of the sixth generation of the Hexagon processor architecture.

## 1.1 Hexagon V73 processor architecture

### 1.1.1 Memory

The Hexagon processor features a unified byte-addressable memory. This memory has a single 32-bit virtual address space, which holds both instructions and data. It operates in little-endian mode.

The load/store architecture supports a complete set of addressing modes for both compiler code generation and DSP application programming.

#### 1.1.1.1 Cache memory

Memory accesses are cached or uncached. Separate L1 instruction and data caches exist for program code and data. A unified L2 cache is partly or wholly configured as tightly-coupled memory (TCM).

#### 1.1.1.2 Virtual memory

Memory is addressed virtually, with virtual-to-physical memory mapping handled by a resident OS. Virtual memory supports the implementation of memory management and memory protection in a hardware-independent manner.

### 1.1.2 Registers

The Hexagon processor has two sets of registers: [General registers](#) and [Control registers](#).

The general registers include thirty-two 32-bit registers (named R0 through R31), which are accessed either as single registers or as aligned 64-bit register pairs. The general registers contain all data, including pointer, scalar, vector, and accumulator data.

The control registers include special-purpose registers such as program counter, status register, loop registers, and so on.

### 1.1.3 Instruction sequencer

The instruction sequencer processes packets of one to four instructions in each cycle. If a packet contains more than one instruction, the instructions execute in parallel.

The instruction combinations allowed in a packet are limited to the instruction types that can execute in parallel in the four execution units (shown in [Figure 1-1](#)).

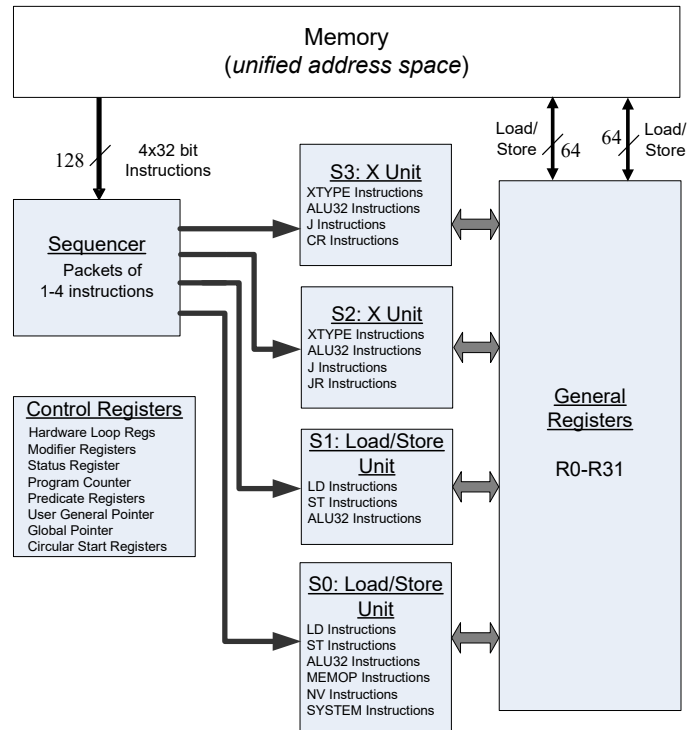


Figure 1-1 Hexagon V73 processor architecture

### 1.1.4 Execution units

The dual execution units are identical: each includes a 64-bit shifter and a vector multiply/accumulate unit with four  $16 \times 16$  multipliers to support both scalar and vector instructions.

These units also perform 32- and 64-bit ALU instructions, and jump and loop instructions.

**NOTE:** Each execution unit supports floating-point instructions.

### 1.1.5 Load/store units

The two load/store units can operate on signed or unsigned bytes, halfwords (16-bit), words (32-bit), or double words (64-bit).

To increase the number of instruction combinations allowed in packets, the load units also support 32-bit ALU instructions.

## 1.2 Instruction set

For the Hexagon processor to achieve large amounts of work per cycle, the instruction set is designed with the following properties:

- Static grouping (VLIW) architecture
- Static fusing of simple dependent instructions
- Extensive compound instructions
- A large set of SIMD and application-specific instructions

To support efficient compilation, the instruction set is designed to be orthogonal with respect to registers, addressing modes, and load/store access size.

### 1.2.1 Addressing modes

The Hexagon processor supports the following memory addressing modes:

- 32-bit absolute
- 32-bit absolute-set
- Absolute with register offset
- Global pointer relative
- Indirect
- Indirect with offset
- Indirect with register offset
- Indirect with auto-increment (immediate or register)
- Circular with auto-increment (immediate or register)
- Bit-reversed with auto-increment register

```
For example:
R2 = memw (##myvariable)
R2 = memw (R3=##myvariable)
R2 = memw (R4<<#3+##myvariable)
R2 = memw (GP+#200)
R2 = memw (R1)
R2 = memw (R3+#100)
R2 = memw (R3+R4<<#2)
R2 = memw (R3++#4)
R2 = memw (R0++M1)
R0 = memw (R2++#8:circ (M0) )
R0 = memw (R2++I:circ (M0) )
R2 = memw (R0++M1:brev)
```

Auto-increment with register addressing uses one of the two dedicated address-modify registers M0 and M1 (which are part of the control registers).

**NOTE:** Atomic memory operations (load locked/store conditional) are supported to implement multi-thread synchronization.

## 1.2.2 Program flow

The Hexagon processor supports zero-overhead hardware loops. For example:

```
loop0(start,#3)      // loop 3 times
start:
  { R0 = mpyi(R0,R0) } :endloop0
```

The loop instructions support nestable loops, with few restrictions on their use.

Software branches use a predicated branch mechanism. Explicit compare instructions generate a predicate bit, which is then tested by conditional branch instructions. For example:

```
P1 = cmp.eq(R2, R3)
if (P1) jump end
```

Jumps and subroutine calls are conditional or unconditional, and support both PC-relative and register indirect addressing modes. For example:

```
jump end
jumpr R1
call function
callr R2
```

The subroutine call instructions store the return address in register R31. Subroutine returns are performed using a jump indirect instruction through this register. For example:

```
jumpr R31      // Subroutine return
```

Two program flow instructions can be grouped into one packet.

## 1.2.3 Instruction pipeline

Pipeline hazards are resolved by the hardware: instruction scheduling is not constrained by pipeline restrictions.

## 1.3 Technical assistance

For assistance or clarification on information in this document, submit a case to Qualcomm Technologies, Inc. (QTI) at <https://createpoint.qti.qualco> For assistance or clarification on information in this document, open a technical support case at <https://support.qualcomm.com/>.

You will need to register for a Qualcomm ID account and your company must have support enabled to access our Case system.

Other systems and support resources are listed on <https://qualcomm.com/support>.

If you need further assistance, you can send an email to [qualcomm.support@qti.qualcomm.com](mailto:qualcomm.support@qti.qualcomm.com).

# 2 Registers

---

General registers are used for general-purpose computation, including address generation, and scalar and vector arithmetic.

Control registers support special-purpose processor features such as hardware loops and predicates.

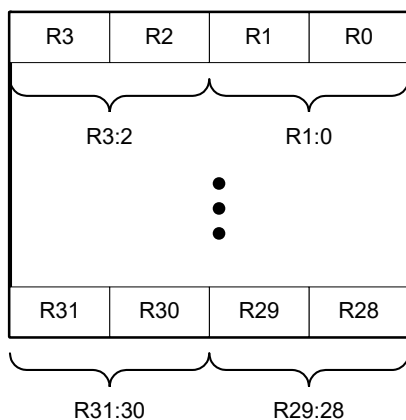
## 2.1 General registers

The Hexagon processor has thirty-two 32-bit general-purpose registers (named R0 through R31). These registers store operands in all of the instructions:

- Memory addresses for load/store instructions
- Data operands for arithmetic/logic instructions
- Vector operands for vector instructions

For example:

```
R1 = memh(R0)           // Load from address R0
R4 = add(R2,R3)         // Add
R28 = vaddh(R11,R10)    // Vector add halfword
```



**Figure 2-1 General registers**

## Aliased registers

Three of the general registers – R29 through R31 – support subroutines (Section 8.3.2) and the [Software Stack](#). The subroutine and stack instructions implicitly modify the registers. They have symbol aliases that indicate when these registers are accessed as subroutine and stack registers.

For example:

```
SP = add(SP, #-8)    // SP is alias of R29
allocframe          // Modifies SP (R29) and FP (R30)
call init           // Modifies LR (R31)
```

**Table 2-1 General register aliases**

Register	Alias	Name	Description
R29	SP	Stack pointer	Points to topmost element of stack in memory.
R30	FP	Frame pointer	Points to current procedure frame on stack.  Used by external debuggers to examine the stack and determine call sequence, parameters, local variables, and so on.
R31	LR	Link register	Stores return address of a subroutine call.

## Register pairs

The general registers can be specified as register pairs that represent a single 64-bit register. For example:

```
R1:0 = memd(R3)           // Load doubleword
R7:6 = valignb(R9:8,R7:6, #2) // Vector align
```

**NOTE:** The first register in a register pair must always be odd-numbered, and the second must be the next lower register.

**Table 2-1 General register pairs**

Register	Register pair
R0	R1:0
R1	
R2	R3:2
R3	
R4	R5:4
R5	
R6	R7:6
R7	
...	
R24	R25:24
R25	
R26	R27:26
R27	

**Table 2-1 General register pairs**

Register	Register pair
R28	R29:28
R29 (SP)	
R30 (FP)	R31:30 (LR:FP)
R31 (LR)	

## 2.2 Control registers

The Hexagon processor includes a set of 32-bit control registers that provide access to processor features such as the program counter, hardware loops, and vector predicates.

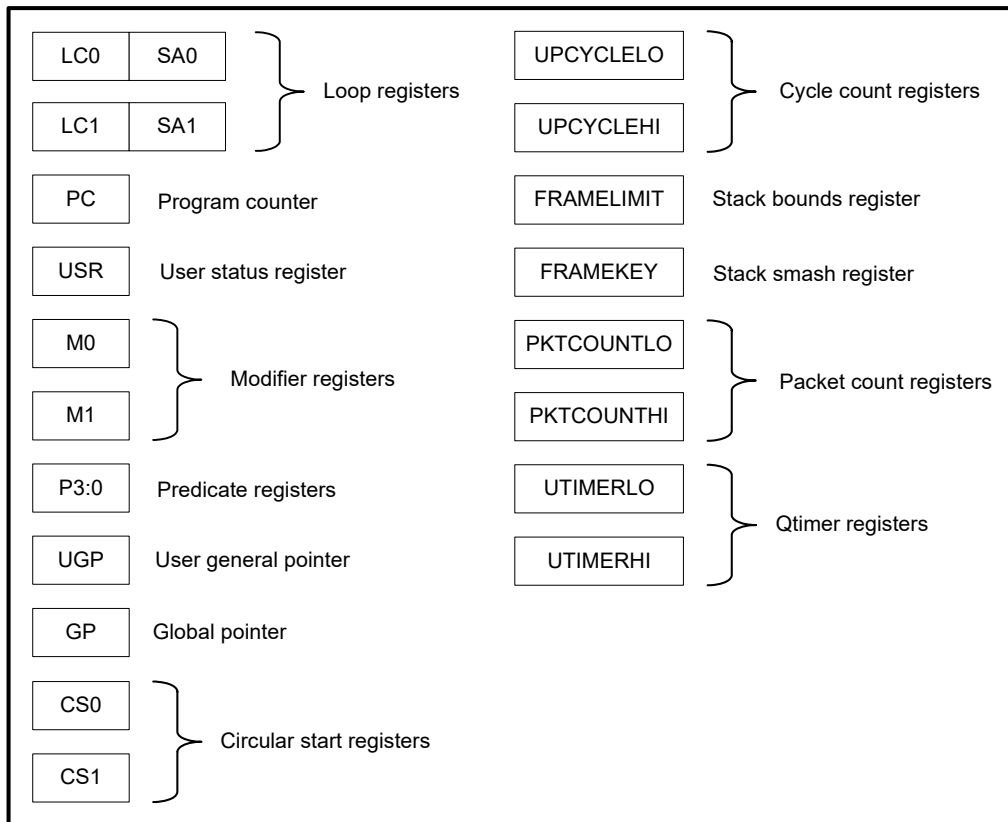
Unlike general registers, control registers are used as instruction operands only in the following cases:

- Instructions that require a specific control register as an operand
- Register transfer instructions

For example:

```
R2 = memw(R0++M1) // Auto-increment addressing mode (M1)
R9 = PC           // Get program counter (PC)
LC1 = R3          // Set hardware loop count (LC1)
```

**NOTE:** When a control register is used in a register transfer, the other operand must be a general register.



**Figure 2-2 Control registers**



## Aliased registers

The control registers have numeric aliases (C0 through C31).

**Table 2-2 Aliased control registers**

Register	Alias	Name
SA0	C0	Loop start address register 0
LC0	C1	Loop count register 0
SA1	C2	Loop start address register 1
LC1	C3	Loop count register 1
P3:0	C4	Predicate registers 3:0
reserved	C5	–
M0	C6	Modifier register 0
M1	C7	Modifier register 1
USR	C8	User status register
PC	C9	Program counter
UGP	C10	User general pointer
GP	C11	Global pointer
CS0	C12	Circular start register 0
CS1	C13	Circular start register 1
UPCYCLELO	C14	Cycle count register (low)
UPCYCLEHI	C15	Cycle count register (high)
UPCYCLE	C15:14	Cycle count register
FRAMELIMIT	C16	Frame limit register
FRAMEKEY	C17	Frame key register
PKTCOUNTLO	C18	Packet count register (low)
PKTCOUNTHI	C19	Packet count register (high)
PKTCOUNT	C19:18	Packet count register
reserved	C20-29	–
UTIMERLO	C30	Qtimer register (low)
UTIMERHI	C31	Qtimer register (high)
UTIMER	C31:30	Qtimer register

**NOTE:** The control register numbers (0 through 31) specify the control registers in [Instruction Encodings](#).

## Register pairs

The control registers can be specified as register pairs that represent a single 64-bit register. Control registers specified as pairs must use their numeric aliases. For example:

```
C1:0 = R5:4      // C1:0 specifies the LC0/SA0 register pair
```

**NOTE:** The first register in a control register pair must always be odd-numbered, and the second must be the next lower register.

**Table 2-2 Control register pairs**

Register	Register pair
C0	C1:0
C1	
C2	C3:2
C3	
C4	C5:4
C5	
C6	C7:6
C7	
...	
C30	C31:30
C31	

### 2.2.1 Program counter

The program counter (PC) register points to the next instruction packet to execute ([Section 3.3](#)). It is modified implicitly by instruction execution, but can be read directly. For example:

```
R7 = PC          // Get program counter
```

**NOTE:** The PC register is read-only: writing to it has no effect.

## 2.2.2 Loop registers

The Hexagon processor includes two sets of loop registers to support nested [Hardware loops](#). Each hardware loop is implemented with a pair of registers containing the loop count and loop start address. The loop registers are modified implicitly by the loop instruction, but can also be accessed directly. For example:

```
loop0(start, R4) // Modifies LC0 and SA0 (LC0=R4, SA0=&start)
LC1 = R22       // Set loop1 count
R9 = SA1        // Get loop1 start address
```

**Table 2-3 Loop registers**

Register	Name	Description
LC0, LC1	Loop count	Number of loop iterations to execute.
SA0, SA1	Loop start address	Address of first instruction in loop.

## 2.2.3 User status register

The user status register (USR) stores processor status and control bits that user programs can access. The status bits contain the status results of certain instructions, while the control bits contain user-settable processor modes for hardware prefetching. For example:

```
R9:8 = vaddw(R9:8, R3:2):sat // Vector add words
R6 = USR // Get saturation status
```

USR stores the following status and control values:

- [Cache prefetch](#) enable
- Cache prefetch status
- [Floating point](#) modes
- Floating point status
- Hardware loop configuration ([Section 8.2](#))
- Sticky [Saturation](#) overflow

**NOTE:** A user control register transfer to USR cannot be grouped in an instruction packet with a [Floating point](#) instruction.

When a transfer to USR changes the enable trap bits [29:25], an isync instruction ([Section 5.11](#)) must execute before the new exception programming can take effect.

**Table 2-4 User status register**

Name	RW	Bits	Field	Description
USR		32		User status register
	R	31	PFA	L2 prefetch active. 1: I2fetch instruction in progress 0: I2fetch finished (or inactive) Set when nonblocking I2fetch instruction is prefetching requested data. Remains set until I2fetch prefetch operation completes (or inactive).
	R	30	reserved	Return 0 if read. Reserved for future expansion. To remain compatible with future processor versions, software should always write this field with the same value read from the field.
	RW	29	FPINEE	Enable trap on IEEE inexact.
	RW	28	FPUNFE	Enable trap on IEEE underflow.
	RW	27	FPOVFE	Enable trap on IEEE overflow.
	RW	26	FPDBZE	Enable trap on IEEE divide-by-zero.
	RW	25	FPINVE	Enable trap on IEEE invalid.
	R	24	reserved	Reserved
	RW	23:22	FPRND	Rounding mode for floating point instructions. 00: Round to nearest, ties to even (default) 01: Toward zero 10: Downward (toward negative infinity) 11: Upward (toward positive infinity)
	R	21:20	reserved	Return 0 if read. Reserved for future expansion. To remain compatible with future processor versions, software should always write this field with the same value read from the field.
	R	19:18	reserved	Reserved
	R	17	reserved	Return 0 if read. Reserved for future expansion. To remain compatible with future processor versions, software should always write this field with the same value read from the field.
	RW	16:15	HFI	L1 instruction prefetch. 00: Disable 01: Enable (1 line) 10: Enable (2 lines)

**Table 2-4 User status register (cont.)**

Name	RW	Bits	Field	Description
	RW	14:13	HFD	L1 data cache prefetch. Four levels are defined from disabled to aggressive. Implementation defines how to interpret these levels. 00: disable 01: conservative 10: moderate 11: aggressive
	RW	12	PCMME	Enable packet counting in Monitor mode.
	RW	11	PCGME	Enable packet counting in Guest mode.
	RW	10	PCUME	Enable packet counting in User mode.
	RW	9:8	LPCFGE	Hardware loop configuration. Number of loop iterations (0 to 3) remaining before pipeline predicate should be set.
	R	7:6	reserved	Return 0 if read. Reserved for future expansion. To remain compatible with future processor versions, software should always write this field with the same value read from the field.
	RW	5	FPINPF	Floating-point IEEE inexact sticky flag.
	RW	4	FPUNFF	Floating-point IEEE underflow sticky flag.
	RW	3	FPOVFF	Floating-point IEEE overflow sticky flag.
	RW	2	FPDBZF	Floating-point IEEE divide-by-zero sticky flag.
	RW	1	FPINVF	Floating-point IEEE invalid sticky flag.
	RW	0	OVF	Sticky saturation overflow. 1: saturation occurred 0: no saturation Set when saturation occurs while executing an instruction that specifies optional saturation. Remains set until explicitly cleared by a <code>USR = Rs</code> instruction.

## 2.2.4 Modifier registers

The modifier registers (M0 to M1) are used in the following addressing modes.

### Indirect auto-increment

In [Indirect with auto-increment register](#) addressing the modifier registers store a signed 32-bit value that specifies the increment (or decrement) value.

**Table 2-5** Modifier registers used in indirect auto-increment addressing

Register	Name	Description
M0, M1	Increment	Signed auto-increment value.

### Circular

In circular addressing ([Section 5.8.10](#)) the modifier registers store the circular buffer length and related “I” values.

**Table 2-6** Modifier registers as used in circular addressing

Name	RW	Bits	Field	Description
M0, M1		32		Circular buffer specifier.
	RW	31:28	I[10:7]	I value (MSB - see <a href="#">Section 5.8.11</a> )
	RW	27:24		0x0
	RW	23:17	I[6:0]	I value (LSB)
	RW	16:0	Length	Circular buffer length

### Bit-reversed

In bit-reversed addressing ([Section 5.8.12](#)) the modifier registers store a signed 32-bit value that specifies the increment (or decrement) value.

**Table 2-7** Modifier registers as used in bit-reversed addressing

Register	Name	Description
M0, M1	Increment	Signed auto-increment value.

## 2.2.5 Predicate registers

The predicate registers (P0 through P3) store the status results of the scalar and vector compare instructions ([Chapter 6](#)). For example:

```
P1 = cmp.eq(R2, R3)      // Scalar compare
if (P1) jump end       // Jump to address (conditional)
R8 = P1                 // Get compare status (P1 only)
P3:0 = R4               // Set compare status (P0-P3)
```

The four predicate registers can be specified as a register quadruple (P3:0) that represents a single 32-bit register.

**NOTE:** Unlike the other control registers, the predicate registers are only 8 bits wide because vector compares return a maximum of eight status results.

**Table 2-3 Predicate registers**

Register	Bits	Description
P0, P1,P2, P3	8	Compare status results.
P3:0	32	Compare status results.
	31:24	P3 register
	23:16	P2 register
	15:8	P1 register
	7:0	P0 register

## 2.2.6 Circular start registers

The circular start registers (CS0 through CS1) store the start address of a circular buffer in circular addressing ([Section 5.8.10](#)). For example:

```
CS0 = R5 // Set circ start register
M0 = R7 // Set modifier register
R0 = memb(R2++#4:circ(M0)) // Load from circ buffer pointed
// to by CS0 with size/K vals in M0
```

**Table 2-8 Circular start registers**

Register	Name	Description
CS0, CS1	Circular start	Circular buffer start address.

## 2.2.7 User general pointer register

The user general pointer (UGP) register is a general-purpose control register. For example:

```
R9 = UGP // Get UGP
UGP = R3 // Set UGP
```

**NOTE:** UGP typically stores the address of thread local storage.

**Table 2-4 User general pointer register**

Register	Name	Description
UGP	User general pointer	General-purpose control register.

## 2.2.8 Global pointer

The global pointer (GP) is used in GP-relative addressing. For example:

```
GP = R7 // Set GP
R2 = memw(GP+#200) // GP-relative load
```

**Table 2-9 Global pointer register**

Name	R/W	Bits	Field	Description
GP		32		Global pointer register
	R/W	31:6	GDP	Global data pointer ( <a href="#">Section 5.8.4</a> ).
	R	5:0	reserved	Return 0 if read. Reserved for future expansion. To remain forward-compatible with future processor versions, software should always write this field with the same value read from the field.



## 2.2.9 Cycle count registers

The cycle count registers (UPCYCLELO to UPCYCLEHI) store a 64-bit value containing the current number of processor cycles executed since the Hexagon processor was last reset. For example:

```
R5 = UPCYCLEHI    // Get cycle count (high)
R4 = UPCYCLELO    // Get cycle count (low)
R5:4 = UPCYCLE    // Get cycle count
```

**NOTE:** The RTOS must grant permission to access these registers. Without this permission, reading these registers from user code returns zero.

**Table 2-5** Cycle count registers

Register	Name	Description
UPCYCLELO	Cycle count (low)	Processor cycle count (low 32 bits)
UPCYCLEHI	Cycle count (high)	Processor cycle count (high 32 bits)
UPCYCLE	Cycle count	Processor cycle count (64 bits)

## 2.2.10 Frame limit register

The frame limit register (FRAMELIMIT) stores the low address of the memory area reserved for the software stack ([Section 7.3.1](#)). For example:

```
R9 = FRAMELIMIT      // Get frame limit register
FRAMELIMIT = R3      // Set frame limit register
```

**Table 2-10** Frame limit register

Register	Name	Description
FRAMELIMIT	Frame limit	Low address of software stack area.

## 2.2.11 Frame key register

The frame key register (FRAMEKEY) stores the key value that XOR-scrambles return addresses when they are stored on the software stack ([Section 7.3.2](#)). For example:

```
R2 = FRAMEKEY        // Get frame key register
FRAMEKEY = R1        // Set frame key register
```

**Table 2-11** Frame key register

Register	Name	Description
FRAMEKEY	Frame key	Key to scramble return addresses stored on software stack.

## 2.2.12 Packet count registers

The packet count registers (PKTCOUNTLO to PKTCOUNTHI) store a 64-bit value containing the current number of instruction packets executed since a PKTCOUNT register was last written to. For example:

```
R9 = PKTCOUNTHI      // Get packet count (high)
R8 = PKTCOUNTLO      // Get packet count (low)
R9:8 = PKTCOUNT      // Get packet count
```

Packet counting can be configured to operate only in specific sets of processor modes (for example, User mode only, or Guest and Monitor modes only). Bits [12:10] in the [User status register](#) control the configuration for each mode.

Packets with exceptions are not counted as committed packets.

**NOTE:** Each hardware thread has its own set of packet count registers.

The RTOS must grant permission to access these registers. Without this permission, reading these registers from user code returns zero.

When a value is written to a PKTCOUNT register, the 64-bit packet count value is incremented before the value is stored in the register.

**Table 2-12 Packet count registers**

Register	Name	Description
PKTCOUNTLO	Packet count (low)	Processor packet count (low 32 bits)
PKTCOUNTHI	Packet count (high)	Processor packet count (high 32 bits)
PKTCOUNT	Cycle count	Processor packet count (64 bits)

## 2.2.13 Qtimer registers

The Qtimer registers (UTIMERLO to UTIMERHI) provide access to the Qtimer global reference count value. They enable Hexagon software to read the 64-bit time value without having to perform an expensive advanced high-performance bus (AHB) load. For example:

```
R5 = UTIMERHI    // Get Qtimer reference count (high)
R4 = UTIMERLO    // Get Qtimer reference count (low)
R5:4 = UTIMER    // Get Qtimer reference count
```

These registers are read-only – hardware automatically updates these registers to contain the current Qtimer value.

**NOTE:** The RTOS must grant permission to access these registers. Without this permission, reading these registers from user code returns zero.

**Table 2-13 Qtimer registers**

Register	Name	Description
UTIMERLO	Qtimer (low)	Qtimer global reference count (low 32 bits)
UTIMERHI	Qtimer (high)	Qtimer global reference count (high 32 bits)
UTIMER	Qtimer	Qtimer global reference count (64 bits)

# 3 Instructions

---

Instruction encoding is described in [Chapter 10](#). For detailed descriptions of the Hexagon processor instructions, see [Chapter 11](#).

## 3.1 Hexagon processor instruction syntax

**NOTE:** The notation described here does not appear in actual assembly language instructions. It is used only to specify the instruction syntax and behavior.

Most Hexagon processor instructions have the following syntax:

```
dest = instr_name(source1,source2,...)[:option1][:option2]...
```

The item specified on the left-hand side (LHS) of the equation is assigned the value specified by the right-hand side (RHS). For example:

```
R2 = add(R3,R1) // Add R3 and R1, assign result to R2
```

- `Courier font` is used for instructions
- Square brackets enclose optional items (for example, `[:sat]`, means that saturation is optional)
- Braces indicate a choice of items (for example, `{Rs,#s16}`, means that either `Rs` or a signed 16-bit immediate can be used)

**Table 3-1 Instruction symbols**

Symbol	Example	Meaning
=	R2 = R3	Assignment of RHS to LHS
#	R1 = #1	Immediate constant value
##	##2147483647	32-bit immediate constant value
0x	0xBABE	Hexadecimal number prefix
MEMxx	R2 = MEMxx(R3)	Access memory. xx specifies the size and type of access.
;	R2 = R3; R4 = R5;	Instruction delimiter, or end of instruction
{ ... }	{R2 = R3; R5 = R6}	Instruction packet delimiter; indicates a group of parallel instructions
( ... )	R2 = memw(R0 + #100)	Source list delimiter
:endloopX	:endloop0	Loop end X specifies loop instruction (0 or 1)

**Table 3-1 Instruction symbols (cont.)**

Symbol	Example	Meaning
:t	if (P0.new) jump:t target	Direction hint (jump taken)
:nt	if (!P1.new) jump:nt target	Direction hint (jump not taken)
:sat	R2 = add(R1,R2):sat	Saturate result
:rnd	R2 = mpy(R1.H,R2.H):rnd	Round result
:carry	R5:4=add(R1:0,R3:2,P1):carry	Predicate used as carry input and output
:<<16	R2 = add(R1.L,R2.L):<<16	Shift result left by halfword
:mem_noshuf	{memw(R5) = R2; R3 = memh(R6)}:mem_noshuf	Inhibit load/store reordering ( <a href="#">Section 5.5</a> )

### 3.1.1 Numeric operands

[Table 3-2](#) lists the notation that describes numeric operands in the syntax and behavior of instructions:

**Table 3-2 Instruction operands**

Symbol	Meaning	Min	Max	Example
#uN	Unsigned N-bit immediate value	0	$2^N-1$	–
#sN	Signed N-bit immediate value	$-2^{N-1}$	$2^{N-1}-1$	–
#mN	Signed N-bit immediate value	$-(2^{N-1}-1)$	$2^{N-1}-1$	–
#uN:S	Unsigned N-bit immediate value representing integral multiples of $2^S$ in specified range	0	$(2^N-1) \times 2^S$	–
#sN:S	Signed N-bit immediate value representing integral multiples of $2^S$ in specified range	$(-2^{N-1}) \times 2^S$	$(2^{N-1}-1) \times 2^S$	–
#rN:S	Same as #sN:S, but value is offset from PC of current packet	$(-2^{N-1}) \times 2^S$	$(2^{N-1}-1) \times 2^S$	–
usat <sub>N</sub>	Saturate value to unsigned N-bit number	0	$2^N-1$	usat <sub>16</sub> (Rs)
sat <sub>N</sub>	Saturate value to signed N-bit number	$-2^{N-1}$	$2^{N-1}-1$	sat <sub>16</sub> (Rs)
sxt x->y	Sign-extend value from x to y bits	–	–	sxt32->64(Rs)
zxt x->y	Zero-extend value from x to y bits	–	–	zxt32->64(Rs)
>>>	Logical right shift	–	–	Rss >>> offset

The #uN, #sN, and #mN symbols specify immediate operands in instructions. The # symbol appears in the actual instruction to indicate the immediate operand.

The #rN symbol specifies loop and branch destinations in instructions. The # symbol does not appear in the actual instruction; instead, the entire #rN symbol (including its :s suffix) is expressed as a loop or branch symbol whose numeric value is determined by the assembler and linker. For example:

```
call my_proc          // instructi
ple
```

The `:s` suffix indicates that the `s` least-significant bits in a value are implied zero bits and therefore not encoded in the instruction. The implied zero bits are called scale bits.

For example, `#s4:2` denotes a signed immediate operand represented by four bits encoded in the instruction, and two scale bits. The possible values for this operand are -32, -28, -24, -20, -16, -12, -8, -4, 0, 4, 8, 12, 16, 20, 24, and 28.

The `##` symbol specifies a 32-bit immediate operand in an instruction (including a loop or branch destination). The `##` symbol indicates the operand in the actual instruction.

Examples of operand symbols:

```
Rd = add(Rs,#s16)      // #s16   -> signed 16-bit imm value
Rd = memw(Rs++#s4:2)  // #s4:2  -> scaled signed 4-bit imm value
call #r22:2           // #r22:2 -> scaled 22-bit PC-rel addr value
Rd = ##u32           // ##u32  -> unsigned 32-bit imm value
```

**NOTE:** When an instruction contains more than one immediate operand, the operand symbols are specified in upper and lower case (for example, `#uN` and `#UN`) to indicate where they appear in the instruction encodings

## 3.1.2 Terminology

[Table 3-3](#) lists the symbols Hexagon processor instruction names use to specify the supported data types.

**Table 3-3 Data symbols**

Size	Symbol	Type
8-bit	B	Byte
8-bit	UB	Unsigned byte
16-bit	H	Half word
16-bit	UH	Unsigned half word
32-bit	W	Word
32-bit	UW	Unsigned word
64-bit	D	Double word

### 3.1.3 Register operands

The following notation describes register operands in the syntax and behavior of instructions:

```
Rds[.elst]
```

The ds field indicates the register operand type and bit size (as defined in [Table 3-4](#)).

**Table 3-4 Register symbols**

Symbol	Operand type	Size (in bits)
d	Destination	32
dd		64
s	First source	32
ss		64
t	Second source	32
tt		64
u	Third source	32
uu		64
x	Source <i>and</i> destination	32
xx		64

Examples of ds field (describing instruction syntax):

```
Rd = neg(Rs)           // Rd -> 32-bit dest, Rs 32-bit source
Rd = xor(Rs,Rt)        // Rt -> 32-bit second source
Rx = insert(Rs,Rtt)   // Rx -> both source and dest
```

Examples of ds field (describing instruction behavior):

```
Rdd = Rss + Rtt       // Rdd, Rss, Rtt -> 64-bit registers
```

The optional elst (element size and type) field specifies parts of a register when the register is used as a vector. It can specify the following values:

- A signed or unsigned byte, halfword, or word within the register (as defined in [Figure 3-1](#))
- A bit field within the register (as defined in [Table 3-5](#))

Examples of elst field:

```
EA = Rt.h[1]          // .h[1] -> bit field 31:16 in Rt
Pd = (Rss.u64 > Rtt.u64) // .u64 -> unsigned 64-bit value
Rd = mpyu(Rs.L,Rt.H)  // .L/.H -> low/high 16-bit fields
```

**NOTE:** The control and predicate registers use the same notation as the general registers, but are written as Cx and Px (respectively) instead of Rx.

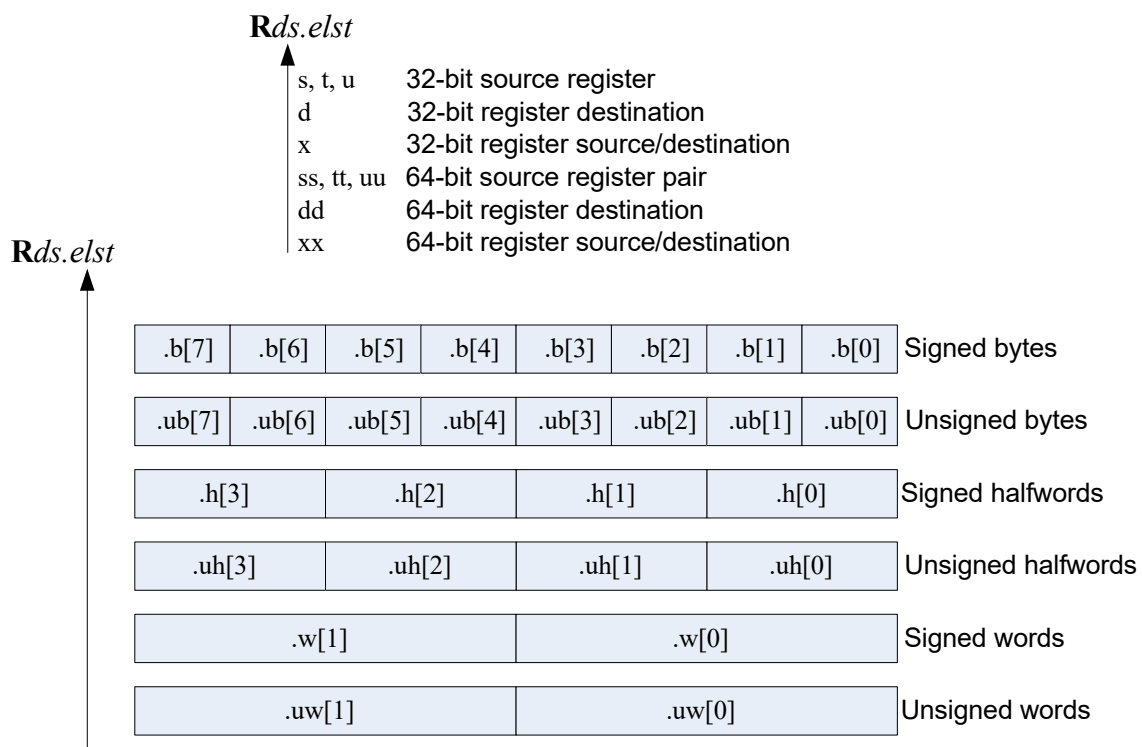


Figure 3-1 Register field symbols

Table 3-5 Register bit field symbols

Symbol	Meaning
.sN	Bits [N-1:0] are treated as a N-bit signed number. For example, R0.s16 means that the least significant 16-bits of R0 are treated as a 16-bit signed number.
.uN	Bits [N-1:0] are treated as a N-bit unsigned number.
.H	The most-significant 16 bits of a 32-bit register.
.L	The least-significant 16 bits of a 32-bit register.



## 3.2 Instruction classes

The Hexagon processor instructions are assigned to specific instruction classes. Classes determine the combinations of instructions that can be written in parallel (Section 3.3). presents an overview of the instruction classes and how they can be grouped together.

Instruction classes logically correspond with instruction types, so they serve as mnemonics for looking up specific instructions. For instance, the ALU32 class contains ALU instructions that operate on 32-bit operands.

**Table 3-6 Instruction classes and subclasses**

Class	Subclass	Description
XTYPE	–	Various operations
	XTYPE ALU	64-bit ALU operations
	XTYPE BIT	Bit operations
	XTYPE COMPLEX	Complex math (using real and imaginary numbers)
	XTYPE FP	Floating point operations
	XTYPE MPY	Multiply operations
	XTYPE PERM	Vector permute and format conversion (pack, splat, swizzle)
	XTYPE PRED	Predicate operations
ALU32	–	32-bit ALU operations
	ALU32 ALU	Arithmetic and logical
	ALU32 PERM	Permute
	ALU32 PRED	Predicate operations
CR	–	Control register access, loops
JR	–	Jumps (register indirect addressing mode)
J	–	Jumps (PC-relative addressing mode)
LD	–	Memory load operations
MEMOP	–	Memory operations
NV	–	New-value operations
	NV J	New-value jumps
	NV ST	New-value stores
ST	–	Memory store operations; allocate stack frame
SYSTEM	–	Operating system access
	SYSTEM USER	Application-level access

## 3.3 Instruction packets

Instructions can be grouped into very long instruction word (VLIW) packets for parallel execution, with each packet containing from one to four instructions. Packets of varying length can be freely mixed in a program.

Vector instructions operate on single instruction multiple data (SIMD) vectors.

Instruction packets must be explicitly specified in software. They are expressed in assembly language by enclosing groups of instructions in curly braces.

For example, two instructions grouped in a packet:

```
{ R0 = R1; R2 = R3 }
```

Four instructions grouped in a packet:

```
{ R8 = memh(R3++#2);  
  R12 = memw(R1++#4);  
  R = mpy(R10, R6) :<<1:sat;  
  R7 = add(R9, #2);  
}
```

Packets have various restrictions on the allowable instruction combinations. The primary restriction is determined by the instruction class of the instructions in a packet. In particular, packet formation is subject to the following constraints:

- **Resource constraints** determine how many instructions of a specific type can appear in a packet. The Hexagon processor has a fixed number of execution units: each instruction executes on a particular type of unit, and each unit can process at most one instruction at a time. Thus, for example, because the Hexagon processor contains only two load units, an instruction packet with three load instructions is invalid.
- **Grouping constraints** are a small set of rules that apply above and beyond the resource constraints.
- **Dependency constraints** ensure that no write-after-write hazards exist in a packet.
- **Ordering constraints** dictate the ordering of instructions within a packet.
- **Alignment constraints** dictate the placement of packets in memory.

**NOTE:** The Hexagon processor executes individual instructions (which are not explicitly grouped in packets) as packets containing a single instruction.

### 3.3.1 Packet execution semantics

Packets are defined to have parallel execution semantics. The execution behavior of a packet is defined as follows:

- First, instructions in the packet read their source registers in parallel.
- Next, instructions in the packet execute.
- Finally, instructions in the packet write their destination registers in parallel.

For example, consider the following packet:

```
{ R2 = R3; R3 = R2; }
```

In the first phase, registers R3 and R2 are read from the register file. Then, after execution, R2 is written with the old value of R3 and R3 is written with the old value of R2. The result of this packet is the swap of the values of R2 and R3.

**NOTE:** [Dual stores](#), [Dual jumps](#), [New-value stores](#), [New-value compare jumps](#), and [Dot-new predicates](#) have non-parallel execution semantics.

### 3.3.2 Sequencing semantics

Packets of any length can freely mix in code. A packet is considered an atomic unit: in essence, a single large instruction. From the program perspective, a packet either executes to completion or not at all; it never partially executes. For example, if a packet causes a memory exception, the exception point is established before the packet.

A packet containing multiple load/store instructions can require service from the external system. For instance, consider a packet that performs two load operations that both miss in the cache. The packet requires the memory system to supply the data:

- From the memory system perspective, the two resulting load requests are processed serially.
- From the program perspective, however, both load operations must complete before the packet can complete.

Thus, the packet is atomic from the program perspective.

Packets have a single PC address, which is the address of the start of the packet. Branches cannot be performed into the middle of a packet.

Architecturally, packets execute to completion – including updating all registers and memory – before the next packet begins. As a result, application programs are not exposed to any pipeline artifacts.

### 3.3.3 Resource constraints

A packet cannot use more hardware resources than are physically available on the processor. For instance, because the Hexagon processor has only two load units, a packet with three load instructions is invalid. The behavior of such a packet is undefined. The assembler automatically rejects packets that oversubscribe the hardware resources.

The processor supports up to four parallel instructions. The instructions are executed in four parallel pipelines, which are referred to as slots. The four slots are named Slot 0, Slot 1, Slot 2, and Slot 3.

**NOTE:** The `endloopN` instructions ([Section 8.2.2](#)) do not use any slots.

Each instruction belongs to specific [Instruction classes](#). For example, jumps belong to instruction class J, while loads belong to instruction class LD. An instruction's class determines which slot it can execute in.

Figure 3-2 shows which instruction classes can be assigned to each of the four slots.

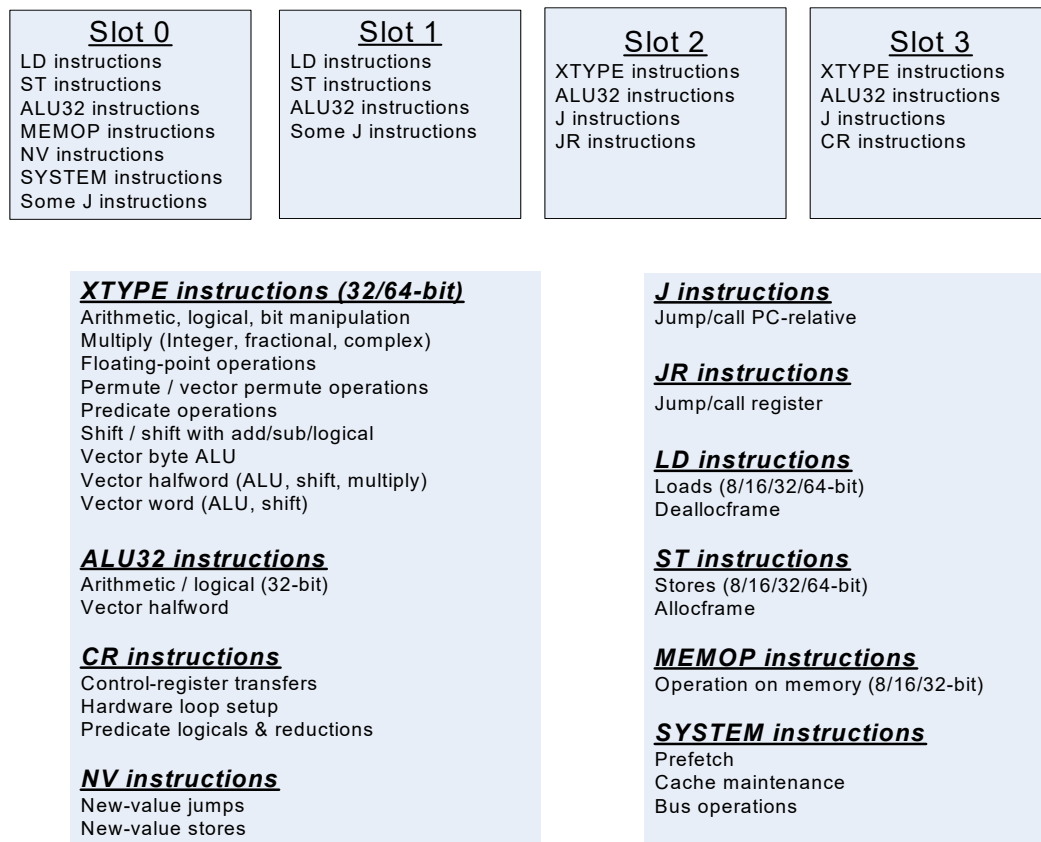


Figure 3-2 Packet grouping combinations

### 3.3.4 Grouping constraints

A small number of restrictions determines what constitutes a valid packet. The assembler ensures that packets follow valid grouping rules. If a packet executes that violates a grouping rule, the behavior is undefined. The following rules must be followed:

- Dot-new conditional instructions ([Section 6.1.4](#)) must be grouped in a packet with an instruction that generates dot-new predicates.
- ST-class instructions can be placed in Slot 1. In this case Slot 0 normally must contain a second ST-class instruction ([Section 5.4](#)).
- J-class instructions can be placed in Slots 2 or 3. However, only certain combinations of program flow instructions (J or JR) can be grouped together in a packet ([Section 8.7](#)). Otherwise, at most one program flow instruction is allowed in a packet. Some Jump and compare-Jump instructions can execute on slots 0 or 1, excluding calls, such as the following:
  - Instructions of the form `"Pd=cmp.xx(); if(Pd.new) jump:hint <target>"`
  - Instructions of the form `"If(Pd[.new]) jump[:hint] <target>"`
  - The `"jump<target>"` instruction

- JR-class instructions can be placed in Slot 2. However, when encoded in a duplex jump instruction, R31 can be placed in Slot 0 (Section 10.3).
- Restrictions limit the instructions that can appear in a packet at the setup or end of a hardware loop (Section 8.2.4).
- A user control register transfer to the control register `USR` cannot be grouped with a floating point instruction (Section 2.2.3).
- The SYSTEM-class instructions include prefetch, cache operations, bus operations, load locked, and store conditional instructions (Section 5.10). These instructions have the following grouping rules:
  - The `brkpt`, `trap`, `pause`, `icinva`, `isync`, and `syncht` instructions are solo instructions. They must not be grouped with other instructions in a packet.
  - The `memw_locked`, `memd_locked`, `l2fetch`, and `trace` instructions must execute on Slot 0. They must be grouped only with ALU32 or (non-FP) XTYPE instructions.
  - The `dccleana`, `dcinva`, `dccleaninva`, and `dczeroa` instructions must execute on Slot 0. Slot 1 must be empty or an ALU32 instruction.

### 3.3.5 Dependency constraints

Instructions in a packet cannot write to the same destination register. The assembler automatically flags such packets as invalid. If the processor executes a packet with two writes to the same general register, an error exception is raised.

If the processor executes a packet that performs multiple writes to the same predicate or control register, the behavior is undefined. Three special cases exist for this rule:

- Conditional writes are allowed to target the same destination register only if at most one of the writes is actually performed (Section 6.1.5).
- The overflow flag in the status register has defined behavior when multiple instructions write to it (Section 2.2.3). Do not group instructions that write to the entire user status register (for example, `USR = R2`) in a packet with any instruction that writes to a bit in the user status register.
- Multiple compare instructions are allowed to target the same predicate register to perform a logical AND of the results (Section 6.1.3).

### 3.3.6 Ordering constraints

In assembly code, instructions can appear in a packet in any order (with the exception of [Dual jumps](#)). The assembler automatically encodes instructions in the packet in the proper order.

In the binary encoding of a packet, the instructions must be ordered from Slot 3 down to Slot 0. If the packet contains less than four instructions, any unused slot is skipped – a NOP is unnecessary as the hardware handles the proper spacing of the instructions.

In memory, instructions in a packet must appear in strictly decreasing slot order. Additionally, if an instruction can go in a higher-numbered slot, and that slot is empty, it must be moved into the higher-numbered slot.

For example, if a packet contains three instructions and slot 1 is not used, encode the instructions in the packet as follows:

- Slot 3 instruction at lowest address
- Slot 2 instruction follows Slot 3 instruction
- Slot 0 instructions at the last (highest) address

If a packet contains a single load or store instruction, that instruction must go in Slot 0, which is the highest address. As an example, a packet containing both LD and ALU32 instructions must be ordered so the LD is in Slot 0 and the ALU32 in another slot.

### 3.3.7 Alignment constraints

Packets have the following constraints on their placement or alignment in memory:

- Packets must be word-aligned (32-bit). If the processor executes an improperly aligned packet, it raises an error exception ([Section 8.10](#)).
- Packets should not wrap the 4 GB address space. If address wraparound occurs, the processor behavior is undefined.

No other core-based restrictions exist for code placement or alignment.

If the processor branches to a packet that crosses a 16-byte address boundary, the resulting instruction fetch stalls for one cycle. Packets that are jump targets or loop body entries can be explicitly aligned to ensure this does not occur ([Section 8.3.5](#)).

## 3.4 Instruction intrinsics

To support efficient coding of the time-critical sections of a program (without resorting to assembly language), the C compilers support intrinsics that directly express Hexagon processor instructions from within C code.

For example:

```
int main()
{
    long long v1 = 0xFFFF0000FFFF0000LL;
    long long v2 = 0x0000FFFF0000FFFFLL;
    long long result;

    // Find the minimum for each half-word in 64-bit vector
    result = Q6_P_vminh_PP(v1,v2);
}
```

Intrinsics are defined for most of the Hexagon processor instructions.

## 3.5 Compound instructions

The Hexagon processor supports compound instructions, which encode pairs of common operations in a single instruction. For example, each of the following is a single compound instruction:

```
dealloc_return          // Deallocate frame and return
R2 &= and(R1, R0)      // And and and
R7 = add(R4, sub(#15, R3)) // Subtract and add
R3 = sub(#20, asl(R3, #16)) // Shift and subtract
R5 = add(R2, mpyi(#8, R4)) // Multiply and add
{
    P0 = cmp.eq (R2, R5) // Compare and jump
    if (P0.new) jump:nt target
}
{
    R2 = #15 // Register transfer and jump
    jump target
}
```

Compound instructions reduce code size and improves code performance.

**NOTE:** Compound instructions (with the exception of X-and-jump, as shown above) have distinct assembly syntax from the instructions they are composed of.

## 3.6 Duplex instructions

To reduce code size the Hexagon processor supports duplex instructions, which encode pairs of common instructions in a 32-bit instruction container.

Unlike [Compound instructions](#), duplex instructions do not have distinctive syntax – in assembly code they appear identical to the instructions they are composed of. The assembler is responsible for recognizing when a pair of instructions can be encoded as a single duplex rather than a pair of regular instruction words.

To fit two instructions into a single 32-bit word, [Duplexes](#) are limited to a subset of the most common instructions (load, store, branch, ALU), and the most common register operands.

# 4 Data Processing

---

The Hexagon processor provides a rich set of operations for processing scalar and vector data. Instructions can perform a wide variety of operations on fixed-point or floating-point data. The fixed-point operations support scalar and vector data in a variety of sizes. The floating-point operations support single-precision data.

This chapter presents an overview of the operations provided by the following Hexagon processor instruction classes:

- [XTYPE](#) – General-purpose data operations
- [ALU32](#) – Arithmetic/logical operations on 32-bit data

## 4.1 Data types

The Hexagon processor provides operations for processing the following data types.

### 4.1.1 Fixed-point data

The Hexagon processor provides operations to process 8-, 16-, 32-, or 64-bit fixed-point data. The data is either integer or fractional, and in signed or unsigned format.

#### 4.1.1.1 Scalar operations

The Hexagon processor includes the following scalar operations on fixed-point data:

- Multiplication of 16-bit, 32-bit, and complex data
- Addition and subtraction of 16-, 32-, and 64-bit data (with and without saturation)
- Logical operations on 32- and 64-bit data (AND, OR, XOR, NOT)
- Shifts on 32- and 64-bit data (arithmetic and logical)
- Min/max, negation, absolute value, parity, norm, swizzle
- Compares of 8-, 16-, 32-, and 64-bit data
- Sign and zero extension (8- and 16- to 32-bit, 32- to 64-bit)
- Bit manipulation
- Predicate operations



### 4.1.1.2 Vector operations

The Hexagon processor includes the following vector operations on fixed-point data:

- Multiplication (halfwords, word by half, vector reduce, dual multiply)
- Addition and subtraction of word and halfword data
- Shifts on word and halfword data (arithmetic and logical)
- Min/max, average, negative average, absolute difference, absolute value
- Compares of word, halfword, and byte data
- Reduce, sum of absolute differences on unsigned bytes
- Special-purpose data arrangement (such as pack, splat, shuffle, align, saturate, splice, truncate, complex conjugate, complex rotate, zero extend)

**NOTE:** Certain vector operations support automatic scaling, saturation, and rounding.

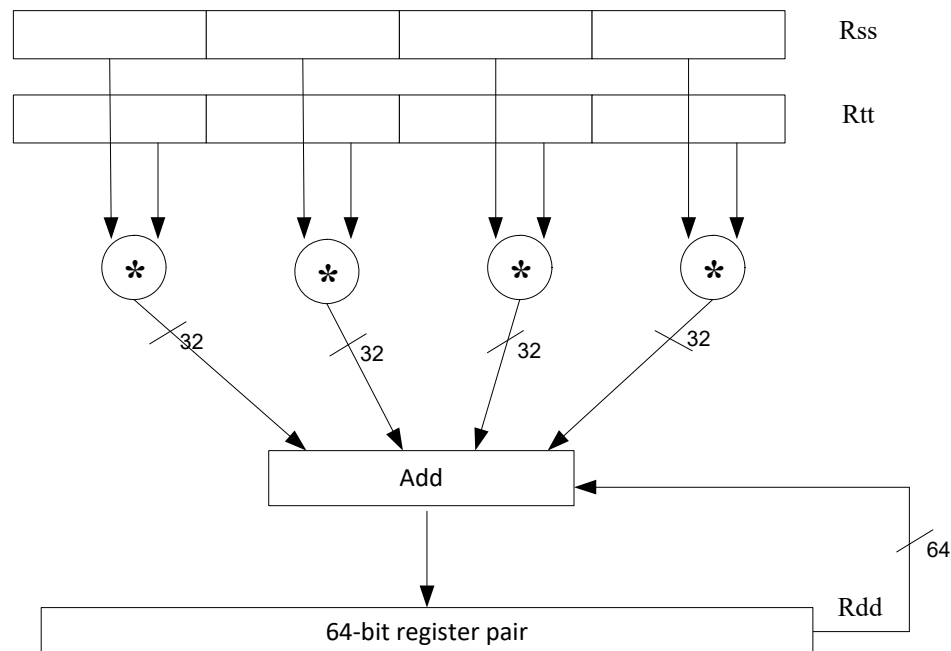
For example, the following instruction performs a vector operation:

```
R1:0 += vrmphy(R3:2, R5:4)
```

It is defined to perform the following operations in one cycle:

```
R1:0 += ((R2.L * R4.L) +
         (R2.H * R4.H) +
         (R3.L * R5.L) +
         (R3.H * R5.H)
        )
```

Figure 4-1 shows a schematic of this instruction type.



**Figure 4-1** Vector instruction example

## 4.1.2 Floating-point data

The Hexagon processor provides operations to process 32-bit floating-point numbers. The numbers are stored in IEEE single-precision floating-point format.

Per the IEEE standard, certain floating-point values are defined to represent positive or negative infinity, as well as Not-a-Number (NaN), which represents values that have no mathematical meaning.

Floating-point numbers can be held in a general register.

### 4.1.2.1 Floating-point operations

The Hexagon processor includes the following operations on floating-point data:

- Addition and subtraction
- Multiplication (with optional scaling)
- Min/max/compare
- Reciprocal/square root approximation
- Format conversion

## 4.1.3 Complex data

The Hexagon processor provides operations to process 32- or 64-bit complex data.

Complex numbers include a signed real portion and a signed imaginary portion. Given two complex numbers  $(a + bi)$  and  $(c + di)$ , the complex multiply operations computes both the real portion  $(ac - bd)$  and the imaginary portion  $(ad + bc)$  in a single instruction.

Complex numbers can be packed in a general register or register pair. When packed, the imaginary portion occupies the most-significant portion of the register or register pair.

## 4.1.4 Vector data

The Hexagon processor provides operations to process 64-bit vector data.

Vector data types pack multiple data items – bytes, halfwords, or words – into 64-bit registers. Vector data operations are common in video and image processing.

Eight 8-bit bytes can be packed into a 64-bit register.

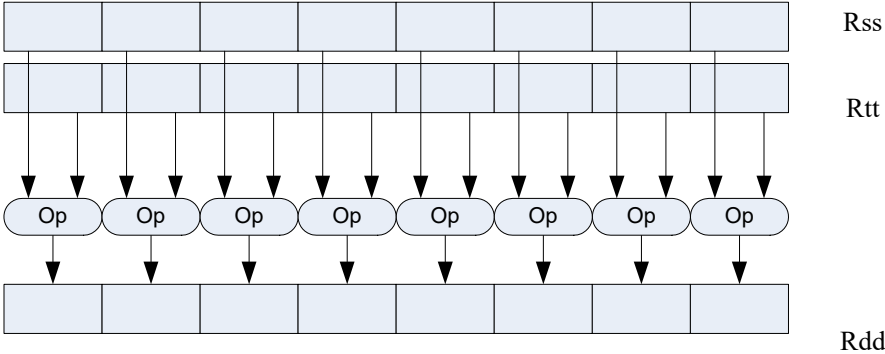


Figure 4-1 Vector byte operation example

Four 16-bit halfword values can be packed in a single 64-bit register pair.

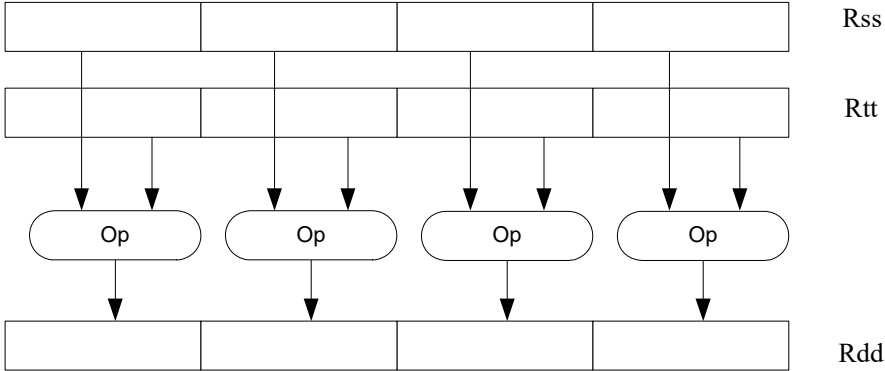


Figure 4-2 Vector halfword operation example

Two 32-bit word values can be packed in a single 64-bit register pair.

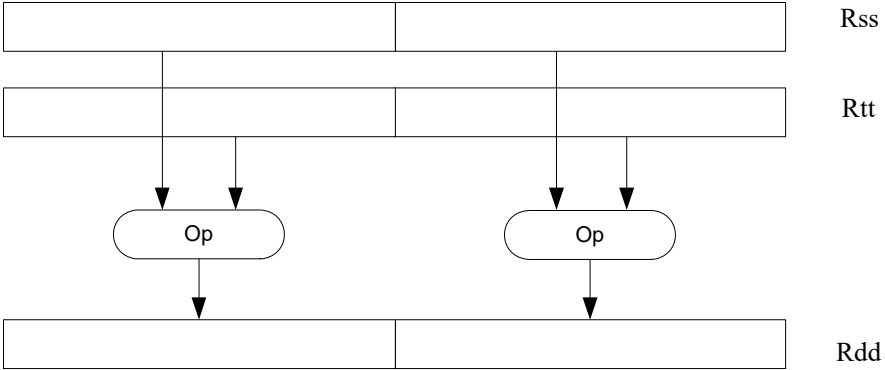


Figure 4-3 Vector word operation example

## 4.2 Instruction options

Some instructions support optional scaling, saturation, and rounding. There are no mode bits controlling these options – instead, they are explicitly specified as part of the instruction name. The options are described in this section.

### 4.2.1 Fractional scaling

In fractional data format, data is treated as fixed-point fractional values whose range is determined by the word length and radix point position.

Fractional scaling is specified in an instruction by adding the `<<1` specifier. For example:

```
R3:2 = cmpy(R0,R1):<<1:sat
```

When two fractional numbers are multiplied, the product must be scaled to restore the original fractional data format. The Hexagon processor allows specification of the fractional scaling of the product in the instruction for shifts of 0 and 1. Perform a shift of 1 for Q1.15 numbers, perform a shift of 0 for integer multiplication.

### 4.2.2 Saturation

Certain instructions are available in saturating form. If a saturating arithmetic instruction has a result which is smaller than the minimum value, the result is set to the minimum value. Similarly, if the operation has a result which is greater than the maximum value, the result is set to the maximum value.

Saturation is specified in an instruction by adding the `:sat` specifier. For example:

```
R2 = abs(R1):sat
```

The open virtualization format (OVF) bit in the [User status register](#) is set whenever a saturating operation saturates to the maximum or minimum value. It remains set until explicitly cleared by a control register transfer to USR. For vector-type saturating operations, if any of the individual elements of the vector saturate, OVF is set.

### 4.2.3 Arithmetic rounding

Certain signed multiply instructions support optional arithmetic rounding (also known as biased rounding). The arithmetic rounding operation takes a double precision fractional value and adds 0x8000 to the low 16-bits (least significant 16-bit halfword).

Rounding is specified in an instruction by adding the `:rnd` specifier. For example:

```
R2 = mpy(R1.h,R2.h):rnd
```

**NOTE:** Arithmetic rounding can accumulate numerical errors, especially when the number to round is exactly 0.5. This happens most frequently when dividing by 2 or averaging.

## 4.2.4 Convergent rounding

To address the problem of error accumulation in [Convergent rounding](#), the Hexagon processor includes four instructions that support positive and negative averaging with a convergent rounding option.

These instructions work as follows:

1. Compute  $(A + B)$  or  $(A - B)$  for AVG and NAVG respectively.
2. Based on the two least-significant bits of the result, add a rounding constant as follows:
  - If the two LSBs are 00, add 0
  - If the two LSBs are 01, add 0
  - If the two LSBs are 10, add 0
  - If the two LSBs are 11, add 1
3. Shift the result right by one bit.

## 4.2.5 Scaling for divide and square-root

On the Hexagon processor, floating point divide and square-root operations are implemented in software using library functions. To enable the efficient implementation of these operations, the processor supports special variants of the multiply-accumulate instruction, named scale FMA.

Scale FMA supports optional scaling of the product generated by the floating-point fused multiply-add instruction.

Scaling is specified in the instruction by adding the `:scale` specifier and a predicate register operand. For example:

```
R3 += sfmpy(R0,R1,P2):scale
```

For single precision, the scaling factor is two raised to the power specified by the contents of the predicate register (which is treated as an 8-bit two's complement value). For double precision, the predicate register value is doubled before use as a power of two.

**NOTE:** Do not use scale FMA instructions outside of divide and square-root library routines. No guarantee is provided that future versions of the Hexagon processor will implement these instructions using the same semantics. Future versions assume only that compatibility for scale FMA is limited to the needs of divide and square-root library routines.

## 4.3 XTYPE operations

The XTYPE instruction class includes most of the data-processing operations performed by the Hexagon processor. These operations are categorized by their operation type.

### 4.3.1 ALU

**XTYPE ALU** operations modify 8-, 16-, 32-, and 64-bit data. These operations include:

- Add and subtract with and without saturation
- Add and subtract with accumulate
- Absolute value
- Logical operations
- Min, max, negate instructions
- Register transfers of 64-bit data
- Word to doubleword sign extension
- Comparisons

### 4.3.2 Bit manipulation

**XTYPE BIT** manipulation operations modify bit fields in a register or register pair. These operations include:

- Bit field insert
- Bit field signed and unsigned extract
- Count leading and trailing bits
- Compare bit masks
- Set/clear/toggle bit
- Test bit operation
- Interleave/deinterleave bits
- Bit reverse
- Split bit field
- Masked parity and linear feedback shift
- Table index formation

### 4.3.3 Complex

**XTYPE COMPLEX** operations manipulate complex numbers. These operations include:

- Complex add and subtract
- Complex multiply with optional round and pack
- Vector complex multiply
- Vector complex conjugate
- Vector complex rotate
- Vector reduce complex multiply real or imaginary

### 4.3.4 Floating point

**XTYPE FP** operations manipulate single-precision floating point numbers. These operations include:

- Addition and subtraction
- Multiplication (with optional scaling)
- Min/max/compare
- Format conversion

The Hexagon floating-point operations are defined to support the IEEE floating-point standard. However, certain IEEE-required operations – such as divide and square root – are not supported directly. Instead, special instructions are defined to support the implementation of the required operations as library routines. These instructions include:

- A special version of the fused multiply-add instruction (designed specifically for use in library routines)
- Reciprocal/square root approximations (which compute the approximate initial values used in reciprocal and reciprocal-square-root routines)
- Extreme value assistance (which adjusts input values if they cannot produce correct results using convergence algorithms)

**NOTE:** The special floating-point instructions are not intended for use directly in user code – use the special floating-point instructions only in the floating point library.

#### Format conversion

The floating-point conversion instructions `sfmake` and `dfmake` convert an unsigned 10-bit immediate value into the corresponding floating-point value.

The immediate value must be encoded so bits [5:0] contain the significand, and bits [9:6] the exponent. The exponent value is added to the initial exponent value (`bias - 6`).

For example, to generate the single-precision floating point value 2.0, bits [5:0] must be set to 0, and bits [9:6] set to 7. Performing the `sfmake` operation on this immediate value yields the floating point value `0x40000000`, which is 2.0.

**NOTE:** The conversion instructions are designed to handle common floating point values, including most integers and many basic fractions ( $1/2$ ,  $3/4$ , and so on).

## Rounding

The Hexagon [User status register](#) includes the FPRND field, which specifies the IEEE-defined floating-point rounding mode.

## Exceptions

The Hexagon user status register includes five status fields, which work as sticky flags for the five IEEE-defined exception conditions: inexact, overflow, underflow, divide by zero, and invalid. A sticky flag is set when the corresponding exception occurs, and remains set until explicitly cleared.

The user status register also includes five mode fields which specify whether to perform an operating-system trap if one of the floating-point exceptions occur. For every instruction packet containing a floating-point operation, if a floating-point sticky flag and the corresponding trap-enable bit are both set, a floating-point trap is generated. After the packet commits, the Hexagon processor then automatically traps to the operating system.

**NOTE:** Non-floating-point instructions never generate a floating-point trap, regardless of the state of the sticky flag and trap-enable bits.

## 4.3.5 Multiply

Multiply operations support fixed-point multiplication, including both single- and double-precision multiplication, and polynomial multiplication.

### Single precision

In single-precision arithmetic a 16-bit value is multiplied by another 16-bit value. These operands can come from the high portion or low portion of any register. Depending on the instruction, the result of the  $16 \times 16$  operation can optionally be accumulated, saturated, rounded, or shifted left by 0 to 1 bits.

The instruction set supports operations on signed  $\times$  signed, unsigned  $\times$  unsigned, and signed  $\times$  unsigned data.



**Table 4-1** summarizes the options available for  $16 \times 16$  single precision multiplications. The symbols used in the table are as follows:

- SS – Perform signed  $\times$  signed multiply
- UU – Perform unsigned  $\times$  unsigned multiply
- SU – Perform signed  $\times$  unsigned multiply
- A+ – Result added to accumulator
- A- – Result subtracted from accumulator
- 0 – Result not added to accumulator

**Table 4-1 Single-precision multiply options**

Multiply	Result	Sign	Accumulate	Sat	Rnd	Scale
$16 \times 16$	32	SS	A+, A-	Yes	No	0-1
$16 \times 16$	32	SS	0	Yes	Yes	0-1
$16 \times 16$	64	SS	A+, A-	No	No	0-1
$16 \times 16$	64	SS	0	No	Yes	0-1
$16 \times 16$	32	UU	A+, A-, 0	No	No	0-1
$16 \times 16$	64	UU	A+, A-, 0	No	No	0-1
$16 \times 16$	32	SU	A+, 0	Yes	No	0-1

### Double precision

Double precision instructions are available for both  $32 \times 32$  and  $32 \times 16$  multiplication:

- For  $32 \times 32$  multiplication the result is either 64 or 32 bits. The 32-bit result is either the high or low portion of the 64-bit product.
- For  $32 \times 16$  multiplication the result is always taken as the upper 32 bits.

The operands are either signed or unsigned.

**Table 4-2 Double precision multiply options**

Multiply	Result	Sign	Accumulate	Sat	Rnd	Scale
$32 \times 32$	64	SS, UU	A+, A-, 0	No	No	0
$32 \times 32$	32 (upper)	SS, UU	0	No	Yes	0
$32 \times 32$	32 (low)	SS, UU	A+, 0	No	No	0
$32 \times 16$	32 (upper)	SS, UU	A+, 0	Yes	Yes	0-1
$32 \times 32$	32 (upper)	SU	0	No	No	0

## Polynomial

Polynomial **XTYPE MPY** instructions are available for both words and vector halfwords.

These instructions are useful for many algorithms including scramble code generation, cryptographic algorithms, convolutional, and Reed Solomon code.

### 4.3.6 Permute

**XTYPE PERM** operations perform various operations on vector data, including arithmetic, format conversion, and rearrangement of vector elements. Many types of conversions are supported:

- Swizzle bytes
- Vector shuffle
- Vector align
- Vector saturate and pack
- Vector splat bytes
- Vector splice
- Vector sign extend halfwords
- Vector zero extend bytes
- Vector zero extend halfwords
- Scalar saturate to byte, halfword, word
- Vector pack high and low halfwords
- Vector round and pack
- Vector splat halfwords

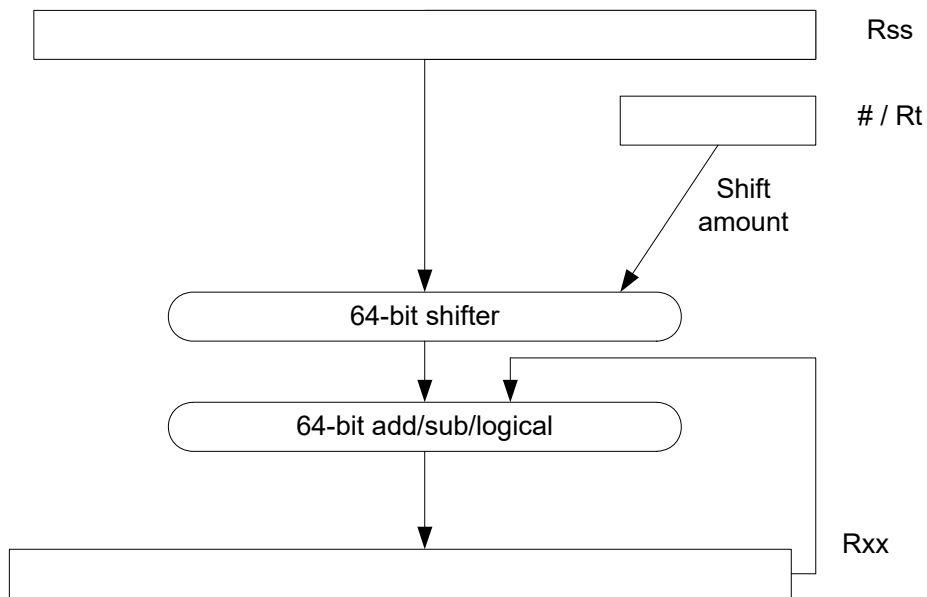
### 4.3.7 Predicate

**XTYPE PRED** operations modify predicate source data. The categories of instructions available include:

- Vector mask generation
- Predicate transfers
- Viterbi packing

### 4.3.8 Shift

Scalar **XTYPE SHIFT** operations perform a variety of 32 and 64-bit shifts followed by an optional add/sub or logical operation. **Figure 4-4** shows the general operation.



**Figure 4-4 64-bit shift and add/sub/logical**

Four shift types are supported:

- ASR – Arithmetic shift right
- ASL – Arithmetic shift left
- LSR – Logical shift right
- LSL – Logical shift left

In register-based shifts, the Rt register is a signed two's-complement number. If this value is positive, the instruction opcode tells the direction of shift (right or left). If this value is negative, the shift direction indicated by the opcode is reversed.

When arithmetic right shifts are performed, the sign bit is shifted in, whereas logical right shifts shift in zeros. Left shifts always shift in zeros.

Some shifts are available with saturation and rounding options.

## 4.4 ALU32 operations

The [ALU32](#) instruction class includes general arithmetic/logical operations on 32-bit data:

- Add, subtract, negate without saturation on 32-bit data
- Logical operations such as AND, OR, XOR, AND with immediate, and OR with immediate
- Scalar 32-bit compares
- Combine halfwords, combine words, combine with immediates, shift halfwords, and Mux
- Conditional add, combine, logical, subtract, and transfer.
- NOP
- Sign and zero-extend bytes and halfwords
- Transfer immediates and registers
- Vector add, subtract, and average halfwords

**NOTE:** ALU32 instructions can execute on any slot ([Section 3.3.3](#)).

[Chapter 6](#) describes the conditional execution and compare instructions.

## 4.5 Vector operations

Vector operations support arithmetic operations on vectors of bytes, halfwords, and words.

The vector operations belong to the XTYPE instruction class (except for vector add, subtract, and average halfwords, which are ALU32).

### Vector byte operations

The vector byte operations process packed vectors of signed or unsigned bytes. They include the following operations:

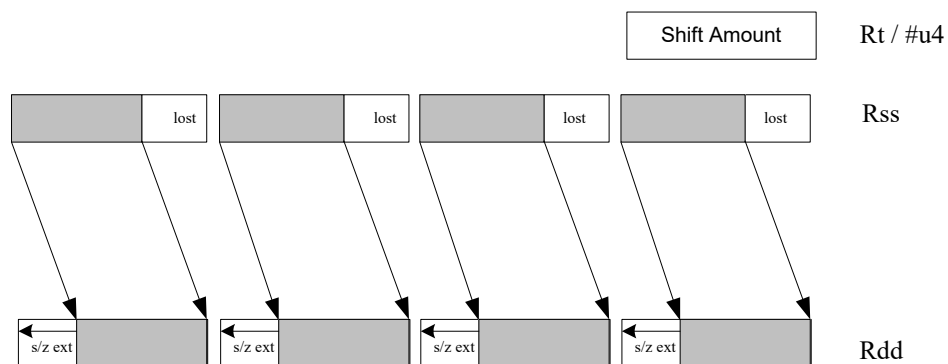
- Vector add and subtract signed or unsigned bytes
- Vector min and max signed or unsigned bytes
- Vector compare signed or unsigned bytes
- Vector average unsigned bytes
- Vector reduce add unsigned bytes
- Vector sum of absolute differences unsigned bytes

## Vector halfword operations

The vector halfword operations process packed 16-bit halfwords. They include the following operations:

- Vector add and subtract halfwords
- Vector average halfwords
- Vector compare halfwords
- Vector min and max halfwords
- Vector shift halfwords
- Vector dual multiply
- Vector dual multiply with round and pack
- Vector multiply even halfwords with optional round and pack
- Vector multiply halfwords
- Vector reduce multiply halfwords

For example, [Figure 4-5](#) shows the operation of the vector arithmetic shift right halfword (vasrh) instruction. In this instruction, each 16-bit half-word is shifted right by the same amount which is specified in a register or with an immediate value. Because the shift is arithmetic, the bits shifted in are copies of the sign bit.



**Figure 4-5 Vector halfword shift right**

## Vector word operations

The vector word operations process packed vectors of two words. They include the following operations:

- Vector add and subtract words
- Vector average words
- Vector compare words
- Vector min and max words
- Vector shift words with optional truncate and pack

For more information on vector operations see [Section 11.1.1](#) and [Section 11.10.1](#).

## 4.6 CR operations

The **CR** instruction class includes operations that access the [Control registers](#).

**Table 4-3 Control register transfer instructions**

Syntax	Operation
Rd = Cs Cd = Rs	Move control register to / from a general register.  <b>NOTE:</b> PC is not a valid destination register.
Rdd = Css Cdd = Rss	Move control register pair to / from a general register pair.  <b>NOTE:</b> PC is not a valid destination register.

**NOTE:** In register-pair transfers, control registers must be specified using their numeric alias names – see [Section 2.2](#) for details.

## 4.7 Compound operations

The instruction set includes a number of instructions that perform multiple logical or arithmetic operations in a single instruction. They include the following operations:

- AND/OR with inverted input
- Compound logical register
- Compound logical predicate
- Compound add-subtract with immediates
- Compound shift-operation with immediates (arithmetic or logical)
- Multiply-add with immediates

For more information see [Section 11.10.1](#).

## 4.8 Special operations

The instruction set includes a number of special-purpose instructions to support specific applications.

### 4.8.1 H.264 CABAC processing

H.264/AVC is adopted in a diverse range of multimedia applications:

- HD-DVDs
- HDTV broadcasting
- Internet video streaming

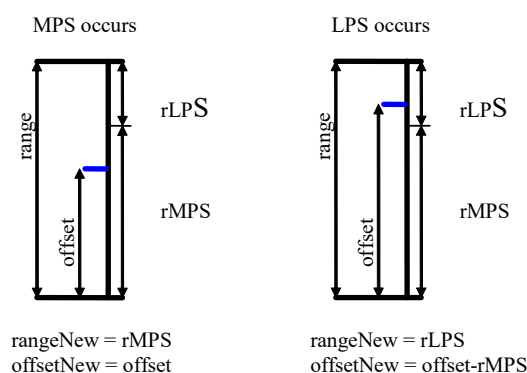
Context Adaptive Binary Arithmetic Coding (CABAC) is one of the two alternative entropy coding methods specified in the H.264 main profile. CABAC offers superior coding efficiency at the expense of greater computational complexity. The Hexagon processor includes a dedicated instruction (decbin) to support CABAC decoding.

Binary arithmetic coding is based on the principle of recursive interval subdivision, and its state is characterized by two quantities:

- The current interval range
- The current offset in the current code interval

The offset is read from the encoded bit stream. When decoding a bin, the interval range is subdivided in two intervals based on the estimation of the probability  $p_{LPS}$  of least probable symbol (LPS): one interval with width of  $rLPS = \text{range} \times p_{LPS}$ , and another with width of  $rMPS = \text{range} \times p_{MPS} = \text{range} - rLPS$ , where MPS stands for most probable symbol.

Depending on which subinterval the offset falls into, the decoder decides whether the bin is decoded as MPS or LPS, after which the two quantities are iteratively updated, as shown in [Figure 4-1](#).

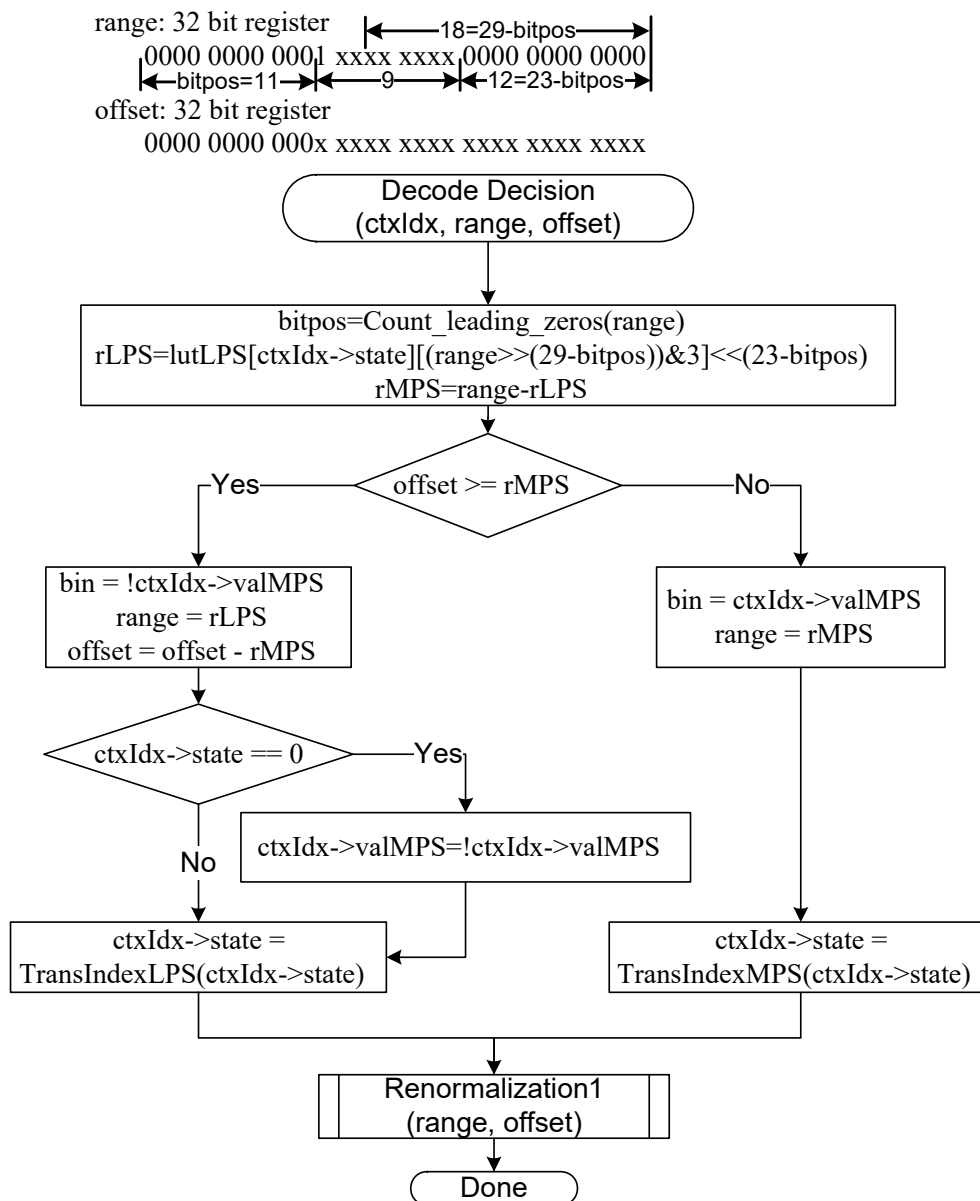


**Figure 4-1**      **Arithmetic decoding for one bin**

### 4.8.1.1 CABAC implementation

In H.264 range is a 9-bit quantity, and offset is 9 bits in regular mode and 10 bits in bypass mode during the whole decoding process. The calculation of rLPS is approximated by a  $64 \times 4$  table of 256 bytes, where the range and the context state (selected for the bin to decode) address the lookup table. To maintain the precision of the whole decoding process, the new range must be renormalized to ensure that the most significant bit is always 1, and that the offset is synchronously refilled from the bit stream.

To simplify the renormalization/refilling process, the decoding scheme shown in [Figure 4-2](#) significantly reduces the frequency of renormalization and refilling bits from the bit-stream, while also being suitable for DSP implementation.



**Figure 4-2 CABAC decoding engine for regular bin**



The Hexagon processor can use the `decbin` instruction to decode one regular bin in two cycles (not counting the bin refilling process).

For more information on the `decbin` instruction see [Section 11.10.6](#).

For example:

```
Rdd = decbin(Rss,Rtt)
```

INPUT: Rss and Rtt register pairs as:

```
Rtt.w1[5:0] = state
Rtt.w1[8] = valMPS
Rtt.w0[4:0] = bitpos
Rss.w0 = range
Rss.w1 = offset
```

OUTPUT: Rdd register pair is packed as

```
Rdd.w0[5:0] = state
Rdd.w0[8] = valMPS
Rdd.w0[31:23] = range
Rdd.w0[22:16] = '0'
Rdd.w1 = offset (normalized)
```

OUTPUT: P0

```
P0 = (bin)
```

#### 4.8.1.2 Code example

```
H264CabacGetBinNC:
/*****
* Non-conventional call:
* Input: R1:0 = offset : range , R2 = dep, R3 = ctxIdx,
*        R4 = (*ctxIdx), R5 = bitpos
*
* Return:
*        R1: 0 - offset : range
*        P0 - (bin)
*****/

// Cycle #1
{ R1:0= decbin(R1:0,R5:4) // Decoding one bin
  R6 = asl(R22,R5) // Where R22 = 0x100
}

// Cycle #2
{ memb(R3) = R0 // Save context to *ctxIdx
  R1:0 = vlsrw(R1:0,R5) // Re-align range and offset
  P1 = cmp.gtu(R6,R1) // Need refill? i.e., P1= (range<0x100)
  IF (!P1.new) jumpr:t LR // Return
}
RENORM_REFILL:
...
```

## 4.8.2 IP Internet checksum

The key features of the Internet checksum<sup>1</sup> include:

- The checksum can be summed in any order
- Carries can be accumulated using an accumulator larger than size being added, and added back in at any time

Using standard data-processing instructions, the Internet checksum can be computed at 8 bytes per cycle in the main loop, by loading words and accumulating into doublewords. After the loop, the upper word is added to the lower word; then the upper halfword is added to the lower halfword, and any carries are added back in.

The Hexagon processor supports a dedicated instruction (`vradduh`) that computes the Internet checksum at a rate of 16 bytes per cycle.

The `vradduh` instruction accepts the halfwords of the two input vectors, adds them all together, and places the result in a 32-bit destination register. This operation can both compute the sum of 16 bytes of input while preserving the carries, and accumulate carries at the end of computation.

For more information on the `vradduh` instruction, see [Vector reduce add halfwords](#).

**NOTE:** This operation utilizes the maximum load bandwidth available in the Hexagon processor.

---

<sup>1</sup> See RFC 1071 (<http://www.faqs.org/rfcs/rfc1071.html>)

### 4.8.2.1 Code example

```

.text
.global fast_ip_check
// Assumes data is 8-byte aligned
// Assumes data is padded at least 16 bytes afterwords with 0's.
// input R0 points to data
// input R1 is length of data
// returns IP checksum in R0

fast_ip_check:
{
    R1 = lsr(R1,#4)           // 16-byte chunks, rounded down, +1
    R9:8 = combine(#0,#0)
    R3:2 = combine(#0,#0)
}
{
    loop0(1f,R1)
    R7:6 = memd(R0+#8)
    R5:4 = memd(R0++#16)
}
.falign
1:
{
    R7:6 = memd(R0+#8)
    R5:4 = memd(R0++#16)
    R2 = vradduh(R5:4,R7:6)   // Accumulate 8 halfwords
    R8 = vradduh(R3:2,R9:8)  // Accumulate carries
}:endloop0
// Drain pipeline
{
    R2 = vradduh(R5:4,R7:6)
    R8 = vradduh(R3:2,R9:8)
    R5:4 = combine(#0,#0)
}
{
    R8 = vradduh(R3:2,R9:8)
    R1 = #0
}
// May have some carries to add back in
{
    R0 = vradduh(R5:4,R9:8)
}
// Possible for one more to pop out
{
    R0 = vradduh(R5:4,R1:0)
}
{
    R0 = not(R0)
    jumpr LR
}

```

## 4.8.3 Software-defined radio

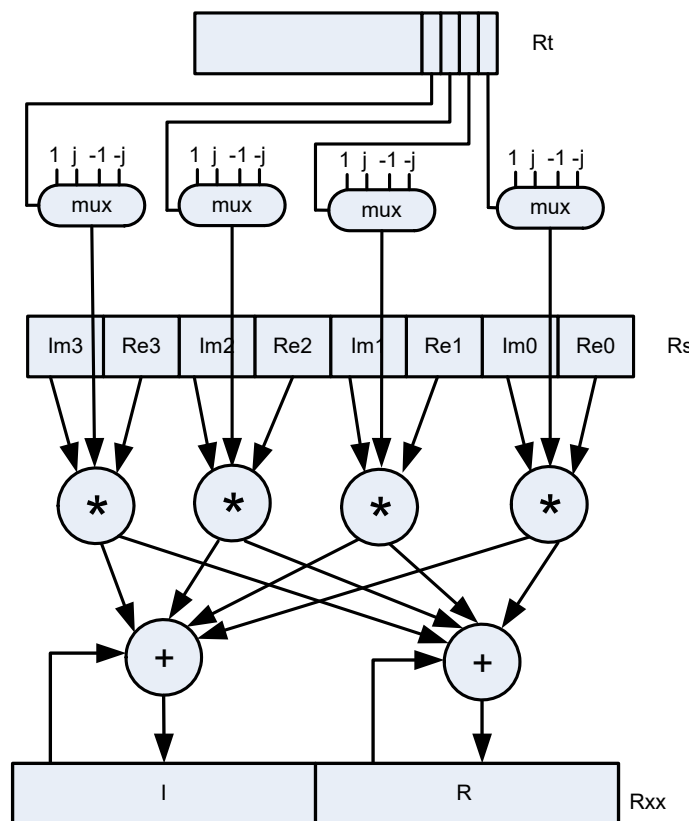
The Hexagon processor includes six special-purpose instructions which support the implementation of software-defined radio. The instructions greatly accelerate the following algorithms.

### 4.8.3.1 Rake despreading

A fundamental operation in despreading is the PN multiply operation. In this operation the received complex chips are compared against a pseudo-random sequence of QAM constellation points and accumulated.

Figure 4-3 shows the `vrrotate` instruction that performs this operation. The products are summed to form a soft 32-bit complex symbol. The instruction has both accumulating and non-accumulating versions.

```
xx += vrrotate(Rss,Rt,#0)
```



**Figure 4-3** Vector reduce complex rotate

For more information on the `vrrotate` instruction, see [Vector reduce complex rotate](#).

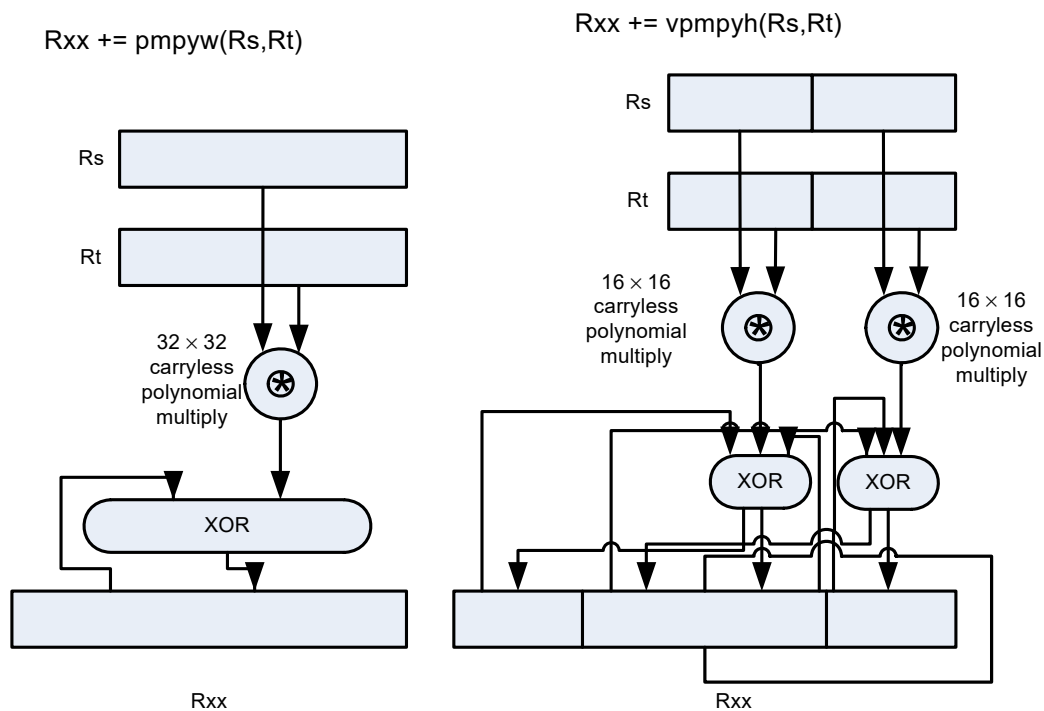
**NOTE:** Using this instruction the Hexagon processor can process 5.3 chips per cycle, and a 12-finger WCDMA user requires only 15 MHz.

### 4.8.3.2 Polynomial operations

The polynomial multiply instructions support the following operations:

- Scramble code generation (at a rate of 8 symbols per cycle for WCDMA)
- Cryptographic algorithms (such as elliptic curve)
- CRC checks (at a rate of 21 bits per cycle)
- Convolutional encoding
- Reed-Solomon codes

The four versions of this instruction support  $32 \times 32$  and vector  $16 \times 16$  multiplication both with and without accumulation, as shown in [Figure 4-4](#).



**Figure 4-4** Polynomial multiply

For more information on the pmpy instructions, see [Polynomial multiply words](#).

# 5 Memory

---

The Hexagon processor features a load/store architecture, where numeric and logical instructions operate on registers. Explicit load instructions move operands from memory to registers, while store instructions move operands from registers to memory. A small number of instructions (known as mem-ops) perform numeric and logical operations directly on memory.

The address space is unified: all accesses target the same linear address space, which contains both instructions and data.

## 5.1 Memory model

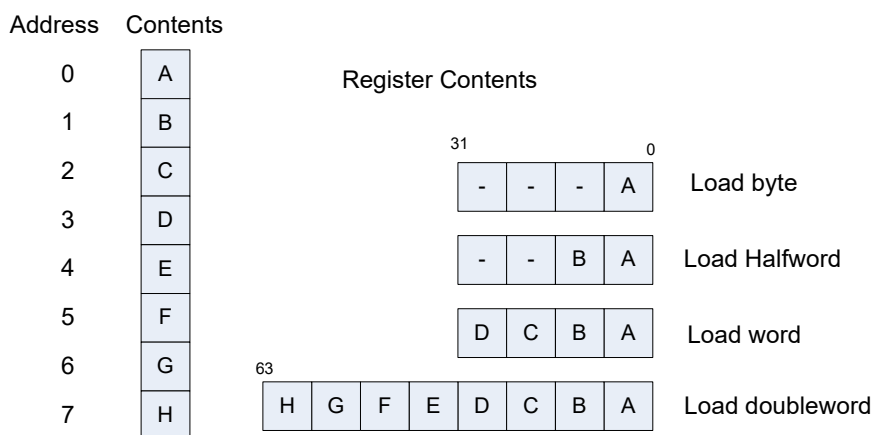
This section describes the memory model for the Hexagon processor.

### 5.1.1 Address space

The Hexagon processor has a 32-bit byte-addressable memory address space. The entire 4G linear address space is addressable by the user application. A virtual-to-physical address translation mechanism is provided.

### 5.1.2 Byte order

The Hexagon processor is a little-endian machine: the lowest address byte in memory is held in the least significant byte of a register, as shown in [Figure 5-1](#).



**Figure 5-1 Hexagon processor byte order**

### 5.1.3 Alignment

Even though the Hexagon processor memory is byte-addressable, instructions and data must be aligned in memory on specific address boundaries:

- Instructions and instruction packets must be 32-bit aligned
- Data must be aligned to its native access size.

Any unaligned memory access causes a memory-alignment exception.

Use the [Permute](#) instructions in applications that must reference unaligned vector data. The loads and stores still must be memory-aligned; however, the permute instructions enable easy rearrangement of the data in registers.

**Table 5-1 Memory alignment restrictions**

Data type	Size (bits)	Exception when
Byte Unsigned byte	8	Never
Halfword Unsigned halfword	16	LSB[0] != 0 <sup>1</sup>
Word Unsigned word	32	LSB[1:0] != 00
Doubleword	64	LSB[2:0] != 000
Instruction Instruction packet	32	LSB[1:0] != 00

<sup>1</sup> LSB = Least significant bits of address

## 5.2 Memory loads

Memory can be loaded in byte, halfword, word, or doubleword sizes. The data types supported are signed or unsigned. The syntax used is `memXX`, where `XX` denotes the data type.

**Table 5-2 Load instructions**

Syntax	Source size (bits)	Destination size (bits)	Data placement	Comment
<code>Rd = memub (Rs)</code>	8	32	Low 8 bits	Zero-extend 8 to 32 bits
<code>Rd = memb (Rs)</code>	8	32	Low 8 bits	Sign-extend 8 to 32 bits
<code>Rd = memuh (Rs)</code>	16	32	Low 16 bits	Zero-extend 16 to 32 bits
<code>Rd = memh (Rs)</code>	16	32	Low 16 bits	Sign-extend 16 to 32 bits
<code>Rd = memubh (Rs)</code>	16	32	Bytes 0 and 2	Bytes 1 and 3 zeroed <sup>1</sup>
<code>Rd = membh (Rs)</code>	16	32	Bytes 0 and 2	Bytes 1 and 3 sign-extended
<code>Rd = memw (Rs)</code>	32	32	All 32 bits	Load word
<code>Rdd = memubh (Rs)</code>	32	64	Bytes 0,2,4,6	Bytes 1,3,5,7 zeroed
<code>Rdd = membh (Rs)</code>	32	64	Bytes 0,2,4,6	Bytes 1,3,5,7 sign-extended
<code>Rdd = memd (Rs)</code>	64	64	All 64 bits	Load doubleword
<code>Ryy = memh_fifo (Rs)</code>	16	64	High 16 bits	Shift vector and load halfword
<code>deallocframe</code>	64	64	All 64 bits	See <a href="#">Chapter 7</a>
<code>dealloc_return</code>	64	64	All 64 bits	See <a href="#">Chapter 7</a>

<sup>1</sup> The `memubh` and `membh` instructions load contiguous bytes from memory (either 2 or 4 bytes) and unpack these bytes into a vector of halfwords. The instructions are useful when bytes are used as input into halfword vector operations, which is common in video and image processing..

**NOTE:** The memory load instructions belong to instruction class LD, and can execute only in slots 0 or 1.



## 5.3 Memory stores

Memory can be stored in byte, halfword, word, or doubleword sizes. The syntax used is `memX`, where `X` denotes the data type.

**Table 5-3 Store instructions**

Syntax	Source size (bits)	Destination size (bits)	Comment
<code>memb (Rs) = Rt</code>	32	8	Store byte (bits 7:0)
<code>memb (Rs) = #s8</code>	8	8	Store byte
<code>memh (Rs) = Rt</code>	32	16	Store lower half (bits 15:0)
<code>memh (Rs) = Rt.H</code>	32	16	Store upper half (bits 31:16)
<code>memh (Rs) = #s8</code>	8	16	Sign-extend 8 to 16 bits
<code>memw (Rs) = Rt</code>	32	32	Store word
<code>memw (Rs) = #s8</code>	8	32	Sign-extend 8 to 32 bits
<code>memd (Rs) = Rtt</code>	64	64	Store doubleword
<code>allocframe (#u11)</code>	64	64	See <a href="#">Chapter 7</a>

**NOTE:** The memory store instructions belong to instruction class ST, and can execute only in slot 0 or – when part of a dual store – slot 1.

## 5.4 Dual stores

Two memory store instructions can appear in the same instruction packet. The resulting operation is considered a dual store. For example:

```
{
    memw(R5) = R2      // Dual store
    memh(R6) = R3
}
```

Unlike most packetized operations, dual stores do not execute in parallel ([Section 3.3.1](#)). Instead, the store instruction in Slot 1 effectively executes first, followed by the store instruction in Slot 0.

**NOTE:** The store instructions in a dual store must belong to instruction class [ST](#), and can execute only in Slots 0 and 1.

## 5.5 Slot 1 store with slot 0 load

A slot 1 store operation with a slot 0 load operation can appear in a packet. The packet attribute `:mem_noshuf` inhibits the instruction reordering that is otherwise done by the assembler. For example:

```
{
    memw(R5) = R2      // Slot 1 store
    R3 = memh(R6)     // Slot 0 load
}:mem_noshuf
```

Unlike most packetized operations, these memory operations do not execute in parallel ([Section 3.3.1](#)). Instead, the store instruction in Slot 1 effectively executes first, followed by the load instruction in Slot 0. If the addresses of the two operations are overlapping, the load receives the newly stored data.

## 5.6 New-value stores

A memory store instruction can store a register that is assigned a new value in the same instruction packet ([Section 3.3](#)). This feature is expressed in assembly language by appending the suffix `.new` to the source register. For example:

```
{
    R2 = memh(R4+#8)   // load halfword
    memw(R5) = R2.new // store newly-loaded value
}
```

New-value store instructions have the following restrictions:

- If an instruction uses auto-increment or absolute-set addressing mode ([Section 5.8](#)), its address register cannot be used as the new-value register.
- If an instruction produces a 64-bit result, its result registers cannot be used as the new-value register.
- If the instruction that sets a new-value register is conditional ([Section 6.1.2](#)), it must always execute.

NOTE: The new-value store instructions belong to instruction class NV, and can execute only in Slot 0.

## 5.7 Mem-ops

Mem-ops perform basic arithmetic, logical, and bit operations directly on memory operands, without the need for a separate load or store. Mem-ops can be performed on byte, halfword, or word sizes.

**Table 5-4 Mem-ops**

Syntax	Operation
<code>memXX (Rs+#u6) [+ -   &amp;] = Rt</code>	Arithmetic/logical on memory
<code>memXX (Rs+#u6) [+ -] = #u5</code>	Arithmetic on memory
<code>memXX (Rs+#u6) = clrbit (#u5)</code>	Clear bit in memory
<code>memXX (Rs+#u6) = setbit (#u5)</code>	Set bit in memory

NOTE: The mem-op instructions belong to instruction class MEMOP, and can execute only in slot 0.

## 5.8 Addressing modes

**Table 5-5 Addressing modes**

Mode	Syntax	Operation <sup>1</sup>
Absolute	<code>memXX (##address)</code>	EA = address
Absolute-set	<code>memXX (Re=##address)</code>	EA = address Re = address
Absolute with register offset	<code>memXX (Ru&lt;&lt;#u2+##U32)</code>	EA = imm + (Ru << #u2)
Global pointer relative	<code>memXX (GP+#immediate)</code> <code>memXX (#immediate)</code>	EA = GP + immediate
Indirect	<code>memXX (Rs)</code>	EA = Rs
Indirect with offset	<code>memXX (Rs+#s11)</code>	EA = Rs + imm
Indirect with register offset	<code>memXX (Rs+Ru&lt;&lt;#u2)</code>	EA = Rs + (Ru << #u2)
Indirect with auto-increment immediate	<code>memXX (Rx++#s4)</code>	EA = Rx; Rx += (imm)
Indirect with auto-increment register	<code>memXX (Rx++Mu)</code>	EA = Rx; Rx += Mu
Circular with auto-increment immediate	<code>memXX (Rx++#s4:circ (Mu) )</code>	EA = Rx; Rx = circ_add (Rx, imm, Mu)
Circular with auto-increment register	<code>memXX (Rx++I:circ (Mu) )</code>	EA = Rx; Rx = circ_add (Rx, I, Mu)
Bit-reversed with auto-increment register	<code>memXX (Rx++Mu:brev)</code>	EA = Rx.H + bit_reverse (Rx.L) Rx += Mu

<sup>1</sup> EA (effective address) is equivalent to VA (virtual address).

## 5.8.1 Absolute

The absolute addressing mode uses a 32-bit constant value as the effective memory address. For example:

```
R2 = memw(##100000) // Load R2 with word from addr 100000
memw(##200000) = R4 // Store R4 to word at addr 200000
```

## 5.8.2 Absolute-set

The absolute-set addressing mode assigns a 32-bit constant value to the specified general register, then uses the assigned value as the effective memory address. For example:

```
R2 = memw(R1=##400000) // Load R2 with word from addr 400000
                        // and load R1 with value 400000
memw(R3=##600000) = R4 // Store R4 to word at addr 600000
                        // and load R3 with value 600000
```

## 5.8.3 Absolute with register offset

The absolute with register offset addressing mode performs an arithmetic left shift of a 32-bit general register value by the amount specified in a 2-bit unsigned immediate value, and then adds the shifted result to an unsigned 32-bit constant value to create the 32-bit effective memory address. For example:

```
R2 = memh(R3 << #3 + ##100000) // load R2 with signed halfword
                                // from addr [100000 + (R3 << 3)]
```

The 32-bit constant value is the base address, and the shifted result is the byte offset.

**NOTE:** This addressing mode is useful for loading an element from a global table, where the immediate value is the name of the table, and the register holds the index of the element.

## 5.8.4 Global pointer relative

The global pointer relative addressing mode adds an unsigned offset value to the Hexagon processor global data pointer GP to create the 32-bit effective memory address. This addressing mode accesses global and static data in C.

Global pointer relative addresses can be expressed two ways in assembly language:

- By explicitly adding an unsigned offset value to register GP
- By specifying only an immediate value as the instruction operand

For example:

```
R2 = memh(GP+#100) // Load R2 with signed halfword
                  // from [GP + 100 bytes]

R3 = memh(#2000) // Load R3 with signed halfword
                // from [GP + #2000 - _SDA_BASE]
```

Specifying only an immediate value causes the assembler and linker to automatically subtract the value of the special symbol `_SDA_BASE_` from the immediate value, and use the result as the effective offset from GP.

The global data pointer is programmed in the GDP field of register GP (Section 2.2.8). This field contains an unsigned 26-bit value that specifies the most significant 26 bits of the 32-bit global data pointer. The least significant 6 bits of the pointer are always defined as zero.

The memory area referenced by the global data pointer is known as the global data area. It can be up to 512 KB in length, and – because of the way the global data pointer is defined – must be aligned to a 64-byte boundary in virtual memory.

When expressed in assembly language, the offset values used in global pointer relative addressing always specify byte offsets from the global data pointer. The offsets must be integral multiples of the size of the instruction data type.

**Table 5-6 Offset ranges (global pointer relative)**

Data type	Offset range	Offset must be multiple of
doubleword	0 ... 524280	8
word	0 ... 262140	4
halfword	0 ... 131070	2
byte	0 ... 65535	1

**NOTE:** When using global pointer relative addressing, the immediate operand should be a symbol in the `.sdata` or `.sbss` section to ensure that the offset is valid.

## 5.8.5 Indirect

The indirect addressing mode uses a 32-bit value stored in a general register as the effective memory address. For example:

```
R2 = memub(R1)    // load R2 with unsigned byte from addr R1
```

## 5.8.6 Indirect with offset

The indirect with offset addressing mode adds a signed offset value to a general register value to create the 32-bit effective memory address. For example:

```
R2 = memh(R3 + #100)    // load R2 with signed halfword
                        // from [R3 + 100 bytes]
```

When expressed in assembly language, the offset values always specify byte offsets from the general register value. The offsets must be integral multiples of the size of the instruction data type.

**Table 5-7 Offset ranges (indirect with offset)**

Data type	Offset range	Offset must be multiple of
doubleword	-8192 ... 8184	8
word	-4096 ... 4092	4
halfword	-2048 ... 2046	2
byte	-1024 ... 1023	1

**NOTE:** The offset range is smaller for conditional instructions ([Section 5.9](#)).

## 5.8.7 Indirect with register offset

The indirect with register offset addressing mode adds a 32-bit general register value to the result created by performing an arithmetic left shift of a second 32-bit general register value by the amount specified in a 2-bit unsigned immediate value, forming the 32-bit effective memory address. For example:

```
R2 = memh(R3+R4<<#1)    // load R2 with signed halfword
                        // from [R3 + (R4 << 1)]
```

The register values always specify byte addresses.

## 5.8.8 Indirect with auto-increment immediate

The indirect with auto-increment immediate addressing mode uses a 32-bit value stored in a general register to specify the effective memory address. However, after the address is accessed, a signed value (known as the increment) is added to the register so it specifies a different memory address (which is accessed in a subsequent instruction). For example:

```
R2 = memw(R3++#4) // R3 contains the effective address
                    // R3 is then incremented by 4
```

When expressed in assembly language, the increment values always specify byte offsets from the general register value. The offsets must be integral multiples of the size of the instruction data type.

**Table 5-8 Increment ranges (indirect with auto-increment immediate)**

Data type	Increment range	Increment must be multiple of
doubleword	-64 ... 56	8
word	-32 ... 28	4
halfword	-16 ... 14	2
byte	-8 ... 7	1

## 5.8.9 Indirect with auto-increment register

The indirect with auto-increment register addressing mode is functionally equivalent to indirect with auto-increment immediate, but uses a modifier register Mx ([Section 2.2.4](#)) instead of an immediate value to hold the increment. For example:

```
R2 = memw(R0++M1) // The effective addr is the value of R0.
                   // Next, M1 is added to R0 and the result
                   // is stored in R0.
```

When auto-incrementing with a modifier register, the increment is a signed 32-bit value which is added to the general register. This offers two advantages over auto-increment immediate:

- A larger increment range
- Variable increments (since the modifier register can be programmed at runtime)

The increment value always specifies a byte offset from the general register value.

**NOTE:** The signed 32-bit increment range is identical for all instruction data types (doubleword, word, halfword, byte).

## 5.8.10 Circular with auto-increment immediate

The circular with auto-increment immediate addressing mode is a variant of indirect with auto-increment addressing – it accesses data buffers in a modulo wrap-around fashion. Circular addressing is commonly used in data stream processing.

Circular addressing is expressed in assembly language with the address modifier “:circ(Mx)”, where Mx specifies a modifier register that is programmed to specify the circular buffer (Section 2.2.4). For example:

```
R0 = memb(R2++#4:circ(M0)) // Load from R2 in circ buf specified
                          // by M0
memw(R2++#8:circ(M1)) = R0 // Store to R2 in circ buf specified
                          // by M1
```

Circular addressing is set up by programming the following elements:

- The Length field of the M<sub>x</sub> register is set to the length (in bytes) of the circular buffer to access. A circular buffer can be from 4 to (128K-1) bytes long.
- Bits 27:24 of the M<sub>x</sub> register are always set to 0.
- The circular start register CS<sub>x</sub> that corresponds to M<sub>x</sub> (CS0 for M0, CS1 for M1) is set to the start address of the circular buffer.

In circular addressing, after memory is accessed at the address specified in the general register, the general register is incremented by the immediate increment value and then modulo'd by the circular buffer length to implement wrap-around access of the buffer.

When expressed in assembly language, the increment values always specify byte offsets from the general register value. The offsets must be integral multiples of the size of the instruction data type.

**Table 5-9 Increment ranges (circular with auto-increment immediate addressing)**

Data type	Increment range	Increment must be multiple of
doubleword	-64 ... 56	8
word	-32 ... 28	4
halfword	-16 ... 14	2
byte	-8 ... 7	1

When programming a circular buffer, the following rules apply:

- The start address must be aligned to the native access size of the buffer elements.
- $ABS(\text{Increment}) < \text{Length}$ . The absolute value of the increment must be less than the buffer length.
- $\text{Access size} < (\text{Length}-1)$ . The memory access size (1 for byte, 2 for halfword, 4 for word, 8 for doubleword) must be less than (Length-1).
- Buffers must not wrap around in the 32-bit address space.

**NOTE:** If any of these rules are not followed, the execution result is undefined.



The following example sets up and accesses a 150-byte circular buffer:

```
R4.H = #0                // M0[27:24]= 0x0
R4.L = #150             // length = 150
M0 = R4
R2 = #cbuf              // start addr = cbuf
CS0 = R2
R0 = memb(R2++#4:circ(M0)) // Load byte from circ buf
                          // specified by M0/CS0
                          // inc R2 by 4 after load
                          // wrap R2 around if >= 150
```

The following C function describes the behavior of the circular add function:

```
unsigned int
fcircadd(unsigned int pointer, int offset,
          unsigned int M_reg, unsigned int CS_reg)
{
    unsigned int length;
    int new_pointer, start_addr, end_addr;

    length = (M_reg&0x01ffff); // Lower 17 bits gives buffer size
    new_pointer = pointer+offset;
    start_addr = CS_reg;
    end_addr = CS_reg + length;
    if (new_pointer >= end_addr) {
        new_pointer -= length;
    } else if (new_pointer < start_addr) {
        new_pointer += length;
    }
    return (new_pointer);
}
```

### 5.8.11 Circular with auto-increment register

The circular with auto-increment register addressing mode is functionally equivalent to circular with auto-increment immediate, but uses a register instead of an immediate value to hold the increment.

Register increments are specified in circular addressing instructions by using the symbol  $I$  as the increment (instead of an immediate value). For example:

```
R0 = memw(R2++I:circ(M1)) // Load byte with incr of I*4 from
                          // circ buf specified by M1/CS1
```

When auto-incrementing with a register, the increment is a signed 11-bit value that is added to the general register. This offers two advantages over circular addressing with immediate increments:

- Larger increment ranges
- Variable increments (since the increment register can be programmed at runtime)

The circular register increment value is programmed in the  $I$  field of the modifier register  $M_x$  (Section 2.2.4) as part of setting up the circular data access. This register field holds the signed 11-bit increment value.

Increment values are expressed in units of the buffer element data type, and are automatically scaled at runtime to the proper data access size.

**Table 5-10 Increment ranges (circular with auto-increment register addressing)**

Data type	Increment range	Increment must be multiple of
doubleword	-8192 ... 8184	8
word	-4096 ... 4092	4
halfword	-2048 ... 2046	2
byte	-1024 ... 1023	1

When programming a circular buffer (with either a register or immediate increment), all the rules that apply to circular addressing must be followed – for details see [Section 5.8.10](#).

**NOTE:** If any of these rules are not followed, the execution result is undefined.

## 5.8.12 Bit-reversed with auto-increment register

The bit-reversed with auto-increment register addressing mode is a variant of indirect with auto-increment addressing – it accesses data buffers using an address value which is the bit-wise reversal of the value stored in the general register. Bit-reversed addressing is used in fast Fourier transforms (FFT) and Viterbi encoding.

The bit-wise reversal of a 32-bit address value is defined as follows:

- The lower 16 bits are transformed by exchanging bit 0 with bit 15, bit 1 with bit 14, and so on.
- The upper 16 bits remain unchanged.

Bit-reversed addressing is expressed in assembly language with the address modifier “:brev”. For example:

```
R2 = memub(R0++M1:brev) // The address is (R0.H | bitrev(R0.L))
                        // The original R0 (not reversed) is added
                        // to M1 and written back to R0
```

The initial values for the address and increment must be set in bit-reversed form, with the hardware bit-reversing the bit-reversed address value to form the effective address.

The buffer length for a bit-reversed buffer must be an integral power of 2, with a maximum length of 64K bytes.

To support bit-reversed addressing, buffers must be properly aligned in memory. A bit-reversed buffer is properly aligned when its starting byte address is aligned to a power of 2 greater than or equal to the buffer size (in bytes). For example:

```
int bitrev_buf[256] __attribute__((aligned(1024)));
```

The bit-reversed buffer declared above is aligned to 1024 bytes because the buffer size is 1024 bytes (256 integer words  $\times$  4 bytes), and 1024 is an integral power of 2.

The buffer location pointer for a bit-reversed buffer must be initialized so the least-significant 16 bits of the address value are bit-reversed.

The increment value must be initialized to the following value:

```
bitreverse(buffer_size_in_bytes / 2)
```

...where `bitreverse` is defined as bit-reversing the least-significant 16 bits while leaving the remaining bits unchanged.

**NOTE:** To simplify the initialization of the bit-reversed pointer, bit-reversed buffers can be aligned to a 64K byte boundary. This initializes the bit-reversed pointer to the base address of the bit-reversed buffer, with no bit-reversing required for the least-significant 16 bits of the pointer value (which are set to 0 by the 64K alignment).

Because buffers allocated on the stack only have an alignment of 8 bytes or less, in most cases bit-reversed buffers should not be declared on the stack.

After a bit-reversed memory access is complete, the general register is incremented by the register increment value. The value in the general register is never affected by the bit-reversal that is performed as part of the memory access.

**NOTE:** The Hexagon processor supports only register increments for bit-reversed addressing – it does not support immediate increments.

## 5.9 Conditional load/stores

Some load and store instructions can be executed conditionally based on predicate values which were set in a previous instruction. The compiler generates conditional loads and stores to increase instruction-level parallelism.

Conditional loads and stores are expressed in assembly language with the instruction prefix “`if (pred_expr)`”, where `pred_expr` specifies a predicate register expression ([Section 6.1](#)). For example:

```
if (P0) R0 = memw(R2)           // Conditional load
if (!P2) memh(R3 + #100) = R1   // Conditional store
if (P1.new) R3 = memw(R3++#4)   // Conditional load
```

Not all addressing modes are supported in conditional loads and stores. [Table 5-11](#) shows which modes are supported.

**Table 5-11 Addressing modes (conditional load/store)**

Addressing mode	Conditional
Absolute	Yes
Absolute-set	No
Absolute with register offset	No

**Table 5-11 Addressing modes (conditional load/store)**

Addressing mode	Conditional
Global pointer relative	No
Indirect	Yes
Indirect with offset	Yes
Indirect with register offset	Yes
Indirect with auto-increment immediate	Yes
Indirect with auto-increment register	No
Circular with auto-increment immediate	No
Circular with auto-increment register	No
Bit-reversed with auto-increment register	No

When a conditional load or store instruction uses [Indirect with offset](#) addressing mode, the offset range is smaller than the range normally defined for indirect-with-offset addressing.

**Table 5-12 Conditional and normal offset ranges (indirect with offset addressing)**

Data type	Offset range (conditional)	Offset range (normal)	Offset must be multiple of
doubleword	0 ... 504	-8192 ... 8184	8
word	0 ... 252	-4096 ... 4092	4
halfword	0 ... 126	-2048 ... 2046	2
byte	0 ... 63	-1024 ... 1023	1

**NOTE:** For more information on conditional execution see [Chapter 6](#).

## 5.10 Cache memory

The Hexagon processor has a cache-based memory architecture:

- A level 1 instruction cache holds recently-fetched instructions.
- A level 1 data cache holds recently-accessed data memory.

Load/store operations that access memory through the level 1 caches are referred to as cached accesses.

Load/stores that bypass the level 1 caches are referred to as uncached accesses.

The memory management unit (MMU) of the Hexagon processor can configure specific memory areas to perform cached or uncached accesses. The operating system is responsible for programming the MMU.

Two types of caching are supported (as cache modes):

- Write-through caching keep the cache data consistent with external memory by always writing to the memory any data that is stored in the cache.

- Write-back caching stores data in the cache without being immediately written to external memory. Cached data that is inconsistent with external memory is referred to as dirty.

The Hexagon processor includes dedicated cache maintenance instructions that push dirty data out to external memory.

### 5.10.1 Uncached memory

In some cases, load/store operations need to bypass the cache memories and be serviced externally (for example, when accessing memory-mapped I/O, registers, and peripheral devices, or other system defined entities). The operating system is responsible for configuring the MMU to generate uncached memory accesses.

Uncached memory is categorized into two distinct types:

- Device-type is for accessing memory that has side-effects (such as a memory-mapped FIFO peripheral). The hardware ensures that interrupts do not cancel a pending device access. The hardware does not reorder device accesses. Mark peripheral control registers as device-type.
- Uncached-type is for memory-like memory. No side effects are associated with an access. The hardware can load from uncached memory multiple times. The hardware can reorder uncached accesses.

For instruction accesses, device-type memory is functionally identical to uncached-type memory. For data accesses, they are different.

Code can execute directly from the L2 cache, bypassing the L1 cache.

### 5.10.2 Tightly coupled memory

The Hexagon processor supports tightly-coupled instruction memory at Level 1, which is defined as memory with similar access properties to the instruction cache.

Tightly-coupled memory is also supported at level 2, which is defined as backing store to the primary caches.

### 5.10.3 Cache maintenance operations

The Hexagon processor includes dedicated cache maintenance instructions that invalidate cache data or push dirty data out to external memory.

The cache maintenance instructions operate on specific memory addresses. If the instruction causes an address error (due to a privilege violation), the processor raises an exception.

**NOTE:** The exception to this rule is the `dcfetch` operation, which never causes a processor exception.

Whenever maintenance operations are performed on the instruction cache, the `isync` instruction ([Section 5.11](#)) must execute immediately afterwards. This instruction ensures that subsequent instructions observe the maintenance operations.

**Table 5-13 Cache instructions (user-level)**

Syntax	Permitted In packet	Operation
<code>icinva (Rs)</code>	Solo <sup>1</sup>	Instruction cache invalidate. Look up instruction cache at address Rs. If the address is in the cache, invalidate it.
<code>dccleaninva (Rs)</code>	Slot 1 empty or ALU32 only	Data cache clean and invalidate. Look up data cache at address Rs. If the address is in the cache and has dirty data, flush that data out to memory. The cache line is then invalidated, whether or not dirty data was written.
<code>dccleana (Rs)</code>	Slot 1 empty or ALU32 only	Data cache clean. Look up data cache at address Rs. If the address is in the cache and has dirty data, flush that data out to memory.
<code>dcinva (Rs)</code>	Slot 1 empty or ALU32 only	Equivalent to <code>dccleaninva(Rs)</code> .
<code>dcfetch (Rs)</code>	Normal <sup>2</sup>	Data cache prefetch. Prefetch data at address Rs into the data cache. NOTE - This instruction does not cause an exception.
<code>l2fetch (Rs, Rt)</code>	ALU32 or XTYPE only	L2 cache prefetch. Prefetch data from memory specified by Rs and Rt into L2 cache.

<sup>1</sup> *Solo* means that the instruction must not be grouped with other instructions in a packet.

<sup>2</sup> *Normal* means that the normal instruction-grouping constraints apply.

## 5.10.4 L2 cache operations

Cache maintenance operations operate on both the L1 and L2 caches.

The data cache coherency operations (including clean, invalidate, and clean and invalidate) affect both the L1 and L2 caches, and ensure that the memory hierarchy remains coherent.

However, the instruction cache invalidate operation affects only the L1 cache. Therefore, invalidating instructions that might be in the L1 or L2 caches requires a two-step procedure:

1. Use `icinva` to invalidate instructions from the L1 cache.
2. Use `dcinva` separately to invalidate instructions from the L2 cache.

## 5.10.5 Cache line zero

The Hexagon processor includes the instruction `dczeroa`. This instruction allocates a line in the L1 data cache and clears it (by storing all zeros). The behavior is as follows:

- The Rs register value must be 32-byte aligned. If it is unaligned, the processor raises an unaligned error exception.

- For a cache hit, the specified cache line is cleared (written with all zeros) and made dirty.
- For a cache miss, the specified cache line is not fetched from external memory. Instead, the line is allocated in the data cache, cleared, and made dirty.

This instruction is useful in optimizing write-only data. It allows for the use of write-back pages – which are the most power and performance efficient – without the need to initially fetch the line to write. This removes unnecessary read bandwidth and latency.

**NOTE:** The `dczeroa` operation has the same exception behavior as write-back stores.

A packet with `dczeroa` must have slot 1 either empty or containing an ALU32 instruction.

## 5.10.6 Cache prefetch

The Hexagon processor supports the following types of cache prefetching.

### 5.10.6.1 Hardware-based instruction cache prefetching

L1 and L2 instruction cache prefetching can be enabled or disabled on a per-thread basis by setting the HFI field in the [User status register](#).

### 5.10.6.2 Software-based data cache prefetching

The Hexagon processor includes the `dcfetch` instruction, which queries the L1 data cache based on the address specified in the instruction:

- If the address is present in the cache, no action is taken.
- If the cache line for the address is missing, the processor attempts to fill the cache line from the next level of memory. The thread does not stall, but rather continues executing while the cache line fill occurs in the background.
- If the address is invalid, no exception is generated and the `dcfetch` instruction is treated as a NOP.

### 5.10.6.3 Software-based L2fetch

More powerful L2 prefetching – of data or instructions – is provided by the `l2fetch` instruction, which specifies an area of memory that is prefetched by the Hexagon processor's hardware prefetch engine. `l2fetch` specifies two registers (`Rs` and `Rt`) as operands. `Rs` contains the 32-bit virtual start address of the memory area to prefetch. `Rt` contains three bit fields that further specify the memory area:

- `Rt[15:8]` – `Width`, specifies the width (in bytes) of a block of memory to fetch.
- `Rt[7:0]` – `Height`, specifies the number of `Width`-sized blocks to fetch.
- `Rt[31:16]` – `Stride`, specifies an unsigned byte offset that increments the pointer after each `Width`-sized block is fetched.

The I2fetch instruction is nonblocking: it initiates a prefetch operation that is performed in the background by the prefetch engine while the thread continues to execute Hexagon processor instructions.

The prefetch engine requests all lines in the specified memory area. If the line(s) of interest are already resident in the L2 cache, the prefetch engine performs no action. If the lines are not in the L2 cache, the prefetch engine attempts to fetch them.

The prefetch engine makes a best effort to prefetch the requested data, and attempts to perform prefetching at a lower priority than demand fetches. This prevents the prefetch engine from adding bus traffic when the system is under a heavy load.

If a program executes an I2fetch instruction while the prefetch operation from a previous I2fetch is still active, the prefetch engine halts the current prefetch operation.

**NOTE:** Executing I2fetch with any bit field operand programmed to zero cancels prefetch activity.

The status of the current prefetch operation is maintained in the PFA field of the user status register. This field can determine whether a prefetch operation has completed.

With respect to MMU permissions and error checking, the I2fetch instruction behaves similarly to a load instruction. If the virtual address causes a processor exception, the exception is taken. This differs from the dcfetch instruction, which is treated as a NOP in the presence of a translation/protection error.

**NOTE:** Prefetches are dropped when the generated prefetch address resides on a different page than the start address. The programmer must use sufficiently large pages to ensure this does not occur.



Figure 5-2 shows two examples of using the l2fetch instruction. The first shows a box prefetch, where a 2D range of memory is defined within a larger frame. The second example shows a prefetch for a large linear memory area of size (Lines \* 128).

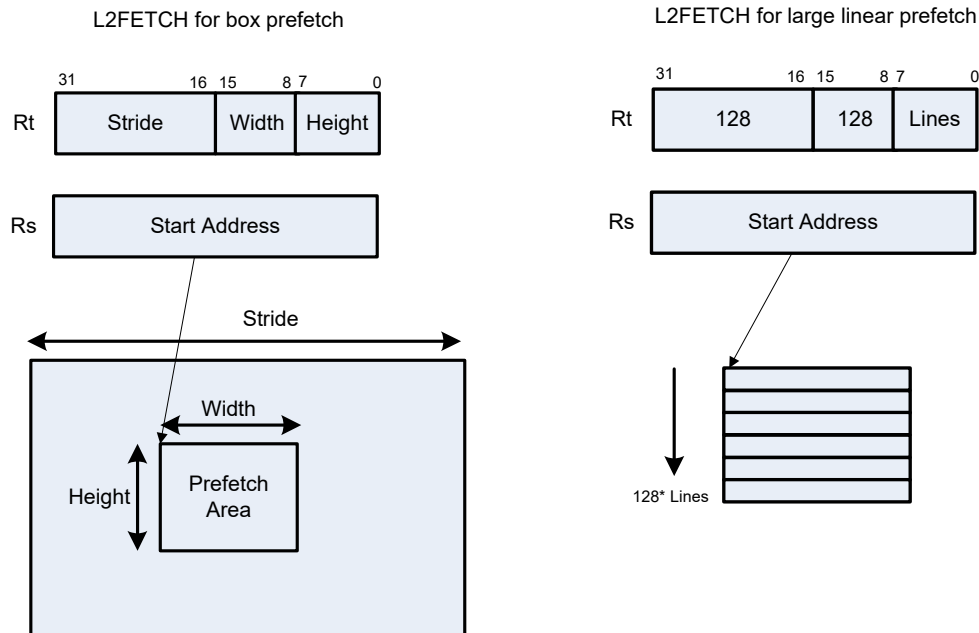


Figure 5-2 L2fetch instruction

#### 5.10.6.4 Hardware-based data cache prefetching

L1 data cache prefetching is enabled or disabled on a per-thread basis by setting the HFD field in the [User status register](#).

When data cache prefetching is enabled, the Hexagon processor observes patterns of data cache misses, and attempts to predict future misses based on any recurring patterns of misses where the addresses are separated by a constant stride. If such patterns are found, the processor attempts to automatically prefetch future cache lines.

Data cache prefetching is user-enabled at four levels of aggressiveness:

- HFD = 00: No prefetching
- HFD = 01: Prefetch up to four lines for misses originating from a load, with a post-update addressing mode that occurs within a hardware loop
- HFD = 10: Prefetch up to four lines for misses originating from loads that occur within a hardware loop
- HFD = 11: Prefetch up to eight lines for misses originating from loads

## 5.11 Memory ordering

Some devices might require synchronization of stores and loads when they are accessed. In this case, a set of processor instructions enable programmer control of the synchronization and ordering of memory accesses.

**Table 5-14 Memory ordering instructions**

Syntax	Operation
<code>isync</code>	Instruction synchronize. This instruction should execute after any instruction cache maintenance operation.
<code>syncht</code>	Synchronize transactions. Perform heavyweight synchronization. Ensure that previous program transactions (for example, <code>memw_locked</code> , <code>cached</code> and <code>uncached load/store</code> ) are complete before execution resumes past this instruction. The <code>syncht</code> instruction ensures that outstanding memory operations from all threads are complete before the <code>syncht</code> instruction is committed.
<code>barrier</code>	Set memory barrier. Ensure proper ordering between the program accesses performed before the instruction and those performed after the instruction. All accesses before the barrier are globally observable before any access occurring after the barrier can be observed. The barrier instruction ensures that all outstanding memory operations from the thread executing the barrier are complete before the instruction is committed.

Data memory accesses and program memory accesses are treated separately and held in separate caches. Software should ensure coherency between data and program code if necessary.

For example, with generated or self-modified code, the modified code is placed in the data cache and can be inconsistent with program cache. The software must explicitly force modified data cache lines to memory (either by using a write-through policy, or through explicit cache clean instructions). Use a barrier instruction to ensure completion of the stores. Finally, invalidate relevant instruction cache contents so the new instructions can be refetched.

Here is the recommended code sequence to change and execute an instruction:

```

ICINVA(R1)      // Clear code from instruction cache
ISYNC           // Ensure that ICINVA is finished
MEMW(R1)=R0    // Write the new instruction
DCCLEANINVA(R1) // Force data out of data cache
SYNCHT         // Ensure that it is in memory
JUMPR R1       // Can now execute code at R1

```

**NOTE:** The memory-ordering instructions must not be grouped with other instructions in a packet, otherwise the behavior is undefined.

This code sequence differs from the one used in previous processor versions.

## 5.12 Atomic operations

The Hexagon processor includes a load locked / store conditional (LL/SC) mechanism to provide the atomic read-modify-write operation that is necessary to implement synchronization primitives such as semaphores and mutexes.

These primitives synchronize the execution of different software programs running concurrently on the Hexagon processor. They can also provide atomic memory support between the Hexagon processor and external blocks.

**Table 5-15 Atomic instructions**

Syntax	Description
<code>Rd = memw_locked(Rs)</code>	Load locked word. Reserve lock on word at address Rs.
<code>memw_locked(Rs, Pd) = Rt</code>	Store conditional word. If no other atomic operation has been performed at the address (that is, atomicity is ensured), perform the store to the word at address Rs and return TRUE in Pd; otherwise return FALSE. TRUE indicates that the LL and SC operations were performed atomically.
<code>Rdd = memd_locked(Rs)</code>	Load locked doubleword. Reserve lock on doubleword at address Rs.
<code>memd_locked(Rs, Pd) = Rtt</code>	Store conditional doubleword. If no other atomic operation has been performed at the address (that is, atomicity is ensured), perform the store to the doubleword at address Rs and return TRUE in Pd; otherwise return FALSE. TRUE indicates that the LL and SC operations have been performed atomically.

Here is the recommended code sequence to acquire a mutex:

```
// Assume mutex address is held in R0
// Assume R1,R3,P0,P1 are scratch

lockMutex:
    R3 = #1
lock_test_spin:
    R1 = memw_locked(R0)           // Do normal test to wait
    P1 = cmp.eq(R1,#0)           // for lock to be available
    if (!P1) jump lock_test_spin
    memw_locked(R0,P0) = r3      // Do store conditional (SC)
    if (!P0) jump lock_test_spin // was LL and SC done atomically?
```

Here is the recommended code sequence to release a mutex:

```
// Assume mutex address is held in R0
// Assume R1 is scratch

R1 = #0
memw(R0) = R1
```

Atomic memX\_locked operations are supported for external accesses that use the advanced extensible interface (AXI) bus and support atomic operations. To perform load-locked operations with external memory, the operating system must define the memory page as uncacheable, otherwise the processor behavior is undefined.

If a load locked operation is performed on an address that does not support atomic operations, the behavior is undefined.

For atomic operations on cacheable memory, the page attributes must be set to cacheable and write-back, otherwise the behavior is undefined. Cacheable memory must be used when threads must synchronize with each other.

**NOTE:** External memX\_locked operations are not supported on the AHB. If they are performed on the AHB, the behavior is undefined.

# 6 Conditional Execution

---

The Hexagon processor uses a conditional execution model based on compare instructions that set predicate bits in one of four 8-bit predicate registers (P0 through P3). These predicate bits can conditionally execute certain instructions.

Conditional scalar operations examine only the least-significant bit in a predicate register, while conditional vector operations examine multiple bits in the register.

Branch instructions are the main consumers of the predicate registers.

## 6.1 Scalar predicates

Scalar predicates are 8-bit values in conditional instructions to represent truth values:

- 0xFF represents true
- 0x00 represents false

The Hexagon processor provides the four 8-bit predicate registers P0 through P3 to hold scalar predicates ([Section 2.2.5](#)). These registers are assigned values by the predicate-generating instructions, and examined by the predicate-consuming instructions.

### 6.1.1 Generating scalar predicates

The following instructions generate scalar predicates:

- Compare byte, halfword, word, doubleword
- Compare single- and double-precision floating point
- Classify floating-point value
- Compare bitmask
- Bounds check
- TLB match
- Store conditional

**Table 6-1 Scalar predicate-generating instructions**

Syntax	Operation
Pd = cmpb.eq(Rs, {Rt, #u8}) Pd = cmph.eq(Rs, {Rt, #s8}) Pd = [!]cmp.eq(Rs, {Rt, #s10}) Pd = cmp.eq(Rss, Rtt) Pd = sfcmp.eq(Rs, Rt) Pd = dfcmp.eq(Rss, Rtt)	Equal (signed). Compare register Rs to Rt or a signed immediate for equality. Assign Pd the resulting truth value.
Pd = cmpb.gt(Rs, {Rt, #s8}) Pd = cmph.gt(Rs, {Rt, #s8}) Pd = [!]cmp.gt(Rs, {Rt, #s10}) Pd = cmp.gt(Rss, Rtt) Pd = sfcmp.gt(Rs, Rt) Pd = dfcmp.gt(Rss, Rtt)	Greater than (signed). Compare register Rs to Rt or a signed immediate for signed greater than. Assign Pd the resulting truth value.
Pd = cmpb.gtu(Rs, {Rt, #u7}) Pd = cmph.gtu(Rs, {Rt, #u7}) Pd = [!]cmp.gtu(Rs, {Rt, #u9}) Pd = cmp.gtu(Rss, Rtt)	Greater than (unsigned). Compare register Rs to Rt or an unsigned immediate for unsigned greater than. Assign Pd the resulting truth value.
Pd = cmp.ge(Rs, #s8) Pd = sfcmp.ge(Rs, Rt) Pd = dfcmp.ge(Rss, Rtt)	Greater than or equal (signed). Compare register Rs to Rt or a signed immediate for signed greater than or equal. Assign Pd the resulting truth value.
Pd = cmp.geu(Rs, #u8)	Greater than or equal (unsigned). Compare register Rs to an unsigned immediate for unsigned greater than or equal. Assign Pd the resulting truth value.
Pd = cmp.lt(Rs, Rt)	Less than (signed). Compare register Rs to Rt for signed less than. Assign Pd the resulting truth value.
Pd = cmp.ltu(Rs, Rt)	Less than (unsigned). Compare register Rs to Rt for unsigned less than. Assign Pd the resulting truth value.
Pd = sfcmp.uo(Rs, Rt) Pd = dfcmp.uo(Rss, Rtt)	Unordered (signed). Determine if register Rs or Rt is set to the value NaN. Assign Pd the resulting truth value.
Pd=sfclass(Rs, #u5) Pd=dfclass(Rss, #u5)	Classify value (signed). Determine if register Rs is set to any of the specified classes. Assign Pd the resulting truth value.
Pd = [!]tstbit(Rs, {Rt, #u5})	Test if bit set. Rt or an unsigned immediate specifies a bit position. Test if the bit in Rs that is specified by the bit position is set. Assign Pd the resulting truth value.
Pd = [!]bitsclr(Rs, {Rt, #u6})	Test if bits clear. Rt or an unsigned immediate specifies a bitmask. Test if the bits in Rs that are specified by the bitmask are all clear. Assign Pd the resulting truth value.

**Table 6-1 Scalar predicate-generating instructions (cont.)**

<code>Pd = [!]bitsset(Rs,Rt)</code>	Test if bits set. Rt specifies a bitmask. Test if the bits in Rs that are specified by the bitmask are all set. Assign Pd the resulting truth value.
<code>memw_locked(Rs,Pd) = Rt</code> <code>memd_locked(Rs,Pd) = Rtt</code>	Store conditional. If no other atomic operation has been performed at the address (atomicity is ensured), perform the store to the word at address Rs. Assign Pd the resulting truth value.
<code>Pd = boundscheck(Rs,Rtt)</code>	Bounds check. Determine if Rs falls in the numeric range defined by Rtt. Assign Pd the resulting truth value.
<code>Pd = tlbmatch(Rss,Rt)</code>	Determine if TLB entry in Rss matches the ASID:PPN specified in Rt. Assign Pd the resulting truth value.

**NOTE:** One of the compare instructions (`cmp.eq`) includes a variant that stores a binary predicate value (0 or 1) in a general register not a predicate register.

## 6.1.2 Consuming scalar predicates

Certain instructions can be conditionally executed based on the value of a scalar predicate (or alternatively specify a scalar predicate as an input to their operation).

The conditional instructions that consume scalar predicates examine only the least-significant bit of the predicate value. In the simplest case, this bit value directly determines whether the instruction executes:

- 1 indicates that the instruction executes
- 0 indicates that the instruction does not execute

If a conditional instruction includes the operator `!` in its predicate expression, the logical negation of the bit value determines whether the instruction is executed.

Conditional instructions are expressed in assembly language with the instruction prefix “`if` (*pred\_expr*)”, where *pred\_expr* specifies the predicate expression. For example:

```
if (P0) jump target           // Jump if P0 is true
if (!P2) R2 = R5              // Assign register if !P2 is true
if (P1) R0 = sub(R2,R3)      // Conditionally subtract if P1
if (P2) R0 = memw(R2)        // Conditionally load word if P2
```

The following instructions can be used as conditional instructions:

- Jumps and calls ([Section 8.3](#))
- Many load and store instructions ([Section 5.9](#))
- Logical instructions (including AND/OR/XOR)
- Shift halfword
- 32-bit add/subtract by register or short immediate

- Sign and zero extend
- 32-bit register transfer and 64-bit combine word
- Register transfer immediate
- Deallocate frame and return

When a conditional load or store executes and the predicate expression is false, the instruction is canceled (including any exceptions that might occur). For example, if a conditional load uses an address with a memory permission violation, and the predicate expression is false, the load does not execute and the exception is not raised.

The mux instruction accepts a predicate as one of its basic operands:

```
Rd = mux (Ps, Rs, Rt)
```

The mux instruction selects either Rs or Rt based on the least significant bit in Ps. If the least-significant bit in Ps is a 1, Rd is set to Rs, otherwise it is set to Rt.

### 6.1.3 Auto-AND predicates

If multiple compare instructions in a packet write to the same predicate register, the result is the logical AND of the individual compare results. For example:

```
{
  P0 = cmp (A)                // If A && B,  jump
  P0 = cmp (B)
  if (P0.new) jump:T taken_path
}
```

To perform the corresponding OR operation, the following instructions can compute the negation of an existing compare (using De Morgan's law):

- Pd = !cmp.{eq,gt}(Rs, {#s10,Rt} )
- Pd = !cmp.gtu(Rs, {#u9,Rt} )
- Pd = !tstbit(Rs, {#u5,Rt} )
- Pd = !bitsclr(Rs, {#u6,Rt} )
- Pd = !bitsset(Rs,Rt)

Auto-AND predicates have the following restrictions:

- If a packet contains endloopN, it cannot perform an auto-AND with predicate register P3.
- If a packet contains a register transfer from a general register to a predicate register, no other instruction in the packet can write to the same predicate register. As a result, a register transfer to P3:0 or C5:4 cannot be grouped with any other predicate-writing instruction.
- The instructions spNloop0, decbin, tlbmatch, memw\_locked, memd\_locked, add:carry, sub:carry, sfcmp, and dfcmp cannot be grouped with another instruction that sets the same predicate register.

**NOTE:** A register transfer from a predicate register to a predicate register has the same auto-AND behavior as a compare instruction.



## 6.1.4 Dot-new predicates

The Hexagon processor can generate and use a scalar predicate in the same instruction packet (Section 3.3). This feature is expressed in assembly language by appending the suffix “.new” to the specified predicate register. For example:

```
if (P0.new) R3 = memw(R4)
```

The following C statement and the corresponding assembly code that is generated from it by the compiler is an example of how to use dot-new predicates.

### C statement

```
if (R2 == 4)
    R3 = *R4;
else
    R5 = 5;
```

### Assembly code

```
{
    P0 = cmp.eq(R2,#4)
    if (P0.new) R3 = memw(R4)
    if (!P0.new) R5 = #5
}
```

In the assembly code, a scalar predicate is generated and then consumed twice within the same instruction packet.

The following conditions apply to using dot-new predicates:

- The predicate must be generated by an instruction in the same packet. The assembler normally enforces this restriction, but if the processor executes a packet that violates this restriction, the execution result is undefined.
- A single packet can contain both the dot-new and normal forms of predicates. The normal form examines the old value in the predicate register, rather than the newly-generated value. For example:

```
{
    P0 = cmp.eq(R2,#4)
    if (P0.new) R3 = memw(R4) // Use newly-generated P0 value
    if (P0) R5 = #5           // Use previous P0 value
}
```

## 6.1.5 Dependency constraints

Two instructions in an instruction packet should not write to the same destination register (Section 3.3.5). An exception to this rule is when the two instructions are conditional, and only one of them ever has the predicate expression value true when the packet executes.

For example, the following packet is valid as long as P2 and P3 never both evaluate to true when the packet is executed:

```
{
  if (P2) R3 = #4      // P2, P3, or both must be false
  if (P3) R3 = #7
}
```

Because predicate values change at runtime, the programmer is responsible for ensuring that such packets are always valid during program execution. If they are invalid, the processor takes the following actions:

- When writing to general registers, an error exception is raised.
- When writing to predicate or control registers, the result is undefined.

## 6.2 Vector predicates

The predicate registers are also used for conditional vector operations. Unlike scalar predicates, vector predicates contain multiple truth values which are generated by vector predicate-generating operations.

For example, a vector compare instruction compares each element of a vector and assigns the compare results to a predicate register. Each bit in the predicate vector contains a truth value indicating the outcome of a separate compare performed by the vector instruction.

The vector mux instruction uses a vector predicate to selectively merge elements from two separate vectors into a single destination vector. This operation is useful for enabling the vectorization of loops with control flow (branches).

The vector instructions that use predicates are described in the following sections.

### 6.2.1 Vector compare

A vector compare instruction inputs two 64-bit vectors, performs separate compares for each pair of vector elements, and generates a predicate value which contains a bit vector of truth values.

In Figure 6-1 two 64-bit vectors of bytes (contained in Rss and Rtt) are being compared. The result is assigned as a vector predicate to the destination register Pd.

In the example vector predicate shown in Figure 6-1, every other compare result in the predicate is true (for example, 1).

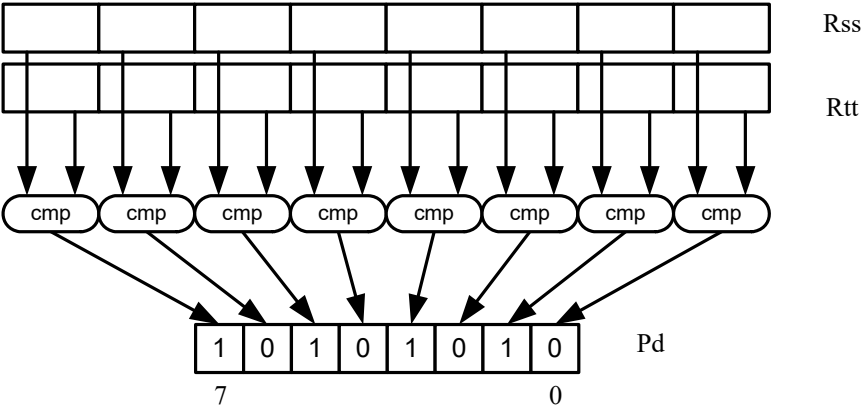


Figure 6-1 Vector byte compare

Figure 6-2 shows how a vector halfword compare generates a vector predicate. Two 64-bit vectors of halfwords are being compared. The result is assigned as a vector predicate to the destination register Pd.

Because a vector halfword compare yields only four truth values, each truth value is encoded as two bits in the generated vector predicate.

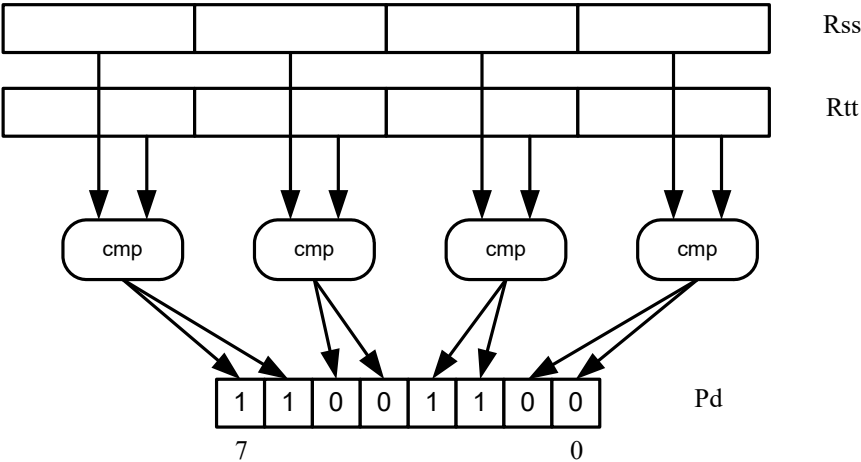


Figure 6-2 Vector halfword compare

## 6.2.2 Vector mux instruction

A vector mux instruction conditionally selects the elements from two vectors. The instruction takes as input two source vectors and a predicate register. For each byte in the vector, the corresponding bit in the predicate register is used to choose from one of the two input vectors. The combined result is written to the destination register.

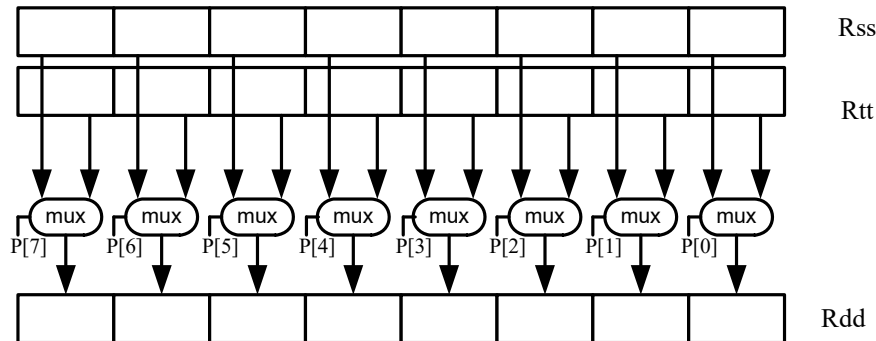


Figure 6-3 Vector mux instruction

Table 6-2 Vector mux instruction

Syntax	Operation
$Rdd = vmux(Ps, Rss, Rtt)$	Select bytes from Rss and Rtt

Changing the order of the source operands in a mux instruction enables formation of both senses of the result. For example:

```
R1:0 = vmux(P0, R3:2, R5:4)    // Choose bytes from R3:2 if true
R1:0 = vmux(P0, R5:4, R3:2)    // Choose bytes from R3:2 if false
```

**NOTE:** By replicating the predicate bits generated by word or halfword compares, the vector mux instruction can select words or halfwords.

## 6.2.3 Using vector conditionals

Vector conditional support is used to vectorize loops with conditional statements.

Consider the following C statement:

```
for (i=0; i<8; i++) {
    if (A[i]) {
        B[i] = C[i];
    }
}
```

Assuming arrays of bytes, this code can be vectorized as follows:

```
R1:0 = memd(R_A)                // R1:0 holds A[7]-A[0]
R3 = #0                          // Clear R3:2
R2 = #0
P0 = vcmpb.eq(R1:0, R3:2)        // Compare bytes in A to zero
```

```

R5:4 = memd(R_B)           // R5:4 holds B[7]-B[0]
R7:6 = memd(R_C)           // R7:6 holds C[7]-C[0]
R3:2 = vmux(P0,R7:6,R5:4)  // if (A[i]) B[i]=C[i]
memd(R_B) = R3:2           // store B[7]-B[0]

```

## 6.3 Predicate operations

The Hexagon processor provides a set of operations for manipulating and moving predicate registers.

**Table 6-3 Predicate register instructions**

Syntax	Operation
Pd = Ps	Transfer predicate Ps to Pd
Pd = Rs	Transfer register Rs to predicate Pd
Rd = Ps	Transfer predicate Ps to register Rd
Pd = and(Ps, [!]Pt)	Set Pd to bitwise AND of Ps and [NOT] Pt
Pd = or(Ps, [!]Pt)	Set Pd to bitwise OR of Ps and [NOT] Pt
Pd = and(Ps, and(Pt, [!]Pu)	Set Pd to AND of Ps and (AND of Pt and [NOT] Pu)
Pd = and(Ps, or(Pt, [!]Pu)	Set Pd to AND of Ps and (OR of Pt and [NOT] Pu)
Pd = or(Ps, and(Pt, [!]Pu)	Set Pd to OR of Ps and (AND of Pt and [NOT] Pu)
Pd = or(Ps, or(Pt, [!]Pu)	Set Pd to OR of Ps and (OR of Pt and [NOT] Pu)
Pd = not(Ps)	Set Pd to bitwise inversion of Ps
Pd = xor(Ps, Pt)	Set Pd to bitwise exclusive OR of Ps and Pt
Pd = any8(Ps)	Set Pd to 0xFF if any bit in Ps is 1, 0x00 otherwise
Pd = all8(Ps)	Set Pd to 0x00 if any bit in Ps is 0, 0xFF otherwise

**NOTE:** These instructions belong to instruction class CR.

Predicate registers can be transferred to and from the general registers either individually or as register quadruples ([Section 2.2.5](#)).

# 7 Software Stack

---

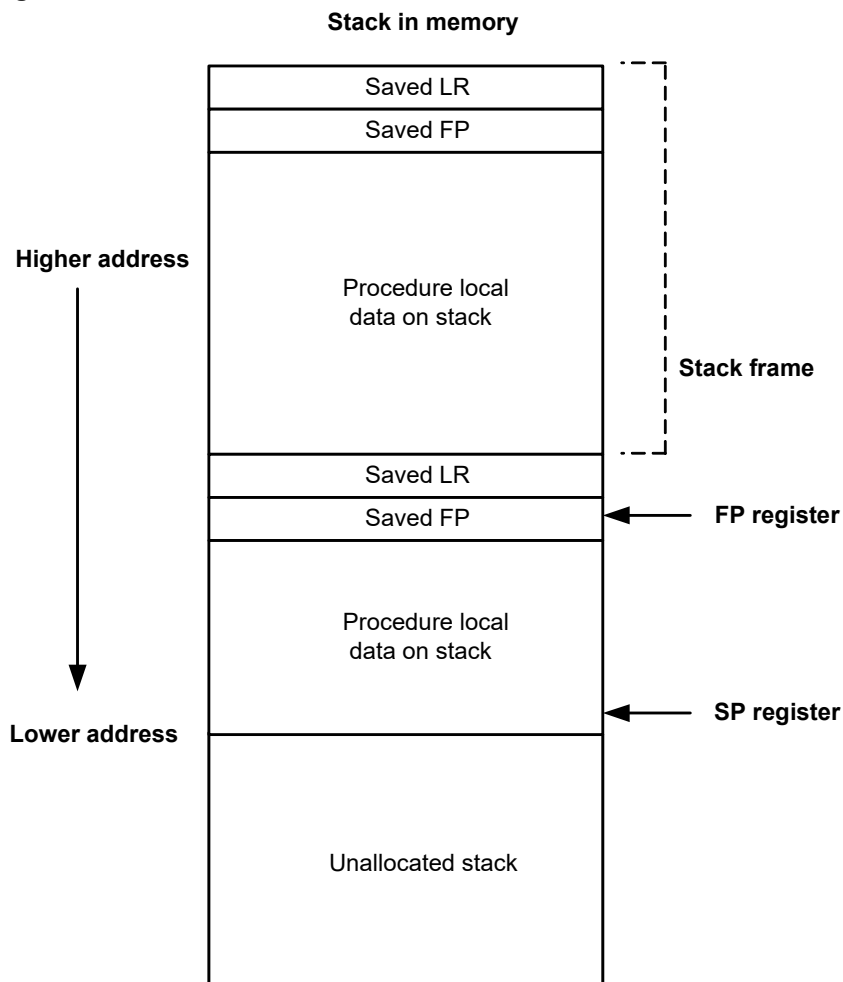
The Hexagon processor includes dedicated registers and instructions to support a call stack for subroutine execution.

The stack structure follows standard C conventions.

## 7.1 Stack structure

The stack is defined to grow from high addresses to low addresses. The stack pointer register SP points to the data element that is currently on the top of the stack.

**Figure 7-1 Stack structure**



**NOTE:** The Hexagon processor supports three dedicated stack instructions: `allocframe`, `deallocframe`, and `dealloc_return` ([Section 7.5](#)).

The SP address must always remain 8-byte aligned for the stack instructions to work properly.

## 7.2 Stack frames

The stack stores stack frames, which are data structures that store state information on the active subroutines in a program (for example, those that were called but have not yet returned). Each stack frame corresponds to an active subroutine in the program.

A stack frame contains the following elements:

- The local variables and data used by the subroutine
- The return address for the subroutine call (pushed from the link register LR)
- The address of the previous stack frame allocated on the stack (pushed from the frame pointer register FP)

The frame pointer register FP always contains the address of the saved frame pointer in the current stack frame. It facilitates debugging by enabling a debugger to examine the stack in memory and easily determine the call sequence, function parameters, and so on.

**NOTE:** For leaf functions it is often unnecessary to save FP and LR. In this case, FP contains the frame pointer of the calling function, not the current function.

## 7.3 Stack protection

The Hexagon processor supports the following features to protect the integrity of the software stack.

### 7.3.1 Stack bounds checking

Stack bounds checking prevents a stack frame from being allocated past the lower boundary of the software stack.

FRAMELIMIT is a 32-bit control register that stores a memory address that specifies the lower bound of the memory area reserved for the software stack. When the `allocframe` instruction allocates a new stack frame, it compares the new stack pointer value in SP with the stack bound value in FRAMELIMIT. If SP is less than FRAMELIMIT, the Hexagon processor raises exception 0x27 ([Section 8.10](#)).

**NOTE:** Stack bounds checking is performed when the processor is in User and Guest modes, but not in Monitor mode.

## 7.3.2 Stack smashing protection

Stack smashing is a technique malicious code uses to gain control over an executing program. Malicious code causes buffer overflows to occur in a procedure's local data, with the goal of modifying the subroutine return address stored in a stack frame so it points to the malicious code instead of the intended return code.

Stack smashing protection prevents this from occurring by scrambling the subroutine return address when a new stack frame is allocated, and then unscrambling the return address when the frame is deallocated. Because the value in FRAMEKEY changes regularly and varies from device to device, it becomes difficult to precalculate a malicious return address.

FRAMEKEY is a 32-bit control register which scrambles return addresses stored on the stack:

- In the allocframe instruction, the 32-bit return address in link register LR is XOR-scrambled with the value in FRAMEKEY before it is stored in the new stack frame.
- In deallocframe and dealloc\_return, the return address loaded from the stack frame is unscrambled with the value in FRAMEKEY before it is stored in LR.

After a processor reset, the default value of FRAMEKEY is 0. If this value is not changed, stack smashing protection is effectively disabled.

**NOTE:** Each hardware thread has its own instance of the FRAMEKEY register.

## 7.4 Stack registers

**Table 7-1 Stack registers**

Register	Name	Description	Alias
SP	Stack pointer	Points to topmost stack element in memory	R29
FP	Frame pointer	Points to previous stack frame on stack	R30
LR	Link register	Contains return address of subroutine call	R31
FRAMELIMIT	Frame limit register	Contains lowest address of stack area	C16
FRAMEKEY	Frame key register	Contains scrambling key for return addresses	C17

**NOTE:** SP, FP, and LR are aliases of three [General registers](#). These general registers are conventionally dedicated for use as stack registers.



## 7.5 Stack instructions

The Hexagon processor includes the `allocframe` and `deallocframe` instructions to efficiently allocate and deallocate stack frames on the call stack.

**Table 7-2 Stack instructions**

Syntax	Operation
<code>allocframe(#u11:3)</code>	<p>Allocate stack frame.</p> <p>This instruction is used after a call. It first XORs the values in LR and FRAMEKEY, and pushes the resulting scrambled return address and FP to the top of stack.</p> <p>Next, it subtracts an unsigned immediate from SP to allocate room for local variables. If the resulting SP is less than FRAMELIMIT, the processor raises exception 0x27. Otherwise, SP is set to the new value, and FP is set to the address of the old frame pointer on the stack.</p> <p>The immediate operand as expressed in assembly syntax specifies the byte offset. This value must be 8-byte aligned. The valid range is from 0 to 16 KB.</p>
<code>deallocframe</code>	<p>Deallocate stack frame.</p> <p>Use this instruction before a return to free a stack frame. It first loads the saved FP and LR values from the address at FP, and XORs the restored LR with the value in FRAMEKEY to unscramble the return address. SP is then pointed back to the previous frame.</p>
<code>dealloc_return</code>	<p>Subroutine return with stack frame deallocate.</p> <p>Perform <code>deallocframe</code> operation, and then perform subroutine return (<a href="#">Section 8.3.3</a>) to the target address loaded from LR by <code>deallocframe</code>.</p>

**NOTE:** The `allocframe` and `deallocframe` instructions load and store the LR and FP registers on the stack as a single aligned 64-bit register pair (LR:FP).

# 8 Program Flow

---

The Hexagon processor supports the following program flow facilities.

## 8.1 Conditional instructions

Many Hexagon processor instructions can conditionally execute. For example:

```
if (P0) R0 = memw(R2)    // Conditionally load word if P0
if (!P1) jump label     // Conditionally jump if not P1
```

The following instructions can be specified as conditional:

- Jumps and calls
- Many load and store instructions
- Logical instructions (including AND/OR/XOR)
- Shift halfword
- 32-bit add/subtract by register or short immediate
- Sign and zero extend
- 32-bit register transfer and 64-bit combine word
- Register transfer immediate
- Deallocate frame and return

For more information, see [Section 5.9](#) and [Chapter 6](#).

## 8.2 Hardware loops

The Hexagon processor includes hardware loop instructions that perform loop branches with zero overhead. For example:

```
loop0(start,#3)          // loop 3 times
start:
  { R0 = mpyi(R0,R0) } :endloop0
```

Two sets of hardware loop instructions are provided – loop0 and loop1 – to nest hardware loops one level deep. For example:

```
// Sum the rows of a 100x200 matrix.

loop1(outer_start,#100)
outer_start:
```

```
R0 = #0
loop0(inner_start, #200)
inner_start:
    R3 = memw(R1++#4)
    { R0 = add(R0, R3) }:endloop0
    { memw(R2++#4) = R0 }:endloop1
```

Use the hardware loop instructions as follows:

- For non-nested loops, loop0 is used.
- For nested loops, loop0 is used for the inner loop, and loop1 for the outer loop.

**NOTE:** If a program must create loops nested more than one level deep, the two innermost loops can be implemented as hardware loops, with the remaining outer loops implemented as software branches.

Each hardware loop is associated with a pair of dedicated loop registers:

- The loop start address register SAn is set to the address of the first instruction in the loop (which is typically expressed in assembly language as a label).
- The loop count register LCn is set to a 32-bit unsigned value which specifies the number of loop iterations to perform. When the PC reaches the end of the loop, LCn is examined to determine whether to repeat or exit the loop.

The hardware loop setup instruction sets both of these registers at once – typically there is no need to set them individually. However, because the loop registers completely specify the hardware loop state, saving and restoring the registers (either automatically by a processor interrupt or manually by the programmer) enables a suspended hardware loop to resume normally once its loop registers are reloaded with the saved values.

The Hexagon processor provides two sets of loop registers for the two hardware loops:

- SA0 and LC0 are used by loop0
- SA1 and LC1 are used by loop1

**Table 8-1 Hardware loop instructions**

Syntax	Description
<code>loopN(start, Rs)</code>	Hardware loop with register loop count. Set registers SAn and LCn for hardware loop N: <ul style="list-style-type: none"> <li>■ SAn is assigned the specified start address of the loop.</li> <li>■ LCn is assigned the value of general register Rs.</li> </ul> <b>NOTE:</b> The loop start operand is encoded as a PC-relative immediate value.
<code>loopN(start, #count)</code>	Hardware loop with immediate loop count. Set registers SAn and LCn for hardware loop N: <ul style="list-style-type: none"> <li>■ SAn is assigned the specified start address of the loop.</li> <li>■ LCn is assigned the specified immediate value (0 to 1023).</li> </ul> <b>NOTE:</b> The loop start operand is encoded as a PC-relative immediate value.
<code>:endloopN</code>	Hardware loop end instruction. Performs the following operation: <pre>if (LCn &gt; 1) {PC = SAn; LCn = LCn-1}</pre> <b>NOTE:</b> This instruction appears in assembly as a suffix appended to the last packet in the loop. It is encoded in the last packet.
<code>SAn = Rs</code>	Set loop start address to general register Rs
<code>LCn = Rs</code>	Set loop count to general register Rs

**NOTE:** The loop instructions are assigned to instruction class CR.

## 8.2.1 Loop setup

To set up a hardware loop, the loop registers SAn and LCn must be set to the proper values. This is done in two ways:

- A `loopN` instruction
- Register transfers to SAn and LCn

The `loopN` instruction performs all the work of setting SAn and LCn. For example:

```
loop0(start,#3)           // SA0=&start, LC0=3
start:
  { R0 = mpyi(R0,R0) } :endloop0
```

In this example, the hardware loop (consisting of a single multiply instruction) executes three times. The `loop0` instruction sets register SA0 to the address value of label `start`, and LC0 to 3.

Loop counts are limited to the range 0 to 1023 when they are expressed as immediate values in `loopN`. If the desired loop count exceeds this range, it must be specified as a register value. For example:

Using `loopN`:

```
R1 = #20000;
loop0(start,R1)           // LC0=20000, SA0=&start
```

```
start:
    { R0 = mpyi(R0,R0) } :endloop0
```

#### Using register transfers:

```
R1 = #20000
LC0 = R1           // LC0=20000
R1 = #start
SA0 = R1           // SA0=&start
start:
    { R0 = mpyi(R0,R0) } :endloop0
```

If a loopN instruction is located too far from its loop start address, the PC-relative offset value that specifies the start address can exceed the maximum range of the instruction's start-address operand. If this occurs, either move the loopN instruction closer to the loop start, or specify the loop start address as a 32-bit constant ([Section 10.9](#)).

For example, using 32-bit constants:

```
R1 = #20000;
loop0(##start,R1) // LC0=20000, SA0=&start
...
```

## 8.2.2 Loop end

The loop end instruction indicates the last packet in a hardware loop. It is expressed in assembly language by appending the packet with the symbol “:endloopN”, where N specifies the hardware loop (0 or 1). For example:

```
loop0(start,#3)
start:
    { R0 = mpyi(R0,R0) } :endloop0 // last packet in loop
```

The last instruction in the loop must always be expressed in assembly language as a packet (using curly braces), even if it is the only instruction in the packet.

Nested hardware loops can specify the same instruction as the end of both the inner and outer loops. For example:

```
// Sum the rows of a 100x200 matrix.
// Software pipeline the outer loop.

p0 = cmp.gt(R0,R0) // p0 = false
loop1(outer_start,#100)
outer_start:
    { if (p0) memw(R2++#4) = R0
      p0 = cmp.eq(R0,R0) // p0 = true
      R0 = #0
      loop0(inner_start,#200) }
inner_start:
    R3 = memw(R1++#4)
    { R0 = add(R0,R3) } :endloop0:endloop1
    memw(R2++#4) = R0
```

Though `endloopN` behaves like a regular instruction (by implementing the loop test and branch), it does not execute in any instruction slot, and does not count as an instruction in the packet. Therefore a single instruction packet which is marked as a loop end can perform up to six operations:

- Four regular instructions (the normal limit for an instruction packet)
- The `endloop0` test and branch
- The `endloop1` test and branch

**NOTE:** The `endloopN` instruction is encoded in the instruction packet ([Section 10.6](#)).

### 8.2.3 Loop execution

After a hardware loop is set up, the loop body always executes at least once regardless of the specified loop count (because the loop count is not examined until the last instruction in the loop). Therefore, if a loop must be optionally executed zero times, it must be preceded with an explicit conditional branch. For example:

```

loop0(start,R1)
  P0 = cmp.eq(R1,#0)
  if (P0) jump skip
start:
  { R0 = mpyi(R0,R0) } :endloop0
skip:

```

In this example a hardware loop is set up with the loop count in R1, but if the value in R1 is zero a software branch skips over the loop body.

After the loop end instruction of a hardware loop is executed, the Hexagon processor examines the value in the corresponding loop count register:

- If the value is greater than 1, the processor decrements the loop count register and performs a zero-cycle branch to the loop start address.
- If the value is less than or equal to 1, the processor resumes program execution at the instruction immediately following the loop end instruction.

**NOTE:** Because nested hardware loops can share the same loop end instruction, the processor can examine both loop count registers in a single operation.

### 8.2.4 Pipelined hardware loops

Software pipelined loops are common for VLIW architectures such as the Hexagon processor. They offer increased code performance in loops by overlapping multiple loop iterations.

A software pipeline has three sections:

- A prologue in which the loop is primed
- A kernel (or steady-state) portion
- An epilogue which drains the pipeline

A simple example is shown in [Table 8-2](#).

**Table 8-2 Software pipelined loop**

<pre>int foo(int *A, int *result) {     int i;     for (i=0;i&lt;100;i++) {         result[i]= A[i]*A[i];     } }</pre>
<pre>foo: {     R3 = R1     loop0(.kernel,#98)        // Decrease loop count by 2 }  {     R1 = memw(R0++#4)        // First prologue stage     R1 = memw(R0++#4)        // Second prologue stage     R2 = mpyi(R1,R1) }  .falign .kernel: {     R1 = memw(R0++#4)        // Kernel     R2 = mpyi(R1,R1)     memw(R3++#4) = R2 }:endloop0 {     R2 = mpyi(R1,R1)        // First epilogue stage     memw(R3++#4) = R2 }      memw(R3++#4) = R2        // Second epilogue stage     jumpr lr</pre>

In [Table 8-2](#), the kernel section of the pipelined loop performs three iterations of the loop in parallel:

- The load for iteration N+2
- The multiply for iteration N+1
- The store for iteration N

One drawback to software pipelining is the extra code necessary for the prologue and epilogue sections of a pipelined loop.

To address this issue, the Hexagon processor provides the `spNloop0` instruction, where the “N” in the instruction name indicates a digit in the range 1 to 3. For example:

```
P3 = sp2loop0(start,#10)    // Set up pipelined loop
```

The `spNloop0` instruction is a variant of the `loop0` instruction: it sets up a normal hardware loop using SA0 and LCO, but also performs the following additional operations:

- When the `spNloop0` instruction executes, it assigns the truth value false to the predicate register P3.

- After the associated loop executes N times, P3 is automatically set to true.

This feature (known as automatic predicate control) enables the store instructions in the kernel section of a pipelined loop to conditionally execute by P3 and thus – because of the way `spNloop0` controls P3 – not execute during the pipeline warm-up. This can reduce the code size of many software pipelined loops by eliminating the need for prologue code.

The `spNloop0` instruction cannot be used to eliminate the epilogue code from a pipelined loop; however, in some cases it is possible to do this through the use of programming techniques.

Typically, the issue affecting the removal of epilogue code is load safety. If the kernel section of a pipelined loop can safely access past the end of its arrays – either because it is known as safe, or because the arrays have been padded at the end – epilogue code is unnecessary. However, if load safety cannot be ensured, explicit epilogue code is required to drain the software pipeline.

[Table 8-3](#) shows how `spNloop0` and load safety simplify the code shown in [Table 8-2](#).

**Table 8-3 Software pipelined loop (using `spNloop0`)**

<pre>int foo(int *A, int *result) {     int i;     for (i=0;i&lt;100;i++) {         result[i]= A[i]*A[i];     } }</pre>
<pre>foo: { // load safety assumed     P3 = sp2loop0(.kernel,#102)    // Set up pipelined loop     R3 = R1 } .falign .kernel: {     R1 = memw(R0++#4)              // Kernel     R2 = mpyi(R1,R1)     if (P3) memw(R3++#4) = R2 }:endloop0      jumpr lr</pre>

**NOTE:** The count value that `spNloop0` uses to control the P3 setting is stored in the user status register `USR.LPCFG`.

## 8.2.5 Loop restrictions

Hardware loops have the following restrictions:

- The loop setup packet in `loopN` or `spNloop0` ([Section 8.2.4](#)) cannot contain a speculative indirect jump, new-value compare jump, or `dealloc_return`.



- The last packet in a hardware loop cannot contain any program flow instructions (including jumps or calls).
- The loop end packet in loop0 cannot contain any instruction that changes SA0 or LC0. Similarly, the loop end packet in loop1 cannot contain any instruction that changes SA1 or LC1.
- The loop end packet in spNloop0 cannot contain any instruction that changes P3.

**NOTE:** SA1 and LC1 can be changed at the end of loop0, while SA0 and LC0 can be changed at the end of loop1.

## 8.3 Software branches

Unlike hardware loops, software branches use an explicit instruction to perform a branch operation. Software branches include jumps, calls, and returns.

The target address for branch instructions is specified as register indirect or PC-relative offsets. PC-relative offsets are normally less than 32 bits, but can be specified as 32 bits by using the appropriate syntax in the target operand ([Section 8.3.4](#)).

Branch instructions are unconditional or conditional, with the execution of conditional instructions controlled by a predicate expression.

**Table 8-4 Software branch instructions**

Syntax	Operation
[if (pred_expr)] jump label [if (pred_expr)] jumpr Rs	Branch to address specified by register Rs or PC-relative offset. Can be conditionally executed.
[if (pred_expr)] call label [if (pred_expr)] callr Rs	Branch to address specified by register Rs or PC-relative offset. Store subroutine return address in link register LR. Can be conditionally executed.
[if (pred_expr)] jumpr LR	Branch to subroutine return address contained in link register LR. Can be conditionally executed.

### 8.3.1 Jumps

Jump instructions change the program flow to a target address, which are specified by either a register or a PC-relative immediate value. Jump instructions can be conditional based on the value of a predicate expression.

**Table 8-5 Jump instructions**

Syntax	Operation
<code>jump label</code>	Direct jump. Branch to address specified by label. Label is encoded as PC-relative signed immediate value.
<code>jumpR Rs</code>	Indirect jump. Branch to address contained in general register Rs.
<code>if ([!]Ps) jump label</code> <code>if ([!]Ps) jumpR Rs</code>	Conditional jump. Perform jump if predicate expression evaluates to true.

**NOTE:** Conditional jumps can be specified as speculative ([Section 8.4](#)).

### 8.3.2 Calls

Call instructions jump to subroutines. The instruction performs a jump to the target address and also stores the return address in the link register LR.

The forms of call are functionally similar to jump instructions and include both PC-relative and register indirect in both unconditional and conditional forms.

**Table 8-6 Call instructions**

Syntax	Operation
<code>call label</code>	Direct subroutine call. Branch to address specified by label, and store return address in register LR. Label is encoded as PC-relative signed immediate value.
<code>callR Rs</code>	Indirect subroutine call. Branch to address contained in general register Rs, and store return address in register LR.
<code>if ([!]Ps) call label</code> <code>if ([!]Ps) callR Rs</code>	Conditional call. If predicate expression evaluates to true, perform subroutine call to specified target address.

### 8.3.3 Returns

Return instructions return from a subroutine. The instruction performs an indirect jump to the subroutine return address stored in link register LR.

Returns are implemented as jump register indirect instructions, and support both unconditional and conditional forms.

**Table 8-7 Return instructions**

Syntax	Operation
<code>jump<sub>r</sub> LR</code>	Subroutine return. Branch to subroutine return address contained in link register LR.
<code>if ([!]Ps) jump<sub>r</sub> LR</code>	Conditional subroutine return. If predicate expression evaluates to true, perform subroutine return to specified target address.
<code>dealloc<sub>r</sub> return</code>	Subroutine return with stack frame deallocate. Perform <code>dealloc<sub>r</sub>frame</code> operation ( <a href="#">Section 7.5</a> ) and then perform subroutine return to the target address loaded by <code>dealloc<sub>r</sub>frame</code> from the link register.
<code>if ([!]Ps) dealloc<sub>r</sub> return</code>	Conditional subroutine return with stack frame deallocate. If predicate expression evaluates to true, perform <code>dealloc<sub>r</sub>frame</code> and subroutine return to the target address loaded by <code>dealloc<sub>r</sub>frame</code> from the link register.

**NOTE:** The link register LR is an alias of general register R31. Therefore subroutine returns can be performed with the instruction `jumpr R31`.

The conditional subroutine returns (including `deallocr return`) can be specified as speculative ([Section 8.4](#)).

### 8.3.4 Extended branches

When a jump or call instruction specifies a PC-relative offset as the branch target, the offset value is normally encoded in significantly less than 32 bits. This can limit the ability for programs to specify “long” branches, which span a large range of the processor’s memory address space.

To support long branches, the jump and call instructions have special versions that encode a full 32-bit value as the PC-relative offset.

**NOTE:** Such instructions use an extra word to store the 32-bit offset ([Section 10.9](#)).

The size of a PC-relative branch offset is expressed in assembly language by optionally prefixing the target label with the symbol “##” or “#”:

- “##” specifies that the assembler *must* use a 32-bit offset
- “#” specifies that the assembler must *not* use a 32-bit offset
- No “#” specifies that the assembler use a 32-bit offset only if necessary

For example:

```
jump ##label    // 32-bit offset
call #label     // Non 32-bit offset
jump label      // Offset size determined by assembler
```

### 8.3.5 Branches to and from packets

Instruction packets are atomic: even if they contain multiple instructions, they are referenced only by the address of the first instruction in the packet. Therefore, branches to a packet can target only the packet’s first instruction.

Packets can contain up to two branches ([Section 8.7](#)). The branch destination can target the current packet or the beginning of another packet.

A branch does not interrupt the execution of the current packet: all of the instructions in the packet execute, even if they appear in the assembly source after the branch instruction.

If a packet is at the end of a hardware loop, it cannot contain a branch instruction.

## 8.4 Speculative jumps

Conditional instructions normally depend on predicates that are generated in a previous instruction packet. However, [Dot-new predicates](#) enable conditional instructions to use a predicate generated in the same packet that contains the conditional instruction.

When dot-new predicates are used with a conditional jump, the resulting instruction is called a speculative jump. For example:

```
{
  P0 = cmp.eq(R9,#16)           // single-packet compare-and-jump
  IF (P0.new) jumpr:t R11       // ... enabled by use of P0.new
}
```

Speculative jumps require the programmer to specify a direction hint in the jump instruction, indicating whether the conditional jump is expected.

The hint initializes the dynamic branch predictor of the Hexagon processor. Whenever the predictor is wrong, the speculative jump instruction takes two cycles to execute instead of one (due to a pipeline stall).

Hints can improve program performance by indicating how speculative jumps are expected to execute over the course of a program: the more often the specified hint indicates how the instruction actually executes, the better the performance.

Hints are expressed in assembly language by appending the suffix “:t” or “:nt” to the jump instruction symbol. For example:

- `jump:t` – The jump instruction is most often taken
- `jump:nt` – The jump instruction is most often not taken

In addition to dot-new predicates, speculative jumps also accept conditional arithmetic expressions (`=0`, `!=0`, `>=0`, `<=0`) involving the general register `Rs`.

**Table 8-8 Speculative jump instructions**

Syntax	Operation
<code>if ([!]Ps.new) jump:t label</code> <code>if ([!]Ps.new) jump:nt label</code>	Speculative direct jump. If predicate expression evaluates to true, jump to address specified by label.
<code>if ([!]Ps.new) jumpr:t Rs</code> <code>if ([!]Ps.new) jumpr:nt Rs</code>	Speculative indirect jump. If predicate expression evaluates to true, jump to address in register <code>Rs</code> .
<code>if (Rs == #0) jump:t label</code> <code>if (Rs == #0) jump:nt label</code>	Speculative direct jump. If predicate <code>Rs = 0</code> is true, jump to address specified by label.
<code>if (Rs != #0) jump:t label</code> <code>if (Rs != #0) jump:nt label</code>	Speculative direct jump. If predicate <code>Rs != 0</code> is true, jump to address specified by label.
<code>if (Rs &gt;= #0) jump:t label</code> <code>if (Rs &gt;= #0) jump:nt label</code>	Speculative direct jump. If predicate <code>Rs &gt;= 0</code> is true, jump to address specified by label.
<code>if (Rs &lt;= #0) jump:t label</code> <code>if (Rs &lt;= #0) jump:nt label</code>	Speculative direct jump. If predicate <code>Rs &lt;= 0</code> is true, jump to address specified by label.

**NOTE:** The hints `:t` and `:nt` interact with the predicate value to determine the instruction cycle count.

Speculative indirect jumps are not supported with register `Rs` predicates.

## 8.5 Compare jumps

To reduce code size, the Hexagon processor supports a compound instruction that combines a compare with a speculative jump in a single 32-bit instruction.

For example:

```
{
    p0 = cmp.eq (R2,R5)           // Single-instr compare-and-jump
    if (p0.new) jump:nt target   // Enabled by compound instr
}
```

The register operands used in a compare jump are limited to R0 through R7 or R16 through R23 (Table 10-3).

The compare and jump instructions that are used in a compare jump are limited to the instructions listed in Table 8-9. The compare can use predicate P0 or P1, while the jump must specify the same predicate that is set in the compare.

A compare jump instruction is expressed in assembly source as two independent compare and jump instructions in a packet. The assembler translates the two instructions into a single compound instruction.

**Table 8-9 Compare jump instructions**

Compare Instruction	Jump Instruction
<code>Pd = cmp.eq (Rs, Rt)</code>	<code>IF (Pd.new) jump:t label</code>
<code>Pd = cmp.gt (Rs, Rt)</code>	<code>IF (Pd.new) jump:nt label</code>
<code>Pd = cmp.gtu (Rs, Rt)</code>	<code>IF (!Pd.new) jump:t label</code>
<code>Pd = cmp.eq (Rs, #U5)</code>	<code>IF (!Pd.new) jump:nt label</code>
<code>Pd = cmp.gt (Rs, #U5)</code>	
<code>Pd = cmp.gtu (Rs, #U5)</code>	
<code>Pd = cmp.eq (Rs, #-1)</code>	
<code>Pd = cmp.gt (Rs, #-1)</code>	
<code>Pd = tstbit (Rs, #0)</code>	

### 8.5.1 New-value compare jumps

A compare jump instruction can access a register that is assigned a new value in the same instruction packet (Section 3.3). This feature is expressed in assembly language by the following changes:

- Appending the suffix `.new` to the new-value register in the compare
- Rewriting the compare jump so its constituent compare and jump operations appear as a single conditional instruction

For example:

```
// load-compare-and-jump packet enabled by new-value compare jump
```

```

{
  R0 = memw(R2+#8)
  if (cmp.eq(R0.new,#0)) jump:nt target
}

```

New-value compare jump instructions have the following restrictions:

- They are limited to the instruction forms listed in [Table 8-10](#).
- They cannot be combined with another jump instruction in the same packet.
- If an instruction produces a 64-bit result or performs a floating-point operation, its result registers cannot be used as the new-value register.
- If an instruction uses auto-increment or absolute-set addressing mode ([Section 5.8](#)), its address register cannot be used as the new-value register.
- If the instruction that sets a new-value register is conditional ([Section 6.1.2](#)), it must always execute.

If the specified jump direction hint is wrong ([Section 8.4](#)), a new-value compare jump takes three cycles to execute instead of one. While this penalty is one cycle longer than in a regular speculative jump, the overall performance is still better than using a regular speculative jump (which must execute an extra packet in all cases).

**NOTE:** New-value compare jump instructions are assigned to instruction class NV, which can execute only in Slot 0. The instruction that assigns the new value must execute in Slot 1, 2, or 3.

**Table 8-10 New-value compare jump instructions**

<pre> if ([!]cmp.eq (Rs.new, Rt) jump:[hint] label if ([!]cmp.gt (Rs.new, Rt) jump:[hint] label if ([!]cmp.gtu (Rs.new, Rt) jump:[hint] label  if ([!]cmp.gt (Rs, Rt.new) jump:[hint] label if ([!]cmp.gtu (Rs, Rt.new) jump:[hint] label </pre>
<pre> if ([!]cmp.eq (Rs.new, #u5) jump:[hint] label if ([!]cmp.gt (Rs.new, #u5) jump:[hint] label if ([!]cmp.gtu (Rs.new, #u5) jump:[hint] label  if ([!]cmp.eq (Rs.new, #-1) jump:[hint] label if ([!]cmp.gt (Rs.new, #-1) jump:[hint] label  if ([!]tstbit (Rs.new, #0) jump:[hint] label </pre>

## 8.6 Register transfer jumps

To reduce code size, the Hexagon processor supports a compound instruction that combines a register transfer with an unconditional jump in a single 32-bit instruction.

For example:

```
{
  jump target    // Jump to label "target"
  R1 = R2        // Assign contents of reg R2 to R1
}
```

The source and target register operands in the register transfer are limited to R0 through R7 or R16 through R23 (Table 10-3).

The target address in the jump is a scaled 9-bit PC-relative address value (as opposed to the 22-bit value in the regular unconditional jump instruction).

A register transfer jump instruction is expressed in assembly source as two independent instructions in a packet. The assembler translates the instructions into a single compound instruction.

**Table 8-11 Register transfer jump instructions**

Syntax	Operation
<pre>jump label; Rd=Rs</pre>	<p>Register transfer jump.</p> <p>Perform register transfer and branch to address specified by label. Label is encoded as PC-relative 9-bit signed immediate value.</p>
<pre>jump label; Rd=#u6</pre>	<p>Register transfer immediate jump.</p> <p>Perform register transfer (of 6-bit unsigned immediate value) and branch to address specified by label. Label is encoded as PC-relative 9-bit signed immediate value.</p>



## 8.7 Dual jumps

Two software branch instructions (referred to here as “jumps”) can appear in the same instruction packet, under the conditions listed in [Table 8-12](#).

The first jump is defined as the jump instruction at the lower address, and the second jump as the jump instruction at the higher address.

Unlike most packetized operations, dual jumps are not executed in parallel ([Section 3.3.1](#)). Instead, the two jumps are processed in a well-defined order in a packet:

1. The predicate in the first jump is evaluated.
2. If the first jump is taken, the second jump is ignored.
3. If the first jump is not taken, the second jump is performed.

**Table 8-12 Dual jump instructions**

Instruction	Description	First jump in packet?	Second jump in packet?
jump	Direct jump	No	Yes
if ([!]Ps[.new]) jump	Conditional jump	Yes	Yes
call if ([!]Ps) call	Direct calls	No	Yes
Pd=cmp.xx ; if ([!]Pd.new) jump	Compare jump	Yes	Yes
if ([!]cmp.xx(Rs.new, Rt)) jump	New-value compare jump	No	No
jumpr if ([!]Ps[.new]) jumpr callr if ([!]Ps) callr dealloc_return if ([!]Ps[.new]) dealloc_return	Indirect jumps Indirect calls dealloc_return	No	No
endloopN	Hardware loop end	No	No

**NOTE:** If a call is ignored in a dual jump, the link register LR is not changed.

## 8.8 Hint indirect jump target

Because it obtains the jump target address from a register, the jumpr instruction ([Section 8.3.1](#)) normally causes the processor to stall for one cycle.

To avoid the stall penalty caused by a jumpr instruction, the Hexagon processor supports the hintjr jump hint instruction, which can be specified before the jumpr instruction.

The `hintjr` instruction indicates that the program is about to execute a `jumpr` to the address contained in the specified register.

**Table 8-13 Speculative jump hint instruction**

Syntax	Operation
<code>hintjr(Rs)</code>	Informs the processor that the <code>jumpr(Rs)</code> instruction is about to be performed.

**NOTE:** To prevent a stall, the `hintjr` instruction must execute at least 2 packets before the corresponding `jumpr` instruction.

The `hintjr` instruction is not needed for `jumpr` instructions used as returns ([Section 8.3.3](#)), because in this case the Hexagon processor automatically predicts the jump targets based on the most recent nested call instructions.

## 8.9 Pauses

Pauses suspend the execution of a program for a period of time, and put it into low-power mode. The program remains suspended for the duration specified in the instruction.

The pause instruction accepts an unsigned 8-bit immediate operand which specifies the pause duration in terms of cycles. The maximum possible duration is 263 cycles (255+8).

Hexagon processor interrupts cause a program to exit the paused state before its specified duration has elapsed.

The pause instruction is useful for implementing user-level low-power synchronization operations (such as spin locks).

**Table 8-14 Pause instruction**

Syntax	Operation
<code>pause (#u8)</code>	Suspend program in low-power mode for specified cycle duration.

## 8.10 Exceptions

Exceptions are internally-generated disruptions to the program flow.

The Hexagon processor OS handles fatal exceptions by terminating the execution of the application system. The user is responsible for fixing the problem and recompiling their applications.

The error messages generated by exceptions include the following information to assist in locating the problem:

- Cause code – Hexadecimal value indicating the type of exception
- User IP – PC value indicating the instruction executed when the exception occurred
- Bad VA – Virtual address indicating the data accessed when the exception occurred

**NOTE:** The cause code, user IP, and Bad VA values are stored in the Hexagon processor system control registers SSR[7:0], ELR, and BADVA respectively.

If multiple exceptions occur simultaneously, the exception with the lowest error code value has the highest exception priority.

If a packet contains multiple loads, or a load and a store, and both operations have an exception of any type, all slot 1 exceptions process before any slot 0 exception is processed.

**Table 8-15 V73 exceptions**

Cause code	Event type	Event description	Notes
0x0	Reset	Software thread reset.	Non-maskable, highest priority
0x01	Precise, unrecoverable	Unrecoverable BIU error (bus error, timeout, L2 parity error, and so on).	Non-maskable
0x03	Precise, unrecoverable	Double exception (exception occurs while SSR[EX]=1).	Non-maskable
0x11	Precise	Privilege violation: User/Guest mode execute to page with no execute permissions (X=0).	Non-maskable
0x12	Precise	Privilege violation: User mode execute to a page with no user permissions (X=1, U=0).	Non-maskable
0x15	Precise	Invalid packet.	Non-maskable
0x16	Precise	Illegal execution of coprocessor instruction.	Non-maskable
0x17	Precise	Instruction cache error.	Non-maskable
0x1A	Precise	Privilege violation: Executing a guest mode instruction in user mode.	Non-maskable
0x1B	Precise	Privilege violation: Executing a supervisor instruction in User/Guest mode.	Non-maskable
0x1D	Precise, unrecoverable	Packet with multiple writes to the same destination register.	Non-maskable
0x1E	Precise, unrecoverable	Program counter values that are not properly aligned.	Non-maskable
0x20	Precise	Load to misaligned address.	Non-maskable
0x21	Precise	Store to misaligned address.	Non-maskable
0x22	Precise	Privilege violation: User/Guest mode read to page with no read permission (R=0).	Non-maskable
0x23	Precise	Privilege violation: User/Guest mode write to page with no write permissions (W=0).	Non-maskable
0x24	Precise	Privilege violation: User mode read to page with no user permission (R=1, U=0).	Non-maskable
0x25	Precise	Privilege violation: User mode write to page with no user permissions (W=1, U=0).	Non-maskable
0x26	Precise	Coprocessor VMEM address error.	Non-maskable
0x27	Precise	Stack overflow: Allocframe instruction exceeded FRAMELIMIT.	Non-maskable,

**Table 8-15 V73 exceptions (cont.)**

Cause code	Event type	Event description	Notes
0x42	Imprecise	Data abort.	Non-maskable
0x43	Imprecise	NMI.	Non-maskable
0x44	Imprecise	Multiple TLB match.	Non-maskable
0x45	Imprecise	Livelock exception.	Non-maskable
0x60	TLB miss-X	Missing fetch address on PC-page.	Non-maskable
0x61	TLB miss-X	Missing fetch on second page from packet that spans pages.	Non-maskable
0x62	TLB miss-X	icinva.	Non-maskable
	Reserved		
0x70	TLB miss-RW	Memory read.	Non-maskable
0x71	TLB miss-RW	Memory write.	Non-maskable
	Reserved		
#u8	Trap0	Software Trap0 instruction.	Non-maskable
#u8	Trap1	Software Trap1 instruction.	Non-maskable
	Reserved		
0x80	Debug	Single-step debug exception.	
	Reserved		
0xBF	Floating-point	Execution of floating-point instruction resulted in exception.	Non-maskable
0xC0	Interrupt0	General external interrupt.	Maskable, highest priority general interrupt
0xC1	Interrupt 1	General external interrupt	Maskable
0xC2	Interrupt 2	General external interrupt	VIC0 interface
0xC3	Interrupt 3	General external interrupt	VIC1 interface
0xC4	Interrupt 4	General external interrupt	VIC2 interface
0xC5	Interrupt 5	General external interrupt	VIC3 interface
0xC6	Interrupt 6	General external interrupt	
0xC7	Interrupt 7	General external interrupt	Lowest-priority interrupt

# 9 PMU Events

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The Hexagon processor can collect execution statistics on the applications it executes. The statistics summarize the various types of Hexagon processor events that occurred while the application was running.

Execution statistics are collected in hardware or software:

- Statistics are collected in hardware with the performance monitor unit (PMU), which is defined as part of the Hexagon processor architecture.
- Statistics are collected in software using the Hexagon simulator. The simulator statistics are presented in the same format used by the PMU.

Execution statistics are expressed in terms of processor events. This chapter defines the event symbols, along with their associated numeric codes.

**NOTE:** Because the types of execution events vary across the Hexagon processor versions, different types of statistics are collected for each version. This chapter lists the event symbols defined for version V73.

## 9.1 V73 processor event symbols

Table 9-1 defines the symbols that represent processor events for the V73 Hexagon processor.

**Table 9-1 V73 processor events symbols**

Event	Symbol	Definition
0x1	COUNTER0_OVERFLOW	Event detected by counter1 to build an effective 64-bit counter.
0x2	COUNTER2_OVERFLOW	Event detected by counter3 to build an effective 64-bit counter.
0x3	COMMITTED_PKT_ANY	Number of packets that are committed by any thread. Packets are executed.
0x4	COMMITTED_PKT_BSB	Number of packets that are committed two cycles after an earlier packet in the same thread.
0x5	COUNTER4_OVERFLOW	Event detected by counter5 to build an effective 64-bit counter.
0x6	COUNTER6_OVERFLOW	Event detected by counter7 to build an effective 64-bit counter.
0x7	COMMITTED_PKT_B2B	Number of packets that are committed one cycle after the earlier packet in the same thread.
0x8	COMMITTED_PKT_SMT	Number of packets that are committed on the SMT threads. Includes the second, third, and fourth packets that are committed in one cycle.
0xa	CYCLES_5_THREAD_RUNNING	Processor cycles that exactly five threads are running. Running means the threads are not in the Wait or Stop state.

**Table 9-1 V73 processor events symbols**

Event	Symbol	Definition
0xb	CYCLES_6_THREAD_RUNNING	Processor cycles that exactly six threads are running. Running means the threads are not in the Wait or Stop state.
0xc	COMMITTED_PKT_T0	Number of packets that are committed by thread 0. Packets are executed.
0xd	COMMITTED_PKT_T1	Number of packets that are committed by thread 1. Packets are executed.
0xe	COMMITTED_PKT_T2	Number of packets that are committed by thread 2. Packets are executed.
0xf	COMMITTED_PKT_T3	Number of packets that are committed by thread 3. Packets are executed.
0x10	COMMITTED_PKT_T4	Number of packets that are committed by thread 4. Packets are executed.
0x11	COMMITTED_PKT_T5	Number of packets that are committed by thread 5. Packets are executed.
0x12	ICACHE_DEMAND_MISS	Number of I-cache cacheable demand primary or secondary misses. Includes secondary misses.
0x13	DCACHE_DEMAND_MISS	Number of D-cache cacheable demand primary or secondary misses. Includes dczero stalls. Excludes uncachable, prefetches, and no-allocate store misses.
0x14	DCACHE_STORE_MISS	Number of D-cache cacheable store misses.
0x15	COMMITTED_PKT_T6	Thread 6 committed a packet. Packets are executed.
0x16	COMMITTED_PKT_T7	Thread 7 committed a packet. Packets are executed.
0x17	CU_PKT_READY_NOT_DISPATCHED	Packets were ready at the CU scheduler but were not scheduled because either the scheduler's thread was not picked or there was an inter-cluster resource conflict.
0x18	COMMITTED_PKT_5_THREAD_RUNNING	Number of committed packets with five threads running. Running means the threads are not in Wait or Stop mode.
0x19	COMMITTED_PKT_6_THREAD_RUNNING	Number of committed packets with six threads running. Running means the threads are not in Wait or Stop mode.
0x1a	COMMITTED_PKT_7_THREAD_RUNNING	Number of committed packets with seven threads running. Running means the threads are not in Wait or Stop mode.
0x1b	COMMITTED_PKT_8_THREAD_RUNNING	Number of committed packets with eight threads running. Running means the threads are not in Wait or Stop mode.
0x1c	IU_L1S_ACCESS	Number of IU L1S loads. Includes demands or prefetches.
0x1d	IU_L1S_PREFETCH	Number of IU L1S prefetches.
0x1e	IU_L1S_AXIS_STALL	Number of IU L1S stalls due to an AXI slave.
0x1f	IU_L1S_NO_GRANT	IU request to L1S, and no grant from the vector unit.
0x20	ANY_IU_REPLAY	Any IU stall other than an I-cache miss. Includes a jump register stall, fetchcross stall, ITLB miss stall, and so on. Excludes a CU replay.
0x21	ANY_DU_REPLAY	Any DU replay. Includes a bank conflict, store buffer full, and so on. Excludes a stall due to a cache miss.

**Table 9-1 V73 processor events symbols**

Event	Symbol	Definition
0x22	CYCLES_7_THREAD_RUNNING	Processor cycles that exactly seven threads are running. Running means the threads are not in Wait or Stop mode.
0x23	ISSUED_PACKETS	Speculatively issued packets were delivered from an IU.
0x24	LOOPCACHE_PACKETS	Committed packets were cloned from the packet queue during a pinned hardware loop.
0x25	COMMITTED_PKT_1_THREAD_RUNNING	Number of committed packets with one thread running. Running means the thread is not in Wait or Stop mode.
0x26	COMMITTED_PKT_2_THREAD_RUNNING	Number of committed packets with two threads running. Running means the threads are not in Wait or Stop mode.
0x27	COMMITTED_PKT_3_THREAD_RUNNING	Number of committed packets with three threads running. Running means the threads are not in Wait or Stop mode.
0x28	THREAD_LMH_THROTTLE	For a specific thread, the sustained power exceeds the limits management threshold and limits budget threshold. Results in throttling that is based on thread priority.
0x29	LMH_THROTTLE	Throttling is based on the value of the peak current that is over the current limits of LMH.
0x2a	COMMITTED_INSTS	Number of committed instructions. Increments by up to eight per cycle. Duplex of two instructions counts as two instructions. Does not include end loops.
0x2b	COMMITTED_TC1_INSTS	Number of committed TC1 class instructions. Increments by up to eight per cycle. Duplex of two TC1 instructions counts as two separate TC1 instructions. Does not include NOP instructions.
0x2c	COMMITTED_PRIVATE_INSTS	Number of committed instructions that have per-cluster (private) execution resources. Increments by up to eight per cycle. Duplex of two private instructions counts as two private instructions.
0x2d	GLOBAL_POWERLIMITS_OVER	Sustained global power that exceeds the overall global limits management threshold and limits the budget threshold. Causes the thread-specific LMH to engage.
0x2e	CYCLES_8_THREAD_RUNNING	Processor cycles that exactly eight threads are running. Running means the threads are not in Wait or Stop mode.
0x2f	COMMITTED_PKT_4_THREAD_RUNNING	Number of committed packets with four threads running. Running means the threads are not in Stop or Wait mode.
0x30	COMMITTED_LOADS	Number of committed load instructions. Includes cached and uncached. Increments by two for dual loads. Excludes prefetches, memory operations, and coprocessor loads.
0x31	COMMITTED_STORES	Number of committed store instructions. Includes cached and uncached. Increments by two for dual stores. Excludes memory operations and coprocessor stores.
0x32	COMMITTED_MEMOPS	Number of committed memory operations instructions. Cached or uncached.
0x33	COMMITTED_NOPs	Number of committed NOPs.
0x34	ISSUED_INSTS	Speculatively issued instructions delivered from the IU.

**Table 9-1 V73 processor events symbols**

Event	Symbol	Definition
0x35	DISPATCHED_PACKETS	Number of packets that the CU dispatched. NOP instructions are squashed.
0x36	DISPATCHED_INSTS	Number of instructions that the CU dispatched.
0x37	COMMITTED_PROGRAM_FLOW_INSTS	Number of committed packets that contain a program flow instruction. Includes CR jumps, endloop, J, JR, dealloc_return, system/trap, superset of event 56. Dual jumps count as two jumps.
0x38	COMMITTED_PKT_CHANGED_FLOW	Number of committed packets that resulted in a change of flow. Any taken jump. Includes endloop and dealloc_return.
0x39	COMMITTED_PKT_ENDLOOP	Number of committed packets containing an end loop that was taken.
0x3a	PST_USED_P0P1BUSY	Number of times a store port was used when p0 and p1 were both occupied. Only increments when store port is present.
0x3b	CYCLES_1_THREAD_RUNNING	Processor cycles that exactly one thread is running. Running means the thread is not in Wait or Stop mode.
0x3c	CYCLES_2_THREAD_RUNNING	Processor cycles that exactly two threads are running. Running means the threads are not in Wait or Stop mode.
0x3d	CYCLES_3_THREAD_RUNNING	Processor cycles that exactly three threads are running. Running means the threads are not in Wait or Stop mode.
0x3e	CYCLES_4_THREAD_RUNNING	Processor cycles that exactly four threads are running. Running means the threads are not in Wait or Stop mode.
0x3f	AXI_LINE128_READ_REQUEST	Number of 128-byte line read requests issued by the primary AXI master. Includes all interleaved requests.
0x40	AXI_READ_REQUEST	All read requests issued by the primary AXI master. Includes full lines, partial lines, and all interleaved requests.
0x41	AXI_LINE32_READ_REQUEST	Number of 32-byte line read requests issued by the primary AXI master. Includes all interleaved requests.
0x42	AXI_WRITE_REQUEST	All write requests issued by the primary AXI master. Includes full lines, partial lines, and all interleaved requests.
0x43	AXI_LINE32_WRITE_REQUEST	Number of 32-byte line write requests issued by the primary AXI master. Includes all interleaved requests. All bytes are valid.
0x44	AHB_READ_REQUEST	Number of read requests issued by the AHB master.
0x45	AHB_WRITE_REQUEST	Number of write requests issued by the AHB master.
0x46	AXI_LINE128_WRITE_REQUEST	Number of 128-byte line write requests issued by the primary AXI master. Includes all interleaved requests. All bytes are valid.
0x47	AXI_SLAVE_MULTI_BEAT_ACCESS	Number of AXI slave multi-beat accesses.
0x48	AXI_SLAVE_SINGLE_BEAT_ACCESS	Number of AXI slave single-beat accesses.
0x49	AXI2_READ_REQUEST	All read requests issued by the secondary AXI master. Includes full lines and partial lines.
0x4a	AXI2_LINE32_READ_REQUEST	Number of 32-byte line read requests issued by the secondary AXI master.



**Table 9-1 V73 processor events symbols**

Event	Symbol	Definition
0x4b	AXI2_WRITE_REQUEST	All write requests issued by the secondary AXI master. Includes full lines and partial lines.
0x4c	AXI2_LINE32_WRITE_REQUEST	Number of 32-byte line write requests issued by the secondary AXI master.
0x4d	AXI2_CONGESTION	Secondary AXI command or data queue is full. An operation is stuck at the head of the secondary AXI master command queue.
0x50	COMMITTED_FPS	Number of committed floating point instructions. Increments by two for dual floating-point operations. Excludes conversions.
0x51	REDIRECT_BIMODAL_MISPREDICT	Mispredicted bimodal branch direction caused a control flow redirect.
0x52	REDIRECT_TARGET_MISPREDICT	Mispredicted branch target caused a control flow redirect. Includes an RAS mispredict, and HintJR mispredict. Excludes indirect jumps and calls other than JUMPR R31 returns. Excludes direction mispredicts.
0x53	REDIRECT_LOOP_MISPREDICT	Mispredicted hardware loop end caused a control flow redirect. Can only happen when the loop has few packets and the loop count is 2 or less.
0x54	REDIRECT_MISC	Control flow is redirected for a reason other than events 81, 82, and 83. Includes exceptions, traps, interrupts, non-R31 jumps, multiple initialization loops in flight, and so on.
0x55	AXI_LINE256_WRITE_REQUEST	Number of 256-byte line write requests issued by the AXI master. All bytes are valid.
0x56	NUM_PACKET_CRACKED	Number of packets that cracked.
0x58	JTLB_MISS	Instruction or data address translation request was missed in the JTLB.
0x5a	COMMITTED_PKT_RETURN	Number of committed return instructions. Includes canceled returns.
0x5b	COMMITTED_PKT_INDIRECT_JUMP	Number of committed indirect jumps or call instructions. Includes canceled instructions. Does not include JUMPR R31 returns.
0x5c	COMMITTED_BIMODAL_BRANCH_INSTS	Number of committed bimodal branches. Includes *.old and *.new. Increments by two for dual jumps.
0x5f	VTCM_FIFO_FULL_CYCLES	Cycles cluster can be issued if the VTCM FIFO queue is full.
0x60	DU_STORE_BUFFER_COALESCED	Number of times an incoming store was coalesced into an existing valid store buffer entry. Each valid store buffer entry is dword-aligned.
0x61	DU_L1S_LOAD_ACCESS	Number of scalar load accesses to L1S.
0x62	ICACHE_ACCESS	Number of I-cacheline fetches.
0x63	BTB_HIT	Number of branch target buffer hits.
0x64	BTB_MISS	Number of branch target buffer misses.
0x65	IU_DEMAND_SECONDARY_MISS	Number of I-cache secondary misses.

**Table 9-1 V73 processor events symbols**

Event	Symbol	Definition
0x67	FAST_FETCH_KILLED	Number of fast fetches that were killed (after an l-cache access).
0x69	FETCHED_PACKETS_DROPPED	Number of packets that are dropped because the IU cannot deliver them to the CU.
0x6b	IU_PREFETCHES_SENT_TO_L2	Number of IU prefetches sent to the L2 cache. Includes cache-lines not dropped by the L2 cache. Excludes replayed prefetches and only counts prefetches the L2 accepts. Excludes IU prefetches that are sent to L2 ITCM.
0x6c	ITLB_MISS	Number of ITLB misses that go to JTLB.
0x72	FETCH_2_CYCLE	Number of two-cycle fetches in an IU (returns, loop end, fall through, BTB).
0x73	FETCH_3_CYCLE	Number of three-cycle fetches in an IU.
0x75	L2_IU_SECONDARY_MISS	Number of L2 secondary misses from an IU.
0x76	L2_IU_ACCESS	Number of L2 cacheable access from an IU. Includes any access to the L2 cache that was the result of an IU command, either demand or L1 prefetch access. Excludes any prefetches generated in the L2 cache. Excludes L2fetch, TCM accesses, and uncacheables. Address must target the primary AXI master.
0x77	L2_IU_MISS	Number of L2 misses from an IU. Of the events qualified by 0x76, the event that resulted in an L2 miss (demand miss or L1 prefetch miss). An L2 miss is any condition that prevents the immediate return of data to the IU, excluding pipeline conflicts.
0x78	L2_IU_PREFETCH_ACCESS	Number of prefetches from an IU to the L2 cache. Any IU prefetch access sent to the L2 cache. Access must be L2 cacheable and target the primary AXI. Does not include L2 fetch-generated accesses.
0x79	L2_IU_PREFETCH_MISS	Number of L2 misses that were IU prefetches. Of the events qualified by 0x78, the events that resulted in an L2 miss.
0x7c	L2_DU_READ_ACCESS	Number of L2 cacheable read accesses from a DU. Any read access from the DU that might cause a lookup in the L2 cache. Includes loads, L1 prefetches, dcfetches. Excludes the initial L2fetch command, uncacheables, TCM accesses, and coprocessor loads. Must target the primary AXI master.
0x7d	L2_DU_READ_MISS	Number of L2 read misses from a DU. Of the events qualified by 0x7C, any event that resulted in an L2 miss (that is, the line was not previously allocated in the L2 cache and will be fetched from the backing memory).
0x7e	L2FETCH_ACCESS	Number of L2 fetch accesses from a DU. Any access to the L2 cache from the L2 prefetch engine that was initiated by programming the L2Fetch engine.
0x7f	L2FETCH_MISS	Number of L2 fetch misses from a programmed inquiry. Of the events qualified by 0x7E, the event that resulted in an L2 miss (that is, the line was not previously allocated in the L2 cache and will be fetched from the backing memory).

**Table 9-1 V73 processor events symbols**

Event	Symbol	Definition
0x81	L2_ACCESS	All requests to the L2 cache. Does not include internally generated accesses like L2 fetch, however the programming of the L2Fetch engine is counted. All accesses to odd interleave or even interleave are counted. Can be L2 cacheable or TCM.
0x82	L2_PIPE_CONFLICT_STALL	Request is not taken by the L2 cache due to a pipe conflict. The conflict can be a tag bank, data bank, or other pipeline conflict.
0x83	L2_TAG_ARRAY_CONFLICT	Of the items in event 130, the items caused by a conflict with the tag array.
0x87	TCM_DU_ACCESS	Number of TCM accesses from a DU. DU access to the L2 TCM space. Excludes HVX requests.
0x88	TCM_DU_READ_ACCESS	Number of TCM read accesses from a DU. DU read access to the L2 TCM space. Includes HVX requests.
0x89	TCM_IU_ACCESS	Number of TCM accesses from an IU. IU access to the L2 TCM space.
0x8a	L2_CASTOUT	L2 cache evicts a dirty line due to an allocation. This event is not triggered on cache operations.
0x8b	L2_DU_STORE_ACCESS	Number of L2 cacheable store access from a DU. Any store access from the DU that might cause a lookup in the L2 cache. Excludes cache operations, uncacheables, TCM, and coprocessor stores. Must target the primary AXI master.
0x8c	L2_DU_STORE_MISS	Number of L2 misses from a DU. Of the events qualified by 0x8B, the events that resulted in a miss. Specifically, the cases where the line is not in the cache or a coalesce buffer.
0x8d	L2_DU_PREFETCH_ACCESS	Number of L2 prefetch accesses from a DU. Of the events qualified by 0x7C, the events that are dcfetch and dhwprefetch. These L2 cacheable events target the primary AXI master.
0x8e	L2_DU_PREFETCH_MISS	Number of L2 prefetch misses from a DU. Of the events qualified by 0x8D, the events that missed the L2 cache.
0x90	L2_DU_LOAD_SECONDARY_MISS	Number of L2 load secondary misses from a DU. Hit a busy line in the scoreboard, which prevented a return. A busy condition can include pipeline bubbles caused by back-to-back loads, like L1 UC loads.
0x91	L2FETCH_COMMAND	Number of L2fetch commands. Excludes L2 fetch stop commands.
0x92	L2FETCH_COMMAND_KILLED	L2 fetch command was killed because a stop command was issued. Increments once for each L2 fetch command that is killed. If multiple commands are queued to the L2Fetch engine, the kill of each command will be recorded.
0x93	L2FETCH_COMMAND_OVERWRITE	L2 fetch command was overwritten. Kills an old L2 fetch command and replaces it with a new command.
0x94	L2FETCH_ACCESS_CREDIT_FAIL	L2 fetch access cannot get a credit. L2 fetch is blocked due to a missing L2 fetch or L2 evict credit.
0x95	AXI_SLAVE_READ_BUSY	AXI slave read access hit a busy line.
0x96	AXI_SLAVE_WRITE_BUSY	AXI slave write access hit a busy line.

**Table 9-1 V73 processor events symbols**

Event	Symbol	Definition
0x97	L2_ACCESS_EVEN	Of the events in 0x81, number of accesses made to the even L2 cache.
0x98	CLADE_HIGH_PRIO_L2_ACCESS	Number of IU or DU requests for a high-priority CLADE region. Not counted for an L2 fetch.
0x99	CLADE_LOW_PRIO_L2_ACCESS	Number of IU or DU requests for a low-priority CLADE region. Not counted for an L2 fetch.
0x9a	CLADE_HIGH_PRIO_L2_MISS	Number of CLADE high-priority L2 accesses that missed in the L2 cache.
0x9b	CLADE_LOW_PRIO_L2_MISS	Number of CLADE low-priority L2 accesses that missed in the L2 cache.
0x9c	CLADE_HIGH_PRIO_EXCEPTION	CLADE high-priority decode that had an exception.
0x9d	CLADE_LOW_PRIO_EXCEPTION	CLADE low-priority decode that had an exception.
0x9e	AXI2_SLAVE_READ_BUSY	AXI secondary slave read access hit a busy line.
0x9f	AXI2_SLAVE_WRITE_BUSY	AXI secondary slave write access hit a busy line.
0xa0	ANY_DU_STALL	Any DU stall. Increments once when the thread has a DU stall (D-cache miss or DTLB miss).
0xa1	DU_BANK_CONFLICT_REPLAY	DU bank conflict replay. Dual memory access to same bank, but different lines.
0xa2	DU_CREDIT_REPLAY	Number of times a packet took a replay because insufficient QoS DU credits were available.
0xa3	L2_FIFO_FULL_REPLAY	Number of L2 even or odd FIFO full replays.
0xa4	DU_STORE_BUFFER_FULL_REPLAY	Number of DU replays because a demand load access hit in the store buffer.
0xa7	DU_SNOOP_REQUEST	Number of DU snoop requests that were accepted.
0xa8	DU_FILL_REPLAY	Fill has an index conflict with an instruction from the same thread in the pipeline. Fills and demands might be from different threads if there is a prefetch from the deferral queue, or if a fill has not be acknowledged for very long and forces itself into the pipeline.
0xa9	PST_3STORETYPE_SBCONF_REPLAY	Number of times a packet on the lower priority cluster had to replay because one cluster had a dual store and the other cluster had a single store or a memop. Only increments when store port is present.
0xac	DU_READ_TO_L2	Number of DU reads to L2 cache. Total of everything that brings data from the L2 array. Includes prefetches (dcfetch and hwprefetch). Excludes coprocessor loads.
0xad	DU_WRITE_TO_L2	Number of DU writes to L2 cache. Total of everything that is written out of the DU to the L2 array. Includes dczeroa. Excludes dcclean, dccleaninv, tag writes, and coprocessor stores.
0xae	PST_3LDST_L2FIFOCONF_REPLAY	Number of times a packet on lower priority cluster had to replay because one cluster had a dual uncacheable loads and the other cluster had a single store or the one cluster had a dual store and other cluster had a load miss. Only increments when store port is present.

**Table 9-1 V73 processor events symbols**

Event	Symbol	Definition
0xaf	DCZERO_COMMITTED	Dczeroa instruction was committed.
0xb0	L2ITCM_IU_READ	Number of ITCM accesses from an IU. Includes IU demand fetches and IU prefetches. This event is not included in any other L2 events.
0xb1	L2ITCM_DU_READ	Number of L2 ITCM read accesses from a DU. Includes all demands and prefetches. This event is not included in any other L2 events.
0xb2	L2ITCM_DU_WRITE	Number of L2 ITCM write accesses from a DU. Includes stores and dczeroa events. Does not include any cache operations. This event is not included in any other L2 events.
0xb3	DTLB_MISS	DTLB miss that goes to JTLB. When both slots miss to different pages, increments by two. When both slots miss to the same page, only counts S1 because S1 goes first and fills for S0.
0xb4	L2ITCM_BIMODAL_WRITES_SUCCESS	Number of successful bimodal writes into L2 ITCM.
0xb6	STORE_BUFFER_HIT_REPLAY	Store buffer hit is replayed because a packet with two stores is going to the same bank but different cachelines, followed by a load from an address that was pushed into the store buffer.
0xb7	STORE_BUFFER_FORCE_REPLAY	Store buffer must drain, forcing the current packet to replay. Typically occurs on a cache index match between the current packet and store buffer. Can also a store buffer timeout.
0xb8	TAG_WRITE_CONFLICT_REPLAY	Number of inter-cluster tag write conflicts.
0xb9	SMT_BANK_CONFLICT	Number of inter-thread SMT bank conflicts.
0xba	PORT_CONFLICT_REPLAY	Number of all port conflict replays, including the same cluster replays caused by high-priority fills and store buffer force drains. Includes inter-cluster replays.
0xbb	L2ITCM_BIMODAL_WRITES_DROPPED	Number of bimodal writes into L2 ITCM that were dropped.
0xbc	L2ITCM_IU_PREFETCH_READ	Number of ITCM accesses from IU prefetches. Includes only prefetches from an IU. This event is included in event 176. It is not included in any other L2 events.
0xbd	PAGE_CROSS_REPLAY	Page cross from a valid packet that caused a replay. Excludes pkill packets. Counts twice if both slots cause a page cross.
0xbe	PST_STORE_SENTON_OTHPORT	Number of times a slot 0 store was sent to the other store buffer because the other cluster had a slot 1 store or memop. Only increments when store port is present.
0xbf	DU_DEMAND_SECONDARY_MISS	Number of DU demand secondary misses.
0xc0	DU_MISC_REPLAY	All DU replays not counted by other replay events. This event counts every time ANY_DU_REPLAY counts and no other DU replay event counts.
0xc1	GUARDBUF_SETMATCH_CRACKING_REPLAY	Number of replays taken by a younger access due to an index match with a guard buffer entry. If the younger access hits in D\$, only the guard buffer entry of the other thread is checked. Otherwise, an index match with either thread's guard buffer entry results in this replay.

**Table 9-1 V73 processor events symbols**

Event	Symbol	Definition
0xc2	DU_STATE_REPLAY	Number of times an access replayed because the access one cycle ahead of it was allocating or invalidating a way in the same index.
0xc3	DCFETCH_COMMITTED	Number of dcfetches that were committed. Includes hits and drops. Does not include convert-to-prefetches.
0xc4	DCFETCH_HIT	Number of dcfetch hits in D-cache. Includes hitting valid or reserved lines.
0xc5	DCFETCH_MISS	Number of dcfetches missed in L1 cache. Counts the dcfetches issued to L2 FIFO.
0xc6	DCACHE_EVICTION_IN_PIPE_REPLAY	Number of replays taken by a packet because an eviction in progress is occupying the pipe.
0xc7	STBUF_MATCH_PARTIAL_CRACK_REPLAY	Number of replays taken by a partial crack store due to a dword match with an existing store buffer entry.
0xc8	DU_LOAD_UNCACHEABLE	Load instructions with addresses uncacheable in the L1 cache
0xc9	DU_DUAL_LOAD_UNCACHEABLE	Packets where both loads have addresses uncacheable in the L1 cache
0xca	DU_STORE_UNCACHEABLE	Store instructions with addresses uncacheable in the L1 cache
0xcb	DU_STORE_RELEASE_CREDIT_STALL	Stall occurs because there are not enough credits from the store release.
0xcd	AXI_LINE256_READ_REQUEST	Number of 256-byte line read requests issued by the AXI master. All bytes are valid.
0xce	AXI_LINE64_READ_REQUEST	Number of 64-byte line read requests issued by the primary AXI master. Includes all interleaved requests.
0xcf	AXI_LINE64_WRITE_REQUEST	Number of 64-byte line write requests issued by the primary AXI master. Includes all interleaved requests. All bytes are valid.
0xd1	AHB_8_READ_REQUEST	Number of 8-byte read requests issued by the AHB.
0xd3	L2FETCH_COMMAND_PAGE_TERMINATION	L2fetch command terminated because it cannot get a page translation from VA to PA. Includes terminations due to permission errors. That is, an address translation can fail because the VA to PA is not in the TLB, or the properties in the translation are not acceptable and the command terminates.
0xd5	L2_DU_STORE_COALESCE	Number of events from 139 that were coalesced
0xd6	L2_STORE_LINK	Number of times a new store links to something else in the scoreboard.
0xd7	L2_SCOREBOARD_70_PERCENT_FULL	Increments by one for every cycle where the L2 scoreboard is at least 70% full. For a 32-entry scoreboard, 23 or more entries are consumed. This event continues to count even if the scoreboard is more than 80% full. For more than one interleave, this event considers only the scoreboard that has the most entries consumed.

**Table 9-1 V73 processor events symbols**

Event	Symbol	Definition
0xd8	L2_SCOREBOARD_80_PERCENT_FULL	Increments by one for every cycle where the L2 scoreboard is at least 80% full. For a 32-entry scoreboard, 26 or more entries are consumed. This event continues to count even if the scoreboard is more than 90% full. For more than one interleave, this event considers only the scoreboard that has the most entries consumed.
0xd9	L2_SCOREBOARD_90_PERCENT_FULL	Increments by one for every cycle where the L2 scoreboard is at least 90% full. For a 32-entry scoreboard, 29 or more entries are consumed. For more than one interleave, this event considers only the scoreboard that has the most entries consumed.
0xda	L2_SCOREBOARD_FULL_REJECT	L2 scoreboard is too full to accept a selector request, and the selector has a request.
0xdc	L2_EVICTION_BUFFERS_FULL	Counts every cycle when all eviction buffers in any interleave are occupied.
0xdd	AHB_MULTI_BEAT_READ_REQUEST	Number of 32-byte multi-beat read requests issued by the AHB.
0xdf	L2_DU_LOAD_SECONDARY_MISS_ON_SW_PREFETCH	Of the events in 0x90, the events where the primary miss was a DC fetch or L2 fetch.
0xe0	L2FETCH_DROP	L2 fetch data dropped because a previous eviction has not completed.
0xe5	THREAD_OFF_PVIEW_CYCLES	Cycles cluster cannot commit because a thread is in the Off or Wait state.
0xe6	ARCH_LOCK_PVIEW_CYCLES	Cycles cluster cannot commit due to a kernel lock or TLB lock.
0xe7	REDIRECT_PVIEW_CYCLES	Cycles cluster cannot commit because of redirects such as branch mispredicts.
0xe8	IU_NO_PKT_PVIEW_CYCLES	Cycles cluster cannot commit because the interrupt queue is empty.
0xe9	DU_CACHE_MISS_PVIEW_CYCLES	Cycles cluster cannot commit due to a D-cache cacheable miss.
0xea	DU_BUSY_OTHER_PVIEW_CYCLES	Cycles cluster cannot commit due to a DU replay, DU bubble, or DTLB miss.
0xeb	CU_BUSY_PVIEW_CYCLES	Cycles cluster cannot commit due to a register interlock, register port conflict, bubbles due to a timing class such as tc_3stall, no B2B HVX, or HVX FIFO is full.
0xec	SMT_DU_CONFLICT_PVIEW_CYCLES	Cycles cluster cannot commit due to a DU resource conflict.
0xed	COPROC_BUSY_PVIEW_CYCLES	Cycles cluster cannot commit due to a D-cache uncacheable access.
0xee	DU_UNCACHED_PVIEW_CYCLES	Cycles cluster cannot commit due to a DU resource conflict.
0xef	SYSTEM_BUSY_PVIEW_CYCLES	Cycles cluster cannot commit due to system level stalls, including DMA synchronization, ETM is full, Qtimer read is not ready, AXI bus is busy, and global cache operations synchronization.
0xf1	AXI_LINE128_READ_REQUEST_EVEN	Number of 128-byte line read requests issued by the even interleaved AXI master.

**Table 9-1 V73 processor events symbols**

Event	Symbol	Definition
0xf2	AXI_READ_REQUEST_EVEN	All read requests issued by the even interleaved AXI master.
0xf3	AXI_LINE32_READ_REQUEST_EVEN	Number of 32-byte line read requests issued by the even-interleaved AXI master.
0xf4	AXI_WRITE_REQUEST_EVEN	All write requests issued by the even-interleaved AXI master.
0xf5	AXI_LINE32_WRITE_REQUEST_EVEN	Number of 32-byte line write requests issued by the even-interleaved AXI master. All bytes are valid.
0xf6	AXI_LINE128_WRITE_REQUEST_EVEN	Number of 128-byte line write requests issued by the even-interleaved AXI master. All bytes are valid.
0xf8	AXI_LINE64_READ_REQUEST_EVEN	Number of 64-byte line read requests issued by the even-interleaved AXI master.
0xf9	AXI_LINE64_WRITE_REQUEST_EVEN	Number of 64-byte line write requests issued by the even-interleaved AXI master. All bytes are valid.
0xfa	AXI_WR_CONGESTION_EVEN	Even-interleaved AXI write command or data queue is full, and an operation is stuck at the head of the even interleaved AXI master command queue.
0xfb	AXI_INCOMPLETE_WRITE_REQUEST_EVEN	L2 line-sized write was made to the even-interleaved AXI master, but not all bytes were valid. Includes segmented writes. Excludes WT stores. This event captures the number of writes coalesced at a line level.
0xfc	AXI_LINE256_READ_REQUEST_EVEN	Number of 256-byte line read requests issued by an even-interleaved AXI master. All bytes are valid.
0xfd	AXI_LINE256_WRITE_REQUEST_EVEN	Number of 256-byte line write requests issued by an even-interleaved AXI master. All bytes are valid.
0xfe	CYCLES_3_COPROC_THREADS_ONE_CLUSTER	Number of processor cycles during which a cluster has three threads in Run mode with the coprocessor bit (SSR.XE) enabled.
0x11b	HVXLD_L2_SECONDARY_MISS	Of the events in 0xFB, the events where the load cannot be returned due to the immediately prior access for the line being a pending load or pending L2Fetch
0x2fa	L2_CLEAN_CASTOUT	Number of clean line evictions from L2 cache. Triggers when L2 cache evicts a line due to an allocation. Not triggered on cache operations.
0x2fb	AXI3_READ_REQUEST	All read requests issued by the tertiary AXI master. Includes full lines and partial lines.
0x2fc	AXI3_LINE32_READ_REQUEST	Number of 32-byte line read requests issued by the tertiary AXI master.
0x2fd	AXI3_WRITE_REQUEST	All write requests issued by the tertiary AXI master. Includes full lines and partial lines.
0x2fe	AXI3_LINE32_WRITE_REQUEST	Number of 32-byte line write requests issued by the tertiary AXI master.
0x2ff	AXI3_RD_CONGESTION	Tertiary AXI read command queue is full, and an operation is stuck at the head of the primary AXI master command queue. Includes all interleaved requests.
0x300	CYCLES_1_PACKET_COMMITTED	Number of cycles when one packet is committed.



**Table 9-1 V73 processor events symbols**

Event	Symbol	Definition
0x301	CYCLES_2_PACKET_COMMITTED	Number of cycles when two packets are committed.
0x302	CYCLES_3_PACKET_COMMITTED	Number of cycles when three packets are committed.
0x303	CYCLES_4_PACKET_COMMITTED	Number of cycles when four packets are committed.
0x304	SMT_CLUSTER0	Number of cycles when more than one packet is committed in cluster 0.
0x305	SMT_CLUSTER1	Number of cycles when more than one packet is committed in cluster 1.
0x306	SMT_INTERCLUSTER	Number of cycles when packets are committed on both clusters.
0x307	SMT_CONFLICT_FOR_REG_READ_OR_CU_FWD	Number of cases when a packet is SMT-able without slot resource conflicts, but it cannot go due to s0/s1 register read/CU forwarding.
0x308	COMMITTED_PKT_2_THREAD_RUNNING_2T_PLUS_0T	Number of committed packets with two threads running on the same cluster. Running means the threads are not in the Wait or Stop state.
0x309	COMMITTED_PKT_2_THREAD_RUNNING_1T_PLUS_1T	Number of committed packets with two threads running on the different cluster. Running means the threads are not in the Wait or Stop state.
0x30a	COMMITTED_PKT_3_THREAD_RUNNING_3T_PLUS_0T	Number of committed packets with three threads running on the same cluster. Running means the threads are not in the Wait or Stop state.
0x30b	COMMITTED_PKT_3_THREAD_RUNNING_2T_PLUS_1T	Number of committed packets with three threads running, two threads on one cluster and 1 on another cluster. Running means the threads are not in the Wait or Stop state.
0x30c	COMMITTED_PKT_4_THREAD_RUNNING_4T_PLUS_0T	Number of committed packets with four threads running on same cluster. Running means the threads are not in the Wait or Stop state.
0x30d	COMMITTED_PKT_4_THREAD_RUNNING_3T_PLUS_1T	Number of committed packets with four threads running, three threads on one cluster and one thread running on another cluster. Running means the threads are not in the Wait or Stop state.
0x30e	COMMITTED_PKT_4_THREAD_RUNNING_2T_PLUS_2T	Number of committed packets with four threads running, two threads on one cluster and two threads running on another cluster. Running means the threads are not in the Wait or Stop state.
0x30f	COMMITTED_PKT_5_THREAD_RUNNING_4T_PLUS_1T	Number of committed packets with 5 threads running, four threads on one cluster and one thread running on another cluster. Running means the threads are not in the Wait or Stop state.
0x310	COMMITTED_PKT_5_THREAD_RUNNING_3T_PLUS_2T	Number of committed packets with 5 threads running, three threads on one cluster and two threads running on another cluster. Running means the threads are not in the Wait or Stop state.

**Table 9-1 V73 processor events symbols**

Event	Symbol	Definition
0x311	COMMITTED_PKT_6_THREAD_RUNNING_4T_PLUS_2T	Number of committed packets with six threads running, four threads on one cluster and two threads running on another cluster. Running means the threads are not in the Wait or Stop state.
0x312	COMMITTED_PKT_6_THREAD_RUNNING_3T_PLUS_3T	Number of committed packets with six threads running, three threads on one cluster and three threads running on another cluster. Running means the threads are not in the Wait or Stop state.
0x313	ICACHE_DEMAND_MISS_PREFETCH_MISS	Number of iprfetches initiated on demand misses.
0x314	SIMPLE_PACKET	Number of committed simple packets, which can be dispatched on in-cluster SMT threads. Includes eligible packets that are committed on both primary and in-cluster SMT threads.
0x315	AXI3_LINE64_WRITE_REQUEST	Number of 64-byte line write requests issued by the primary AXI3 master. Includes all interleaved requests. All bytes are valid.
0x316	AXI3_LINE64_READ_REQUEST	Number of 64-byte line read requests issued by the primary AXI3 master. Includes all interleaved requests
0x317	AXI3_WR_CONGESTION	Tertiary AXI write command or data queue is full, and an operation is stuck at the head of the primary AXI3 master command queue. Includes all interleaved requests.
0x318	AXI3_INCOMPLETE_WRITE_REQUEST	L2 line-sized write was made to the AXI3 master, but not all bytes were valid. Includes segmented writes. Excludes WT stores.
0x319	ICACHE_DATA_REPLAY	Number of I-cache data replays due to incorrect way predictions.
0x31c	SMT_PKT_PICKED_BUT_NOT_COMMIT_PVIEW_CYCLES	In-cluster SMT thread is picked, but not committed.
0x31d	SMT_PKT_IQ_NO_PKT_PVIEW_CYCLES	In-cluster SMT thread is not picked because no packet is in IQ on the SMT thread.
0x31e	SMT_PKT_NOT_SIMPLE_PVIEW_CYCLES	In-cluster SMT thread is not picked because no simple packet is on the SMT thread.
0x31f	SMT_PKT_NOT_READY_PVIEW_CYCLES	In-cluster SMT thread is not picked because no simple packet is ready for dispatch on the SMT thread.
0x320	SMT_PKT_SLOT_CONFLICT_PVIEW_CYCLES	In-cluster SMT thread is not picked due to a slot conflict between the primary thread and SMT thread.
0x321	SMT_PKT_REG_FWD_BLOCK_PVIEW_CYCLES	In-cluster SMT thread is not picked due to a register and forward block on the SMT thread.
0x322	CLADE2_EB_FULL	CLADE2 can use up to two eviction buffer entries. Indicates that both entries are used and CLADE2 is congested.
0x323	CLADE2_RD_REQ	Number of L2 cache read requests in the CLADE2 region.
0x324	CLADE2_RDCACHE_MISS	Number of L2 cache read request misses in the CLADE2 region.
0x325	CLADE2_WR_REQ	Number of L2 cache write requests in the CLADE2 region.

**Table 9-1 V73 processor events symbols**

Event	Symbol	Definition
0x326	CLADE2_WRCACHE_MISS	Number of L2 cache write request misses in the CLADE2 region.
0x327	AXI_EWD_REQUEST	L2 cache eviction of clean data to the AXI master bus.
0x328	AXI_EWD_REQUEST_EVEN	L2 cache eviction of clean data to the even interleaved AXI master bus.
0x329	AXI_CMO_REQUEST	Cache maintenance operation request from the QDSP6 core to AXIM.
0x32a	AXI_CMO_REQUEST_EVEN	Cache maintenance operation request from the QDSP6 core to AXIM Interleave0.
0x32b	ICACHE_DEMAND_MISS_PREFETCH_MISS_IU0	Number of ipfetched initiated on a demand miss in IU0.
0x32c	ICACHE_DEMAND_MISS_PREFETCH_MISS_IU1	Number of ipfetched initiated on a demand miss in IU1.
0x330	VMEM_ST_SMT_DU_PORT_CONFLICT_REPLAY	Number of times any packet takes a replay because CU didn't allocate a port at schedule time and it didn't arbitrate for the port in DU based on the state bit (VMEMSttoVTCM) but needed a port as it mapped to L2.
0x333	DU_SPF_DTLBPGCROSS	Number of stopping prefetching due to page cross
0x334	DU_SPF_DCACHE_HIT	Number of prefetch requests hitting in L1D\$
0x335	DU_SPF_DCACHE_MISS	Number of prefetch requests missing in L1D\$
0x336	DU_SPF_L2FIFOFULL_RETRY	Number of prefetch retry on L2FIFO queue full
0x337	DU_SPF_L2BUFFULL_RETRY	Number of prefetch retry on L2 credits/AQoS busy
0x338	DU_SPF_CONFLICT_RETRY	Number of cycles that prefetches losing arbitration
0x350	DU_NUM_WAY_PREDICTIONS	Number of times DU did way predictions for loads
0x351	DU_WAY_PRED_REPLAYS	Number of times DU replayed due to way misprediction
0x352	DU_BANKCONFLICTREPLAY_INVALID	Number of time bank conflict replay was prevented due to way prediction
0xbbf	DU_CACHE_MISS_L2HIT_PVIEW_CYCLES	Cycles cluster cannot commit due to D-cache cacheable miss hitting in L2.
0xbc0	DU_CACHE_MISS_TCM_PVIEW_CYCLES	Cycles cluster cannot commit due to D-cache cacheable miss going to TCM.
0xbc1	DU_CACHE_MISS_AXI_PVIEW_CYCLES	Cycles cluster cannot commit due to D-cache cacheable miss going to AXI.
0xbc2	DU_CACHE_MISS_AHB_PVIEW_CYCLES	Cycles cluster cannot commit due to D-cache cacheable miss going to AHB.
0xbc3	DU_CACHE_MISS_AXI2_PVIEW_CYCLES	Cycles cluster cannot commit due to D-cache cacheable miss going to AXI2.
0xbc4	ICACHE_DEMAND_MISS_L2HIT_PRI	—
0xbc5	ICACHE_DEMAND_MISS_L2MISS_PRI	—
0xbc6	ICACHE_DEMAND_MISS_L2TCMHIT_PRI	—
0xbc7	ICACHE_DEMAND_MISS_L2ITCMHIT_PRI	—

**Table 9-1 V73 processor events symbols**

<b>Event</b>	<b>Symbol</b>	<b>Definition</b>
0xbc8	ICACHE_IPREFETCHES_SENT_L2HIT	—
0xbc9	ICACHE_IPREFETCHES_SENT_L2MISS	—
0xbca	ICACHE_IPREFETCHES_SENT_L2TCM	—
0xpcb	ICACHE_IPREFETCHES_SENT_L2ITCM	—

# 10 Instruction Encoding

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## 10.1 Instructions

Hexagon processor instructions are encoded in a 32-bit instruction word. The instruction word format varies according to the instruction type.

The instruction words contain two types of bit fields:

- Common fields appear in every processor instruction, and are defined the same in all instructions.
- Instruction-specific fields appear only in some instructions, or vary in definition across the instruction set.

**Table 10-1 Instruction bit fields**

Name	Description	Type
ICLASS	Instruction class	Common
Parse	Packet / loop bits	

**Table 10-1 Instruction bit fields**

Name	Description	Type
MajOp Maj	Major opcode	Instruction-specific
MinOp Min	Minor opcode	
RegType	Register type (32-bit, 64-bit)	
Type	Operand type (byte, halfword, and so on)	
Amode	Addressing mode	
<i>dn</i>	Destination register operand	
<i>sn</i>	Source register operand	
<i>tn</i>	Source register operand #2	
<i>xn</i>	Source and destination register operand	
<i>un</i>	Predicate or modifier register operand	
sH	Source register bit field (Rs.H or Rs.L)	
tH	Source register #2 bit field (Rt.H or Rt.L)	
UN	Unsigned operand	
Rs	No source register read	
P	Predicate expression	
PS	Predicate sense (Pu or !Pu)	
DN	Dot-new predicate	
PT	Predict taken	
sm	Supervisor mode only	

**NOTE:** In some cases, instruction-specific fields encode instruction attributes other than the ones described for the fields in [Table 10-1](#).

### Reserved bits

Some instructions contain reserved bits that do not currently encode instruction attributes. Always set these bits to 0 to ensure compatibility with any future changes in the instruction encoding.

**NOTE:** Reserved bits appear as '-' characters in the instruction encoding tables.

## 10.2 Sub-instructions

To reduce code size, the Hexagon processor supports the encoding of certain pairs of instructions in a single 32-bit container. Instructions encoded this way are sub-instructions, and the containers are duplexes ([Section 10.3](#)).

Sub-instructions are limited to certain commonly-used instructions:

- Arithmetic and logical operations
- Register transfer
- Loads and stores
- Stack frame allocation/deallocation
- Subroutine return

[Table 10-2](#) lists the sub-instructions along with the group identifiers that encode them in duplexes.

Sub-instructions can access only a subset of the general registers (R0 to R7, R16 to R23).

[Table 10-3](#) lists the sub-instruction register encodings.

**NOTE:** Certain sub-instructions implicitly access registers such as SP (R29).

**Table 10-2 Sub-instructions**

Group	Instruction	Description
L1	Rd = memw (Rs+#u4:2)	Word load
L1	Rd = memub (Rs+#u4:0)	Unsigned byte load
Group	Instruction	Instruction
L2	Rd = memh/memuh (Rs+#u3:1)	Halfword loads
L2	Rd = memb (Rs+#u3:0)	Signed byte load
L2	Rd = memw (r29+#u5:2)	Load word from stack
L2	Rdd = memd (r29+#u5:3)	Load pair from stack
L2	deallocframe	Deallocate stack frame
L2	if ([!]P0) dealloc_return if ([!]P0.new) dealloc_return:nt	Deallocate stack frame and return
L2	jumpr R31 if ([!]P0) jumpr R31 if ([!]P0.new) jumpr:nt R31	Return
Group	Instruction	Instruction
S1	memw (Rs+#u4:2) = Rt	Store word
S1	memb (Rs+#u4:0) = Rt	Store byte
Group	Instruction	Instruction
S2	memh (Rs+#u3:1) = Rt	Store halfword
S2	memw (r29+#u5:2) = Rt	Store word to stack
S2	memd (r29+#s6:3) = Rtt	Store pair to stack
S2	memw (Rs+#u4:2) = #U1	Store immediate word #0 or #1

**Table 10-2 Sub-instructions (cont.)**

Group	Instruction	Description
S2	memb(Rs+#u4) = #U1	Store immediate byte #0 or #1
S2	allocframe(#u5:3)	Allocate stack frame
Group	Instruction	Instruction
A	Rx = add(Rx, #s7)	Add immediate
A	Rd = Rs	Transfer
A	Rd = #u6	Set to unsigned immediate
A	Rd = #-1	Set to -1
A	if ([!]P0[.new]) Rd = #0	Conditional clear
A	Rd = add(r29, #u6:2)	Add immediate to stack pointer
A	Rx = add(Rx, Rs)	Register add
A	P0 = cmp.eq(Rs, #u2)	Compare register equal immediate
A	Rdd = combine(#0, Rs)	Combine zero and register into pair
A	Rdd = combine(Rs, #0)	Combine register and zero into pair
A	Rdd = combine(#u2, #U2)	Combine immediates into pair
A	Rd = add(Rs, #1) Rd = add(Rs, #-1)	Add and subtract 1
A	Rd = sxth/sxtb/zxtb/zxth(Rs)	Sign- and zero-extends
A	Rd = and(Rs, #1)	And with 1

**Table 10-3 Sub-instruction registers**

Register	Encoding
Rs, Rt, Rd, Rx	0000 = R0 0001 = R1 0010 = R2 0011 = R3 0100 = R4 0101 = R5 0110 = R6 0111 = R7 1000 = R16 1001 = R17 1010 = R18 1011 = R19 1100 = R20 1101 = R21 1110 = R22 1111 = R23



**Table 10-3 Sub-instruction registers**

Register	Encoding
Rdd, Rtt	000 = R1:0
	001 = R3:2
	010 = R5:4
	011 = R7:6
	100 = R17:16
	101 = R19:18
	110 = R21:20
	111 = R23:22

## 10.3 Duplexes

A duplex is encoded as a 32-bit instruction with bits [15:14] set to 00. The sub-instructions that comprise a duplex are encoded as 13-bit fields in the duplex.

An instruction packet can contain one duplex and up to two other (non-duplex) instructions. The duplex must always appear as the last word in a packet.

The sub-instructions in a duplex always execute in slot 0 and slot 1.

**Table 10-4 Duplex instruction encoding**

Bits	Name	Description
15:14	Parse bits	00 = Duplex type, ends the packet and indicates that word contains two sub-instructions
12:0	Sub-instruction low	Encodes slot 0 sub-instruction
28:16	Sub-instruction high	Encodes slot 1 sub-instruction
31:29, 13	4-bit ICLASS	Indicates which group the low and high sub-instructions belong to.

The duplex ICLASS field values that specify the group of each sub-instruction in a duplex are shown in [Table 10-5](#)

**Table 10-5 Duplex ICLASS field**

ICLASS	Low slot 0 subinsn type	High slot 1 subinsn type
0x0	L1-type	L1-type
0x1	L2-type	L1-type
0x2	L2-type	L2-type
0x3	A-type	A-type
0x4	L1-type	A-type
0x5	L2-type	A-type
0x6	S1-type	A-type
0x7	S2-type	A-type
0x8	S1-type	L1-type
0x9	S1-type	L2-type
0xA	S1-type	S1-type

**Table 10-5 Duplex ICLASS field**

ICLASS	Low slot 0 subinsn type	High slot 1 subinsn type
0xB	S2-type	S1-type
0xC	S2-type	L1-type
0xD	S2-type	L2-type
0xE	S2-type	S2-type
0xF	Reserved	Reserved

Duplexes have the following grouping constraints:

- **Constant extenders** expand the range of an instruction's immediate operand to 32 bits, and can expand the following sub-instructions:
  - $Rx = \text{add}(Rx, \#s7)$
  - $Rd = \#u6$

A duplex can contain only one constant-extended instruction, and it must appear in the slot 1 position.
- When the sub-instructions are treated as 13-bit unsigned integer values for two instructions with the same sub-instruction group in a duplex, the instruction corresponding to the numerically smaller value must be encoded in the slot 1 position of the duplex.<sup>1</sup>
- Sub-instructions must conform to any slot assignment grouping rules that apply to the individual instructions, even if a duplex pattern exists that violates those assignments. One exception to this rule exists:
  - `jumpr R31` must appear in the Slot 0 position

<sup>1</sup> The sub-instruction register and immediate fields are assumed to be 0 when performing this comparison.

## 10.4 Instruction classes

The instruction class ([Section 3.2](#)) is encoded in the four most-significant bits of the instruction word (31:28). These bits are referred to as the instruction's ICLASS field. The Slots column in [Table 10-6](#) indicates which slots can receive the instruction class.

**Table 10-6 Instruction class encoding**

Encoding	Instruction class	Slots
0000	Constant extender ( <a href="#">Section 10.9</a> )	–
0001	J	2,3
0010	J	2,3
0011	LD ST	0,1
0100	LD ST (conditional or GP-relative)	0,1
0101	J	2,3
0110	CR	3
0111	ALU32	0,1,2,3
1000	XTYPE	2,3
1001	LD	0,1
1010	ST	0
1011	ALU32	0,1,2,3
1100	XTYPE	2,3
1101	XTYPE	2,3
1110	XTYPE	2,3
1111	ALU32	0,1,2,3

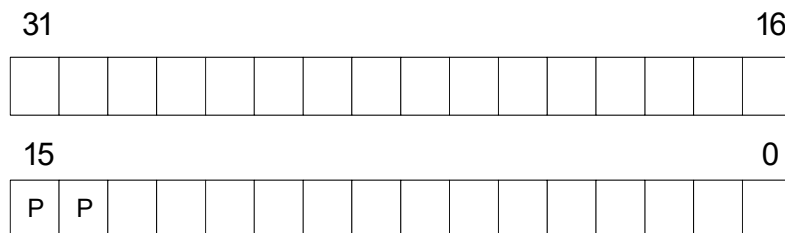
For details on encoding the individual class types, see [Chapter 11](#).

## 10.5 Instruction packets

Instruction packets are encoded using two bits of the instruction word (15:14), which are referred to as the Parse field of the instruction word. The field values have the following definitions:

- '11' indicates that an instruction is the last instruction in a packet (the instruction word at the highest address).
- '01' or '10' indicate that an instruction is not the last instruction in a packet.
- '00' indicates a duplex.

If any sequence of four consecutive instructions occurs without one of them containing '11' or '00', the processor raises an error exception (illegal opcode).



Packet / loop parse bits:

01, 10 = not end of packet

11 = end of packet

00 = duplex

**Figure 10-1 Parse field instruction packet encoding**

The following examples show how to use the Parse field to encode instruction packets:

```
{ A ; B}
 01 11           // Parse fields of instructions A,B

{ A ; B ; C}
 01 01  11      // Parse fields of instructions A,B,C

{ A ; B ; C ; D}
 01 01  01  11  // Parse fields of instructions A,B,C,D
```

## 10.6 Loop packets

In addition to encoding the last instruction in a packet, the Parse field of the instruction word (Section 10.5) encodes the last packet in a hardware loop.

The Hexagon processor supports two **Hardware loops**, labeled 0 and 1. The last packet in these loops is subject to the following restrictions:

- The last packet in a hardware loop 0 must contain two or more instruction words.
- The last packet in a hardware loop 1 must contain three or more instruction words.

If the last packet in a loop is expressed in assembly language with fewer than the required number of words, the assembler automatically adds one or two NOP instructions to the encoded packet so it contains the minimum required number of instruction words.

The Parse fields in a packet's first and second instruction words (the words at the lowest addresses) encode whether the packet is the last packet in a hardware loop.

**Table 10-7 Parse field loop packet encoding**

Packet	Parse field in first Instruction	Parse field in second Instruction
Not last in loop	01 or 11	01 or 11 <sup>1</sup>
Last in loop 0	10	01 or 11
Last in loop 1	01	10
Last in loops 0 & 1	10	10

<sup>1</sup> Not applicable for single-instruction packets.

The following examples show how to use the Parse field to encode loop packets:

```
{ A  B}:endloop0
 10 11 // Parse fields of instrs A,B

{ A  B  C}:endloop0
 10 01 11 // Parse fields of instrs A,B,C

{ A  B  C  D}:endloop0
 10 01 01 11 // Parse fields of instrs A,B,C,D

{ A  B  C}:endloop1
 01 10 11 // Parse fields of instrs A,B,C

{ A  B  C  D}:endloop1
 01 10 01 11 // Parse fields of instrs A,B,C,D

{ A  B  C}:endloop0:endloop1
 10 10 11 // Parse fields of instrs A,B,C

{ A  B  C  D}:endloop0:endloop1
 10 10 01 11 // Parse fields of instrs A,B,C,D
```

## 10.7 Immediate values

To conserve encoding space, the Hexagon processor often stores immediate values in instruction fields that are smaller (in bit size) than the values actually needed in the instruction operation.

When an instruction operates on one of its immediate operands, the processor automatically extends the immediate value to the bit size required by the operation:

- Signed immediate values are sign-extended
- Unsigned immediate values are zero-extended

## 10.8 Scaled immediate values

To minimize the number of bits in instruction words to store certain immediate values, the Hexagon processor stores the values as scaled immediate values. Use scaled immediate values when an immediate value must represent integral multiples of a power of 2 in a specific range.

For example, consider an instruction operand whose possible values are the following:

-32, -28, -24, -20, -16, -12, -8, -4, 0, 4, 8, 12, 16, 20, 24, 28

Encoding the full range of integers -32...28 normally requires 6 bits. However, if the operand is stored as a scaled immediate, it can first be shifted right by two bits, storing only the four remaining bits in the instruction word. When the operand is fetched from the instruction word, the processor automatically shifts the value left by two bits to recreate the original operand value.

**NOTE:** The scaled immediate value in the example above is represented notationally as #s4 : 2.

Scaled immediate values commonly encode address offsets that apply to data types of varying size. For example, [Table 10-8](#) shows how to use the byte offsets in immediate-with-offset addressing mode that are stored as 11-bit scaled immediate values. This enables the offsets to span the same range of data elements regardless of the data type.

**Table 10-8 Scaled immediate encoding (indirect offsets)**

Data type	Offset size (stored)	Scale bits	Offset size (effective)	Offset range (bytes)	Offset range (elements)
byte	11	0	11	-1024 ... 1023	-1024 ... 1023
halfword	11	1	12	-2048 ... 2046	-1024 ... 1023
word	11	2	13	-4096 ... 4092	-1024 ... 1023
doubleword	11	3	14	-8192 ... 8184	-1024 ... 1023

## 10.9 Constant extenders

To support the use of 32-bit operands in a number of instructions, the Hexagon processor defines constant extenders, which are an instruction word that exists solely to extend the bit range of an immediate or address operand that is contained in an adjacent instruction in a packet.

For example, the absolute addressing mode specifies a 32-bit constant value as the effective address. Instructions using this addressing mode are encoded in a single packet containing both the normal instruction word and a second word with a constant extender that increases the range of the instruction's normal constant operand to a full 32 bits.

**NOTE:** Constant extended operands can encode symbols.

A constant extender is encoded as a 32-bit instruction with the 4-bit ICLASS field set to 0 and the 2-bit Parse field set to its usual value (Section 10.5). The remaining 26 bits in the instruction word store the data bits that are prepended to an operand as small as 6 bits to create a full 32-bit value.

**Table 10-9 Constant extender encoding**

Bits	Name	Description
31:28	ICLASS	Instruction class = 0000
27:16	Extender high	High 12 bits of 26-bit constant extension
15:14	Parse	Parse bits
13:0	Extender low	Low 14 bits of 26-bit constant extension

Within a packet, a constant extender must be positioned immediately before the instruction that it extends: in terms of memory addresses, the extender word must reside at address (<instr\_address> - 4).

The constant extender effectively serves as a prefix for an instruction: it does not execute in a slot, nor does it consume any slot resources. All packets must contain four or fewer words, and the constant extender occupies one word.

If the instruction operand to extend is longer than 6 bits, the overlapping bits in the base instruction must be encoded as zeros. The value in the constant extender always supplies the upper 26 bits.

The Regclass field in Table 10-10 lists the values to set bits [27:24] to in the instruction word to identify the instruction as one that might include a constant extender.

**NOTE:** When the base instruction encodes two constant operands, the extended immediate is the one specified in the table.

Constant extenders appear in disassembly listings as Hexagon instructions with the name `immext`.

**NOTE:** If a constant extender is encoded in a packet for an instruction that does not accept a constant extender, the execution result is undefined. The assembler normally ensures that only valid constant extenders are generated.

Table 10-10 Constant extender instructions

ICLASS	Regclass	Instructions
LD	---1	Rd = mem{b,ub,h,uh,w,d} (##U32) if ([!]Pt[.new]) Rd = mem{b,ub,h,uh,w,d} (Rs + ##U32) // predicated loads
LD	----	Rd = mem{b,ub,h,uh,w,d} (Rs + ##U32) Rd = mem{b,ub,h,uh,w,d} (Re=##U32) Rd = mem{b,ub,h,uh,w,d} (Rt<<#u2 + ##U32) if ([!]Pt[.new]) Rd = mem{b,ub,h,uh,w,d} (##U32)
ST	---0	mem{b,h,w,d} (##U32) = Rs[.new] // GP-stores if ([!]Pt[.new]) mem{b,h,w,d} (Rs + ##U32) = Rt[.new] // predicated stores
ST	----	mem{b,h,w,d} (Rs + ##U32) = Rt[.new] mem{b,h,w,d} (Rd=##U32) = Rt[.new] mem{b,h,w,d} (Ru<<#u2 + ##U32) = Rt[.new] if ([!]Pt[.new]) mem{b,h,w,d} (##U32) = Rt[.new]
MEMOP	----	[if [!]Ps] memw(Rs + #u6) = ##U32 // constant store memw(Rs + Rt<<#u2) = ##U32 // constant store
NV	----	if (cmp.xx(Rs.new,##U32)) jump:hint target
ALU32	----	Rd = ##u32 Rdd = combine(Rs,##u32) Rdd = combine(##u32,Rs) Rdd = combine(##u32,#s8) Rdd = combine(#s8,##u32) Rd = mux (Pu, Rs,##u32) Rd = mux (Pu, ##u32, Rs) Rd = mux (Pu, ##u32, #s8) if ([!]Pu[.new]) Rd = add(Rs,##u32) if ([!]Pu[.new]) Rd = ##u32 Pd = [!]cmp.eq (Rs,##u32) Pd = [!]cmp.gt (Rs,##u32) Pd = [!]cmp.gtu (Rs,##u32) Rd = [!]cmp.eq(Rs,##u32) Rd = and(Rs,##u32) Rd = or(Rs,##u32) Rd = sub(##u32,Rs)
ALU32	----	Rd = add(Rs,##s32)
XTYPE	00--	Rd = mpyi(Rs,##u32) Rd += mpyi(Rs,##u32) Rd -= mpyi(Rs,##u32) Rx += add(Rs,##u32) Rx -= add(Rs,##u32)
ALU32	---- 1	Rd = ##u32 Rd = add(Rs,##s32)
J	1---	jump (PC + ##s32) call (PC + ##s32) if ([!]Pu) call (PC + ##s32)
CR	----	Pd = spNloop0(PC+##s32,Rs/#U10) loop0/1 (PC+##s32,##Rs/#U10)



**Table 10-10 Constant extender instructions (cont.)**

ICLASS	Regclass	Instructions
XTYPE	1---	Rd = add (pc, ##s32) Rd = add (##u32, mpyi (Rs, #u6) ) Rd = add (##u32, mpyi (Rs, Rt) ) Rd = add (Rs, add (Rt, ##u32) ) Rd = add (Rs, sub (##u32, Rt) ) Rd = sub (##u32, add (Rs, Rt) ) Rd = or (Rs, and (Rt, ##u32) ) Rx = add/sub/and/or (##u32, asl/asr/lsr (Rx, #U5) ) Rx = add/sub/and/or (##u32, asl/asr/lsr (Rs, Rx) ) Rx = add/sub/and/or (##u32, asl/asr/lsr (Rx, Rs) ) Pd = cmpb/h. {eq, gt, gtu} (Rs, ##u32)

<sup>1</sup> Constant extension is only for a Slot 1 sub-instruction.

### Encoding 32-bit address operands in load/stores

Two methods exist for encoding a 32-bit absolute address in a load or store instruction:

1. For unconditional load/stores, the GP-relative load/store instruction is used. The assembler encodes the absolute 32-bit address as follows:

- The upper 26 bits are encoded in a constant extender
- The lower 6 bits are encoded in the 6 operand bits contained in the GP-relative instruction

In this case the 32-bit value encoded must be a plain address, and the value stored in the GP register is ignored.

**NOTE:** When a constant extender is explicitly specified with a GP-relative load/store, the processor ignores the value in GP and creates the effective address directly from the 32-bit constant value.

2. For conditional load/store instructions that have their base address encoded only by a 6-bit immediate operand, a constant extender must be explicitly specified; otherwise, the execution result is undefined. The assembler ensures that these instructions always include a constant extender.

This case applies also to instructions that use the absolute-set addressing mode or absolute-plus-register-offset addressing mode.

### Encoding 32-bit immediate operands

The immediate operands of certain instructions use scaled immediates ([Section 10.8](#)) to increase their addressable range. When using constant extenders, scaled immediates are not scaled by the processor. Instead, the assembler must encode the full 32-bit unscaled value as follows:

- The upper 26 bits are encoded in the constant extender
- The lower six 6 bits are encoded in the base instruction in the least-significant bit positions of the immediate operand field.
- Any overlapping bits in the base instruction are encoded as zeros.

## Encoding 32-bit jump/call target addresses

When a jump/call has a constant extender, the resulting target address is forced to a 32-bit alignment (bits 1:0 in the address are cleared by hardware). The resulting jump/call operation never causes an alignment violation.

## 10.10 New-value operands

Instructions that include a new-value register operand specify in their encodings which instruction in the packet has its destination register accessed as the new-value register.

New-value consumers include a 3-bit instruction field named `Nt` that specifies this information.

- `Nt[0]` is reserved and must always be encoded as zero. A nonzero value produces undefined results.
- `Nt[2:1]` encodes the distance (in instructions) from the producer to the consumer, as follows:
  - `Nt[2:1] = 00` // reserved
  - `Nt[2:1] = 01` // producer is +1 instruction ahead of consumer
  - `Nt[2:1] = 10` // producer is +2 instructions ahead of consumer
  - `Nt[2:1] = 11` // producer is +3 instructions ahead of consumer

“ahead” is defined here as the instruction encoded at a lower memory address than the consumer instruction, not counting empty slots or constant extenders. For example, the following producer/consumer relationship is encoded with `Nt[2:1]` set to `01`.

```
...
<producer instruction word>
<consumer constant extender word>
<consumer instruction word>
...
```

**NOTE:** Instructions with 64-bit register pair destinations cannot produce new-values. The assembler flags this case with an error, as the result is undefined.

## 10.11 Instruction mapping

Some Hexagon processor instructions are encoded by the assembler as variants of other instructions. This is done for operations that are functionally equivalent to other instructions, but are still defined as separate instructions because of their programming utility as common operations.

**Table 10-11 Instructions mapped to other instructions**

Instruction	Mapping
<code>Rd = not (Rs)</code>	<code>Rd = sub (#-1, Rs)</code>
<code>Rd = neg (Rs)</code>	<code>Rd = sub (#0, Rs)</code>
<code>Rdd = Rss</code>	<code>Rdd = combine (Rss.H32, Rss.L32)</code>

# 11 Instruction Set

---

This chapter describes the instruction set for version 7 of the Hexagon processor.

The instructions are listed alphabetically within instruction categories. The following information is provided for each instruction:

- Instruction name
- A brief description of the instruction
- A high-level functional description (syntax and behavior) with possible operand types
- Instruction class and slot information for grouping instructions in packets
- C intrinsic functions that provide access to the instruction
- Instruction encoding

## 11.1 ALU32

The ALU32 instruction class includes instructions that perform arithmetic and logical operations on 32-bit data.

ALU32 instructions are executable on any slot.

### 11.1.1 ALU32 ALU

The ALU32 ALU instruction subclass includes instructions that perform arithmetic and logical operations on individual 32-bit items.

#### Add

Add a source register either to another source register or to a signed 16-bit immediate value. Store the result in destination register. Source and destination registers are 32 bits. If the result overflows 32 bits, it wraps around. Optionally saturate result to a signed value between 0x80000000 and 0x7fffffff.

For 64-bit versions of this operation, see the XTYPE add instructions.

Syntax	Behavior
<code>Rd=add(Rs, #s16)</code>	<code>apply_extension(#s); Rd=Rs+#s;</code>
<code>Rd=add(Rs, Rt)</code>	<code>Rd=Rs+Rt;</code>
<code>Rd=add(Rs, Rt) : sat</code>	<code>Rd=sat<sub>32</sub>(Rs+Rt);</code>

**Class: ALU32 (slots 0,1,2,3)**

#### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

#### Intrinsics

<code>Rd=add(Rs, #s16)</code>	<code>Word32 Q6_R_add_RI(Word32 Rs, Word32 Is16)</code>
<code>Rd=add(Rs, Rt)</code>	<code>Word32 Q6_R_add_RR(Word32 Rs, Word32 Rt)</code>
<code>Rd=add(Rs, Rt) : sat</code>	<code>Word32 Q6_R_add_RR_sat(Word32 Rs, Word32 Rt)</code>

## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS											s5					Parse												d5					
1	0	1	1	i	i	i	i	i	i	i	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=add(Rs,#s16)
ICLASS				P	MajOp			MinOp			s5					Parse		t5					d5										
1	1	1	1	0	0	1	1	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=add(Rs,Rt)	
1	1	1	1	0	1	1	0	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=add(Rs,Rt):sat	

Field name	Description
MajOp	Major opcode
MinOp	Minor opcode
P	Predicated
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Logical operations

Perform bitwise logical operations (AND, OR, XOR, NOT) either on two source registers or on a source register and a signed 10-bit immediate value. Store result in destination register. Source and destination registers are 32 bits.

For 64-bit versions of these operations, see the [XTYPE](#) logical instructions.

Syntax	Behavior
Rd=and(Rs, #s10)	apply_extension(#s); Rd=Rs&#s;
Rd=and(Rs, Rt)	Rd=Rs&Rt;
Rd=and(Rt, ~Rs)	Rd = (Rt & ~Rs);
Rd=not(Rs)	Assembler mapped to: "Rd=sub(#-1, Rs)"
Rd=or(Rs, #s10)	apply_extension(#s); Rd=Rs #s;
Rd=or(Rs, Rt)	Rd=Rs Rt;
Rd=or(Rt, ~Rs)	Rd = (Rt   ~Rs);
Rd=xor(Rs, Rt)	Rd=Rs^Rt;

### Class: ALU32 (slots 0,1,2,3)

#### Intrinsics

Rd=and(Rs, #s10)	Word32 Q6_R_and_RI(Word32 Rs, Word32 Is10)
Rd=and(Rs, Rt)	Word32 Q6_R_and_RR(Word32 Rs, Word32 Rt)
Rd=and(Rt, ~Rs)	Word32 Q6_R_and_RnR(Word32 Rt, Word32 Rs)
Rd=not(Rs)	Word32 Q6_R_not_R(Word32 Rs)
Rd=or(Rs, #s10)	Word32 Q6_R_or_RI(Word32 Rs, Word32 Is10)
Rd=or(Rs, Rt)	Word32 Q6_R_or_RR(Word32 Rs, Word32 Rt)
Rd=or(Rt, ~Rs)	Word32 Q6_R_or_RnR(Word32 Rt, Word32 Rs)
Rd=xor(Rs, Rt)	Word32 Q6_R_xor_RR(Word32 Rs, Word32 Rt)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		Rs	MajOp		MinOp		s5					Parse					d5															
0	1	1	1	0	1	1	0	0	0	i	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=and(Rs,#s10)
0	1	1	1	0	1	1	0	1	0	i	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=or(Rs,#s10)
ICLASS		P	MajOp		MinOp		s5					Parse					t5					d5										

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
1	1	1	1	0	0	0	1	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=and(Rs,Rt)
1	1	1	1	0	0	0	1	0	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=or(Rs,Rt)
1	1	1	1	0	0	0	1	0	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=xor(Rs,Rt)
1	1	1	1	0	0	0	1	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=and(Rt,~Rs)
1	1	1	1	0	0	0	1	1	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=or(Rt,~Rs)

Field name	Description
MajOp	Major opcode
MinOp	Minor opcode
Rs	No Rs read
P	Predicated
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Negate

Perform arithmetic negation on a source register. Store result in destination register. Source and destination registers are 32 bits.

For 64-bit and saturating versions of this instruction, see the XTYPE-class negate instructions.

Syntax	Behavior
<code>Rd=neg(Rs)</code>	Assembler mapped to: <code>"Rd=sub(#0,Rs)"</code>

**Class:** N/A

### Intrinsics

<code>Rd=neg(Rs)</code>	<code>Word32 Q6_R_neg_R(Word32 Rs)</code>
-------------------------	---



## NOP

Perform no operation. This instruction is used for padding and alignment.

Within a packet, it can be positioned in any slot 0 through 3.

### Syntax

nop

### Behavior

**Class: ALU32 (slots 0,1,2,3)**

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS				Rs		MajOp								Parse																			
0	1	1	1	1	1	1	1	1	-	-	-	-	-	-	-	-	P	P	-	-	-	-	-	-	-	-	-	-	-	-	-	-	nop

Field name	Description
MajOp	Major opcode
Rs	No Rs read
ICLASS	Instruction class
Parse	Packet/loop parse bits

## Subtract

Subtract a source register from either another source register or from a signed 10-bit immediate value. Store the result in the destination register. Source and destination registers are 32 bits. If the result underflows 32 bits, it wraps around. Optionally saturate result to a signed value between 0x8000\_0000 and 0x7fff\_ffff.

For 64-bit versions of this operation, see the XTYPE subtract instructions.

Syntax	Behavior
Rd=sub(#s10,Rs)	apply_extension(#s); Rd=#s-Rs;
Rd=sub(Rt,Rs)	Rd=Rt-Rs;
Rd=sub(Rt,Rs):sat	Rd=sat <sub>32</sub> (Rt - Rs);

### Class: ALU32 (slots 0,1,2,3)

#### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

#### Intrinsics

Rd=sub(#s10,Rs)	Word32 Q6_R_sub_IR(Word32 Is10, Word32 Rs)
Rd=sub(Rt,Rs)	Word32 Q6_R_sub_RR(Word32 Rt, Word32 Rs)
Rd=sub(Rt,Rs):sat	Word32 Q6_R_sub_RR_sat(Word32 Rt, Word32 Rs)

#### Encoding

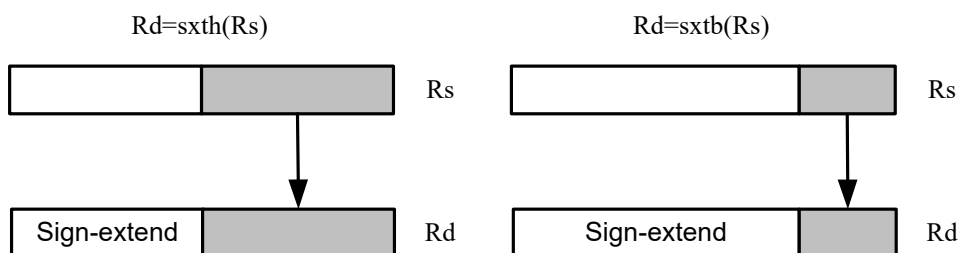
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		Rs		MajOp		MinOp		s5					Parse					d5														
0	1	1	1	0	1	1	0	0	1	i	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=sub(#s10,Rs)
ICLASS		P		MajOp		MinOp		s5					Parse					t5					d5									
1	1	1	1	0	0	1	1	0	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=sub(Rt,Rs)
1	1	1	1	0	1	1	0	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=sub(Rt,Rs):sat

Field name	Description
MajOp	Major opcode
MinOp	Minor opcode
Rs	No Rs read
P	Predicated
ICLASS	Instruction class
Parse	Packet/loop parse bits

<b>Field name</b>	<b>Description</b>
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Sign extend

Sign-extend the least-significant byte or halfword from the source register and place the 32-bit result in the destination register.



### Syntax

`Rd=sxtb(Rs)`

`Rd=sxth(Rs)`

### Behavior

`Rd = sxt8->32(Rs) ;`

`Rd = sxt16->32(Rs) ;`

## Class: ALU32 (slots 0,1,2,3)

### Intrinsics

`Rd=sxtb(Rs)`

`Word32 Q6_R_sxtb_R(Word32 Rs)`

`Rd=sxth(Rs)`

`Word32 Q6_R_sxth_R(Word32 Rs)`

### Encoding

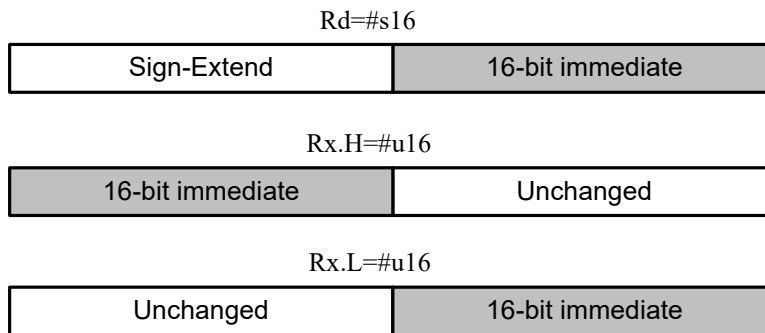
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS			Rs	MajOp			MinOp			s5					Parse		C											d5					
0	1	1	1	0	0	0	0	1	0	1	s	s	s	s	s	P	P	0	-	-	-	-	-	-	-	-	-	d	d	d	d	d	Rd=sxtb(Rs)
0	1	1	1	0	0	0	0	1	1	1	s	s	s	s	s	P	P	0	-	-	-	-	-	-	-	-	-	d	d	d	d	d	Rd=sxth(Rs)

Field name	Description
MajOp	Major opcode
MinOp	Minor opcode
Rs	No Rs read
C	Conditional
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s

## Transfer immediate

Assign an immediate value to a 32-bit destination register.

Two types of assignment are supported. The first sign-extends a 16-bit signed immediate value to 32 bits. The second assigns a 16-bit unsigned immediate value to either the upper or lower 16 bits of the destination register, leaving the other 16 bits unchanged.



Syntax	Behavior
Rd=#s16	apply_extension(#s); Rd=#s;
Rdd=#s8	if ("#s8<0") { Assembler mapped to: "Rdd=combine(#-1,#s8)"; } else { Assembler mapped to: "Rdd=combine(#0,#s8)"; }
Rx.[HL]=#u16	Rx.h[01]=#u;

### Class: ALU32 (slots 0,1,2,3)

#### Intrinsics

Rd=#s16	Word32 Q6_R_equals_I(Word32 Is16)
Rdd=#s8	Word64 Q6_P_equals_I(Word32 Is8)
Rx.H=#u16	Word32 Q6_Rh_equals_I(Word32 Rx, Word32 Iu16)
Rx.L=#u16	Word32 Q6_Rl_equals_I(Word32 Rx, Word32 Iu16)

## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				Rs	MajOp			MinOp			x5					Parse																
0	1	1	1	0	0	0	1	i	i	1	x	x	x	x	x	P	P	i	i	i	i	i	i	i	i	i	i	i	i	i	i	Rx.L=#u16
0	1	1	1	0	0	1	0	i	i	1	x	x	x	x	x	P	P	i	i	i	i	i	i	i	i	i	i	i	i	i	i	Rx.H=#u16
ICLASS				Rs	MajOp			MinOp			Parse												d5									
0	1	1	1	1	0	0	0	i	i	-	i	i	i	i	i	P	P	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=#s16

Field name	Description
MajOp	Major opcode
MinOp	Minor opcode
Rs	No Rs read
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
x5	Field to encode register x

## Transfer register

Transfer a source register to a destination register. Source and destination registers are either 32 bits or 64 bits.

Syntax	Behavior
Rd=Rs	Rd=Rs;
Rdd=Rss	Assembler mapped to: "Rdd=combine(Rss.H32,Rss.L32)"

**Class: ALU32 (slots 0,1,2,3)**

### Intrinsics

Rd=Rs `Word32 Q6_R_equals_R(Word32 Rs)`

Rdd=Rss `Word64 Q6_P_equals_P(Word64 Rss)`

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				Rs	MajOp			MinOp			s5					Parse	C	d5														
0	1	1	1	0	0	0	0	0	1	1	s	s	s	s	s	P	P	0	-	-	-	-	-	-	-	-	d	d	d	d	d	Rd=Rs

Field name	Description
MajOp	Major opcode
MinOp	Minor opcode
Rs	No Rs read
C	Conditional
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s

## Vector add halfwords

Add the two 16-bit halfwords of Rs to the two 16-bit halfwords of Rt. The results are optionally saturated to signed or unsigned 16-bit values.

Syntax	Behavior
<code>Rd=vaddh(Rs,Rt)[:sat]</code>	<pre>for (i=0;i&lt;2;i++) {     Rd.h[i]=[sat<sub>16</sub>](Rs.h[i]+Rt.h[i]); }</pre>
<code>Rd=vadduh(Rs,Rt):sat</code>	<pre>for (i=0;i&lt;2;i++) {     Rd.h[i]=usat<sub>16</sub>(Rs.uh[i]+Rt.uh[i]); }</pre>

**Class: ALU32 (slots 0,1,2,3)**

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the Status Register is set. OVF remains set until explicitly cleared by a transfer to the status register.

### Intrinsics

<code>Rd=vaddh(Rs,Rt)</code>	Word32 Q6_R_vaddh_RR(Word32 Rs, Word32 Rt)
<code>Rd=vaddh(Rs,Rt):sat</code>	Word32 Q6_R_vaddh_RR_sat(Word32 Rs, Word32 Rt)
<code>Rd=vadduh(Rs,Rt):sat</code>	Word32 Q6_R_vadduh_RR_sat(Word32 Rs, Word32 Rt)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		P	MajOp		MinOp			s5					Parse		t5					d5												
1	1	1	1	0	1	1	0	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=vaddh(Rs,Rt)
1	1	1	1	0	1	1	0	0	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=vaddh(Rs,Rt):sat
1	1	1	1	0	1	1	0	0	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=vadduh(Rs,Rt):sat

Field name	Description
MajOp	Major opcode
MinOp	Minor opcode
P	Predicated
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t



## Vector average halfwords

The `vavgh` instruction adds the two 16-bit halfwords of `Rs` to the two 16-bit halfwords of `Rd`, and shifts the result right by 1 bit. Optionally, a rounding constant is added before shifting.

The `vnavgh` instruction subtracts the two 16-bit halfwords of `Rt` from the two 16-bit halfwords of `Rs`, and shifts the result right by 1 bit. For vector negative average with rounding, see the XTYPE `VNAVGH` instruction ([Vector average halfwords](#)).

Syntax	Behavior
<code>Rd=vavgh(Rs,Rt)</code>	<pre>for (i=0;i&lt;2;i++) {     Rd.h[i]=((Rs.h[i]+Rt.h[i])&gt;&gt;1); }</pre>
<code>Rd=vavgh(Rs,Rt):rnd</code>	<pre>for (i=0;i&lt;2;i++) {     Rd.h[i]=((Rs.h[i]+Rt.h[i]+1)&gt;&gt;1); }</pre>
<code>Rd=vnavgh(Rt,Rs)</code>	<pre>for (i=0;i&lt;2;i++) {     Rd.h[i]=((Rt.h[i]-Rs.h[i])&gt;&gt;1); }</pre>

**Class: ALU32 (slots 0,1,2,3)**

### Intrinsics

<code>Rd=vavgh(Rs,Rt)</code>	<code>Word32 Q6_R_vavgh_RR(Word32 Rs, Word32 Rt)</code>
<code>Rd=vavgh(Rs,Rt):rnd</code>	<code>Word32 Q6_R_vavgh_RR_rnd(Word32 Rs, Word32 Rt)</code>
<code>Rd=vnavgh(Rt,Rs)</code>	<code>Word32 Q6_R_vnavgh_RR(Word32 Rt, Word32 Rs)</code>

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				P	MajOp			MinOp		s5					Parse		t5					d5										
1	1	1	1	0	1	1	1	-	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=vavgh(Rs,Rt)
1	1	1	1	0	1	1	1	-	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=vavgh(Rs,Rt):rnd
1	1	1	1	0	1	1	1	-	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=vnavgh(Rt,Rs)

Field name	Description
MajOp	Major opcode
MinOp	Minor opcode
P	Predicated
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector subtract halfwords

Subtract each of the two halfwords in 32-bit vector Rs from the corresponding halfword in vector Rt. Optionally, saturate each 16-bit addition to either a signed or unsigned 16-bit value.

Applying saturation to the vsubh instruction clamps the result to the signed range 0x8000 to 0x7fff, whereas applying saturation to vsubuh ensures that the unsigned result is in the range 0 to 0xffff. When saturation is not needed, use vsubh.

Syntax	Behavior
<code>Rd=vsubh(Rt,Rs)[:sat]</code>	<pre>for (i=0;i&lt;2;i++) {   Rd.h[i]=[sat<sub>16</sub>](Rt.h[i]-Rs.h[i]); }</pre>
<code>Rd=vsubuh(Rt,Rs):sat</code>	<pre>for (i=0;i&lt;2;i++) {   Rd.h[i]=usat<sub>16</sub>(Rt.uh[i]-Rs.uh[i]); }</pre>

**Class: ALU32 (slots 0,1,2,3)**

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

### Intrinsics

<code>Rd=vsubh(Rt,Rs)</code>	<code>Word32 Q6_R_vsubh_RR(Word32 Rt, Word32 Rs)</code>
<code>Rd=vsubh(Rt,Rs):sat</code>	<code>Word32 Q6_R_vsubh_RR_sat(Word32 Rt, Word32 Rs)</code>
<code>Rd=vsubuh(Rt,Rs):sat</code>	<code>Word32 Q6_R_vsubuh_RR_sat(Word32 Rt, Word32 Rs)</code>

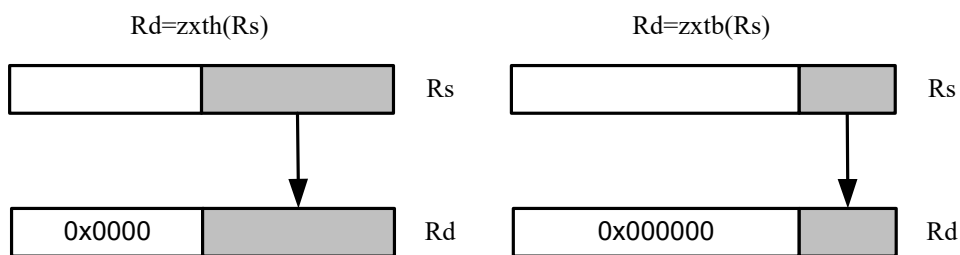
### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				P	MajOp			MinOp			s5					Parse		t5					d5									
1	1	1	1	0	1	1	0	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=vsubh(Rt,Rs)
1	1	1	1	0	1	1	0	1	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=vsubh(Rt,Rs):sat
1	1	1	1	0	1	1	0	1	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=vsubuh(Rt,Rs):sat

Field name	Description
MajOp	Major opcode
MinOp	Minor opcode
P	Predicated
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Zero extend

Zero-extend the least significant byte or halfword from Rs and place the 32-bit result in Rd.



### Syntax

Rd=zxtb(Rs)

Rd=zxtb(Rs)

### Behavior

Assembler mapped to: "Rd=and(Rs, #255)"

Rd = zxt<sub>16->32</sub>(Rs);

## Class: ALU32 (slots 0,1,2,3)

### Intrinsics

Rd=zxtb(Rs)

Word32 Q6\_R\_zxtb\_R(Word32 Rs)

Rd=zxtb(Rs)

Word32 Q6\_R\_zxtb\_R(Word32 Rs)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0							
ICLASS			Rs	MajOp		MinOp		s5					Parse		C											d5												
0	1	1	1	0	0	0	0	1	1	0	s	s	s	s	s	P	P	0	-	-	-	-	-	-	-	-	-	-	-	-	-	-	d	d	d	d	d	Rd=zxtb(Rs)

Field name	Description
MajOp	Major opcode
MinOp	Minor opcode
Rs	No Rs read
C	Conditional
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s

## 11.1.2 ALU32 PERM

The ALU32 PERM instruction subclass includes instructions that rearrange or perform format conversion on vector data types.

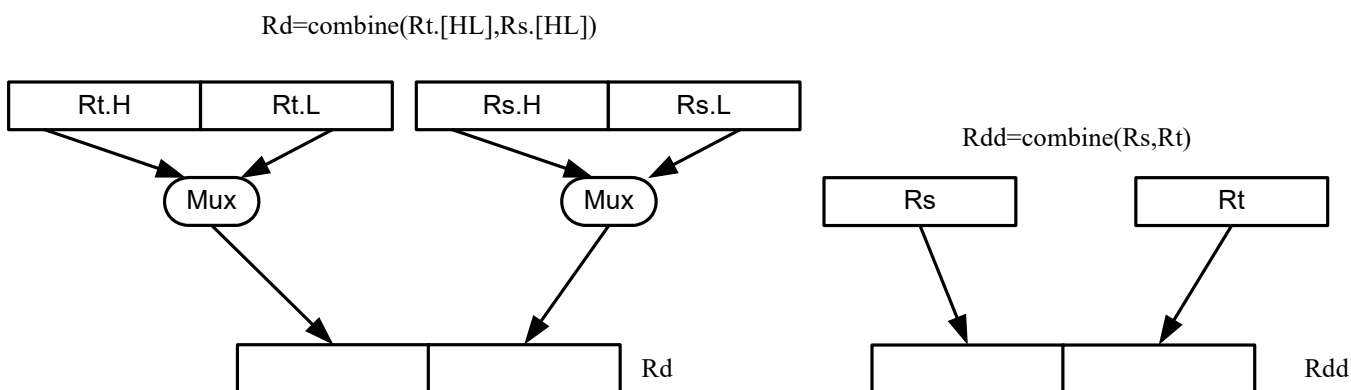
### Combine words into doubleword

Combine halfwords or words into larger values.

In a halfword combine, either the high or low halfword of the first source register is transferred to the most-significant halfword of the destination register, while either the high or low halfword of the second source register is transferred to the least-significant halfword of the destination register. Source and destination registers are 32 bits.

In a word combine, the first source register is transferred to the most-significant word of the destination register, while the second source register is transferred to the least-significant word of the destination register. Source registers are 32 bits and the destination register is 64 bits.

In a variant of word combine, signed 8-bit immediate values (instead of registers) are transferred to the most- and least-significant words of the 64-bit destination register. Optionally, one of the immediate values can be 32 bits.



#### Syntax

#### Behavior

$Rd = \text{combine}(Rt.[HL], Rs.[HL])$	$Rd = (Rt.uh[01] \ll 16)   Rs.uh[01];$
$Rdd = \text{combine}(\#s8, \#s8)$	<pre> apply_extension(#s); Rdd.w[0]=#s; Rdd.w[1]=#s; </pre>
$Rdd = \text{combine}(\#s8, \#U6)$	<pre> apply_extension(#U); Rdd.w[0]=#U; Rdd.w[1]=#s; </pre>
$Rdd = \text{combine}(\#s8, Rs)$	<pre> apply_extension(#s); Rdd.w[0]=Rs; Rdd.w[1]=#s; </pre>
$Rdd = \text{combine}(Rs, \#s8)$	<pre> apply_extension(#s); Rdd.w[0]=#s; Rdd.w[1]=Rs; </pre>

**Syntax**

Rdd=combine (Rs, Rt)

**Behavior**

Rdd.w[0]=Rt;  
Rdd.w[1]=Rs;

**Class: ALU32 (slots 0,1,2,3)**

**Intrinsics**

Rd=combine (Rt.H, Rs.H)	Word32 Q6_R_combine_RhRh (Word32 Rt, Word32 Rs)
Rd=combine (Rt.H, Rs.L)	Word32 Q6_R_combine_RhRl (Word32 Rt, Word32 Rs)
Rd=combine (Rt.L, Rs.H)	Word32 Q6_R_combine_RlRh (Word32 Rt, Word32 Rs)
Rd=combine (Rt.L, Rs.L)	Word32 Q6_R_combine_RlRl (Word32 Rt, Word32 Rs)
Rdd=combine (#s8, #S8)	Word64 Q6_P_combine_II (Word32 Is8, Word32 IS8)
Rdd=combine (#s8, Rs)	Word64 Q6_P_combine_IR (Word32 Is8, Word32 Rs)
Rdd=combine (Rs, #s8)	Word64 Q6_P_combine_RI (Word32 Rs, Word32 Is8)
Rdd=combine (Rs, Rt)	Word64 Q6_P_combine_RR (Word32 Rs, Word32 Rt)

**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		Rs		MajOp		MinOp		s5					Parse					d5														
0	1	1	1	0	0	1	1	-	0	0	s	s	s	s	s	P	P	1	i	i	i	i	i	i	i	i	d	d	d	d	Rdd=combine(Rs,#s8)	
0	1	1	1	0	0	1	1	-	0	1	s	s	s	s	s	P	P	1	i	i	i	i	i	i	i	i	d	d	d	d	Rdd=combine(#s8,Rs)	
ICLASS		Rs		MajOp		MinOp		s5					Parse					d5														
0	1	1	1	1	1	0	0	0	I	I	I	I	I	I	I	P	P	I	i	i	i	i	i	i	i	i	d	d	d	d	Rdd=combine(#s8,#S8)	
0	1	1	1	1	1	0	0	1	-	-	I	I	I	I	I	P	P	I	i	i	i	i	i	i	i	i	d	d	d	d	Rdd=combine(#s8,#U6)	
ICLASS		P		MajOp		MinOp		s5					Parse					t5			d5											
1	1	1	1	0	0	1	1	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	Rd=combine(Rt.H,Rs.H)	
1	1	1	1	0	0	1	1	1	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	Rd=combine(Rt.H,Rs.L)	
1	1	1	1	0	0	1	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	Rd=combine(Rt.L,Rs.H)	
1	1	1	1	0	0	1	1	1	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	Rd=combine(Rt.L,Rs.L)	
1	1	1	1	0	1	0	1	0	-	-	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	Rdd=combine(Rs,Rt)	

Field name	Description
MajOp	Major opcode
MinOp	Minor opcode
Rs	No Rs read
P	Predicated
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Mux

Select between two source registers based on the least-significant bit of a predicate register. If the bit is 1, transfer the first source register to the destination register; otherwise, transfer the second source register. Source and destination registers are 32 bits.

In a variant of mux, signed 8-bit immediate values are used instead of registers for either or both source operands.

For 64-bit versions of this instruction, see the XTYPE vmux ([Vector mux](#)) instruction.

Syntax	Behavior
Rd=mux (Pu, #s8, #S8)	PREDUSE_TIMING; apply_extension(#s); Rd = (Pu[0] ? #s : #S);
Rd=mux (Pu, #s8, Rs)	PREDUSE_TIMING; apply_extension(#s); Rd = (Pu[0] ? #s : Rs);
Rd=mux (Pu, Rs, #s8)	PREDUSE_TIMING; apply_extension(#s); Rd = (Pu[0] ? Rs : #s);
Rd=mux (Pu, Rs, Rt)	PREDUSE_TIMING; Rd = (Pu[0] ? Rs : Rt);

### Class: ALU32 (slots 0,1,2,3)

#### Intrinsics

Rd=mux (Pu, #s8, #S8)	Word32 Q6_R_mux_pII(Byte Pu, Word32 Is8, Word32 IS8)
Rd=mux (Pu, #s8, Rs)	Word32 Q6_R_mux_pIR(Byte Pu, Word32 Is8, Word32 Rs)
Rd=mux (Pu, Rs, #s8)	Word32 Q6_R_mux_pRI(Byte Pu, Word32 Rs, Word32 Is8)
Rd=mux (Pu, Rs, Rt)	Word32 Q6_R_mux_pRR(Byte Pu, Word32 Rs, Word32 Rt)

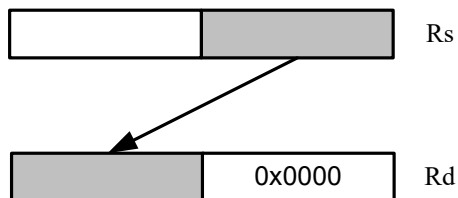
#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				Rs	MajOp				u2			s5					Parse					d5										
0	1	1	1	0	0	1	1	0	u	u	s	s	s	s	s	P	P	0	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=mux(Pu,Rs,#s8)
0	1	1	1	0	0	1	1	1	u	u	s	s	s	s	s	P	P	0	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=mux(Pu,#s8,Rs)
ICLASS				Rs		u1			Parse					d5																		
0	1	1	1	1	0	1	u	u	l	l	l	l	l	l	l	P	P	l	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=mux(Pu,#s8,#S8)
ICLASS				P	MajOp				s5					Parse					t5			u2		d5								
1	1	1	1	0	1	0	0	-	-	-	s	s	s	s	s	P	P	-	t	t	t	t	t	-	u	u	d	d	d	d	d	Rd=mux(Pu,Rs,Rt)

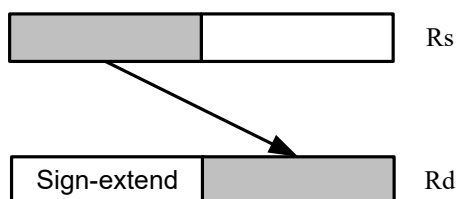
<b>Field name</b>	<b>Description</b>
MajOp	Major opcode
MinOp	Minor opcode
Rs	No Rs read
P	Predicated
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
u1	Field to encode register u
u2	Field to encode register u

## Shift word by 16

ASLH performs an arithmetic left shift of the 32-bit source register by 16 bits (one halfword). The lower 16 bits of the destination are zero-filled.



ASRH performs an arithmetic right shift of the 32-bit source register by 16 bits (one halfword). The upper 16 bits of the destination are sign-extended.



### Syntax

$Rd=aslh(Rs)$

$Rd=asrh(Rs)$

### Behavior

$Rd=R_s \ll 16;$

$Rd=R_s \gg 16;$

## Class: ALU32 (slots 0,1,2,3)

### Intrinsics

$Rd=aslh(Rs)$

Word32 Q6\_R\_aslh\_R(Word32 Rs)

$Rd=asrh(Rs)$

Word32 Q6\_R\_asrh\_R(Word32 Rs)

### Encoding

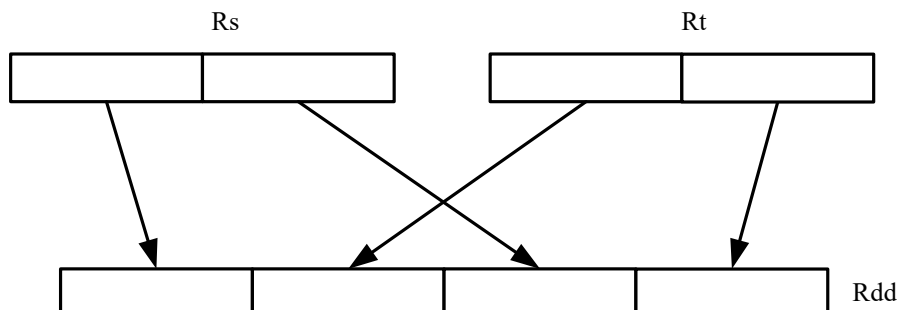
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				Rs	MajOp			MinOp			s5					Parse	C	d5														
0	1	1	1	0	0	0	0	0	0	0	s	s	s	s	s	P	P	0	-	-	-	-	-	-	-	-	d	d	d	d	d	Rd=aslh(Rs)
0	1	1	1	0	0	0	0	0	0	1	s	s	s	s	s	P	P	0	-	-	-	-	-	-	-	-	d	d	d	d	d	Rd=asrh(Rs)

Field name	Description
MajOp	Major opcode
MinOp	Minor opcode
Rs	No Rs read
C	Conditional
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s



## Pack high and low halfwords

Pack together the most-significant halfwords from Rs and Rt into the most-significant word of register pair Rdd, and the least-significant halfwords from Rs and Rt into the least-significant halfword of Rdd.



### Syntax

```
Rdd=packhl (Rs, Rt)
```

### Behavior

```
Rdd.h[0]=Rt.h[0];
Rdd.h[1]=Rs.h[0];
Rdd.h[2]=Rt.h[1];
Rdd.h[3]=Rs.h[1];
```

**Class: ALU32 (slots 0,1,2,3)**

### Intrinsics

```
Rdd=packhl (Rs, Rt)
```

```
Word64 Q6_P_packhl_RR(Word32 Rs, Word32 Rt)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				P	MajOp		MinOp		s5					Parse		t5					d5											
1	1	1	1	0	1	0	1	1	-	-	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rdd=packhl(Rs,Rt)

Field name	Description
MajOp	Major opcode
MinOp	Minor opcode
P	Predicated
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

### 11.1.3 ALU32 PRED

The ALU32 PRED instruction subclass includes instructions that perform conditional arithmetic and logical operations based on the values stored in a predicate register, and which produce predicate results. They are executable on any slot.

#### Conditional add

If the least-significant bit of predicate Pu is set, add a 32-bit source register to either another register or an immediate value. The result is placed in 32-bit destination register. If the predicate is false, the instruction does nothing.

Syntax	Behavior
<pre>if ([!]Pu[.new]) Rd=add(Rs, #s8)</pre>	<pre>if ([!]Pu[.new][0]) {     apply_extension(#s);     Rd=Rs+#s; } else {     NOP; }</pre>
<pre>if ([!]Pu[.new]) Rd=add(Rs, Rt)</pre>	<pre>if ([!]Pu[.new][0]) {     Rd=Rs+Rt; } else {     NOP; }</pre>

**Class: ALU32 (slots 0,1,2,3)**

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS		Rs		MajOp		PS		u2		s5					Parse		D N											d5					
0	1	1	1	0	1	0	0	0	u	u	s	s	s	s	s	P	P	0	i	i	i	i	i	i	i	i	d	d	d	d	d	if (Pu) Rd=add(Rs,#s8)	
0	1	1	1	0	1	0	0	0	u	u	s	s	s	s	s	P	P	1	i	i	i	i	i	i	i	i	d	d	d	d	d	if (Pu.new) Rd=add(Rs,#s8)	
0	1	1	1	0	1	0	0	1	u	u	s	s	s	s	s	P	P	0	i	i	i	i	i	i	i	i	d	d	d	d	d	if (!Pu) Rd=add(Rs,#s8)	
0	1	1	1	0	1	0	0	1	u	u	s	s	s	s	s	P	P	1	i	i	i	i	i	i	i	i	d	d	d	d	d	if (!Pu.new) Rd=add(Rs,#s8)	
ICLASS		P	MajOp		MinOp		s5					Parse		D N	t5					PS	u2		d5										
1	1	1	1	1	0	1	1	0	-	0	s	s	s	s	s	P	P	0	t	t	t	t	t	t	0	u	u	d	d	d	d	d	if (Pu) Rd=add(Rs,Rt)
1	1	1	1	1	0	1	1	0	-	0	s	s	s	s	s	P	P	0	t	t	t	t	t	t	1	u	u	d	d	d	d	d	if (!Pu) Rd=add(Rs,Rt)
1	1	1	1	1	0	1	1	0	-	0	s	s	s	s	s	P	P	1	t	t	t	t	t	t	0	u	u	d	d	d	d	d	if (Pu.new) Rd=add(Rs,Rt)
1	1	1	1	1	0	1	1	0	-	0	s	s	s	s	s	P	P	1	t	t	t	t	t	t	1	u	u	d	d	d	d	d	if (!Pu.new) Rd=add(Rs,Rt)

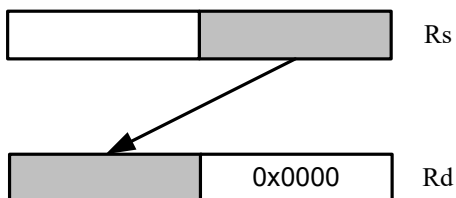
Field name	Description
MajOp	Major opcode
MinOp	Minor opcode
Rs	No Rs read
DN	Dot-new
PS	Predicate sense
P	Predicated

<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
u2	Field to encode register u

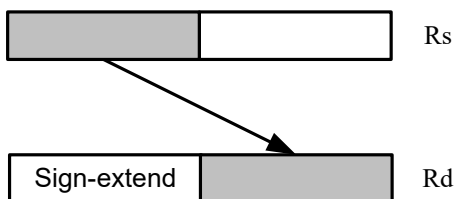
## Conditional shift halfword

Conditionally shift a halfword.

The aslh instruction performs an arithmetic left shift of the 32-bit source register by 16 bits (one halfword). The lower 16 bits of the destination are zero-filled.



The asrh instruction performs an arithmetic right shift of the 32-bit source register by 16 bits (one halfword). The upper 16 bits of the destination are sign-extended.



### Syntax

```
if ([!]Pu[.new]) Rd=aslh(Rs)
```

```
if ([!]Pu[.new]) Rd=asrh(Rs)
```

### Behavior

```
if ([!]Pu[.new][0]) {
    Rd=Rs<<16;
} else {
    NOP;
}
```

```
if ([!]Pu[.new][0]) {
    Rd=Rs>>16;
} else {
    NOP;
}
```

**Class: ALU32 (slots 0,1,2,3)**

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		Rs	MajOp			MinOp			s5					Parse	C	S	dn	u2	d5													
0	1	1	1	0	0	0	0	0	0	0	s	s	s	s	s	P	P	1	-	0	0	u	u	-	-	-	d	d	d	d	d	if (Pu) Rd=aslh(Rs)
0	1	1	1	0	0	0	0	0	0	0	s	s	s	s	s	P	P	1	-	0	1	u	u	-	-	-	d	d	d	d	d	if (Pu.new) Rd=aslh(Rs)
0	1	1	1	0	0	0	0	0	0	0	s	s	s	s	s	P	P	1	-	1	0	u	u	-	-	-	d	d	d	d	d	if (!Pu) Rd=aslh(Rs)
0	1	1	1	0	0	0	0	0	0	0	s	s	s	s	s	P	P	1	-	1	1	u	u	-	-	-	d	d	d	d	d	if (!Pu.new) Rd=aslh(Rs)
0	1	1	1	0	0	0	0	0	0	1	s	s	s	s	s	P	P	1	-	0	0	u	u	-	-	-	d	d	d	d	d	if (Pu) Rd=asrh(Rs)
0	1	1	1	0	0	0	0	0	0	1	s	s	s	s	s	P	P	1	-	0	1	u	u	-	-	-	d	d	d	d	d	if (Pu.new) Rd=asrh(Rs)
0	1	1	1	0	0	0	0	0	0	1	s	s	s	s	s	P	P	1	-	1	0	u	u	-	-	-	d	d	d	d	d	if (!Pu) Rd=asrh(Rs)
0	1	1	1	0	0	0	0	0	0	1	s	s	s	s	s	P	P	1	-	1	1	u	u	-	-	-	d	d	d	d	d	if (!Pu.new) Rd=asrh(Rs)

<b>Field name</b>	<b>Description</b>
MajOp	Major opcode
MinOp	Minor opcode
Rs	No Rs read
C	Conditional
S	Predicate sense
dn	Dot-new
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
u2	Field to encode register u

## Conditional combine

If the least-significant bit of predicate Pu is set, the most-significant word of destination Rdd is taken from the first source register Rs, while the least-significant word is taken from the second source register Rt. If the predicate is false, this instruction does nothing.

### Syntax

```
if ([!]Pu[.new])
Rdd=combine(Rs,Rt)
```

### Behavior

```
if ([!]Pu[.new][0]) {
  Rdd.w[0]=Rt;
  Rdd.w[1]=Rs;
} else {
  NOP;
}
```

**Class: ALU32 (slots 0,1,2,3)**

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				P	MajOp			MinOp			s5					Parse		<sup>D</sup> <sub>N</sub>	t5					PS	u2		d5					
1	1	1	1	1	1	0	1	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	u	u	d	d	d	d	d	if (Pu) Rdd=combine(Rs,Rt)
1	1	1	1	1	1	0	1	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	u	u	d	d	d	d	d	if (!Pu) Rdd=combine(Rs,Rt)
1	1	1	1	1	1	0	1	0	0	0	s	s	s	s	s	P	P	1	t	t	t	t	t	0	u	u	d	d	d	d	d	if (Pu.new) Rdd=combine(Rs,Rt)
1	1	1	1	1	1	0	1	0	0	0	s	s	s	s	s	P	P	1	t	t	t	t	t	1	u	u	d	d	d	d	d	if (!Pu.new) Rdd=combine(Rs,Rt)

Field name	Description
DN	Dot-new
MajOp	Major opcode
MinOp	Minor opcode
P	Predicated
PS	Predicate sense
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
u2	Field to encode register u

## Conditional logical operations

If the least-significant bit of predicate Pu is set, do a logical operation on the source values. The result is placed in 32-bit destination register. If the predicate is false, the instruction does nothing.

Syntax	Behavior
<pre>if ([!]Pu[.new]) Rd=and(Rs,Rt)</pre>	<pre>if ([!]Pu[.new][0]) {     Rd=Rs&amp;Rt; } else {     NOP; }</pre>
<pre>if ([!]Pu[.new]) Rd=or(Rs,Rt)</pre>	<pre>if ([!]Pu[.new][0]) {     Rd=Rs Rt; } else {     NOP; }</pre>
<pre>if ([!]Pu[.new]) Rd=xor(Rs,Rt)</pre>	<pre>if ([!]Pu[.new][0]) {     Rd=Rs^Rt; } else {     NOP; }</pre>

**Class: ALU32 (slots 0,1,2,3)**

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		P	MajOp		MinOp		s5					Parse	DN	t5				PS	u2	d5												
1	1	1	1	1	0	0	1	-	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	u	u	d	d	d	d	d	if (Pu) Rd=and(Rs,Rt)
1	1	1	1	1	0	0	1	-	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	u	u	d	d	d	d	d	if (!Pu) Rd=and(Rs,Rt)
1	1	1	1	1	0	0	1	-	0	0	s	s	s	s	s	P	P	1	t	t	t	t	t	0	u	u	d	d	d	d	d	if (Pu.new) Rd=and(Rs,Rt)
1	1	1	1	1	0	0	1	-	0	0	s	s	s	s	s	P	P	1	t	t	t	t	t	1	u	u	d	d	d	d	d	if (!Pu.new) Rd=and(Rs,Rt)
1	1	1	1	1	0	0	1	-	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	u	u	d	d	d	d	d	if (Pu) Rd=or(Rs,Rt)
1	1	1	1	1	0	0	1	-	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	u	u	d	d	d	d	d	if (!Pu) Rd=or(Rs,Rt)
1	1	1	1	1	0	0	1	-	0	1	s	s	s	s	s	P	P	1	t	t	t	t	t	0	u	u	d	d	d	d	d	if (Pu.new) Rd=or(Rs,Rt)
1	1	1	1	1	0	0	1	-	0	1	s	s	s	s	s	P	P	1	t	t	t	t	t	1	u	u	d	d	d	d	d	if (!Pu.new) Rd=or(Rs,Rt)
1	1	1	1	1	0	0	1	-	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	u	u	d	d	d	d	d	if (Pu) Rd=xor(Rs,Rt)
1	1	1	1	1	0	0	1	-	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	u	u	d	d	d	d	d	if (!Pu) Rd=xor(Rs,Rt)
1	1	1	1	1	0	0	1	-	1	1	s	s	s	s	s	P	P	1	t	t	t	t	t	0	u	u	d	d	d	d	d	if (Pu.new) Rd=xor(Rs,Rt)
1	1	1	1	1	0	0	1	-	1	1	s	s	s	s	s	P	P	1	t	t	t	t	t	1	u	u	d	d	d	d	d	if (!Pu.new) Rd=xor(Rs,Rt)

Field name	Description
DN	Dot-new
MajOp	Major opcode
MinOp	Minor opcode
P	Predicated
PS	Predicate sense
ICLASS	Instruction class
Parse	Packet/loop parse bits

<b>Field name</b>	<b>Description</b>
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
u2	Field to encode register u



## Conditional subtract

If the least-significant bit of predicate Pu is set, subtract a 32-bit source register Rt from register Rs. The result is placed in a 32-bit destination register. If the predicate is false, the instruction does nothing.

### Syntax

```
if ([!]Pu[.new])
Rd=sub(Rt, Rs)
```

### Behavior

```
if ([!]Pu[.new][0]) {
    Rd=Rt-Rs;
} else {
    NOP;
}
```

**Class: ALU32 (slots 0,1,2,3)**

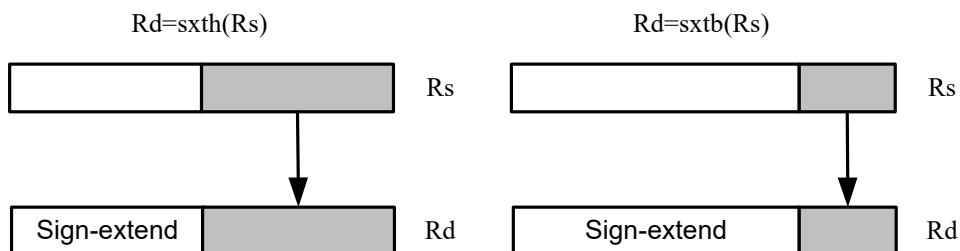
### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS				P	MajOp		MinOp		s5					Parse		DN	t5					PS	u2		d5								
1	1	1	1	1	0	1	1	0	-	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	u	u	d	d	d	d	d	d	if (Pu) Rd=sub(Rt,Rs)
1	1	1	1	1	0	1	1	0	-	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	u	u	d	d	d	d	d	d	if (!Pu) Rd=sub(Rt,Rs)
1	1	1	1	1	0	1	1	0	-	1	s	s	s	s	s	P	P	1	t	t	t	t	t	0	u	u	d	d	d	d	d	d	if (Pu.new) Rd=sub(Rt,Rs)
1	1	1	1	1	0	1	1	0	-	1	s	s	s	s	s	P	P	1	t	t	t	t	t	1	u	u	d	d	d	d	d	d	if (!Pu.new) Rd=sub(Rt,Rs)

Field name	Description
DN	Dot-new
MajOp	Major opcode
MinOp	Minor opcode
P	Predicated
PS	Predicate sense
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
u2	Field to encode register u

## Conditional sign extend

Conditionally sign-extend the least-significant byte or halfword from Rs and put the 32-bit result in Rd.



### Syntax

```
if ([!]Pu[.new]) Rd=sxtb(Rs)

if ([!]Pu[.new]) Rd=sxth(Rs)
```

### Behavior

```
if ([!]Pu[.new][0]) {
    Rd=sxt8->32(Rs);
} else {
    NOP;
}

if ([!]Pu[.new][0]) {
    Rd=sxt16->32(Rs);
} else {
    NOP;
}
```

**Class: ALU32 (slots 0,1,2,3)**

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				Rs	MajOp				MinOp				s5				Parse	C	S	dn	u2	d5										
0	1	1	1	0	0	0	0	1	0	1	s	s	s	s	s	P	P	1	-	0	0	u	u	-	-	-	d	d	d	d	d	if (Pu) Rd=sxtb(Rs)
0	1	1	1	0	0	0	0	1	0	1	s	s	s	s	s	P	P	1	-	0	1	u	u	-	-	-	d	d	d	d	d	if (Pu.new) Rd=sxtb(Rs)
0	1	1	1	0	0	0	0	1	0	1	s	s	s	s	s	P	P	1	-	1	0	u	u	-	-	-	d	d	d	d	d	if (!Pu) Rd=sxtb(Rs)
0	1	1	1	0	0	0	0	1	0	1	s	s	s	s	s	P	P	1	-	1	1	u	u	-	-	-	d	d	d	d	d	if (!Pu.new) Rd=sxtb(Rs)
0	1	1	1	0	0	0	0	1	1	1	s	s	s	s	s	P	P	1	-	0	0	u	u	-	-	-	d	d	d	d	d	if (Pu) Rd=sxth(Rs)
0	1	1	1	0	0	0	0	1	1	1	s	s	s	s	s	P	P	1	-	0	1	u	u	-	-	-	d	d	d	d	d	if (Pu.new) Rd=sxth(Rs)
0	1	1	1	0	0	0	0	1	1	1	s	s	s	s	s	P	P	1	-	1	0	u	u	-	-	-	d	d	d	d	d	if (!Pu) Rd=sxth(Rs)
0	1	1	1	0	0	0	0	1	1	1	s	s	s	s	s	P	P	1	-	1	1	u	u	-	-	-	d	d	d	d	d	if (!Pu.new) Rd=sxth(Rs)

Field name	Description
MajOp	Major opcode
MinOp	Minor opcode
Rs	No Rs read
C	Conditional
S	Predicate sense
dn	Dot-new
ICLASS	Instruction class

<b>Field name</b>	<b>Description</b>
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
u2	Field to encode register u

## Conditional transfer

If the LSB of predicate Pu is set, transfer register Rs or a signed immediate into destination Rd. If the predicate is false, this instruction does nothing.

Syntax	Behavior
if ([!]Pu[.new]) Rd=#s12	apply_extension(#s); if ([!]Pu[.new][0]) Rd=#s; else NOP;
if ([!]Pu[.new]) Rd=Rs	Assembler mapped to: "if ([!]Pu[.new]) Rd=add(Rs,#0) "
if ([!]Pu[.new]) Rdd=Rss	Assembler mapped to: "if ([!]Pu[.new]) Rdd=combine(Rss.H32,Rss.L32) "

### Class: ALU32 (slots 0,1,2,3)

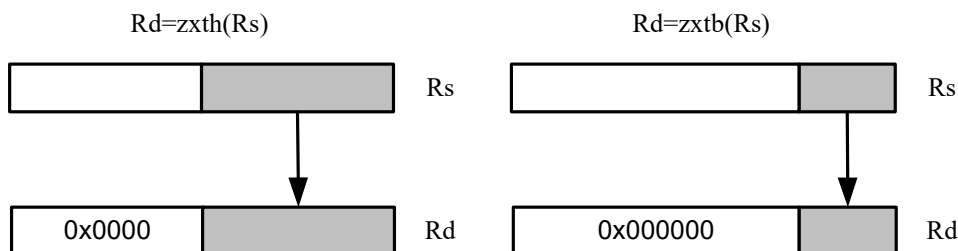
#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				Rs	MajOp		PS	u2						Parse	D N					d5												
0	1	1	1	1	1	1	0	0	u	u	0	i	i	i	i	P	P	0	i	i	i	i	i	i	i	i	d	d	d	d	d	if (Pu) Rd=#s12
0	1	1	1	1	1	1	0	0	u	u	0	i	i	i	i	P	P	1	i	i	i	i	i	i	i	i	d	d	d	d	d	if (Pu.new) Rd=#s12
0	1	1	1	1	1	1	0	1	u	u	0	i	i	i	i	P	P	0	i	i	i	i	i	i	i	i	d	d	d	d	d	if (!Pu) Rd=#s12
0	1	1	1	1	1	1	0	1	u	u	0	i	i	i	i	P	P	1	i	i	i	i	i	i	i	i	d	d	d	d	d	if (!Pu.new) Rd=#s12

Field name	Description
MajOp	Major opcode
MinOp	Minor opcode
Rs	No Rs read
DN	Dot-new
PS	Predicate sense
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
u2	Field to encode register u

## Conditional zero extend

Conditionally zero-extend the least-significant byte or halfword from Rs and put the 32-bit result in Rd.



### Syntax

```
if ([!]Pu[.new]) Rd=zxth(Rs)
```

```
if ([!]Pu[.new]) Rd=zxth(Rs)
```

### Behavior

```
if ([!]Pu[.new][0]) {
    Rd=zxth8->32(Rs);
} else {
    NOP;
}
```

```
if ([!]Pu[.new][0]) {
    Rd=zxth16->32(Rs);
} else {
    NOP;
}
```

**Class: ALU32 (slots 0,1,2,3)**

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		Rs		MajOp		MinOp		s5					Parse		C	S	dn	u2		d5												
0	1	1	1	0	0	0	0	1	0	0	s	s	s	s	s	P	P	1	-	0	0	u	u	-	-	-	d	d	d	d	d	if (Pu) Rd=zxth(Rs)
0	1	1	1	0	0	0	0	1	0	0	s	s	s	s	s	P	P	1	-	0	1	u	u	-	-	-	d	d	d	d	d	if (Pu.new) Rd=zxth(Rs)
0	1	1	1	0	0	0	0	1	0	0	s	s	s	s	s	P	P	1	-	1	0	u	u	-	-	-	d	d	d	d	d	if (!Pu) Rd=zxth(Rs)
0	1	1	1	0	0	0	0	1	0	0	s	s	s	s	s	P	P	1	-	1	1	u	u	-	-	-	d	d	d	d	d	if (!Pu.new) Rd=zxth(Rs)
0	1	1	1	0	0	0	0	1	1	0	s	s	s	s	s	P	P	1	-	0	0	u	u	-	-	-	d	d	d	d	d	if (Pu) Rd=zxth(Rs)
0	1	1	1	0	0	0	0	1	1	0	s	s	s	s	s	P	P	1	-	0	1	u	u	-	-	-	d	d	d	d	d	if (Pu.new) Rd=zxth(Rs)
0	1	1	1	0	0	0	0	1	1	0	s	s	s	s	s	P	P	1	-	1	0	u	u	-	-	-	d	d	d	d	d	if (!Pu) Rd=zxth(Rs)
0	1	1	1	0	0	0	0	1	1	0	s	s	s	s	s	P	P	1	-	1	1	u	u	-	-	-	d	d	d	d	d	if (!Pu.new) Rd=zxth(Rs)

Field name	Description
MajOp	Major opcode
MinOp	Minor opcode
Rs	No Rs read
C	Conditional
S	Predicate sense
dn	Dot-new
ICLASS	Instruction class
Parse	Packet/loop parse bits

<b>Field name</b>	<b>Description</b>
d5	Field to encode register d
s5	Field to encode register s
u2	Field to encode register u

## Compare

The register form compares two 32-bit registers for unsigned greater than, greater than, or equal.

The immediate form compares a register against a signed or unsigned immediate value. The 8-bit predicate register Pd is set to all 1's or all 0's depending on the result. For 64-bit versions of this instruction, see the XTYPE compare instructions.

Syntax	Behavior
Pd=[!]cmp.eq(Rs, #s10)	apply_extension(#s); Pd=Rs[!]=#s ? 0xff : 0x00;
Pd=[!]cmp.eq(Rs, Rt)	Pd=Rs[!]=Rt ? 0xff : 0x00;
Pd=[!]cmp.gt(Rs, #s10)	apply_extension(#s); Pd=Rs<=#s ? 0xff : 0x00;
Pd=[!]cmp.gt(Rs, Rt)	Pd=Rs<=Rt ? 0xff : 0x00;
Pd=[!]cmp.gtu(Rs, #u9)	apply_extension(#u); Pd=Rs.uw[0]<=#u.uw[0] ? 0xff : 0x00;
Pd=[!]cmp.gtu(Rs, Rt)	Pd=Rs.uw[0]<=Rt.uw[0] ? 0xff : 0x00;
Pd=cmp.ge(Rs, #s8)	Assembler mapped to: "Pd=cmp.gt(Rs, #s8-1)"
Pd=cmp.geu(Rs, #u8)	if ("#u8==0") { Assembler mapped to: "Pd=cmp.eq(Rs, Rs)"; } else { Assembler mapped to: "Pd=cmp.gtu(Rs, #u8-1)"; }
Pd=cmp.lt(Rs, Rt)	Assembler mapped to: "Pd=cmp.gt(Rt, Rs)"
Pd=cmp.ltu(Rs, Rt)	Assembler mapped to: "Pd=cmp.gtu(Rt, Rs)"

### Class: ALU32 (slots 0,1,2,3)

#### Intrinsics

Pd=!cmp.eq(Rs, #s10)	Byte Q6_p_not_cmp_eq_RI(Word32 Rs, Word32 Is10)
Pd=!cmp.eq(Rs, Rt)	Byte Q6_p_not_cmp_eq_RR(Word32 Rs, Word32 Rt)
Pd=!cmp.gt(Rs, #s10)	Byte Q6_p_not_cmp_gt_RI(Word32 Rs, Word32 Is10)
Pd=!cmp.gt(Rs, Rt)	Byte Q6_p_not_cmp_gt_RR(Word32 Rs, Word32 Rt)
Pd=!cmp.gtu(Rs, #u9)	Byte Q6_p_not_cmp_gtu_RI(Word32 Rs, Word32 Iu9)
Pd=!cmp.gtu(Rs, Rt)	Byte Q6_p_not_cmp_gtu_RR(Word32 Rs, Word32 Rt)
Pd=cmp.eq(Rs, #s10)	Byte Q6_p_cmp_eq_RI(Word32 Rs, Word32 Is10)
Pd=cmp.eq(Rs, Rt)	Byte Q6_p_cmp_eq_RR(Word32 Rs, Word32 Rt)
Pd=cmp.ge(Rs, #s8)	Byte Q6_p_cmp_ge_RI(Word32 Rs, Word32 Is8)
Pd=cmp.geu(Rs, #u8)	Byte Q6_p_cmp_geu_RI(Word32 Rs, Word32 Iu8)
Pd=cmp.gt(Rs, #s10)	Byte Q6_p_cmp_gt_RI(Word32 Rs, Word32 Is10)

Pd=cmp.gt(Rs,Rt)	Byte Q6_p_cmp_gt_RR(Word32 Rs, Word32 Rt)
Pd=cmp.gtu(Rs,#u9)	Byte Q6_p_cmp_gtu_RI(Word32 Rs, Word32 Iu9)
Pd=cmp.gtu(Rs,Rt)	Byte Q6_p_cmp_gtu_RR(Word32 Rs, Word32 Rt)
Pd=cmp.lt(Rs,Rt)	Byte Q6_p_cmp_lt_RR(Word32 Rs, Word32 Rt)
Pd=cmp.ltu(Rs,Rt)	Byte Q6_p_cmp_ltu_RR(Word32 Rs, Word32 Rt)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS		Rs	MajOp		MinOp		s5					Parse												d2									
0	1	1	1	0	1	0	1	0	0	i	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	i	0	0	0	d	d	Pd=cmp.eq(Rs,#s10)
0	1	1	1	0	1	0	1	0	0	i	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	i	1	0	0	d	d	Pd=!cmp.eq(Rs,#s10)
0	1	1	1	0	1	0	1	0	1	i	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	i	0	0	0	d	d	Pd=cmp.gt(Rs,#s10)
0	1	1	1	0	1	0	1	0	1	i	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	i	1	0	0	d	d	Pd=!cmp.gt(Rs,#s10)
0	1	1	1	0	1	0	1	1	0	0	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	i	0	0	0	d	d	Pd=cmp.gtu(Rs,#u9)
0	1	1	1	0	1	0	1	1	0	0	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	i	1	0	0	d	d	Pd=!cmp.gtu(Rs,#u9)
ICLASS		P	MajOp		MinOp		s5					Parse		t5					d2														
1	1	1	1	0	0	1	0	-	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	0	0	0	d	d	Pd=cmp.eq(Rs,Rt)	
1	1	1	1	0	0	1	0	-	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	1	0	0	d	d	Pd=!cmp.eq(Rs,Rt)	
1	1	1	1	0	0	1	0	-	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	0	0	0	d	d	Pd=cmp.gt(Rs,Rt)	
1	1	1	1	0	0	1	0	-	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	1	0	0	d	d	Pd=!cmp.gt(Rs,Rt)	
1	1	1	1	0	0	1	0	-	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	0	0	0	d	d	Pd=cmp.gtu(Rs,Rt)	
1	1	1	1	0	0	1	0	-	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	1	0	0	d	d	Pd=!cmp.gtu(Rs,Rt)	

Field name	Description
MajOp	Major opcode
MinOp	Minor opcode
Rs	No Rs read
P	Predicated
ICLASS	Instruction class
Parse	Packet/loop parse bits
d2	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t



## Compare to general register

The register form compares two 32-bit registers for unsigned greater than, greater than, or equal. The immediate form compares a register against a signed or unsigned immediate value. The resulting zero or one is placed in a general register.

Syntax	Behavior
Rd=[!]cmp.eq(Rs,#s8)	apply_extension(#s); Rd=(Rs[!]=#s);
Rd=[!]cmp.eq(Rs,Rt)	Rd=(Rs[!]=Rt);

### Class: ALU32 (slots 0,1,2,3)

#### Intrinsics

Rd=!cmp.eq(Rs,#s8)	Word32 Q6_R_not_cmp_eq_RI(Word32 Rs, Word32 Is8)
Rd=!cmp.eq(Rs,Rt)	Word32 Q6_R_not_cmp_eq_RR(Word32 Rs, Word32 Rt)
Rd=cmp.eq(Rs,#s8)	Word32 Q6_R_cmp_eq_RI(Word32 Rs, Word32 Is8)
Rd=cmp.eq(Rs,Rt)	Word32 Q6_R_cmp_eq_RR(Word32 Rs, Word32 Rt)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		Rs		MajOp		MinOp		s5					Parse					d5														
0	1	1	1	0	0	1	1	-	1	0	s	s	s	s	s	P	P	1	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=cmp.eq(Rs,#s8)
0	1	1	1	0	0	1	1	-	1	1	s	s	s	s	s	P	P	1	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=!cmp.eq(Rs,#s8)
ICLASS		P		MajOp		MinOp		s5					Parse					t5					d5									
1	1	1	1	0	0	1	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=cmp.eq(Rs,Rt)
1	1	1	1	0	0	1	1	0	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=!cmp.eq(Rs,Rt)

Field name	Description
MajOp	Major opcode
MinOp	Minor opcode
Rs	No Rs read
P	Predicated
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## 11.2 CR

The CR instruction class includes instructions which manage control registers, including hardware looping, modulo addressing, and status flags.

CR instructions are executable on slot 3.

### End loop instructions

The endloop instructions mark the end of a hardware loop. If the loop count (LC) register indicates that a loop should continue to iterate, the LC register is decremented and the program flow changes to the address in the start address (SA) register.

The endloopN instruction is actually a pseudo-instruction encoded in bits 15:14 of each instruction. Therefore, no distinct 32-bit encoding exists for this instruction.

Syntax	Behavior
endloop0	<pre> if (USR.LPCFG) {   if (USR.LPCFG==1) {     P3=0xff;   }   USR.LPCFG=USR.LPCFG-1; } if (LC0&gt;1) {   PC=SA0;   LC0=LC0-1; } </pre>
endloop01	<pre> if (USR.LPCFG) {   if (USR.LPCFG==1) {     P3=0xff;   }   USR.LPCFG=USR.LPCFG-1; } if (LC0&gt;1) {   PC=SA0;   LC0=LC0-1; } else {   if (LC1&gt;1) {     PC=SA1;     LC1=LC1-1;   } } </pre>
endloop1	<pre> if (LC1&gt;1) {   PC=SA1;   LC1=LC1-1; } </pre>

**Class:** N/A

#### Notes

- This instruction cannot be grouped in a packet with any program flow instructions.

- The Next PC value is the address immediately following the last instruction in the packet containing this instruction.
- The PC value is the address of the start of the packet

## Corner detection acceleration

The FASTCORNER9 instruction takes the Ps and Pt values and treats them as a circular bit string. If any contiguous nine bits are set around the circle, the result is true, false otherwise. The sense may be optionally inverted. This instruction is used to accelerate FAST corner detection.

### Syntax

```
Pd=[!]fastcorner9(Ps,Pt)
```

### Behavior

```
PREDUSE_TIMING;
tmp.h[0]=(Ps<<8)|Pt;
tmp.h[1]=(Ps<<8)|Pt;
for (i = 1; i < 9; i++) {
    tmp &= tmp >> 1;
}
Pd = tmp == 0 ? 0xff : 0x00;
```

### Class: CR (slot 2,3)

### Notes

- This instruction may execute on either slot2 or slot3, even though it is a CR-type

### Intrinsics

```
Pd=!fastcorner9(Ps,Pt)
```

```
Byte Q6_p_not_fastcorner9_pp(Byte Ps, Byte Pt)
```

```
Pd=fastcorner9(Ps,Pt)
```

```
Byte Q6_p_fastcorner9_pp(Byte Ps, Byte Pt)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				sm								s2		Parse				t2				d2										
0	1	1	0	1	0	1	1	0	0	0	0	-	-	s	s	P	P	1	-	-	-	t	t	1	-	-	1	-	-	d	d	Pd=fastcorner9(Ps,Pt)
0	1	1	0	1	0	1	1	0	0	0	1	-	-	s	s	P	P	1	-	-	-	t	t	1	-	-	1	-	-	d	d	Pd=!fastcorner9(Ps,Pt)

Field name	Description
sm	Supervisor mode only
ICLASS	Instruction class
Parse	Packet/loop parse bits
d2	Field to encode register d
s2	Field to encode register s
t2	Field to encode register t

## Logical reductions on predicates

The any8 instruction sets a destination predicate register to 0xff if any of the low eight bits in source predicate register Ps are set. Otherwise, the predicate is set to 0x00.

The all8 instruction sets a destination predicate register to 0xff if all of the low eight bits in the source predicate register Ps are set. Otherwise, the predicate is set to 0x00.

Syntax	Behavior
Pd=all8(Ps)	PREDUSE_TIMING; Pd = (Ps == 0xff ? 0xff : 0x00);
Pd=any8(Ps)	PREDUSE_TIMING; Pd = (Ps ? 0xff : 0x00);

### Class: CR (slot 2,3)

#### Notes

- This instruction may execute on either slot2 or slot3, even though it is a CR-type

#### Intrinsics

Pd=all8(Ps)      Byte Q6\_p\_all8\_p(Byte Ps)

Pd=any8(Ps)      Byte Q6\_p\_any8\_p(Byte Ps)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS				sm										s2		Parse								d2									
0	1	1	0	1	0	1	1	1	0	0	0	-	-	s	s	P	P	0	-	-	-	-	-	-	-	-	-	-	-	-	d	d	Pd=any8(Ps)
0	1	1	0	1	0	1	1	1	0	1	0	-	-	s	s	P	P	0	-	-	-	-	-	-	-	-	-	-	-	d	d	Pd=all8(Ps)	

Field name	Description
sm	Supervisor mode only
ICLASS	Instruction class
Parse	Packet/loop parse bits
d2	Field to encode register d
s2	Field to encode register s

## Looping instructions

loopN is a single instruction which sets up a hardware loop. The N in the instruction name indicates the set of loop registers to use. Loop0 is the innermost loop, while loop1 is the outer loop. The loopN instruction first sets the start address (SA) register based on a PC-relative immediate add. The relative immediate is added to the PC and stored in SA. The loop count (LC) register is set to either an unsigned immediate or to a register value.

Syntax	Behavior
loop0(#r7:2,#U10)	<pre> apply_extension(#r); #r=#r &amp; ~PCALIGN_MASK; SA0=PC+#r; LC0=#U; USR.LPCFG=0; </pre>
loop0(#r7:2,Rs)	<pre> apply_extension(#r); #r=#r &amp; ~PCALIGN_MASK; SA0=PC+#r; LC0=Rs; USR.LPCFG=0; </pre>
loop1(#r7:2,#U10)	<pre> apply_extension(#r); #r=#r &amp; ~PCALIGN_MASK; SA1=PC+#r; LC1=#U; </pre>
loop1(#r7:2,Rs)	<pre> apply_extension(#r); #r=#r &amp; ~PCALIGN_MASK; SA1=PC+#r; LC1=Rs; </pre>

### Class: CR (slot 3)

#### Notes

- This instruction cannot execute in the last address of a hardware loop.
- The Next PC value is the address immediately following the last instruction in the packet containing this instruction.
- The PC value is the address of the start of the packet
- A PC-relative address is formed by taking the decoded immediate value and adding it to the current PC value.

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				sm				s5				Parse																				
0	1	1	0	0	0	0	0	0	0	0	s	s	s	s	s	P	P	-	i	i	i	i	i	-	-	-	i	i	-	-	-	loop0(#r7:2,Rs)
0	1	1	0	0	0	0	0	0	0	1	s	s	s	s	s	P	P	-	i	i	i	i	i	-	-	-	i	i	-	-	-	loop1(#r7:2,Rs)
ICLASS				sm								Parse																				
0	1	1	0	1	0	0	1	0	0	0	I	I	I	I	I	P	P	-	i	i	i	i	i	I	I	I	i	i	-	I	I	loop0(#r7:2,#U10)
0	1	1	0	1	0	0	1	0	0	1	I	I	I	I	I	P	P	-	i	i	i	i	i	I	I	I	i	i	-	I	I	loop1(#r7:2,#U10)

<b>Field name</b>	<b>Description</b>
sm	Supervisor mode only
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s

## Add to PC

Add an immediate value to the program counter (PC) and place the result in a destination register. This instruction is typically used with a constant extender to add a 32-bit immediate value to PC.

### Syntax

```
Rd=add(pc, #u6)
```

### Behavior

```
Rd=PC+apply_extension(#u);
```

### Class: CR (slot 3)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS				sm												Parse					d5												
0	1	1	0	1	0	1	0	0	1	0	0	1	0	0	1	P	P	-	i	i	i	i	i	i	i	-	-	d	d	d	d	d	Rd=add(pc,#u6)

Field name	Description
sm	Supervisor mode only
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d



## Pipelined loop instructions

The spNloop0 instruction is a single instruction that sets up a hardware loop with automatic predicate control. This saves code size by enabling many software pipelined loops to generate without prologue code. Upon executing this instruction, the P3 register automatically clears. After the loop executes N times (where N is selectable from 1-3), the P3 register is set. This ensures that store instructions in the loop are predicated with P3 and thus not enabled during the pipeline warm-up.

In the spNloop0 instruction, the loop 0 (inner-loop) registers are used. This instruction sets the start address (SA0) register based on a PC-relative immediate add. The relative immediate is added to the PC and stored in SA0. The loop count (LC0) is set to either an unsigned immediate or to a register value. The predicate P3 is cleared. The USR.LPCFG bits are set based on the N value.

Syntax	Behavior
<code>p3=sp1loop0 (#r7:2, #U10)</code>	<pre> apply_extension(#r); #r=#r &amp; ~PCALIGN_MASK; SA0=PC+#r; LC0=#U; USR.LPCFG=1; P3=0; </pre>
<code>p3=sp1loop0 (#r7:2, Rs)</code>	<pre> apply_extension(#r); #r=#r &amp; ~PCALIGN_MASK; SA0=PC+#r; LC0=Rs; USR.LPCFG=1; P3=0; </pre>
<code>p3=sp2loop0 (#r7:2, #U10)</code>	<pre> apply_extension(#r); #r=#r &amp; ~PCALIGN_MASK; SA0=PC+#r; LC0=#U; USR.LPCFG=2; P3=0; </pre>
<code>p3=sp2loop0 (#r7:2, Rs)</code>	<pre> apply_extension(#r); #r=#r &amp; ~PCALIGN_MASK; SA0=PC+#r; LC0=Rs; USR.LPCFG=2; P3=0; </pre>
<code>p3=sp3loop0 (#r7:2, #U10)</code>	<pre> apply_extension(#r); #r=#r &amp; ~PCALIGN_MASK; SA0=PC+#r; LC0=#U; USR.LPCFG=3; P3=0; </pre>
<code>p3=sp3loop0 (#r7:2, Rs)</code>	<pre> apply_extension(#r); #r=#r &amp; ~PCALIGN_MASK; SA0=PC+#r; LC0=Rs; USR.LPCFG=3; P3=0; </pre>

**Class: CR (slot 3)****Notes**

- The predicate generated by this instruction can not be used as a .new predicate, nor can it be automatically ANDed with another predicate.
- This instruction cannot execute in the last address of a hardware loop.
- The Next PC value is the address immediately following the last instruction in the packet containing this instruction.
- The PC value is the address of the start of the packet
- A PC-relative address is formed by taking the decoded immediate value and adding it to the current PC value.

**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				sm							s5					Parse																
0	1	1	0	0	0	0	0	1	0	1	s	s	s	s	s	P	P	-	i	i	i	i	i	-	-	-	i	i	-	-	-	p3=sp1loop0(#7:2,Rs)
0	1	1	0	0	0	0	0	1	1	0	s	s	s	s	s	P	P	-	i	i	i	i	i	-	-	-	i	i	-	-	-	p3=sp2loop0(#7:2,Rs)
0	1	1	0	0	0	0	0	1	1	1	s	s	s	s	s	P	P	-	i	i	i	i	i	-	-	-	i	i	-	-	-	p3=sp3loop0(#7:2,Rs)
ICLASS				sm												Parse																
0	1	1	0	1	0	0	1	1	0	1	I	I	I	I	I	P	P	-	i	i	i	i	i	I	I	I	i	i	-	I	I	p3=sp1loop0(#7:2,#U10)
0	1	1	0	1	0	0	1	1	1	0	I	I	I	I	I	P	P	-	i	i	i	i	i	I	I	I	i	i	-	I	I	p3=sp2loop0(#7:2,#U10)
0	1	1	0	1	0	0	1	1	1	1	I	I	I	I	I	P	P	-	i	i	i	i	i	I	I	I	i	i	-	I	I	p3=sp3loop0(#7:2,#U10)

Field name	Description
sm	Supervisor mode only
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s

## Logical operations on predicates

Perform bitwise logical operations on predicate registers.

Syntax	Behavior
<code>Pd=Ps</code>	Assembler mapped to: " <code>Pd=or(Ps,Ps)</code> "
<code>Pd=and(Ps, and(Pt, [!]Pu))</code>	PREDUSE_TIMING; <code>Pd = Ps &amp; Pt &amp; (~Pu);</code>
<code>Pd=and(Ps, or(Pt, [!]Pu))</code>	PREDUSE_TIMING; <code>Pd = Ps &amp; (Pt   (~Pu));</code>
<code>Pd=and(Pt, [!]Ps)</code>	PREDUSE_TIMING; <code>Pd=Pt &amp; (~Ps);</code>
<code>Pd=not(Ps)</code>	PREDUSE_TIMING; <code>Pd=~Ps;</code>
<code>Pd=or(Ps, and(Pt, [!]Pu))</code>	PREDUSE_TIMING; <code>Pd = Ps   (Pt &amp; (~Pu));</code>
<code>Pd=or(Ps, or(Pt, [!]Pu))</code>	PREDUSE_TIMING; <code>Pd = Ps   Pt   (~Pu);</code>
<code>Pd=or(Pt, [!]Ps)</code>	PREDUSE_TIMING; <code>Pd=Pt   (~Ps);</code>
<code>Pd=xor(Ps, Pt)</code>	PREDUSE_TIMING; <code>Pd=Ps ^ Pt;</code>

### Class: CR (slot 2,3)

#### Notes

- This instruction may execute on either slot2 or slot3, even though it is a CR-type

#### Intrinsics

<code>Pd=Ps</code>	Byte <code>Q6_p_equals_p(Byte Ps)</code>
<code>Pd=and(Ps, and(Pt, !Pu))</code>	Byte <code>Q6_p_and_and_ppnp(Byte Ps, Byte Pt, Byte Pu)</code>
<code>Pd=and(Ps, and(Pt, Pu))</code>	Byte <code>Q6_p_and_and_ppp(Byte Ps, Byte Pt, Byte Pu)</code>
<code>Pd=and(Ps, or(Pt, !Pu))</code>	Byte <code>Q6_p_and_or_ppnp(Byte Ps, Byte Pt, Byte Pu)</code>
<code>Pd=and(Ps, or(Pt, Pu))</code>	Byte <code>Q6_p_and_or_ppp(Byte Ps, Byte Pt, Byte Pu)</code>
<code>Pd=and(Pt, !Ps)</code>	Byte <code>Q6_p_and_pnp(Byte Pt, Byte Ps)</code>
<code>Pd=and(Pt, Ps)</code>	Byte <code>Q6_p_and_pp(Byte Pt, Byte Ps)</code>
<code>Pd=not(Ps)</code>	Byte <code>Q6_p_not_p(Byte Ps)</code>
<code>Pd=or(Ps, and(Pt, !Pu))</code>	Byte <code>Q6_p_or_and_ppnp(Byte Ps, Byte Pt, Byte Pu)</code>
<code>Pd=or(Ps, and(Pt, Pu))</code>	Byte <code>Q6_p_or_and_ppp(Byte Ps, Byte Pt, Byte Pu)</code>

Pd=or(Ps,or(Pt,!Pu))	Byte Q6_p_or_or_ppnp(Byte Ps, Byte Pt, Byte Pu)
Pd=or(Ps,or(Pt,Pu))	Byte Q6_p_or_or_ppp(Byte Ps, Byte Pt, Byte Pu)
Pd=or(Pt,!Ps)	Byte Q6_p_or_pnp(Byte Pt, Byte Ps)
Pd=or(Pt,Ps)	Byte Q6_p_or_pp(Byte Pt, Byte Ps)
Pd=xor(Ps,Pt)	Byte Q6_p_xor_pp(Byte Ps, Byte Pt)

## Encoding

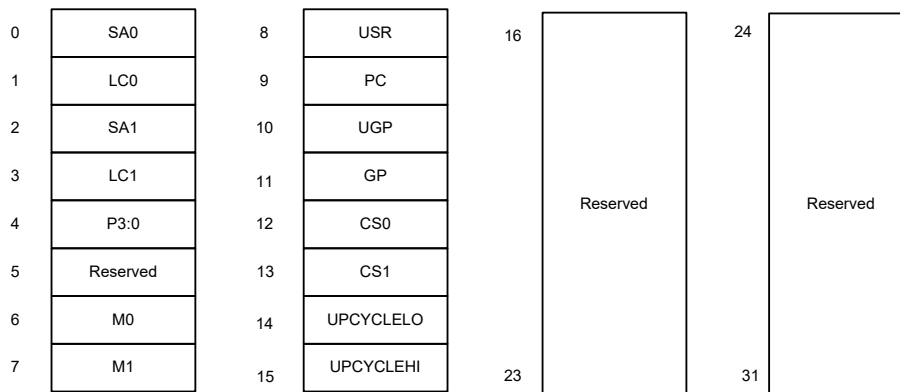
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		sm												s2		Parse				t2				d2								
0	1	1	0	1	0	1	1	0	0	0	0	-	-	s	s	P	P	0	-	-	-	t	t	-	-	-	-	-	-	d	d	Pd=and(Pt,Ps)
ICLASS		sm												s2		Parse				t2		u2		d2								
0	1	1	0	1	0	1	1	0	0	0	1	-	-	s	s	P	P	0	-	-	-	t	t	u	u	-	-	-	-	d	d	Pd=and(Ps,and(Pt,Pu))
ICLASS		sm												s2		Parse				t2				d2								
0	1	1	0	1	0	1	1	0	0	1	0	-	-	s	s	P	P	0	-	-	-	t	t	-	-	-	-	-	-	d	d	Pd=or(Pt,Ps)
ICLASS		sm												s2		Parse				t2		u2		d2								
0	1	1	0	1	0	1	1	0	0	1	1	-	-	s	s	P	P	0	-	-	-	t	t	u	u	-	-	-	-	d	d	Pd=and(Ps,or(Pt,Pu))
ICLASS		sm												s2		Parse				t2				d2								
0	1	1	0	1	0	1	1	0	1	0	0	-	-	s	s	P	P	0	-	-	-	t	t	-	-	-	-	-	-	d	d	Pd=xor(Ps,Pt)
ICLASS		sm												s2		Parse				t2		u2		d2								
0	1	1	0	1	0	1	1	0	1	0	1	-	-	s	s	P	P	0	-	-	-	t	t	u	u	-	-	-	-	d	d	Pd=or(Ps,and(Pt,Pu))
ICLASS		sm												s2		Parse				t2				d2								
0	1	1	0	1	0	1	1	0	1	1	0	-	-	s	s	P	P	0	-	-	-	t	t	-	-	-	-	-	-	d	d	Pd=and(Pt,!Ps)
ICLASS		sm												s2		Parse				t2		u2		d2								
0	1	1	0	1	0	1	1	0	1	1	1	-	-	s	s	P	P	0	-	-	-	t	t	u	u	-	-	-	-	d	d	Pd=or(Ps,or(Pt,Pu))
0	1	1	0	1	0	1	1	1	0	0	1	-	-	s	s	P	P	0	-	-	-	t	t	u	u	-	-	-	-	d	d	Pd=and(Ps,and(Pt,!Pu))
0	1	1	0	1	0	1	1	1	0	1	1	-	-	s	s	P	P	0	-	-	-	t	t	u	u	-	-	-	-	d	d	Pd=and(Ps,or(Pt,!Pu))
ICLASS		sm												s2		Parse				t2				d2								
0	1	1	0	1	0	1	1	1	1	0	0	-	-	s	s	P	P	0	-	-	-	-	-	-	-	-	-	-	-	d	d	Pd=not(Ps)
ICLASS		sm												s2		Parse				t2		u2		d2								
0	1	1	0	1	0	1	1	1	1	0	1	-	-	s	s	P	P	0	-	-	-	t	t	u	u	-	-	-	-	d	d	Pd=or(Ps,and(Pt,!Pu))
ICLASS		sm												s2		Parse				t2				d2								
0	1	1	0	1	0	1	1	1	1	1	0	-	-	s	s	P	P	0	-	-	-	t	t	-	-	-	-	-	-	d	d	Pd=or(Pt,!Ps)
ICLASS		sm												s2		Parse				t2		u2		d2								
0	1	1	0	1	0	1	1	1	1	1	1	-	-	s	s	P	P	0	-	-	-	t	t	u	u	-	-	-	-	d	d	Pd=or(Ps,or(Pt,!Pu))

Field name	Description
sm	Supervisor mode only
ICLASS	Instruction class
Parse	Packet/loop parse bits
d2	Field to encode register d
s2	Field to encode register s
t2	Field to encode register t
u2	Field to encode register u

## User control register transfer

Move 32- or 64-bit values between a user control register and a general register. The user control registers include SA, LC, Predicates, M, USR, PC, UGP, GP, and CS, and UPCYCLE. Registers are moved as singles or as aligned 64-bit pairs.

The PC register is not writable. A program flow instruction must be used to change the PC value.



**Figure 11-1 User control registers and their register field encodings**

Syntax	Behavior
Cd=Rs	Cd=Rs;
Cdd=Rss	Cdd=Rss;
Rd=Cs	Rd=Cs;
Rdd=Css	Rdd=Css;

### Class: CR (slot 3)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				sm		s5					Parse		d5																			
0	1	1	0	0	0	1	0	0	0	1	s	s	s	s	s	P	P	-	-	-	-	-	-	-	-	-	d	d	d	d	d	Cd=Rs
0	1	1	0	0	0	1	1	0	0	1	s	s	s	s	s	P	P	-	-	-	-	-	-	-	-	-	d	d	d	d	d	Cdd=Rss
0	1	1	0	1	0	0	0	0	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	-	-	-	d	d	d	d	d	Rdd=Css
0	1	1	0	1	0	1	0	0	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	-	-	-	d	d	d	d	d	Rd=Cs

Field name	Description
sm	Supervisor mode only
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s

## 11.3 JR

The JR instruction class includes instructions to change the program flow to a new location contained in a register.

JR instructions are executable on slot 2.

### Call subroutine from register

Change the program flow to a subroutine. This instruction first transfers the next program counter (NPC) value into the link register, and then jumps to a target address contained in a register.

This instruction can only appear in slot 2.

Syntax	Behavior
<code>callr Rs</code>	LR=NPC; PC=Rs;
<code>callrh Rs</code>	LR=NPC; PC=Rs;
<code>if ([!]Pu) callr Rs</code>	if ([!]Pu[0]) { LR=NPC; PC=Rs; }

### Class: JR (slot 2)

#### Notes

- This instruction can conditionally execute based on the value of a predicate register. If the instruction is preceded by 'if Pn', the instruction only executes if the least-significant bit of the predicate register is 1. Similarly, if the instruction is preceded by 'if !Pn', the instruction executes only if the least-significant bit of Pn is 0.

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0					
ICLASS											s5					Parse																				
0	1	0	1	0	0	0	0	1	0	1	s	s	s	s	s	s	P	P	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	callr Rs	
0	1	0	1	0	0	0	0	1	1	0	s	s	s	s	s	s	P	P	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	callrh Rs
ICLASS											s5					Parse										u2										
0	1	0	1	0	0	0	1	0	0	0	s	s	s	s	s	s	P	P	-	-	-	-	u	u	-	-	-	-	-	-	-	-	-	-	-	if (Pu) callr Rs
0	1	0	1	0	0	0	1	0	0	1	s	s	s	s	s	s	P	P	-	-	-	-	u	u	-	-	-	-	-	-	-	-	-	-	-	if (!Pu) callr Rs

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s
u2	Field to encode register u

## Hinted call subroutine from register

Change the program flow to a subroutine. This instruction first transfers the next program counter (NPC) value into the link register, and then jumps to a target address contained in a register. This instruction is effective only when a preceding hintjr exists.

This instruction can only appear in slot 2.

Syntax	Behavior
callrh Rs	LR=NPC; PC=Rs;

### Class: JR (slot 2)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS			s5											Parse																			
0	1	0	1	0	0	0	0	0	1	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	-	-	-	-	-	-	-	-	callrh Rs

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s

## Hint an indirect jump address

Provide a hint indicating that there will soon be an indirect call to the address specified in Rs. The indirect call can be either jumprh or callrh.

This instruction can appear in either slot 2 or slot 3.

Syntax	Behavior
<code>hintjr(Rs)</code>	;

**Class: JR (slot 2,3)**

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS					s5											Parse																
0	1	0	1	0	0	1	0	1	0	1	s	s	s	s	s	P	P	-	-	-	-	-	-	-	-	-	-	-	-	-	-	hintjr(Rs)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s



## Jump to address from register

Change the program flow to a target address. This instruction changes the program counter to a target address contained in a register.

This instruction can appear only in slot 2.

Syntax	Behavior
<code>if ([!]Pu) jumpr Rs</code>	Assembler mapped to: <code>"if ([!]Pu) ""jumpr"":nt ""Rs"</code>
<code>if ([!]Pu[.new]) jumpr:&lt;hint&gt; Rs</code>	<pre>{ if ([!]Pu[.new][0]) {     PC=Rs; }</pre>
<code>jumpr Rs</code>	<code>PC=Rs;</code>
<code>jumprh Rs</code>	<code>PC=Rs;</code>

### Class: JR (slot 2)

#### Notes

- This instruction can conditionally execute based on the value of a predicate register. If the instruction is preceded by 'if Pn', the instruction only executes if the least-significant bit of the predicate register is 1. Similarly, if the instruction is preceded by 'if !Pn', the instruction executes only if the least-significant bit of Pn is 0.

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS											s5					Parse																	
0	1	0	1	0	0	1	0	1	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	jumpr Rs
0	1	0	1	0	0	1	0	1	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	jumprh Rs
ICLASS											s5					Parse					u2												
0	1	0	1	0	0	1	1	0	1	0	s	s	s	s	s	P	P	-	0	0	-	u	u	-	-	-	-	-	-	-	-	if (Pu) jumpr:nt Rs	
0	1	0	1	0	0	1	1	0	1	0	s	s	s	s	s	P	P	-	0	1	-	u	u	-	-	-	-	-	-	-	-	if (Pu.new) jumpr:nt Rs	
0	1	0	1	0	0	1	1	0	1	0	s	s	s	s	s	P	P	-	1	0	-	u	u	-	-	-	-	-	-	-	-	if (Pu) jumpr:t Rs	
0	1	0	1	0	0	1	1	0	1	0	s	s	s	s	s	P	P	-	1	1	-	u	u	-	-	-	-	-	-	-	-	if (Pu.new) jumpr:t Rs	
0	1	0	1	0	0	1	1	0	1	1	s	s	s	s	s	P	P	-	0	0	-	u	u	-	-	-	-	-	-	-	-	if (!Pu) jumpr:nt Rs	
0	1	0	1	0	0	1	1	0	1	1	s	s	s	s	s	P	P	-	0	1	-	u	u	-	-	-	-	-	-	-	-	if (!Pu.new) jumpr:nt Rs	
0	1	0	1	0	0	1	1	0	1	1	s	s	s	s	s	P	P	-	1	0	-	u	u	-	-	-	-	-	-	-	-	if (!Pu) jumpr:t Rs	
0	1	0	1	0	0	1	1	0	1	1	s	s	s	s	s	P	P	-	1	1	-	u	u	-	-	-	-	-	-	-	-	if (!Pu.new) jumpr:t Rs	

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s
u2	Field to encode register u

## Hinted jump to address from register

Change the program flow to a target address. This instruction changes the program counter to a target address contained in a register. This instruction is effective only when a preceding hintjr exists.

This instruction can appear only in slot 2.

Syntax	Behavior
jumprh Rs	PC=Rs;

### Class: JR (slot 2)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0				
ICLASS											s5					Parse																			
0	1	0	1	0	0	1	0	1	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	jumprh Rs

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s

## 11.4 J

The J instruction class includes branch instructions (jumps and calls) that obtain the target address from a (PC-relative) immediate address value.

J instructions are executable on slot 2 and slot 3.

### Call subroutine

Change the program flow to a subroutine. This instruction first transfers the next program counter (NPC) value into the link register, and then jumps to the target address.

This instruction can appear in slots 2 or 3.

Syntax	Behavior
<code>call #r22:2</code>	<pre> apply_extension(#r); #r=#r &amp; ~PCALIGN_MASK; LR=NPC; PC=PC+#r; </pre>
<code>if ([!]Pu) call #r15:2</code>	<pre> apply_extension(#r); #r=#r &amp; ~PCALIGN_MASK; if ([!]Pu[0]) {     LR=NPC;     PC=PC+#r; } </pre>

### Class: J (slots 2,3)

#### Notes

- This instruction can conditionally execute based on the value of a predicate register. If the instruction is preceded by 'if Pn', the instruction only executes if the least-significant bit of the predicate register is 1. Similarly, if the instruction is preceded by 'if !Pn', the instruction is executed only if the least-significant bit of Pn is 0.
- The Next PC value is the address immediately following the last instruction in the packet containing this instruction.
- The PC value is the address of the start of the packet
- A PC-relative address is formed by taking the decoded immediate value and adding it to the current PC value.

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS																Parse																
0	1	0	1	1	0	1	i	i	i	i	i	i	i	i	i	P	P	i	i	i	i	i	i	i	i	i	i	i	i	i	0	call #r22:2
ICLASS																Parse											D N	u2				
0	1	0	1	1	1	0	1	i	i	0	i	i	i	i	i	P	P	i	-	0	-	u	u	i	i	i	i	i	i	i	-	if (Pu) call #r15:2
0	1	0	1	1	1	0	1	i	i	1	i	i	i	i	i	P	P	i	-	0	-	u	u	i	i	i	i	i	i	-	if (!Pu) call #r15:2	

<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
DN	Dot-new
Parse	Packet/loop parse bits
u2	Field to encode register u

## Compare and jump

Compare two registers, or a register and immediate value, and write a predicate with the result. Then use the predicate result to conditionally jump to a PC-relative target address.

The registers available as operands are restricted to R0-R7 and R16-R23. The predicate destination is restricted to P0 and P1.

In assembly syntax, this instruction appears as two instructions in the packet: a compare and a separate conditional jump. The assembler may convert adjacent compare and jump instructions into compound compare-jump form.

Syntax	Behavior
<code>p[01]=cmp.eq(Rs,#-1); if (!p[01].new) jump:&lt;hint&gt; #r9:2</code>	<code>P[01]=(Rs==-1) ? 0xff : 0x00 if (!P[01].new[0]) {   apply_extension(#r);   #r=#r &amp; ~PCALIGN_MASK;   PC=PC+#r; }</code>
<code>p[01]=cmp.eq(Rs,#U5); if (!p[01].new) jump:&lt;hint&gt; #r9:2</code>	<code>P[01]=(Rs==#U) ? 0xff : 0x00 if (!P[01].new[0]) {   apply_extension(#r);   #r=#r &amp; ~PCALIGN_MASK;   PC=PC+#r; }</code>
<code>p[01]=cmp.eq(Rs,Rt); if (!p[01].new) jump:&lt;hint&gt; #r9:2</code>	<code>P[01]=(Rs==Rt) ? 0xff : 0x00 if (!P[01].new[0]) {   apply_extension(#r);   #r=#r &amp; ~PCALIGN_MASK;   PC=PC+#r; }</code>
<code>p[01]=cmp.gt(Rs,#-1); if (!p[01].new) jump:&lt;hint&gt; #r9:2</code>	<code>P[01]=(Rs&gt;-1) ? 0xff : 0x00 if (!P[01].new[0]) {   apply_extension(#r);   #r=#r &amp; ~PCALIGN_MASK;   PC=PC+#r; }</code>
<code>p[01]=cmp.gt(Rs,#U5); if (!p[01].new) jump:&lt;hint&gt; #r9:2</code>	<code>P[01]=(Rs&gt;#U) ? 0xff : 0x00 if (!P[01].new[0]) {   apply_extension(#r);   #r=#r &amp; ~PCALIGN_MASK;   PC=PC+#r; }</code>
<code>p[01]=cmp.gt(Rs,Rt); if (!p[01].new) jump:&lt;hint&gt; #r9:2</code>	<code>P[01]=(Rs&gt;Rt) ? 0xff : 0x00 if (!P[01].new[0]) {   apply_extension(#r);   #r=#r &amp; ~PCALIGN_MASK;   PC=PC+#r; }</code>

Syntax	Behavior
<code>p[01]=cmp.gtu(Rs,#U5); if ([!]p[01].new) jump:&lt;hint&gt; #r9:2</code>	<code>P[01]=(Rs.uw[0]&gt;#U) ? 0xff : 0x00 if ([!]P[01].new[0]) { apply_extension(#r); #r=#r &amp; ~PCALIGN_MASK; PC=PC+#r; }</code>
<code>p[01]=cmp.gtu(Rs,Rt); if ([!]p[01].new) jump:&lt;hint&gt; #r9:2</code>	<code>P[01]=(Rs.uw[0]&gt;Rt) ? 0xff : 0x00 if ([!]P[01].new[0]) { apply_extension(#r); #r=#r &amp; ~PCALIGN_MASK; PC=PC+#r; }</code>
<code>p[01]=tstbit(Rs,#0); if ([!]p[01].new) jump:&lt;hint&gt; #r9:2</code>	<code>P[01]=(Rs &amp; 1) ? 0xff : 0x00 if ([!]P[01].new[0]) { apply_extension(#r); #r=#r &amp; ~PCALIGN_MASK; PC=PC+#r; }</code>

**Class: J (slots 0,1,2,3)**

**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS												s4				Parse																	
0	0	0	1	0	0	0	1	1	0	i	i	s	s	s	s	P	P	0	-	-	-	0	0	i	i	i	i	i	i	i	i	-	p0=cmp.eq(Rs,#-1); if (p0.new) jump:nt #r9:2
0	0	0	1	0	0	0	1	1	0	i	i	s	s	s	s	P	P	0	-	-	-	0	1	i	i	i	i	i	i	i	i	-	p0=cmp.gt(Rs,#-1); if (p0.new) jump:nt #r9:2
0	0	0	1	0	0	0	1	1	0	i	i	s	s	s	s	P	P	0	-	-	-	1	1	i	i	i	i	i	i	i	-	p0=tstbit(Rs,#0); if (p0.new) jump:nt #r9:2	
0	0	0	1	0	0	0	1	1	0	i	i	s	s	s	s	P	P	1	-	-	-	0	0	i	i	i	i	i	i	i	-	p0=cmp.eq(Rs,#-1); if (p0.new) jump:t #r9:2	
0	0	0	1	0	0	0	1	1	0	i	i	s	s	s	s	P	P	1	-	-	-	0	1	i	i	i	i	i	i	i	-	p0=cmp.gt(Rs,#-1); if (p0.new) jump:t #r9:2	
0	0	0	1	0	0	0	1	1	0	i	i	s	s	s	s	P	P	1	-	-	-	1	1	i	i	i	i	i	i	i	-	p0=tstbit(Rs,#0); if (p0.new) jump:t #r9:2	
0	0	0	1	0	0	0	1	1	1	i	i	s	s	s	s	P	P	0	-	-	-	0	0	i	i	i	i	i	i	i	-	p0=cmp.eq(Rs,#-1); if (!p0.new) jump:nt #r9:2	
0	0	0	1	0	0	0	1	1	1	i	i	s	s	s	s	P	P	0	-	-	-	0	1	i	i	i	i	i	i	i	-	p0=cmp.gt(Rs,#-1); if (!p0.new) jump:nt #r9:2	
0	0	0	1	0	0	0	1	1	1	i	i	s	s	s	s	P	P	0	-	-	-	1	1	i	i	i	i	i	i	i	-	p0=tstbit(Rs,#0); if (!p0.new) jump:nt #r9:2	
0	0	0	1	0	0	0	1	1	1	i	i	s	s	s	s	P	P	1	-	-	-	0	0	i	i	i	i	i	i	i	-	p0=cmp.eq(Rs,#-1); if (!p0.new) jump:t #r9:2	
0	0	0	1	0	0	0	1	1	1	i	i	s	s	s	s	P	P	1	-	-	-	0	1	i	i	i	i	i	i	i	-	p0=cmp.gt(Rs,#-1); if (!p0.new) jump:t #r9:2	
0	0	0	1	0	0	0	1	1	1	i	i	s	s	s	s	P	P	1	-	-	-	1	1	i	i	i	i	i	i	i	-	p0=tstbit(Rs,#0); if (!p0.new) jump:t #r9:2	
0	0	0	1	0	0	0	0	0	0	i	i	s	s	s	s	P	P	0	l	l	l	l	l	i	i	i	i	i	i	i	-	p0=cmp.eq(Rs,#U5); if (p0.new) jump:nt #r9:2	
0	0	0	1	0	0	0	0	0	0	i	i	s	s	s	s	P	P	1	l	l	l	l	l	i	i	i	i	i	i	i	-	p0=cmp.eq(Rs,#U5); if (p0.new) jump:t #r9:2	
0	0	0	1	0	0	0	0	0	1	i	i	s	s	s	s	P	P	0	l	l	l	l	l	i	i	i	i	i	i	i	-	p0=cmp.eq(Rs,#U5); if (!p0.new) jump:nt #r9:2	
0	0	0	1	0	0	0	0	0	1	i	i	s	s	s	s	P	P	1	l	l	l	l	l	i	i	i	i	i	i	i	-	p0=cmp.eq(Rs,#U5); if (!p0.new) jump:t #r9:2	

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	0	1	0	0	0	0	1	0	i	i	s	s	s	s	P	P	0	l	l	l	l	l	i	i	i	i	i	i	i	-	p0=cmp.gt(Rs,#U5); if (p0.new) jump:nt #r9:2
0	0	0	1	0	0	0	0	1	0	i	i	s	s	s	s	P	P	1	l	l	l	l	l	i	i	i	i	i	i	i	-	p0=cmp.gt(Rs,#U5); if (p0.new) jump:t #r9:2
0	0	0	1	0	0	0	0	1	1	i	i	s	s	s	s	P	P	0	l	l	l	l	l	i	i	i	i	i	i	i	-	p0=cmp.gt(Rs,#U5); if (lp0.new) jump:nt #r9:2
0	0	0	1	0	0	0	0	1	1	i	i	s	s	s	s	P	P	1	l	l	l	l	l	i	i	i	i	i	i	i	-	p0=cmp.gt(Rs,#U5); if (lp0.new) jump:t #r9:2
0	0	0	1	0	0	0	1	0	0	i	i	s	s	s	s	P	P	0	l	l	l	l	l	i	i	i	i	i	i	i	-	p0=cmp.gtu(Rs,#U5); if (p0.new) jump:nt #r9:2
0	0	0	1	0	0	0	1	0	0	i	i	s	s	s	s	P	P	1	l	l	l	l	l	i	i	i	i	i	i	i	-	p0=cmp.gtu(Rs,#U5); if (p0.new) jump:t #r9:2
0	0	0	1	0	0	0	1	0	1	i	i	s	s	s	s	P	P	0	l	l	l	l	l	i	i	i	i	i	i	i	-	p0=cmp.gtu(Rs,#U5); if (lp0.new) jump:nt #r9:2
0	0	0	1	0	0	0	1	0	1	i	i	s	s	s	s	P	P	1	l	l	l	l	l	i	i	i	i	i	i	i	-	p0=cmp.gtu(Rs,#U5); if (lp0.new) jump:t #r9:2
0	0	0	1	0	0	1	1	1	0	i	i	s	s	s	s	P	P	0	-	-	-	0	0	i	i	i	i	i	i	i	-	p1=cmp.eq(Rs,#1); if (p1.new) jump:nt #r9:2
0	0	0	1	0	0	1	1	1	0	i	i	s	s	s	s	P	P	0	-	-	-	0	1	i	i	i	i	i	i	i	-	p1=cmp.gt(Rs,#-1); if (p1.new) jump:nt #r9:2
0	0	0	1	0	0	1	1	1	0	i	i	s	s	s	s	P	P	0	-	-	-	1	1	i	i	i	i	i	i	i	-	p1=tstbit(Rs,#0); if (p1.new) jump:nt #r9:2
0	0	0	1	0	0	1	1	1	0	i	i	s	s	s	s	P	P	1	-	-	-	0	0	i	i	i	i	i	i	i	-	p1=cmp.eq(Rs,#-1); if (p1.new) jump:t #r9:2
0	0	0	1	0	0	1	1	1	0	i	i	s	s	s	s	P	P	1	-	-	-	0	1	i	i	i	i	i	i	i	-	p1=cmp.gt(Rs,#-1); if (p1.new) jump:t #r9:2
0	0	0	1	0	0	1	1	1	0	i	i	s	s	s	s	P	P	1	-	-	-	1	1	i	i	i	i	i	i	i	-	p1=tstbit(Rs,#0); if (p1.new) jump:t #r9:2
0	0	0	1	0	0	1	1	1	1	i	i	s	s	s	s	P	P	0	-	-	-	0	0	i	i	i	i	i	i	i	-	p1=cmp.eq(Rs,#-1); if (lp1.new) jump:nt #r9:2
0	0	0	1	0	0	1	1	1	1	i	i	s	s	s	s	P	P	0	-	-	-	0	1	i	i	i	i	i	i	i	-	p1=cmp.gt(Rs,#-1); if (lp1.new) jump:nt #r9:2
0	0	0	1	0	0	1	1	1	1	i	i	s	s	s	s	P	P	0	-	-	-	1	1	i	i	i	i	i	i	i	-	p1=tstbit(Rs,#0); if (lp1.new) jump:nt #r9:2
0	0	0	1	0	0	1	1	1	1	i	i	s	s	s	s	P	P	1	-	-	-	0	0	i	i	i	i	i	i	i	-	p1=cmp.eq(Rs,#-1); if (lp1.new) jump:t #r9:2
0	0	0	1	0	0	1	1	1	1	i	i	s	s	s	s	P	P	1	-	-	-	0	1	i	i	i	i	i	i	i	-	p1=cmp.gt(Rs,#-1); if (lp1.new) jump:t #r9:2
0	0	0	1	0	0	1	1	1	1	i	i	s	s	s	s	P	P	1	-	-	-	1	1	i	i	i	i	i	i	i	-	p1=tstbit(Rs,#0); if (lp1.new) jump:t #r9:2
0	0	0	1	0	0	1	0	0	0	i	i	s	s	s	s	P	P	0	l	l	l	l	l	i	i	i	i	i	i	i	-	p1=cmp.eq(Rs,#U5); if (p1.new) jump:nt #r9:2
0	0	0	1	0	0	1	0	0	0	i	i	s	s	s	s	P	P	1	l	l	l	l	l	i	i	i	i	i	i	i	-	p1=cmp.eq(Rs,#U5); if (lp1.new) jump:t #r9:2
0	0	0	1	0	0	1	0	0	1	i	i	s	s	s	s	P	P	1	l	l	l	l	l	i	i	i	i	i	i	i	-	p1=cmp.eq(Rs,#U5); if (lp1.new) jump:t #r9:2
0	0	0	1	0	0	1	0	1	0	i	i	s	s	s	s	P	P	0	l	l	l	l	l	i	i	i	i	i	i	i	-	p1=cmp.gt(Rs,#U5); if (p1.new) jump:nt #r9:2
0	0	0	1	0	0	1	0	1	0	i	i	s	s	s	s	P	P	1	l	l	l	l	l	i	i	i	i	i	i	i	-	p1=cmp.gt(Rs,#U5); if (p1.new) jump:t #r9:2
0	0	0	1	0	0	1	0	1	1	i	i	s	s	s	s	P	P	0	l	l	l	l	l	i	i	i	i	i	i	i	-	p1=cmp.gt(Rs,#U5); if (lp1.new) jump:nt #r9:2
0	0	0	1	0	0	1	0	1	1	i	i	s	s	s	s	P	P	1	l	l	l	l	l	i	i	i	i	i	i	i	-	p1=cmp.gtu(Rs,#U5); if (p1.new) jump:nt #r9:2
0	0	0	1	0	0	1	1	0	0	i	i	s	s	s	s	P	P	1	l	l	l	l	l	i	i	i	i	i	i	i	-	p1=cmp.gtu(Rs,#U5); if (p1.new) jump:t #r9:2
0	0	0	1	0	0	1	1	0	1	i	i	s	s	s	s	P	P	0	l	l	l	l	l	i	i	i	i	i	i	i	-	p1=cmp.gtu(Rs,#U5); if (lp1.new) jump:nt #r9:2
0	0	0	1	0	0	1	1	0	1	i	i	s	s	s	s	P	P	0	l	l	l	l	l	i	i	i	i	i	i	i	-	p1=cmp.gtu(Rs,#U5); if (lp1.new) jump:t #r9:2

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
0	0	0	1	0	0	1	1	0	1	i	i	s	s	s	s	P	P	1	l	l	l	l	l	i	i	i	i	i	i	i	i	-	p1=cmp.gtu(Rs,#U5); if (lp1.new) jump:t #r9:2
ICLASS												s4				Parse				t4													
0	0	0	1	0	1	0	0	0	0	i	i	s	s	s	s	P	P	0	0	t	t	t	t	i	i	i	i	i	i	i	i	-	p0=cmp.eq(Rs,Rt); if (p0.new) jump:nt #r9:2
0	0	0	1	0	1	0	0	0	0	i	i	s	s	s	s	P	P	0	1	t	t	t	t	i	i	i	i	i	i	i	i	-	p1=cmp.eq(Rs,Rt); if (p1.new) jump:nt #r9:2
0	0	0	1	0	1	0	0	0	0	i	i	s	s	s	s	P	P	1	0	t	t	t	t	i	i	i	i	i	i	i	i	-	p0=cmp.eq(Rs,Rt); if (p0.new) jump:t #r9:2
0	0	0	1	0	1	0	0	0	0	i	i	s	s	s	s	P	P	1	1	t	t	t	t	i	i	i	i	i	i	i	i	-	p1=cmp.eq(Rs,Rt); if (p1.new) jump:t #r9:2
0	0	0	1	0	1	0	0	0	1	i	i	s	s	s	s	P	P	0	0	t	t	t	t	i	i	i	i	i	i	i	i	-	p0=cmp.eq(Rs,Rt); if (lp0.new) jump:nt #r9:2
0	0	0	1	0	1	0	0	0	1	i	i	s	s	s	s	P	P	0	1	t	t	t	t	i	i	i	i	i	i	i	i	-	p1=cmp.eq(Rs,Rt); if (lp1.new) jump:nt #r9:2
0	0	0	1	0	1	0	0	0	1	i	i	s	s	s	s	P	P	1	0	t	t	t	t	i	i	i	i	i	i	i	i	-	p0=cmp.eq(Rs,Rt); if (!p0.new) jump:t #r9:2
0	0	0	1	0	1	0	0	0	1	i	i	s	s	s	s	P	P	1	1	t	t	t	t	i	i	i	i	i	i	i	i	-	p1=cmp.eq(Rs,Rt); if (lp1.new) jump:t #r9:2
0	0	0	1	0	1	0	0	1	0	i	i	s	s	s	s	P	P	0	0	t	t	t	t	i	i	i	i	i	i	i	i	-	p0=cmp.gt(Rs,Rt); if (p0.new) jump:nt #r9:2
0	0	0	1	0	1	0	0	1	0	i	i	s	s	s	s	P	P	0	1	t	t	t	t	i	i	i	i	i	i	i	i	-	p1=cmp.gt(Rs,Rt); if (p1.new) jump:nt #r9:2
0	0	0	1	0	1	0	0	1	0	i	i	s	s	s	s	P	P	1	0	t	t	t	t	i	i	i	i	i	i	i	i	-	p0=cmp.gt(Rs,Rt); if (lp0.new) jump:t #r9:2
0	0	0	1	0	1	0	0	1	0	i	i	s	s	s	s	P	P	1	1	t	t	t	t	i	i	i	i	i	i	i	i	-	p1=cmp.gt(Rs,Rt); if (lp1.new) jump:t #r9:2
0	0	0	1	0	1	0	0	1	1	i	i	s	s	s	s	P	P	0	1	t	t	t	t	i	i	i	i	i	i	i	i	-	p0=cmp.gt(Rs,Rt); if (!p0.new) jump:nt #r9:2
0	0	0	1	0	1	0	0	1	1	i	i	s	s	s	s	P	P	1	0	t	t	t	t	i	i	i	i	i	i	i	i	-	p1=cmp.gt(Rs,Rt); if (lp1.new) jump:t #r9:2
0	0	0	1	0	1	0	1	0	0	i	i	s	s	s	s	P	P	0	0	t	t	t	t	i	i	i	i	i	i	i	i	-	p0=cmp.gtu(Rs,Rt); if (p0.new) jump:nt #r9:2
0	0	0	1	0	1	0	1	0	0	i	i	s	s	s	s	P	P	0	1	t	t	t	t	i	i	i	i	i	i	i	i	-	p1=cmp.gtu(Rs,Rt); if (lp1.new) jump:nt #r9:2
0	0	0	1	0	1	0	1	0	0	i	i	s	s	s	s	P	P	1	0	t	t	t	t	i	i	i	i	i	i	i	i	-	p0=cmp.gtu(Rs,Rt); if (p0.new) jump:t #r9:2
0	0	0	1	0	1	0	1	0	0	i	i	s	s	s	s	P	P	1	1	t	t	t	t	i	i	i	i	i	i	i	i	-	p1=cmp.gtu(Rs,Rt); if (lp1.new) jump:t #r9:2
0	0	0	1	0	1	0	1	0	1	i	i	s	s	s	s	P	P	0	1	t	t	t	t	i	i	i	i	i	i	i	i	-	p0=cmp.gtu(Rs,Rt); if (lp0.new) jump:nt #r9:2
0	0	0	1	0	1	0	1	0	1	i	i	s	s	s	s	P	P	1	0	t	t	t	t	i	i	i	i	i	i	i	i	-	p1=cmp.gtu(Rs,Rt); if (!p1.new) jump:nt #r9:2
0	0	0	1	0	1	0	1	0	1	i	i	s	s	s	s	P	P	1	1	t	t	t	t	i	i	i	i	i	i	i	i	-	p0=cmp.gtu(Rs,Rt); if (lp0.new) jump:t #r9:2
0	0	0	1	0	1	0	1	0	1	i	i	s	s	s	s	P	P	1	1	t	t	t	t	i	i	i	i	i	i	i	i	-	p1=cmp.gtu(Rs,Rt); if (lp1.new) jump:t #r9:2

<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
Parse	Packet/loop parse bits
s4	Field to encode register s
t4	Field to encode register t



## Jump to address

Change the program flow to a target address. This instruction changes the program counter to a target address that is relative to the PC address. The offset from the current PC address is contained in the instruction encoding.

A speculated jump instruction includes a hint ("taken" or "not taken") that specifies the expected value of the conditional expression. If the actual generated value of the predicate differs from this expected value, the jump instruction incurs a performance penalty.

This instruction can appear in slots 2 or 3.

Syntax	Behavior
<code>if ([!]Pu) jump #r15:2</code>	Assembler mapped to: <code>"if ([!]Pu) ""jump"":nt ""#r15:2"</code>
<code>if ([!]Pu) jump:&lt;hint&gt; #r15:2</code>	<pre>if ([!]Pu[0]) {     apply_extension(#r);     #r=#r &amp; ~PCALIGN_MASK;     PC=PC+#r; }</pre>
<code>jump #r22:2</code>	<pre>apply_extension(#r); #r=#r &amp; ~PCALIGN_MASK; PC=PC+#r;</pre>

### Class: J (slots 0,1,2,3)

#### Notes

- This instruction can conditionally execute based on the value of a predicate register. If the instruction is preceded by 'if Pn', the instruction only executes if the least-significant bit of the predicate register is 1. Similarly, if the instruction is preceded by 'if !Pn', the instruction executes only if the least-significant bit of Pn is 0.

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0								
ICLASS																Parse																							
0	1	0	1	1	0	0	i	i	i	i	i	i	i	i	i	P	P	i	i	i	i	i	i	i	i	i	i	i	i	i	i	-	jump #r22:2						
ICLASS																Parse											PT	DN	u2										
0	1	0	1	1	1	0	0	i	i	0	i	i	i	i	i	P	P	i	0	0	-	u	u	i	i	i	i	i	i	i	i	-	if (Pu) jump:nt #r15:2						
0	1	0	1	1	1	0	0	i	i	0	i	i	i	i	i	P	P	i	1	0	-	u	u	i	i	i	i	i	i	i	-	if (Pu) jump:t #r15:2							
0	1	0	1	1	1	0	0	i	i	1	i	i	i	i	i	P	P	i	0	0	-	u	u	i	i	i	i	i	i	-	if (!Pu) jump:nt #r15:2								
0	1	0	1	1	1	0	0	i	i	1	i	i	i	i	i	P	P	i	1	0	-	u	u	i	i	i	i	i	-	if (!Pu) jump:t #r15:2									

Field name	Description
ICLASS	Instruction class
DN	Dot-new
PT	Predict-taken
Parse	Packet/loop parse bits
u2	Field to encode register u

## Jump to address conditioned on new predicate

Perform speculated jump.

Jump if the LSB of the newly-generated predicate is true. The predicate must be generated in the same packet as the speculated jump instruction.

A speculated jump instruction includes a hint ("taken" or "not taken") that specifies the expected value of the conditional expression. If the actual generated value of the predicate differs from this expected value, the jump instruction incurs a performance penalty.

This instruction can appear in slots 2 or 3.

Syntax	Behavior
<pre>if ([!]Pu.new) jump:&lt;hint&gt; #r15:2</pre>	<pre>{ if ([!]Pu.new[0]){   apply_extension(#r);   #r=#r &amp; ~PCALIGN_MASK;   PC=PC+#r; }</pre>

### Class: J (slots 0,1,2,3)

#### Notes

- This instruction can conditionally execute based on the value of a predicate register. If the instruction is preceded by 'if Pn', the instruction only executes if the least-significant bit of the predicate register is 1. Similarly, if the instruction is preceded by 'if !Pn', the instruction executes only if the least-significant bit of Pn is 0.

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS																Parse		PT	<sup>D</sup> <sub>N</sub>	u2													
0	1	0	1	1	1	0	0	i	i	0	i	i	i	i	i	P	P	i	0	1	-	u	u	i	i	i	i	i	i	i	i	-	if (Pu.new) jump:nt #r15:2
0	1	0	1	1	1	0	0	i	i	0	i	i	i	i	i	P	P	i	1	1	-	u	u	i	i	i	i	i	i	i	i	-	if (Pu.new) jump:t #r15:2
0	1	0	1	1	1	0	0	i	i	1	i	i	i	i	i	P	P	i	0	1	-	u	u	i	i	i	i	i	i	i	i	-	if (!Pu.new) jump:nt #r15:2
0	1	0	1	1	1	0	0	i	i	1	i	i	i	i	i	P	P	i	1	1	-	u	u	i	i	i	i	i	i	i	i	-	if (!Pu.new) jump:t #r15:2

Field name	Description
ICLASS	Instruction class
DN	Dot-new
PT	Predict-taken
Parse	Packet/loop parse bits
u2	Field to encode register u

## Jump to address condition on register value

Perform register-conditional jump.

Jump if the specified register expression is true.

A register-conditional jump includes a hint ("taken" or "not taken") that specifies the expected value of the register expression. If the actual generated value of the expression differs from this expected value, the jump instruction incurs a performance penalty.

This instruction can appear only in slot 3.

Syntax	Behavior
<code>if (Rs!=#0) jump:nt #r13:2</code>	<code>if (Rs != 0) {   PC=PC+#r; }</code>
<code>if (Rs!=#0) jump:t #r13:2</code>	<code>if (Rs != 0) {   PC=PC+#r; }</code>
<code>if (Rs&lt;=#0) jump:nt #r13:2</code>	<code>if (Rs &lt;= 0) {   PC=PC+#r; }</code>
<code>if (Rs&lt;=#0) jump:t #r13:2</code>	<code>if (Rs &lt;= 0) {   PC=PC+#r; }</code>
<code>if (Rs==#0) jump:nt #r13:2</code>	<code>if (Rs == 0) {   PC=PC+#r; }</code>
<code>if (Rs==#0) jump:t #r13:2</code>	<code>if (Rs == 0) {   PC=PC+#r; }</code>
<code>if (Rs&gt;=#0) jump:nt #r13:2</code>	<code>if (Rs &gt;= 0) {   PC=PC+#r; }</code>
<code>if (Rs&gt;=#0) jump:t #r13:2</code>	<code>if (Rs &gt;= 0) {   PC=PC+#r; }</code>

### Class: J (slot 3)

#### Notes

- This instruction will be deprecated in a future version.

## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS			sm			s5					Parse																						
0	1	1	0	0	0	0	1	0	0	i	s	s	s	s	s	P	P	i	0	i	i	i	i	i	i	i	i	i	i	i	i	-	if (Rs!=#0) jump:nt #r13:2
0	1	1	0	0	0	0	1	0	0	i	s	s	s	s	s	P	P	i	1	i	i	i	i	i	i	i	i	i	i	i	i	-	if (Rs!=#0) jump:t #r13:2
0	1	1	0	0	0	0	1	0	1	i	s	s	s	s	s	P	P	i	0	i	i	i	i	i	i	i	i	i	i	i	i	-	if (Rs>=#0) jump:nt #r13:2
0	1	1	0	0	0	0	1	0	1	i	s	s	s	s	s	P	P	i	1	i	i	i	i	i	i	i	i	i	i	i	i	-	if (Rs>=#0) jump:t #r13:2
0	1	1	0	0	0	0	1	1	0	i	s	s	s	s	s	P	P	i	0	i	i	i	i	i	i	i	i	i	i	i	i	-	if (Rs==#0) jump:nt #r13:2
0	1	1	0	0	0	0	1	1	0	i	s	s	s	s	s	P	P	i	1	i	i	i	i	i	i	i	i	i	i	i	i	-	if (Rs==#0) jump:t #r13:2
0	1	1	0	0	0	0	1	1	1	i	s	s	s	s	s	P	P	i	0	i	i	i	i	i	i	i	i	i	i	i	i	-	if (Rs<=#0) jump:nt #r13:2
0	1	1	0	0	0	0	1	1	1	i	s	s	s	s	s	P	P	i	1	i	i	i	i	i	i	i	i	i	i	i	i	-	if (Rs<=#0) jump:t #r13:2

Field name	Description
sm	Supervisor mode only
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s

## Transfer and jump

Move an unsigned immediate or register value into a destination register and unconditionally jump. In assembly syntax, this instruction appears as two instructions in the packet, a transfer and a separate jump. The assembler may convert adjacent transfer and jump instructions into compound transfer-jump form.

Syntax	Behavior
Rd=#U6 ; jump #r9:2	<pre>apply_extension(#r); #r=#r &amp; ~PCALIGN_MASK; Rd=#U; PC=PC+#r;</pre>
Rd=Rs ; jump #r9:2	<pre>apply_extension(#r); #r=#r &amp; ~PCALIGN_MASK; Rd=Rs; PC=PC+#r;</pre>

### Class: J (slots 2,3)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0				
ICLASS												d4				Parse																			
0	0	0	1	0	1	1	0	-	-	i	i	d	d	d	d	P	P	I	I	I	I	I	I	I	I	i	i	i	i	i	i	i	-	Rd=#U6 ; jump #r9:2	
ICLASS												s4				Parse				d4															
0	0	0	1	0	1	1	1	-	-	i	i	s	s	s	s	P	P	-	-	d	d	d	d	d	d	i	i	i	i	i	i	i	-	Rd=Rs ; jump #r9:2	

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d4	Field to encode register d
s4	Field to encode register s

## 11.5 LD

The LD instruction class includes load instructions, which are used to load values into registers.

LD instructions are executable on slot 0 and slot 1.

### Load doubleword

Load a 64-bit doubleword from memory and place in a destination register pair.

Syntax	Behavior
<code>Rdd=memd (Re=#U6)</code>	<code>apply_extension (#U);</code> <code>EA=#U;</code> <code>Rdd = *EA;</code> <code>Re=#U;</code>
<code>Rdd=memd (Rs+#s11:3)</code>	<code>apply_extension (#s);</code> <code>EA=Rs+#s;</code> <code>Rdd = *EA;</code>
<code>Rdd=memd (Rs+Rt&lt;&lt;#u2)</code>	<code>EA=Rs+ (Rt&lt;&lt;#u);</code> <code>Rdd = *EA;</code>
<code>Rdd=memd (Rt&lt;&lt;#u2+#U6)</code>	<code>apply_extension (#U);</code> <code>EA=#U+ (Rt&lt;&lt;#u);</code> <code>Rdd = *EA;</code>
<code>Rdd=memd (Rx++#s4:3)</code>	<code>EA=Rx;</code> <code>Rx=Rx+#s;</code> <code>Rdd = *EA;</code>
<code>Rdd=memd (Rx++#s4:3:circ (Mu))</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, #s, MuV);</code> <code>Rdd = *EA;</code>
<code>Rdd=memd (Rx++I:circ (Mu))</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, I&lt;&lt;3, MuV);</code> <code>Rdd = *EA;</code>
<code>Rdd=memd (Rx++Mu)</code>	<code>EA=Rx;</code> <code>Rx=Rx+MuV;</code> <code>Rdd = *EA;</code>
<code>Rdd=memd (Rx++Mu:brev)</code>	<code>EA=Rx.h[1]   brev (Rx.h[0]);</code> <code>Rx=Rx+MuV;</code> <code>Rdd = *EA;</code>
<code>Rdd=memd (gp+#u16:3)</code>	<code>apply_extension (#u);</code> <code>EA=(Constant_extended ? (0) : GP)+#u;</code> <code>Rdd = *EA;</code>

### Class: LD (slots 0,1)

#### Intrinsics

Rdd=memd(Rx++#s4:3:circ(Mu)) Word32 Q6\_R\_memd\_IM\_circ(void\*\* StartAddress, Word32 Is4\_3, Word32 Mu, void\* BaseAddress)

Rdd=memd(Rx++I:circ(Mu)) Word32 Q6\_R\_memd\_M\_circ(void\*\* StartAddress, Word32 Mu, void\* BaseAddress)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS											s5					Parse		t5					d5										
0	0	1	1	1	0	1	0	1	1	0	s	s	s	s	s	P	P	i	t	t	t	t	t	i	-	-	d	d	d	d	d	Rdd=memd(Rs+Rt<<#u2)	
ICLASS				Type							UN						Parse							d5									
0	1	0	0	1	i	i	1	1	1	0	i	i	i	i	i	P	P	i	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rdd=memd(gp+#u16:3)
ICLASS				Amode			Type			UN	s5					Parse							d5										
1	0	0	1	0	i	i	1	1	1	0	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rdd=memd(Rs+#s11:3)
ICLASS				Amode			Type			UN	x5					Parse		u1						d5									
1	0	0	1	1	0	0	1	1	1	0	x	x	x	x	x	P	P	u	0	-	-	0	i	i	i	i	d	d	d	d	d	Rdd=memd(Rx++#s4:3:circ(Mu))	
1	0	0	1	1	0	0	1	1	1	0	x	x	x	x	x	P	P	u	0	-	-	1	-	0	-	-	d	d	d	d	d	Rdd=memd(Rx++I:circ(Mu))	
ICLASS				Amode			Type			UN	e5					Parse							d5										
1	0	0	1	1	0	1	1	1	1	0	e	e	e	e	e	P	P	0	1	l	l	l	l	-	l	l	d	d	d	d	d	Rdd=memd(Re+#U6)	
ICLASS				Amode			Type			UN	x5					Parse							d5										
1	0	0	1	1	0	1	1	1	1	0	x	x	x	x	x	P	P	0	0	-	-	-	i	i	i	i	d	d	d	d	d	Rdd=memd(Rx++#s4:3)	
ICLASS				Amode			Type			UN	t5					Parse							d5										
1	0	0	1	1	1	0	1	1	1	0	t	t	t	t	t	P	P	i	1	l	l	l	l	i	l	l	d	d	d	d	d	Rdd=memd(Rt<<#u2+#U6)	
ICLASS				Amode			Type			UN	x5					Parse		u1						d5									
1	0	0	1	1	1	0	1	1	1	0	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	d	d	d	d	d	Rdd=memd(Rx++Mu)	
1	0	0	1	1	1	1	1	1	1	0	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	d	d	d	d	d	Rdd=memd(Rx++Mu:brev)	

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
Parse	Packet/loop parse bits
d5	Field to encode register d
e5	Field to encode register e
s5	Field to encode register s
t5	Field to encode register t
u1	Field to encode register u
x5	Field to encode register x

## Load-acquire doubleword

Load a 64-bit doubleword from memory and place in a destination register pair. The load-acquire memory operation is observed before any following memory operations (in program order) have been observed at the local point of serialization. A different order may be observed at the global point of serialization. (see Ordering and Synchronization).

### Syntax

```
Rdd=memd_aq (Rs)
```

### Behavior

```
EA=Rs;
Rdd = *EA
```

**Class: LD (slots 0)**

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS				Amode			Type		U N	s5					Parse												d5						
1	0	0	1	0	0	1	0	0	0	0	s	s	s	s	s	P	P	0	1	1	-	-	-	0	0	0	d	d	d	d	d		Rdd=memd_aq(Rs)

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s



## Load doubleword conditionally

Load a 64-bit doubleword from memory and place in a destination register pair.

This instruction is conditional based on a predicate value. If the predicate is true, the instruction is performed, otherwise it is treated as a NOP.

Syntax	Behavior
<code>if ([!]Pt[.new]) Rdd=memd(#u6)</code>	<pre> apply_extension(#u); EA=#u; if ([!]Pt[.new][0]) {     Rdd = *EA; } else {     NOP; }                     </pre>
<code>if ([!]Pt[.new]) Rdd=memd(Rs+#u6:3)</code>	<pre> apply_extension(#u); EA=Rs+#u; if ([!]Pt[.new][0]) {     Rdd = *EA; } else {     NOP; }                     </pre>
<code>if ([!]Pt[.new]) Rdd=memd(Rx++#s4:3)</code>	<pre> EA=Rx; if ([!]Pt[.new][0]) {     Rx=Rx+#s;     Rdd = *EA; } else {     NOP; }                     </pre>
<code>if ([!]Pv[.new]) Rdd=memd(Rs+Rt&lt;&lt;#u2)</code>	<pre> EA=Rs+(Rt&lt;&lt;#u); if ([!]Pv[.new][0]) {     Rdd = *EA; } else {     NOP; }                     </pre>

**Class: LD (slots 0,1)**

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS											s5					Parse		t5					d5										
0	0	1	1	0	0	0	0	1	1	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	v	v	d	d	d	d	d	if (Pv) Rdd=memd(Rs+Rt<<#u2)
0	0	1	1	0	0	0	1	1	1	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	v	v	d	d	d	d	d	if (!Pv) Rdd=memd(Rs+Rt<<#u2)
0	0	1	1	0	0	1	0	1	1	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	v	v	d	d	d	d	d	if (Pv.new) Rdd=memd(Rs+Rt<<#u2)
0	0	1	1	0	0	1	1	1	1	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	v	v	d	d	d	d	d	if (!Pv.new) Rdd=memd(Rs+Rt<<#u2)
ICLASS			Se	Pr		Type	UN	s5					Parse		t2					d5													

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
0	1	0	0	0	0	0	1	1	1	0	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d		if (Pt) Rdd=memd(Rs+#u6:3)
0	1	0	0	0	0	1	1	1	1	0	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d		if (Pt.new) Rdd=memd(Rs+#u6:3)
0	1	0	0	0	1	0	1	1	1	0	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d		if (!Pt) Rdd=memd(Rs+#u6:3)
0	1	0	0	0	1	1	1	1	1	0	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d		if (!Pt.new) Rdd=memd(Rs+#u6:3)
ICLASS			Amode			Type			UN	x5					Parse		t2					d5											
1	0	0	1	1	0	1	1	1	1	0	x	x	x	x	x	P	P	1	0	0	t	t	i	i	i	i	d	d	d	d	d		if (Pt) Rdd=memd(Rx++#s4:3)
1	0	0	1	1	0	1	1	1	1	0	x	x	x	x	x	P	P	1	0	1	t	t	i	i	i	i	d	d	d	d	d		if (!Pt) Rdd=memd(Rx++#s4:3)
1	0	0	1	1	0	1	1	1	1	0	x	x	x	x	x	P	P	1	1	0	t	t	i	i	i	i	d	d	d	d	d		if (Pt.new) Rdd=memd(Rx++#s4:3)
1	0	0	1	1	0	1	1	1	1	0	x	x	x	x	x	P	P	1	1	1	t	t	i	i	i	i	d	d	d	d	d		if (!Pt.new) Rdd=memd(Rx++#s4:3)
ICLASS			Amode			Type			UN						Parse		t2					d5											
1	0	0	1	1	1	1	1	1	1	0	i	i	i	i	i	P	P	1	0	0	t	t	i	1	-	-	d	d	d	d	d		if (Pt) Rdd=memd(#u6)
1	0	0	1	1	1	1	1	1	1	0	i	i	i	i	i	P	P	1	0	1	t	t	i	1	-	-	d	d	d	d	d		if (!Pt) Rdd=memd(#u6)
1	0	0	1	1	1	1	1	1	1	0	i	i	i	i	i	P	P	1	1	0	t	t	i	1	-	-	d	d	d	d	d		if (Pt.new) Rdd=memd(#u6)
1	0	0	1	1	1	1	1	1	1	0	i	i	i	i	i	P	P	1	1	1	t	t	i	1	-	-	d	d	d	d	d		if (!Pt.new) Rdd=memd(#u6)

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
PredNew	PredNew
Sense	Sense
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t2	Field to encode register t
t5	Field to encode register t
v2	Field to encode register v
x5	Field to encode register x

## Load byte

Load a signed byte from memory. The byte at the effective address in memory is placed in the least-significant 8 bits of the destination register. The destination register is then sign-extended from 8 bits to 32.

Syntax	Behavior
<code>Rd=memb (Re=#U6)</code>	<code>apply_extension (#U);</code> <code>EA=#U;</code> <code>Rd = *EA;</code> <code>Re=#U;</code>
<code>Rd=memb (Rs+#s11:0)</code>	<code>apply_extension (#s);</code> <code>EA=Rs+#s;</code> <code>Rd = *EA;</code>
<code>Rd=memb (Rs+Rt&lt;&lt;#u2)</code>	<code>EA=Rs+ (Rt&lt;&lt;#u);</code> <code>Rd = *EA;</code>
<code>Rd=memb (Rt&lt;&lt;#u2+#U6)</code>	<code>apply_extension (#U);</code> <code>EA=#U+ (Rt&lt;&lt;#u);</code> <code>Rd = *EA;</code>
<code>Rd=memb (Rx++#s4:0)</code>	<code>EA=Rx;</code> <code>Rx=Rx+#s;</code> <code>Rd = *EA;</code>
<code>Rd=memb (Rx++#s4:0:circ (Mu))</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, #s, MuV);</code> <code>Rd = *EA;</code>
<code>Rd=memb (Rx++I:circ (Mu))</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, I&lt;&lt;0, MuV);</code> <code>Rd = *EA;</code>
<code>Rd=memb (Rx++Mu)</code>	<code>EA=Rx;</code> <code>Rx=Rx+MuV;</code> <code>Rd = *EA;</code>
<code>Rd=memb (Rx++Mu:brev)</code>	<code>EA=Rx.h[1]   brev (Rx.h[0]);</code> <code>Rx=Rx+MuV;</code> <code>Rd = *EA;</code>
<code>Rd=memb (gp+#u16:0)</code>	<code>apply_extension (#u);</code> <code>EA=(Constant_extended ? (0) : GP)+#u;</code> <code>Rd = *EA;</code>

### Class: LD (slots 0,1)

#### Intrinsics

<code>Rd=memb (Rx++#s4:0:circ (Mu))</code>	<code>Word32 Q6_R_memb_IM_circ (void** StartAddress, Word32 Is4_0, Word32 Mu, void* BaseAddress)</code>
<code>Rd=memb (Rx++I:circ (Mu))</code>	<code>Word32 Q6_R_memb_M_circ (void** StartAddress, Word32 Mu, void* BaseAddress)</code>

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS											s5					Parse			t5					d5									
0	0	1	1	1	0	1	0	0	0	0	s	s	s	s	s	P	P	i	t	t	t	t	t	i	-	-	d	d	d	d	d	Rd=memb(Rs+Rt<<#u2)	
ICLASS				Type							UN						Parse								d5								
0	1	0	0	1	i	i	1	0	0	0	i	i	i	i	i	P	P	i	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=memb(gp+#u16:0)
ICLASS				Amode			Type			UN	s5					Parse								d5									
1	0	0	1	0	i	i	1	0	0	0	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=memb(Rs+#s11:0)
ICLASS				Amode			Type			UN	x5					Parse			u1						d5								
1	0	0	1	1	0	0	1	0	0	0	x	x	x	x	x	P	P	u	0	-	-	0	i	i	i	i	d	d	d	d	d	Rd=memb(Rx++#s4:0:circ(Mu))	
1	0	0	1	1	0	0	1	0	0	0	x	x	x	x	x	P	P	u	0	-	-	1	-	0	-	-	d	d	d	d	d	Rd=memb(Rx++l:circ(Mu))	
ICLASS				Amode			Type			UN	e5					Parse								d5									
1	0	0	1	1	0	1	1	0	0	0	e	e	e	e	e	P	P	0	1	l	l	l	l	-	l	l	d	d	d	d	d	Rd=memb(Re=#U6)	
ICLASS				Amode			Type			UN	x5					Parse								d5									
1	0	0	1	1	0	1	1	0	0	0	x	x	x	x	x	P	P	0	0	-	-	-	i	i	i	i	d	d	d	d	d	Rd=memb(Rx++#s4:0)	
ICLASS				Amode			Type			UN	t5					Parse								d5									
1	0	0	1	1	1	0	1	0	0	0	t	t	t	t	t	P	P	i	1	l	l	l	l	i	l	l	d	d	d	d	d	Rd=memb(Rt<<#u2+#U6)	
ICLASS				Amode			Type			UN	x5					Parse			u1						d5								
1	0	0	1	1	1	0	1	0	0	0	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	d	d	d	d	d	Rd=memb(Rx++Mu)	
1	0	0	1	1	1	1	1	0	0	0	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	d	d	d	d	d	Rd=memb(Rx++Mu:brev)	

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
Parse	Packet/loop parse bits
d5	Field to encode register d
e5	Field to encode register e
s5	Field to encode register s
t5	Field to encode register t
u1	Field to encode register u
x5	Field to encode register x

## Load byte conditionally

Load a signed byte from memory. The byte at the effective address in memory is placed in the least-significant 8 bits of the destination register. The destination register is then sign-extended from 8 bits to 32.

This instruction is conditional based on a predicate value. If the predicate is true, the instruction is performed, otherwise it is treated as a NOP.

Syntax	Behavior
<code>if ([!]Pt[.new]) Rd=memb(#u6)</code>	<pre> apply_extension(#u); EA=#u; if ([!]Pt[.new][0]) {     Rd = *EA; } else {     NOP; } </pre>
<code>if ([!]Pt[.new]) Rd=memb(Rs+#u6:0)</code>	<pre> apply_extension(#u); EA=Rs+#u; if ([!]Pt[.new][0]) {     Rd = *EA; } else {     NOP; } </pre>
<code>if ([!]Pt[.new]) Rd=memb(Rx++#s4:0)</code>	<pre> EA=Rx; if ([!]Pt[.new][0]) {     Rx=Rx+#s;     Rd = *EA; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) Rd=memb(Rs+Rt&lt;&lt;#u2)</code>	<pre> EA=Rs+(Rt&lt;&lt;#u); if ([!]Pv[.new][0]) {     Rd = *EA; } else {     NOP; } </pre>

**Class: LD (slots 0,1)**

### Encoding

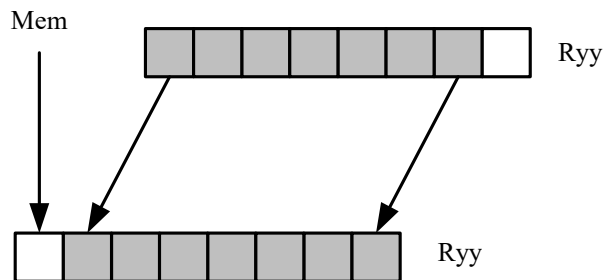
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0			
ICLASS											s5				Parse				t5				d5											
0	0	1	1	0	0	0	0	0	0	0	s	s	s	s	s	P	P	i	t	t	t	t	t	i	v	v	d	d	d	d	d	d	if (Pv) Rd=memb(Rs+Rt<<#u2)	
0	0	1	1	0	0	0	1	0	0	0	s	s	s	s	s	P	P	i	t	t	t	t	t	i	v	v	d	d	d	d	d	d	d	if (!Pv) Rd=memb(Rs+Rt<<#u2)
0	0	1	1	0	0	1	0	0	0	0	s	s	s	s	s	P	P	i	t	t	t	t	t	i	v	v	d	d	d	d	d	d	d	if (Pv.new) Rd=memb(Rs+Rt<<#u2)
0	0	1	1	0	0	1	1	0	0	0	s	s	s	s	s	P	P	i	t	t	t	t	t	i	v	v	d	d	d	d	d	d	d	if (!Pv.new) Rd=memb(Rs+Rt<<#u2)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS					Se ns e	Pr ed Ne w	Type	UN	s5					Parse		t2					d5											
0	1	0	0	0	0	0	1	0	0	0	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d	if (Pt) Rd=memb(Rs+#u6:0)
0	1	0	0	0	0	1	1	0	0	0	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d	if (Pt.new) Rd=memb(Rs+#u6:0)
0	1	0	0	0	1	0	1	0	0	0	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d	if (!Pt) Rd=memb(Rs+#u6:0)
0	1	0	0	0	1	1	1	0	0	0	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d	if (!Pt.new) Rd=memb(Rs+#u6:0)
ICLASS				Amode		Type	UN	x5					Parse		t2					d5												
1	0	0	1	1	0	1	1	0	0	0	x	x	x	x	x	P	P	1	0	0	t	t	i	i	i	i	d	d	d	d	d	if (Pt) Rd=memb(Rx++#s4:0)
1	0	0	1	1	0	1	1	0	0	0	x	x	x	x	x	P	P	1	0	1	t	t	i	i	i	i	d	d	d	d	d	if (!Pt) Rd=memb(Rx++#s4:0)
1	0	0	1	1	0	1	1	0	0	0	x	x	x	x	x	P	P	1	1	0	t	t	i	i	i	i	d	d	d	d	d	if (Pt.new) Rd=memb(Rx++#s4:0)
1	0	0	1	1	0	1	1	0	0	0	x	x	x	x	x	P	P	1	1	1	t	t	i	i	i	i	d	d	d	d	d	if (!Pt.new) Rd=memb(Rx++#s4:0)
ICLASS				Amode		Type	UN						Parse		t2					d5												
1	0	0	1	1	1	1	1	0	0	0	i	i	i	i	i	P	P	1	0	0	t	t	i	1	-	-	d	d	d	d	d	if (Pt) Rd=memb(#u6)
1	0	0	1	1	1	1	1	0	0	0	i	i	i	i	i	P	P	1	0	1	t	t	i	1	-	-	d	d	d	d	d	if (!Pt) Rd=memb(#u6)
1	0	0	1	1	1	1	1	0	0	0	i	i	i	i	i	P	P	1	1	0	t	t	i	1	-	-	d	d	d	d	d	if (Pt.new) Rd=memb(#u6)
1	0	0	1	1	1	1	1	0	0	0	i	i	i	i	i	P	P	1	1	1	t	t	i	1	-	-	d	d	d	d	d	if (!Pt.new) Rd=memb(#u6)

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
PredNew	PredNew
Sense	Sense
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t2	Field to encode register t
t5	Field to encode register t
v2	Field to encode register v
x5	Field to encode register x

## Load byte into shifted vector

Shift a 64-bit vector right by one byte. Insert a byte from memory into the vacated upper byte of the vector.



Syntax	Behavior
<code>Ryy=memb_fifo (Re=#U6)</code>	<pre> apply_extension(#U); EA=#U; {     tmpV = *EA;     Ryy = (((size8u_t)Ryy)&gt;&gt;8)   (tmpV&lt;&lt;56); } Re=#U; </pre>
<code>Ryy=memb_fifo (Rs)</code>	Assembler mapped to: " <code>Ryy=memb_fifo</code> " " <code>(Rs+#0)</code> "
<code>Ryy=memb_fifo (Rs+#s11:0)</code>	<pre> apply_extension(#s); EA=Rs+#s; {     tmpV = *EA;     Ryy = (((size8u_t)Ryy)&gt;&gt;8)   (tmpV&lt;&lt;56); } </pre>
<code>Ryy=memb_fifo (Rt&lt;&lt;#u2+#U6)</code>	<pre> apply_extension(#U); EA=#U+(Rt&lt;&lt;#u); {     tmpV = *EA;     Ryy = (((size8u_t)Ryy)&gt;&gt;8)   (tmpV&lt;&lt;56); } </pre>
<code>Ryy=memb_fifo (Rx++#s4:0)</code>	<pre> EA=Rx; Rx=Rx+#s; {     tmpV = *EA;     Ryy = (((size8u_t)Ryy)&gt;&gt;8)   (tmpV&lt;&lt;56); } </pre>
<code>Ryy=memb_fifo (Rx++#s4:0:circ (Mu))</code>	<pre> EA=Rx; Rx=Rx=circ_add(Rx, #s, MuV); {     tmpV = *EA;     Ryy = (((size8u_t)Ryy)&gt;&gt;8)   (tmpV&lt;&lt;56); } </pre>

Syntax	Behavior
<code>Ryy=memb_fifo(Rx++I:circ(Mu))</code>	<pre>EA=Rx; Rx=Rx=circ_add(Rx,I&lt;&lt;0,MuV); {     tmpV = *EA;     Ryy = (((size8u_t)Ryy)&gt;&gt;8) (tmpV&lt;&lt;56); }</pre>
<code>Ryy=memb_fifo(Rx++Mu)</code>	<pre>EA=Rx; Rx=Rx+MuV; {     tmpV = *EA;     Ryy = (((size8u_t)Ryy)&gt;&gt;8) (tmpV&lt;&lt;56); }</pre>
<code>Ryy=memb_fifo(Rx++Mu:brev)</code>	<pre>EA=Rx.h[1]   brev(Rx.h[0]); Rx=Rx+MuV; {     tmpV = *EA;     Ryy = (((size8u_t)Ryy)&gt;&gt;8) (tmpV&lt;&lt;56); }</pre>

**Class: LD (slots 0,1)****Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS		Amode		Type		UN	s5					Parse		y5																			
1	0	0	1	0	i	i	0	1	0	0	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	y	y	y	y	y	Ryy=memb_fifo(Rs+#s11:0)	
ICLASS		Amode		Type		UN	x5					Parse		u1	y5																		
1	0	0	1	1	0	0	0	1	0	0	x	x	x	x	x	P	P	u	0	-	-	0	i	i	i	i	y	y	y	y	y	Ryy=memb_fifo(Rx++#s4:0:circ(Mu))	
1	0	0	1	1	0	0	0	1	0	0	x	x	x	x	x	P	P	u	0	-	-	1	-	0	-	-	y	y	y	y	y	Ryy=memb_fifo(Rx++I:circ(Mu))	
ICLASS		Amode		Type		UN	e5					Parse		y5																			
1	0	0	1	1	0	1	0	1	0	0	e	e	e	e	e	P	P	0	1	l	l	l	l	l	-	l	l	y	y	y	y	y	Ryy=memb_fifo(Re=#U6)
ICLASS		Amode		Type		UN	x5					Parse		y5																			
1	0	0	1	1	0	1	0	1	0	0	x	x	x	x	x	P	P	0	0	-	-	-	i	i	i	i	y	y	y	y	y	Ryy=memb_fifo(Rx++#s4:0)	
ICLASS		Amode		Type		UN	t5					Parse		y5																			
1	0	0	1	1	1	0	0	1	0	0	t	t	t	t	t	P	P	i	1	l	l	l	l	l	i	l	l	y	y	y	y	y	Ryy=memb_fifo(Rt<<#u2+#U6)
ICLASS		Amode		Type		UN	x5					Parse		u1	y5																		
1	0	0	1	1	1	0	0	1	0	0	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	y	y	y	y	y	Ryy=memb_fifo(Rx++Mu)	
1	0	0	1	1	1	1	0	1	0	0	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	y	y	y	y	y	Ryy=memb_fifo(Rx++Mu:brev)	

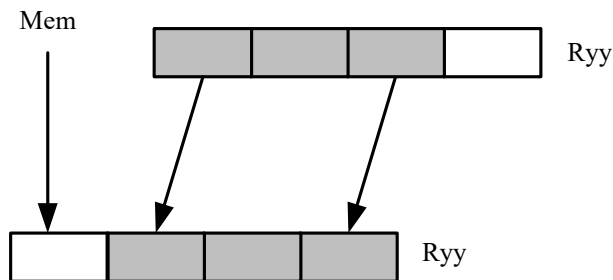
Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type



<b>Field name</b>	<b>Description</b>
UN	Unsigned
Parse	Packet/loop parse bits
e5	Field to encode register e
s5	Field to encode register s
t5	Field to encode register t
u1	Field to encode register u
x5	Field to encode register x
y5	Field to encode register y

## Load half into shifted vector

Shift a 64-bit vector right by one halfword. Insert a halfword from memory into the vacated upper halfword of the vector.



Syntax	Behavior
<code>Ryy=memh_fifo (Re=#U6)</code>	<pre> apply_extension(#U); EA=#U; {     tmpV = *EA;     Ryy = (((size8u_t)Ryy)&gt;&gt;16)   (tmpV&lt;&lt;48); } Re=#U; </pre>
<code>Ryy=memh_fifo (Rs)</code>	Assembler mapped to: "Ryy=memh_fifo" (Rs+#0) "
<code>Ryy=memh_fifo (Rs+#s11:1)</code>	<pre> apply_extension(#s); EA=Rs+#s; {     tmpV = *EA;     Ryy = (((size8u_t)Ryy)&gt;&gt;16)   (tmpV&lt;&lt;48); } </pre>
<code>Ryy=memh_fifo (Rt&lt;&lt;#u2+#U6)</code>	<pre> apply_extension(#U); EA=#U+(Rt&lt;&lt;#u); {     tmpV = *EA;     Ryy = (((size8u_t)Ryy)&gt;&gt;16)   (tmpV&lt;&lt;48); } </pre>
<code>Ryy=memh_fifo (Rx++#s4:1)</code>	<pre> EA=Rx; Rx=Rx+#s; {     tmpV = *EA;     Ryy = (((size8u_t)Ryy)&gt;&gt;16)   (tmpV&lt;&lt;48); } </pre>
<code>Ryy=memh_fifo (Rx++#s4:1:circ (Mu))</code>	<pre> EA=Rx; Rx=Rx=circ_add(Rx, #s, MuV); {     tmpV = *EA;     Ryy = (((size8u_t)Ryy)&gt;&gt;16)   (tmpV&lt;&lt;48); } </pre>

Syntax	Behavior
<code>Ryy=memh_fifo(Rx++I:circ(Mu))</code>	<pre>EA=Rx; Rx=Rx=circ_add(Rx,I&lt;&lt;1,MuV); {     tmpV = *EA;     Ryy = (((size8u_t)Ryy)&gt;&gt;16) (tmpV&lt;&lt;48); }</pre>
<code>Ryy=memh_fifo(Rx++Mu)</code>	<pre>EA=Rx; Rx=Rx+MuV; {     tmpV = *EA;     Ryy = (((size8u_t)Ryy)&gt;&gt;16) (tmpV&lt;&lt;48); }</pre>
<code>Ryy=memh_fifo(Rx++Mu:brev)</code>	<pre>EA=Rx.h[1]   brev(Rx.h[0]); Rx=Rx+MuV; {     tmpV = *EA;     Ryy = (((size8u_t)Ryy)&gt;&gt;16) (tmpV&lt;&lt;48); }</pre>

**Class: LD (slots 0,1)****Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		Amode			Type		U	s5					Parse		y5																	
1	0	0	1	0	i	i	0	0	1	0	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	y	y	y	y	y	Ryy=memh_fifo(Rs+#s11:1)
ICLASS		Amode			Type		U	x5					Parse		u1	y5																
1	0	0	1	1	0	0	0	0	1	0	x	x	x	x	x	P	P	u	0	-	-	0	i	i	i	i	y	y	y	y	y	Ryy=memh_fifo(Rx++#s4:1:circ(Mu))
1	0	0	1	1	0	0	0	0	1	0	x	x	x	x	x	P	P	u	0	-	-	1	-	0	-	-	y	y	y	y	y	Ryy=memh_fifo(Rx++I:circ(Mu))
ICLASS		Amode			Type		U	e5					Parse		y5																	
1	0	0	1	1	0	1	0	0	1	0	e	e	e	e	e	P	P	0	1	l	l	l	l	-	l	l	y	y	y	y	y	Ryy=memh_fifo(Re=#U6)
ICLASS		Amode			Type		U	x5					Parse		y5																	
1	0	0	1	1	0	1	0	0	1	0	x	x	x	x	x	P	P	0	0	-	-	-	i	i	i	i	y	y	y	y	y	Ryy=memh_fifo(Rx++#s4:1)
ICLASS		Amode			Type		U	t5					Parse		y5																	
1	0	0	1	1	1	0	0	0	1	0	t	t	t	t	t	P	P	i	1	l	l	l	l	i	l	l	y	y	y	y	y	Ryy=memh_fifo(Rt<<#u2+#U6)
ICLASS		Amode			Type		U	x5					Parse		u1	y5																
1	0	0	1	1	1	0	0	0	1	0	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	y	y	y	y	y	Ryy=memh_fifo(Rx++Mu)
1	0	0	1	1	1	1	0	0	1	0	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	y	y	y	y	y	Ryy=memh_fifo(Rx++Mu:brev)

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type

<b>Field name</b>	<b>Description</b>
UN	Unsigned
Parse	Packet/loop parse bits
e5	Field to encode register e
s5	Field to encode register s
t5	Field to encode register t
u1	Field to encode register u
x5	Field to encode register x
y5	Field to encode register y

## Load halfword

Load a signed halfword from memory. The 16-bit halfword at the effective address in memory is placed in the least-significant 16 bits of the destination register. The destination register is then sign-extended from 16 bits to 32.

Syntax	Behavior
<code>Rd=memh (Re=#U6)</code>	<code>apply_extension (#U);</code> <code>EA=#U;</code> <code>Rd = *EA;</code> <code>Re=#U;</code>
<code>Rd=memh (Rs+#s11:1)</code>	<code>apply_extension (#s);</code> <code>EA=Rs+#s;</code> <code>Rd = *EA;</code>
<code>Rd=memh (Rs+Rt&lt;&lt;#u2)</code>	<code>EA=Rs+ (Rt&lt;&lt;#u);</code> <code>Rd = *EA;</code>
<code>Rd=memh (Rt&lt;&lt;#u2+#U6)</code>	<code>apply_extension (#U);</code> <code>EA=#U+ (Rt&lt;&lt;#u);</code> <code>Rd = *EA;</code>
<code>Rd=memh (Rx++#s4:1)</code>	<code>EA=Rx;</code> <code>Rx=Rx+#s;</code> <code>Rd = *EA;</code>
<code>Rd=memh (Rx++#s4:1:circ (Mu))</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, #s, MuV);</code> <code>Rd = *EA;</code>
<code>Rd=memh (Rx++I:circ (Mu))</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, I&lt;&lt;1, MuV);</code> <code>Rd = *EA;</code>
<code>Rd=memh (Rx++Mu)</code>	<code>EA=Rx;</code> <code>Rx=Rx+MuV;</code> <code>Rd = *EA;</code>
<code>Rd=memh (Rx++Mu:brev)</code>	<code>EA=Rx.h[1]   brev (Rx.h[0]);</code> <code>Rx=Rx+MuV;</code> <code>Rd = *EA;</code>
<code>Rd=memh (gp+#u16:1)</code>	<code>apply_extension (#u);</code> <code>EA=(Constant_extended ? (0) : GP)+#u;</code> <code>Rd = *EA;</code>

### Class: LD (slots 0,1)

#### Intrinsics

`Rd=memh (Rx++#s4:1:circ (Mu))` `Word32 Q6_R_memh_IM_circ (void** StartAddress, Word32 Is4_1, Word32 Mu, void* BaseAddress)`

`Rd=memh (Rx++I:circ (Mu))` `Word32 Q6_R_memh_M_circ (void** StartAddress, Word32 Mu, void* BaseAddress)`

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS											s5					Parse		t5					d5										
0	0	1	1	1	0	1	0	0	1	0	s	s	s	s	s	P	P	i	t	t	t	t	t	i	-	-	d	d	d	d	d	Rd=memh(Rs+Rt<<#u2)	
ICLASS				Type				UN								Parse							d5										
0	1	0	0	1	i	i	1	0	1	0	i	i	i	i	i	P	P	i	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=memh(gp+#u16:1)
ICLASS				Amode		Type		UN	s5					Parse							d5												
1	0	0	1	0	i	i	1	0	1	0	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=memh(Rs+#s11:1)
ICLASS				Amode		Type		UN	x5					Parse		u1						d5											
1	0	0	1	1	0	0	1	0	1	0	x	x	x	x	x	P	P	u	0	-	-	0	i	i	i	i	d	d	d	d	d	Rd=memh(Rx++#s4:1:circ(Mu))	
1	0	0	1	1	0	0	1	0	1	0	x	x	x	x	x	P	P	u	0	-	-	1	-	0	-	-	d	d	d	d	d	Rd=memh(Rx++l:circ(Mu))	
ICLASS				Amode		Type		UN	e5					Parse							d5												
1	0	0	1	1	0	1	1	0	1	0	e	e	e	e	e	P	P	0	1	l	l	l	l	-	l	l	d	d	d	d	d	Rd=memh(Re=#U6)	
ICLASS				Amode		Type		UN	x5					Parse							d5												
1	0	0	1	1	0	1	1	0	1	0	x	x	x	x	x	P	P	0	0	-	-	-	i	i	i	i	d	d	d	d	d	Rd=memh(Rx++#s4:1)	
ICLASS				Amode		Type		UN	t5					Parse							d5												
1	0	0	1	1	1	0	1	0	1	0	t	t	t	t	t	P	P	i	1	l	l	l	l	i	l	l	d	d	d	d	d	Rd=memh(Rt<<#u2+#U6)	
ICLASS				Amode		Type		UN	x5					Parse		u1						d5											
1	0	0	1	1	1	0	1	0	1	0	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	d	d	d	d	d	Rd=memh(Rx++Mu)	
1	0	0	1	1	1	1	1	0	1	0	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	d	d	d	d	d	Rd=memh(Rx++Mu:brev)	

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
Parse	Packet/loop parse bits
d5	Field to encode register d
e5	Field to encode register e
s5	Field to encode register s
t5	Field to encode register t
u1	Field to encode register u
x5	Field to encode register x

## Load halfword conditionally

Load a signed halfword from memory. The 16-bit halfword at the effective address in memory is placed in the least-significant 16 bits of the destination register. The destination register is then sign-extended from 16 bits to 32.

This instruction is conditional based on a predicate value. If the predicate is true, the instruction is performed, otherwise it is treated as a NOP.

Syntax	Behavior
<code>if ([!]Pt[.new]) Rd=memh(#u6)</code>	<pre> apply_extension(#u); EA=#u; if ([!]Pt[.new][0]) {     Rd = *EA; } else {     NOP; } </pre>
<code>if ([!]Pt[.new]) Rd=memh(Rs+#u6:1)</code>	<pre> apply_extension(#u); EA=Rs+#u; if ([!]Pt[.new][0]) {     Rd = *EA; } else {     NOP; } </pre>
<code>if ([!]Pt[.new]) Rd=memh(Rx++#s4:1)</code>	<pre> EA=Rx; if ([!]Pt[.new][0]) {     Rx=Rx+#s;     Rd = *EA; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) Rd=memh(Rs+Rt&lt;&lt;#u2)</code>	<pre> EA=Rs+(Rt&lt;&lt;#u); if ([!]Pv[.new][0]) {     Rd = *EA; } else {     NOP; } </pre>

### Class: LD (slots 0,1)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS											s5				Parse		t5					d5											
0	0	1	1	0	0	0	0	0	1	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	v	v	d	d	d	d	d	if (Pv) Rd=memh(Rs+Rt<<#u2)
0	0	1	1	0	0	0	1	0	1	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	v	v	d	d	d	d	d	if (!Pv) Rd=memh(Rs+Rt<<#u2)
0	0	1	1	0	0	1	0	0	1	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	v	v	d	d	d	d	d	if (Pv.new) Rd=memh(Rs+Rt<<#u2)
0	0	1	1	0	0	1	1	0	1	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	v	v	d	d	d	d	d	if (!Pv.new) Rd=memh(Rs+Rt<<#u2)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS					Se ns e	Pr ed Ne w	Type	UN	s5					Parse		t2					d5											
0	1	0	0	0	0	0	1	0	1	0	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d	if (Pt) Rd=memh(Rs+#u6:1)
0	1	0	0	0	0	1	1	0	1	0	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d	if (Pt.new) Rd=memh(Rs+#u6:1)
0	1	0	0	0	1	0	1	0	1	0	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d	if (!Pt) Rd=memh(Rs+#u6:1)
0	1	0	0	0	1	1	1	0	1	0	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d	if (!Pt.new) Rd=memh(Rs+#u6:1)
ICLASS				Amode			Type	UN	x5					Parse		t2					d5											
1	0	0	1	1	0	1	1	0	1	0	x	x	x	x	x	P	P	1	0	0	t	t	i	i	i	i	d	d	d	d	d	if (Pt) Rd=memh(Rx++#s4:1)
1	0	0	1	1	0	1	1	0	1	0	x	x	x	x	x	P	P	1	0	1	t	t	i	i	i	i	d	d	d	d	d	if (!Pt) Rd=memh(Rx++#s4:1)
1	0	0	1	1	0	1	1	0	1	0	x	x	x	x	x	P	P	1	1	0	t	t	i	i	i	i	d	d	d	d	d	if (Pt.new) Rd=memh(Rx++#s4:1)
1	0	0	1	1	0	1	1	0	1	0	x	x	x	x	x	P	P	1	1	1	t	t	i	i	i	i	d	d	d	d	d	if (!Pt.new) Rd=memh(Rx++#s4:1)
ICLASS				Amode			Type	UN						Parse		t2					d5											
1	0	0	1	1	1	1	1	0	1	0	i	i	i	i	i	P	P	1	0	0	t	t	i	1	-	-	d	d	d	d	d	if (Pt) Rd=memh(#u6)
1	0	0	1	1	1	1	1	0	1	0	i	i	i	i	i	P	P	1	0	1	t	t	i	1	-	-	d	d	d	d	d	if (!Pt) Rd=memh(#u6)
1	0	0	1	1	1	1	1	0	1	0	i	i	i	i	i	P	P	1	1	0	t	t	i	1	-	-	d	d	d	d	d	if (Pt.new) Rd=memh(#u6)
1	0	0	1	1	1	1	1	0	1	0	i	i	i	i	i	P	P	1	1	1	t	t	i	1	-	-	d	d	d	d	d	if (!Pt.new) Rd=memh(#u6)

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
PredNew	PredNew
Sense	Sense
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t2	Field to encode register t
t5	Field to encode register t
v2	Field to encode register v
x5	Field to encode register x



## Memory copy

Copy  $Mu + 1$  (length) bytes from the address in  $Rt$  (source base) to the address in  $Rs$  (destination base). The source base, destination base, and length values must be aligned to the L2 cache-line size. Behavior is undefined for non-aligned values and for source or destination buffers partially in illegal space. The accesses by the memcopy instruction are noncoherent with the cache-hierarchy of the Q6.

In addition to normal translation exceptions, a coprocessor memory exception occurs if any of the following are true:

- Source or destination base address in illegal space
- Source or destination buffer crosses a page boundary
- Source base address is NOT in AXI space
- Destination base address is NOT in VTCM

This instruction is only available on cores with VTCM.

### Syntax

$Rdd = pmemcopy(Rx, Rtt)$

### Behavior

**Class: LD (slots 0,1)**

### Notes

- This is a solo instruction. It must not be grouped with other instructions in a packet.

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			Amode			Type			UN	t5					Parse		x5					d5										
1	0	0	1	1	0	0	1	1	1	1	t	t	t	t	t	P	P	0	x	x	x	x	x	0	0	0	d	d	d	d	d	Rdd=pmemcopy(Rx,Rtt)

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
Parse	Packet/loop parse bits
d5	Field to encode register d
t5	Field to encode register t
x5	Field to encode register x

## Piecemeal memory copy

Piecemeal memory copy.

Syntax	Behavior
Rd=movlen(Rs,Rtt)	
Rdd=linecpy(Rs,Rtt)	
Rdd=pmemcpy(Rx,Rtt)	

**Class: CR (slot 3)**

### Notes

- This is a solo instruction. It must not be grouped with other instructions in a packet.

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				sm							t5					Parse		s5					d5									
0	1	1	0	1	1	1	1	1	1	1	t	t	t	t	t	P	P	0	s	s	s	s	s	0	1	0	d	d	d	d	d	Rd=movlen(Rs,Rtt)
ICLASS				Amode		Type		UN	t5					Parse		s5					d5											
1	0	0	1	1	0	0	1	1	1	1	t	t	t	t	t	P	P	0	s	s	s	s	s	0	0	1	d	d	d	d	d	Rdd=linecpy(Rs,Rtt)
ICLASS				Amode		Type		UN	t5					Parse		x5					d5											
1	0	0	1	1	0	0	1	1	1	1	t	t	t	t	t	P	P	0	x	x	x	x	x	0	0	0	d	d	d	d	d	Rdd=pmemcpy(Rx,Rtt)

Field name	Description
sm	Supervisor mode only
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Load unsigned byte

Load an unsigned byte from memory. The byte at the effective address in memory is placed in the least-significant 8 bits of the destination register. The destination register is then zero-extended from 8 bits to 32.

Syntax	Behavior
Rd=memub (Re=#U6)	apply_extension (#U); EA=#U; Rd = *EA; Re=#U;
Rd=memub (Rs+#s11:0)	apply_extension (#s); EA=Rs+#s; Rd = *EA;
Rd=memub (Rs+Rt<<#u2)	EA=Rs+ (Rt<<#u); Rd = *EA;
Rd=memub (Rt<<#u2+#U6)	apply_extension (#U); EA=#U+ (Rt<<#u); Rd = *EA;
Rd=memub (Rx++#s4:0)	EA=Rx; Rx=Rx+#s; Rd = *EA;
Rd=memub (Rx++#s4:0:circ (Mu) )	EA=Rx; Rx=Rx=circ_add (Rx, #s, MuV); Rd = *EA;
Rd=memub (Rx++I:circ (Mu) )	EA=Rx; Rx=Rx=circ_add (Rx, I<<0, MuV); Rd = *EA;
Rd=memub (Rx++Mu)	EA=Rx; Rx=Rx+MuV; Rd = *EA;
Rd=memub (Rx++Mu:brev)	EA=Rx.h[1]   brev (Rx.h[0]); Rx=Rx+MuV; Rd = *EA;
Rd=memub (gp+#u16:0)	apply_extension (#u); EA=(Constant_extended ? (0) : GP)+#u; Rd = *EA;

### Class: LD (slots 0,1)

#### Intrinsics

Rd=memub (Rx++#s4:0:circ (Mu) )	Word32 Q6_R_memub_IM_circ (void** StartAddress, Word32 Is4_0, Word32 Mu, void* BaseAddress)
Rd=memub (Rx++I:circ (Mu) )	Word32 Q6_R_memub_M_circ (void** StartAddress, Word32 Mu, void* BaseAddress)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS										s5					Parse		t5					d5											
0	0	1	1	1	0	1	0	0	0	1	s	s	s	s	s	P	P	i	t	t	t	t	t	i	-	-	d	d	d	d	d	Rd=memub(Rs+Rt<<#u2)	
ICLASS										Type	UN	s5					Parse		t5					d5									
0	1	0	0	1	i	i	1	0	0	1	i	i	i	i	i	P	P	i	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=memub(gp+#u16:0)
ICLASS										Amode	Type	UN	s5					Parse		t5					d5								
1	0	0	1	0	i	i	1	0	0	1	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=memub(Rs+#s11:0)	
ICLASS										Amode	Type	UN	x5					Parse	u1	t5					d5								
1	0	0	1	1	0	0	1	0	0	1	x	x	x	x	x	P	P	u	0	-	-	0	i	i	i	i	d	d	d	d	d	Rd=memub(Rx++#s4:0:circ (Mu))	
1	0	0	1	1	0	0	1	0	0	1	x	x	x	x	x	P	P	u	0	-	-	1	-	0	-	-	d	d	d	d	d	Rd=memub(Rx++l:circ(Mu))	
ICLASS										Amode	Type	UN	e5					Parse		t5					d5								
1	0	0	1	1	0	1	1	0	0	1	e	e	e	e	e	P	P	0	1	l	l	l	l	-	l	l	d	d	d	d	d	Rd=memub(Re=#U6)	
ICLASS										Amode	Type	UN	x5					Parse		t5					d5								
1	0	0	1	1	0	1	1	0	0	1	x	x	x	x	x	P	P	0	0	-	-	-	i	i	i	i	d	d	d	d	d	Rd=memub(Rx++#s4:0)	
ICLASS										Amode	Type	UN	t5					Parse		t5					d5								
1	0	0	1	1	1	0	1	0	0	1	t	t	t	t	t	P	P	i	1	l	l	l	l	i	l	l	d	d	d	d	d	Rd=memub(Rt<<#u2+#U6)	
ICLASS										Amode	Type	UN	x5					Parse	u1	t5					d5								
1	0	0	1	1	1	0	1	0	0	1	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	d	d	d	d	d	Rd=memub(Rx++Mu)	
1	0	0	1	1	1	1	1	0	0	1	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	d	d	d	d	d	Rd=memub(Rx++Mu:brev)	

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
Parse	Packet/loop parse bits
d5	Field to encode register d
e5	Field to encode register e
s5	Field to encode register s
t5	Field to encode register t
u1	Field to encode register u
x5	Field to encode register x

## Load unsigned byte conditionally

Load an unsigned byte from memory. The byte at the effective address in memory is placed in the least-significant 8 bits of the destination register. The destination register is then zero-extended from 8 bits to 32.

This instruction is conditional based on a predicate value. If the predicate is true, the instruction is performed, otherwise it is treated as a NOP.

Syntax	Behavior
<pre>if ([!]Pt[.new]) Rd=memub (#u6)</pre>	<pre>apply_extension(#u); EA=#u; if ([!]Pt[.new][0]) {     Rd = *EA; } else {     NOP; }</pre>
<pre>if ([!]Pt[.new]) Rd=memub (Rs+#u6:0)</pre>	<pre>apply_extension(#u); EA=Rs+#u; if ([!]Pt[.new][0]) {     Rd = *EA; } else {     NOP; }</pre>
<pre>if ([!]Pt[.new]) Rd=memub (Rx++#s4:0)</pre>	<pre>EA=Rx; if ([!]Pt[.new][0]) {     Rx=Rx+#s;     Rd = *EA; } else {     NOP; }</pre>
<pre>if ([!]Pv[.new]) Rd=memub (Rs+Rt&lt;&lt;#u2)</pre>	<pre>EA=Rs+(Rt&lt;&lt;#u); if ([!]Pv[.new][0]) {     Rd = *EA; } else {     NOP; }</pre>

### Class: LD (slots 0,1)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0			
ICLASS											s5					Parse		t5					d5											
0	0	1	1	0	0	0	0	0	0	1	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	v	v	d	d	d	d	d	if (Pv) Rd=memub(Rs+Rt<<#u2)	
0	0	1	1	0	0	0	1	0	0	1	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	v	v	d	d	d	d	d	d	if (!Pv) Rd=memub(Rs+Rt<<#u2)
0	0	1	1	0	0	1	0	0	0	1	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	v	v	d	d	d	d	d	d	if (Pv.new) Rd=memub(Rs+Rt<<#u2)
0	0	1	1	0	0	1	1	0	0	1	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	v	v	d	d	d	d	d	d	if (!Pv.new) Rd=memub(Rs+Rt<<#u2)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				Sense	PredNew	Type	UN	s5					Parse		t2				d5													
0	1	0	0	0	0	0	1	0	0	1	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d	if (Pt) Rd=memub(Rs+#u6:0)
0	1	0	0	0	0	1	1	0	0	1	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d	if (Pt.new) Rd=memub(Rs+#u6:0)
0	1	0	0	0	1	0	1	0	0	1	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d	if (!Pt) Rd=memub(Rs+#u6:0)
0	1	0	0	0	1	1	1	0	0	1	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d	if (!Pt.new) Rd=memub(Rs+#u6:0)
ICLASS				Amode		Type	UN	x5					Parse		t2				d5													
1	0	0	1	1	0	1	1	0	0	1	x	x	x	x	x	P	P	1	0	0	t	t	i	i	i	i	d	d	d	d	d	if (Pt) Rd=memub(Rx++#s4:0)
1	0	0	1	1	0	1	1	0	0	1	x	x	x	x	x	P	P	1	0	1	t	t	i	i	i	i	d	d	d	d	d	if (!Pt) Rd=memub(Rx++#s4:0)
1	0	0	1	1	0	1	1	0	0	1	x	x	x	x	x	P	P	1	1	0	t	t	i	i	i	i	d	d	d	d	d	if (Pt.new) Rd=memub(Rx++#s4:0)
1	0	0	1	1	0	1	1	0	0	1	x	x	x	x	x	P	P	1	1	1	t	t	i	i	i	i	d	d	d	d	d	if (!Pt.new) Rd=memub(Rx++#s4:0)
ICLASS				Amode		Type	UN						Parse		t2				d5													
1	0	0	1	1	1	1	1	0	0	1	i	i	i	i	i	P	P	1	0	0	t	t	i	1	-	-	d	d	d	d	d	if (Pt) Rd=memub(#u6)
1	0	0	1	1	1	1	1	0	0	1	i	i	i	i	i	P	P	1	0	1	t	t	i	1	-	-	d	d	d	d	d	if (!Pt) Rd=memub(#u6)
1	0	0	1	1	1	1	1	0	0	1	i	i	i	i	i	P	P	1	1	0	t	t	i	1	-	-	d	d	d	d	d	if (Pt.new) Rd=memub(#u6)
1	0	0	1	1	1	1	1	0	0	1	i	i	i	i	i	P	P	1	1	1	t	t	i	1	-	-	d	d	d	d	d	if (!Pt.new) Rd=memub(#u6)

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
PredNew	PredNew
Sense	Sense
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t2	Field to encode register t
t5	Field to encode register t
v2	Field to encode register v
x5	Field to encode register x

## Load unsigned halfword

Load an unsigned halfword from memory. The 16-bit halfword at the effective address in memory is placed in the least-significant 16 bits of the destination register. The destination register is zero-extended from 16 bits to 32.

Syntax	Behavior
Rd=memuh (Re=#U6)	apply_extension (#U); EA=#U; Rd = *EA; Re=#U;
Rd=memuh (Rs+#s11:1)	apply_extension (#s); EA=Rs+#s; Rd = *EA;
Rd=memuh (Rs+Rt<<#u2)	EA=Rs+ (Rt<<#u); Rd = *EA;
Rd=memuh (Rt<<#u2+#U6)	apply_extension (#U); EA=#U+ (Rt<<#u); Rd = *EA;
Rd=memuh (Rx++#s4:1)	EA=Rx; Rx=Rx+#s; Rd = *EA;
Rd=memuh (Rx++#s4:1:circ (Mu) )	EA=Rx; Rx=Rx=circ_add (Rx, #s, MuV); Rd = *EA;
Rd=memuh (Rx++I:circ (Mu) )	EA=Rx; Rx=Rx=circ_add (Rx, I<<1, MuV); Rd = *EA;
Rd=memuh (Rx++Mu)	EA=Rx; Rx=Rx+MuV; Rd = *EA;
Rd=memuh (Rx++Mu:brev)	EA=Rx.h[1]   brev (Rx.h[0]); Rx=Rx+MuV; Rd = *EA;
Rd=memuh (gp+#u16:1)	apply_extension (#u); EA=(Constant_extended ? (0) : GP)+#u; Rd = *EA;

### Class: LD (slots 0,1)

#### Intrinsics

Rd=memuh (Rx++#s4:1:circ (Mu) )	Word32 Q6_R_memuh_IM_circ (void** StartAddress, Word32 Is4_1, Word32 Mu, void* BaseAddress)
Rd=memuh (Rx++I:circ (Mu) )	Word32 Q6_R_memuh_M_circ (void** StartAddress, Word32 Mu, void* BaseAddress)

## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS											s5					Parse		t5					d5										
0	0	1	1	1	0	1	0	0	1	1	s	s	s	s	s	P	P	i	t	t	t	t	t	i	-	-	d	d	d	d	d	Rd=memuh(Rs+Rt<<#u2)	
ICLASS											Type	UN						Parse							d5								
0	1	0	0	1	i	i	1	0	1	1	i	i	i	i	i	P	P	i	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=memuh(gp+#u16:1)
ICLASS											Amode	Type	UN	s5					Parse							d5							
1	0	0	1	0	i	i	1	0	1	1	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=memuh(Rs+#s11:1)
ICLASS											Amode	Type	UN	x5					Parse		u1						d5						
1	0	0	1	1	0	0	1	0	1	1	x	x	x	x	x	P	P	u	0	-	-	0	i	i	i	i	d	d	d	d	d	Rd=memuh(Rx++#s4:1:circ(Mu))	
1	0	0	1	1	0	0	1	0	1	1	x	x	x	x	x	P	P	u	0	-	-	1	-	0	-	-	d	d	d	d	d	Rd=memuh(Rx++l:circ(Mu))	
ICLASS											Amode	Type	UN	e5					Parse							d5							
1	0	0	1	1	0	1	1	0	1	1	e	e	e	e	e	P	P	0	1	l	l	l	l	-	l	l	d	d	d	d	d	Rd=memuh(Re=#U6)	
ICLASS											Amode	Type	UN	x5					Parse							d5							
1	0	0	1	1	0	1	1	0	1	1	x	x	x	x	x	P	P	0	0	-	-	-	i	i	i	i	d	d	d	d	d	Rd=memuh(Rx++#s4:1)	
ICLASS											Amode	Type	UN	t5					Parse							d5							
1	0	0	1	1	1	0	1	0	1	1	t	t	t	t	t	P	P	i	1	l	l	l	l	i	l	l	d	d	d	d	d	Rd=memuh(Rt<<#u2+#U6)	
ICLASS											Amode	Type	UN	x5					Parse		u1						d5						
1	0	0	1	1	1	0	1	0	1	1	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	d	d	d	d	d	Rd=memuh(Rx++Mu)	
1	0	0	1	1	1	1	1	0	1	1	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	d	d	d	d	d	Rd=memuh(Rx++Mu:brev)	

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
Parse	Packet/loop parse bits
d5	Field to encode register d
e5	Field to encode register e
s5	Field to encode register s
t5	Field to encode register t
u1	Field to encode register u
x5	Field to encode register x



## Load unsigned halfword conditionally

Load an unsigned halfword from memory. The 16-bit halfword at the effective address in memory is placed in the least-significant 16 bits of the destination register. The destination register is zero-extended from 16 bits to 32.

This instruction is conditional based on a predicate value. If the predicate is true, the instruction is performed, otherwise it is treated as a NOP.

Syntax	Behavior
<pre>if ([!]Pt[.new]) Rd=memuh (#u6)</pre>	<pre>apply_extension(#u); EA=#u; if ([!]Pt[.new][0]) {     Rd = *EA; } else {     NOP; }</pre>
<pre>if ([!]Pt[.new]) Rd=memuh (Rs+#u6:1)</pre>	<pre>apply_extension(#u); EA=Rs+#u; if ([!]Pt[.new][0]) {     Rd = *EA; } else {     NOP; }</pre>
<pre>if ([!]Pt[.new]) Rd=memuh (Rx++#s4:1)</pre>	<pre>EA=Rx; if ([!]Pt[.new][0]) {     Rx=Rx+#s;     Rd = *EA; } else {     NOP; }</pre>
<pre>if ([!]Pv[.new]) Rd=memuh (Rs+Rt&lt;&lt;#u2)</pre>	<pre>EA=Rs+(Rt&lt;&lt;#u); if ([!]Pv[.new][0]) {     Rd = *EA; } else {     NOP; }</pre>

**Class: LD (slots 0,1)**

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS											s5					Parse		t5					d5										
0	0	1	1	0	0	0	0	0	1	1	s	s	s	s	s	P	P	i	t	t	t	t	t	i	v	v	d	d	d	d	d	d	if (Pv) Rd=memuh(Rs+Rt<<#u2)
0	0	1	1	0	0	0	1	0	1	1	s	s	s	s	s	P	P	i	t	t	t	t	t	i	v	v	d	d	d	d	d	d	if (!Pv) Rd=memuh(Rs+Rt<<#u2)
0	0	1	1	0	0	1	0	0	1	1	s	s	s	s	s	P	P	i	t	t	t	t	t	i	v	v	d	d	d	d	d	d	if (Pv.new) Rd=memuh(Rs+Rt<<#u2)
0	0	1	1	0	0	1	1	0	1	1	s	s	s	s	s	P	P	i	t	t	t	t	t	i	v	v	d	d	d	d	d	d	if (!Pv.new) Rd=memuh(Rs+Rt<<#u2)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				Sense	PredNew	Type	UN	s5					Parse		t2					d5												
0	1	0	0	0	0	0	1	0	1	1	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d	if (Pt) Rd=memuh(Rs+#u6:1)
0	1	0	0	0	0	1	1	0	1	1	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d	if (Pt.new) Rd=memuh(Rs+#u6:1)
0	1	0	0	0	1	0	1	0	1	1	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d	if (!Pt) Rd=memuh(Rs+#u6:1)
0	1	0	0	0	1	1	1	0	1	1	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d	if (!Pt.new) Rd=memuh(Rs+#u6:1)
ICLASS				Amode		Type	UN	x5					Parse		t2					d5												
1	0	0	1	1	0	1	1	0	1	1	x	x	x	x	x	P	P	1	0	0	t	t	i	i	i	i	d	d	d	d	d	if (Pt) Rd=memuh(Rx++#s4:1)
1	0	0	1	1	0	1	1	0	1	1	x	x	x	x	x	P	P	1	0	1	t	t	i	i	i	i	d	d	d	d	d	if (!Pt) Rd=memuh(Rx++#s4:1)
1	0	0	1	1	0	1	1	0	1	1	x	x	x	x	x	P	P	1	1	0	t	t	i	i	i	i	d	d	d	d	d	if (Pt.new) Rd=memuh(Rx++#s4:1)
1	0	0	1	1	0	1	1	0	1	1	x	x	x	x	x	P	P	1	1	1	t	t	i	i	i	i	d	d	d	d	d	if (!Pt.new) Rd=memuh(Rx++#s4:1)
ICLASS				Amode		Type	UN						Parse		t2					d5												
1	0	0	1	1	1	1	1	0	1	1	i	i	i	i	i	P	P	1	0	0	t	t	i	1	-	-	d	d	d	d	d	if (Pt) Rd=memuh(#u6)
1	0	0	1	1	1	1	1	0	1	1	i	i	i	i	i	P	P	1	0	1	t	t	i	1	-	-	d	d	d	d	d	if (!Pt) Rd=memuh(#u6)
1	0	0	1	1	1	1	1	0	1	1	i	i	i	i	i	P	P	1	1	0	t	t	i	1	-	-	d	d	d	d	d	if (Pt.new) Rd=memuh(#u6)
1	0	0	1	1	1	1	1	0	1	1	i	i	i	i	i	P	P	1	1	1	t	t	i	1	-	-	d	d	d	d	d	if (!Pt.new) Rd=memuh(#u6)

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
PredNew	PredNew
Sense	Sense
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t2	Field to encode register t
t5	Field to encode register t
v2	Field to encode register v
x5	Field to encode register x

## Load word

Load a 32-bit word from memory and place in a destination register.

Syntax	Behavior
<code>Rd=memw (Re=#U6)</code>	<code>apply_extension (#U);</code> <code>EA=#U;</code> <code>Rd = *EA;</code> <code>Re=#U;</code>
<code>Rd=memw (Rs+#s11:2)</code>	<code>apply_extension (#s);</code> <code>EA=Rs+#s;</code> <code>Rd = *EA;</code>
<code>Rd=memw (Rs+Rt&lt;&lt;#u2)</code>	<code>EA=Rs+ (Rt&lt;&lt;#u);</code> <code>Rd = *EA;</code>
<code>Rd=memw (Rt&lt;&lt;#u2+#U6)</code>	<code>apply_extension (#U);</code> <code>EA=#U+ (Rt&lt;&lt;#u);</code> <code>Rd = *EA;</code>
<code>Rd=memw (Rx++#s4:2)</code>	<code>EA=Rx;</code> <code>Rx=Rx+#s;</code> <code>Rd = *EA;</code>
<code>Rd=memw (Rx++#s4:2:circ (Mu))</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, #s, MuV);</code> <code>Rd = *EA;</code>
<code>Rd=memw (Rx++I:circ (Mu))</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, I&lt;&lt;2, MuV);</code> <code>Rd = *EA;</code>
<code>Rd=memw (Rx++Mu)</code>	<code>EA=Rx;</code> <code>Rx=Rx+MuV;</code> <code>Rd = *EA;</code>
<code>Rd=memw (Rx++Mu:brev)</code>	<code>EA=Rx.h[1]   brev (Rx.h[0]);</code> <code>Rx=Rx+MuV;</code> <code>Rd = *EA;</code>
<code>Rd=memw (gp+#u16:2)</code>	<code>apply_extension (#u);</code> <code>EA=(Constant_extended ? (0) : GP)+#u;</code> <code>Rd = *EA;</code>

### Class: LD (slots 0,1)

#### Intrinsics

<code>Rd=memw (Rx++#s4:2:circ (Mu))</code>	<code>Word32 Q6_R_memw_IM_circ (void** StartAddress, Word32 Is4_2, Word32 Mu, void* BaseAddress)</code>
<code>Rd=memw (Rx++I:circ (Mu))</code>	<code>Word32 Q6_R_memw_M_circ (void** StartAddress, Word32 Mu, void* BaseAddress)</code>

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS											s5					Parse		t5					d5										
0	0	1	1	1	0	1	0	1	0	0	s	s	s	s	s	P	P	i	t	t	t	t	t	i	-	-	d	d	d	d	d	Rd=memw(Rs+Rt<<#u2)	
ICLASS				Type							UN						Parse							d5									
0	1	0	0	1	i	i	1	1	0	0	i	i	i	i	i	P	P	i	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=memw(gp+#u16:2)
ICLASS				Amode			Type			UN	s5					Parse							d5										
1	0	0	1	0	i	i	1	1	0	0	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=memw(Rs+#s11:2)
ICLASS				Amode			Type			UN	x5					Parse		u1						d5									
1	0	0	1	1	0	0	1	1	0	0	x	x	x	x	x	P	P	u	0	-	-	0	i	i	i	i	d	d	d	d	d	Rd=memw(Rx++#s4:2:circ(Mu))	
1	0	0	1	1	0	0	1	1	0	0	x	x	x	x	x	P	P	u	0	-	-	1	-	0	-	-	d	d	d	d	d	Rd=memw(Rx++l:circ(Mu))	
ICLASS				Amode			Type			UN	e5					Parse							d5										
1	0	0	1	1	0	1	1	1	0	0	e	e	e	e	e	P	P	0	1	l	l	l	l	-	l	l	d	d	d	d	d	Rd=memw(Re=#U6)	
ICLASS				Amode			Type			UN	x5					Parse							d5										
1	0	0	1	1	0	1	1	1	0	0	x	x	x	x	x	P	P	0	0	-	-	-	i	i	i	i	d	d	d	d	d	Rd=memw(Rx++#s4:2)	
ICLASS				Amode			Type			UN	t5					Parse							d5										
1	0	0	1	1	1	0	1	1	0	0	t	t	t	t	t	P	P	i	1	l	l	l	l	i	l	l	d	d	d	d	d	Rd=memw(Rt<<#u2+#U6)	
ICLASS				Amode			Type			UN	x5					Parse		u1						d5									
1	0	0	1	1	1	0	1	1	0	0	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	d	d	d	d	d	Rd=memw(Rx++Mu)	
1	0	0	1	1	1	1	1	1	0	0	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	d	d	d	d	d	Rd=memw(Rx++Mu:brev)	

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
Parse	Packet/loop parse bits
d5	Field to encode register d
e5	Field to encode register e
s5	Field to encode register s
t5	Field to encode register t
u1	Field to encode register u
x5	Field to encode register x

## Load-acquire word

Load a 32-bit word from memory and place in a destination register. The load-acquire memory operation is observed before any following memory operations (in program order) are observed at the local point of serialization. A different order can be observed at the global point of serialization (see Ordering and Synchronization).

### Syntax

```
Rd=memw_aq(Rs)
```

### Behavior

```
EA=Rs;
Rd = *EA
```

**Class: LD (slots 0)**

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				Amode			Type			U N	s5					Parse		d5														
1	0	0	1	0	0	1	0	0	0	0	s	s	s	s	s	P	P	0	0	1	-	-	-	0	0	0	d	d	d	d	d	Rd=memw_aq(Rs)

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s

## Load word conditionally

Load a 32-bit word from memory and place in a destination register.

This instruction is conditional based on a predicate value. If the predicate is true, the instruction is performed, otherwise it is treated as a NOP.

Syntax	Behavior
<code>if ([!]Pt[.new]) Rd=memw(#u6)</code>	<pre> apply_extension(#u); EA=#u; if ([!]Pt[.new][0]) {     Rd = *EA; } else {     NOP; } </pre>
<code>if ([!]Pt[.new]) Rd=memw(Rs+#u6:2)</code>	<pre> apply_extension(#u); EA=Rs+#u; if ([!]Pt[.new][0]) {     Rd = *EA; } else {     NOP; } </pre>
<code>if ([!]Pt[.new]) Rd=memw(Rx++#s4:2)</code>	<pre> EA=Rx; if ([!]Pt[.new][0]) {     Rx=Rx+#s;     Rd = *EA; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) Rd=memw(Rs+Rt&lt;&lt;#u2)</code>	<pre> EA=Rs+(Rt&lt;&lt;#u); if ([!]Pv[.new][0]) {     Rd = *EA; } else {     NOP; } </pre>

**Class: LD (slots 0,1)**

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS											s5					Parse		t5					d5										
0	0	1	1	0	0	0	0	1	0	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	v	v	d	d	d	d	d	if (Pv) Rd=memw(Rs+Rt<<#u2)
0	0	1	1	0	0	0	1	1	0	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	v	v	d	d	d	d	d	if (!Pv) Rd=memw(Rs+Rt<<#u2)
0	0	1	1	0	0	1	0	1	0	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	v	v	d	d	d	d	d	if (Pv.new) Rd=memw(Rs+Rt<<#u2)
0	0	1	1	0	0	1	1	1	0	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	v	v	d	d	d	d	d	if (!Pv.new) Rd=memw(Rs+Rt<<#u2)
ICLASS					Se ns e	Pr ed Ne w	Type	UN	s5					Parse		t2					d5												

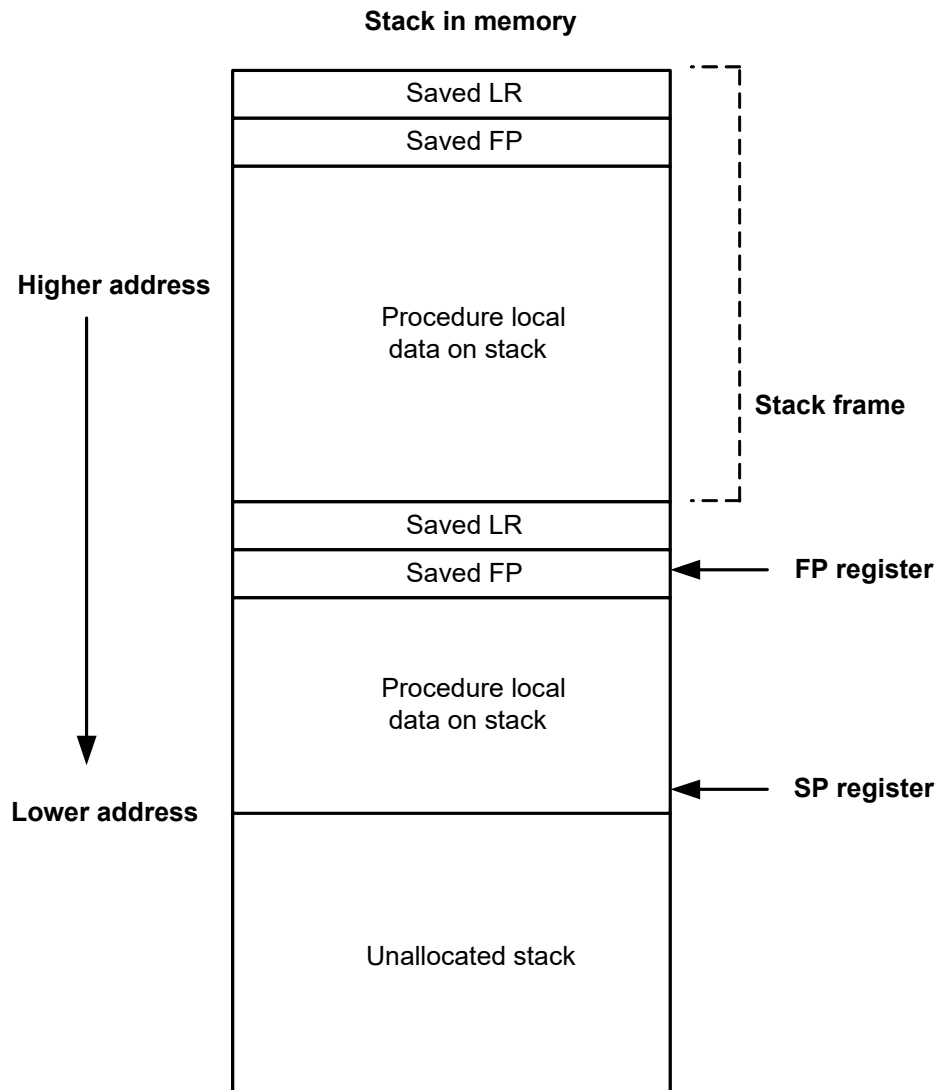
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
0	1	0	0	0	0	0	1	1	0	0	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d		if (Pt) Rd=memw(Rs+#u6:2)
0	1	0	0	0	0	1	1	1	0	0	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d		if (Pt.new) Rd=memw(Rs+#u6:2)
0	1	0	0	0	1	0	1	1	0	0	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d		if (!Pt) Rd=memw(Rs+#u6:2)
0	1	0	0	0	1	1	1	1	0	0	s	s	s	s	s	P	P	0	t	t	i	i	i	i	i	i	d	d	d	d	d		if (!Pt.new) Rd=memw(Rs+#u6:2)
ICLASS			Amode			Type		U N	x5					Parse		t2			d5														
1	0	0	1	1	0	1	1	1	0	0	x	x	x	x	x	P	P	1	0	0	t	t	i	i	i	i	d	d	d	d	d		if (Pt) Rd=memw(Rx+++s4:2)
1	0	0	1	1	0	1	1	1	0	0	x	x	x	x	x	P	P	1	0	1	t	t	i	i	i	i	d	d	d	d	d		if (!Pt) Rd=memw(Rx+++s4:2)
1	0	0	1	1	0	1	1	1	0	0	x	x	x	x	x	P	P	1	1	0	t	t	i	i	i	i	d	d	d	d	d		if (Pt.new) Rd=memw(Rx+++s4:2)
1	0	0	1	1	0	1	1	1	0	0	x	x	x	x	x	P	P	1	1	1	t	t	i	i	i	i	d	d	d	d	d		if (!Pt.new) Rd=memw(Rx+++s4:2)
ICLASS			Amode			Type		U N						Parse		t2			d5														
1	0	0	1	1	1	1	1	1	0	0	i	i	i	i	i	P	P	1	0	0	t	t	i	1	-	-	d	d	d	d	d		if (Pt) Rd=memw(#u6)
1	0	0	1	1	1	1	1	1	0	0	i	i	i	i	i	P	P	1	0	1	t	t	i	1	-	-	d	d	d	d	d		if (!Pt) Rd=memw(#u6)
1	0	0	1	1	1	1	1	1	0	0	i	i	i	i	i	P	P	1	1	0	t	t	i	1	-	-	d	d	d	d	d		if (Pt.new) Rd=memw(#u6)
1	0	0	1	1	1	1	1	1	0	0	i	i	i	i	i	P	P	1	1	1	t	t	i	1	-	-	d	d	d	d	d		if (!Pt.new) Rd=memw(#u6)

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
PredNew	PredNew
Sense	Sense
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t2	Field to encode register t
t5	Field to encode register t
v2	Field to encode register v
x5	Field to encode register x

## Deallocate stack frame

Deallocate a stack frame from the call stack. The instruction first loads the saved FP and saved LR values from the address at FP. It then points SP back to the previous frame.

The following figure shows the stack layout.



### Syntax

```
Rdd=deallocframe(Rs):raw
```

```
deallocframe
```

### Behavior

```
EA=Rs;
tmp = *EA;
Rdd = frame_unscramble(tmp);
SP=EA+8;
```

```
Assembler mapped to: "r31:30=deallocframe(r30):raw"
```



**Class: LD (slots 0,1)****Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS			Amode			Type		U N	s5					Parse												d5							
1	0	0	1	0	0	0	0	0	0	0	s	s	s	s	s	P	P	0	-	-	-	-	-	-	-	-	-	d	d	d	d	d	Rdd=deallocframe(Rs):raw

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s

## Deallocate frame and return

Return from a function with a stack frame. This instruction is equivalent to deallocframe followed by jumpr R31.

Syntax	Behavior
<code>Rdd=dealloc_return(Rs):raw</code>	<pre>EA=Rs; tmp = *EA; Rdd = frame_unscramble(tmp); SP=EA+8; PC=Rdd.w[1];</pre>
<code>dealloc_return</code>	Assembler mapped to: <code>"r31:30=dealloc_return(r30):raw"</code>
<code>if ([!]Pv) Rdd=dealloc_return(Rs):raw</code>	<pre>EA=Rs; if ([!]Pv[0]) {     tmp = *EA;     Rdd = frame_unscramble(tmp);     SP=EA+8;     PC=Rdd.w[1]; } else {     NOP; }</pre>
<code>if ([!]Pv) dealloc_return</code>	Assembler mapped to: <code>"if ([!]Pv"" r31:30=dealloc_return(r30)"":raw"</code>
<code>if ([!]Pv.new) Rdd=dealloc_return(Rs):nt:raw</code>	<pre>EA=Rs; if ([!]Pv.new[0]) {     tmp = *EA;     Rdd = frame_unscramble(tmp);     SP=EA+8;     PC=Rdd.w[1]; } else {     NOP; }</pre>
<code>if ([!]Pv.new) Rdd=dealloc_return(Rs):t:raw</code>	<pre>EA=Rs; if ([!]Pv.new[0]) {     tmp = *EA;     Rdd = frame_unscramble(tmp);     SP=EA+8;     PC=Rdd.w[1]; } else {     NOP; }</pre>
<code>if ([!]Pv.new) dealloc_return:nt</code>	Assembler mapped to: <code>"if ([!]Pv"".new"" r31:30=dealloc_return(r30)"":nt"":raw"</code>
<code>if ([!]Pv.new) dealloc_return:t</code>	Assembler mapped to: <code>"if ([!]Pv"".new"" r31:30=dealloc_return(r30)"":t"":raw"</code>

**Class: LD (slots 0)**

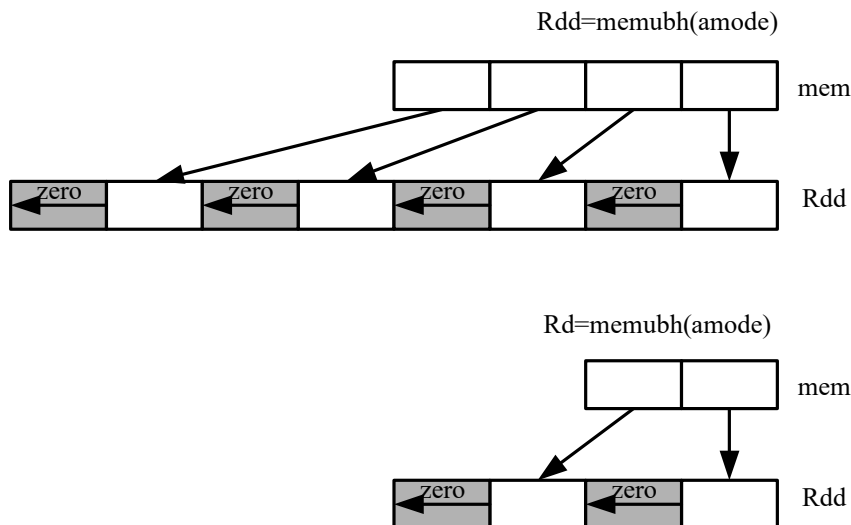
**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			Amode			Type			UN	s5					Parse												d5					
1	0	0	1	0	1	1	0	0	0	0	s	s	s	s	s	P	P	0	0	0	0	-	-	-	-	-	d	d	d	d	d	Rdd=dealloc_return(Rs):raw
1	0	0	1	0	1	1	0	0	0	0	s	s	s	s	s	P	P	0	0	1	0	v	v	-	-	-	d	d	d	d	d	if (Pv.new) Rdd=dealloc_return(Rs):nt:raw
1	0	0	1	0	1	1	0	0	0	0	s	s	s	s	s	P	P	0	1	0	0	v	v	-	-	-	d	d	d	d	d	if (Pv) Rdd=dealloc_return(Rs):raw
1	0	0	1	0	1	1	0	0	0	0	s	s	s	s	s	P	P	0	1	1	0	v	v	-	-	-	d	d	d	d	d	if (Pv.new) Rdd=dealloc_return(Rs):traw
1	0	0	1	0	1	1	0	0	0	0	s	s	s	s	s	P	P	1	0	1	0	v	v	-	-	-	d	d	d	d	d	if (!Pv.new) Rdd=dealloc_return(Rs):nt:raw
1	0	0	1	0	1	1	0	0	0	0	s	s	s	s	s	P	P	1	1	0	0	v	v	-	-	-	d	d	d	d	d	if (!Pv) Rdd=dealloc_return(Rs):raw
1	0	0	1	0	1	1	0	0	0	0	s	s	s	s	s	P	P	1	1	1	0	v	v	-	-	-	d	d	d	d	d	if (!Pv.new) Rdd=dealloc_return(Rs):traw

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
v2	Field to encode register v

## Load and unpack bytes to halfwords

Load contiguous bytes from memory and vector unpack them into halfwords.



Syntax	Behavior
Rd=membh (Re=#U6)	<pre> apply_extension (#U); EA=#U; {     tmpV = *EA;     for (i=0;i&lt;2;i++) {         Rd.h[i]=tmpV.b[i];     } } Re=#U; </pre>
Rd=membh (Rs)	Assembler mapped to: "Rd=membh" (Rs+#0) "
Rd=membh (Rs+#s11:1)	<pre> apply_extension (#s); EA=Rs+#s; {     tmpV = *EA;     for (i=0;i&lt;2;i++) {         Rd.h[i]=tmpV.b[i];     } } </pre>
Rd=membh (Rt<<#u2+#U6)	<pre> apply_extension (#U); EA=#U+ (Rt&lt;&lt;#u); {     tmpV = *EA;     for (i=0;i&lt;2;i++) {         Rd.h[i]=tmpV.b[i];     } } </pre>

Syntax	Behavior
<code>Rd=membh (Rx++#s4:1)</code>	<pre>EA=Rx; Rx=Rx+#s; {     tmpV = *EA;     for (i=0;i&lt;2;i++) {         Rd.h[i]=tmpV.b[i];     } }</pre>
<code>Rd=membh (Rx++#s4:1:circ (Mu) )</code>	<pre>EA=Rx; Rx=Rx=circ_add (Rx, #s, MuV); {     tmpV = *EA;     for (i=0;i&lt;2;i++) {         Rd.h[i]=tmpV.b[i];     } }</pre>
<code>Rd=membh (Rx++I:circ (Mu) )</code>	<pre>EA=Rx; Rx=Rx=circ_add (Rx, I&lt;&lt;1, MuV); {     tmpV = *EA;     for (i=0;i&lt;2;i++) {         Rd.h[i]=tmpV.b[i];     } }</pre>
<code>Rd=membh (Rx++Mu)</code>	<pre>EA=Rx; Rx=Rx+MuV; {     tmpV = *EA;     for (i=0;i&lt;2;i++) {         Rd.h[i]=tmpV.b[i];     } }</pre>
<code>Rd=membh (Rx++Mu:brev)</code>	<pre>EA=Rx.h[1]   brev (Rx.h[0]); Rx=Rx+MuV; {     tmpV = *EA;     for (i=0;i&lt;2;i++) {         Rd.h[i]=tmpV.b[i];     } }</pre>
<code>Rd=memubh (Re=#U6)</code>	<pre>apply_extension (#U); EA=#U; {     tmpV = *EA;     for (i=0;i&lt;2;i++) {         Rd.h[i]=tmpV.ub[i];     } } Re=#U;</pre>

Syntax	Behavior
Rd=memubh (Rs+#s1:1)	<pre> apply_extension(#s); EA=Rs+#s; {     tmpV = *EA;     for (i=0;i&lt;2;i++) {         Rd.h[i]=tmpV.ub[i];     } } </pre>
Rd=memubh (Rt<<#u2+#U6)	<pre> apply_extension(#U); EA=#U+(Rt&lt;&lt;#u); {     tmpV = *EA;     for (i=0;i&lt;2;i++) {         Rd.h[i]=tmpV.ub[i];     } } </pre>
Rd=memubh (Rx++#s4:1)	<pre> EA=Rx; Rx=Rx+#s; {     tmpV = *EA;     for (i=0;i&lt;2;i++) {         Rd.h[i]=tmpV.ub[i];     } } </pre>
Rd=memubh (Rx++#s4:1:circ (Mu) )	<pre> EA=Rx; Rx=Rx=circ_add(Rx,#s,MuV); {     tmpV = *EA;     for (i=0;i&lt;2;i++) {         Rd.h[i]=tmpV.ub[i];     } } </pre>
Rd=memubh (Rx++I:circ (Mu) )	<pre> EA=Rx; Rx=Rx=circ_add(Rx,I&lt;&lt;1,MuV); {     tmpV = *EA;     for (i=0;i&lt;2;i++) {         Rd.h[i]=tmpV.ub[i];     } } </pre>
Rd=memubh (Rx++Mu)	<pre> EA=Rx; Rx=Rx+MuV; {     tmpV = *EA;     for (i=0;i&lt;2;i++) {         Rd.h[i]=tmpV.ub[i];     } } </pre>

Syntax	Behavior
<code>Rd=memubh (Rx++Mu:brev)</code>	<pre>EA=Rx.h[1]   brev(Rx.h[0]); Rx=Rx+MuV; {     tmpV = *EA;     for (i=0;i&lt;2;i++) {         Rd.h[i]=tmpV.ub[i];     } }</pre>
<code>Rdd=membh (Re=#U6)</code>	<pre>apply_extension(#U); EA=#U; {     tmpV = *EA;     for (i=0;i&lt;4;i++) {         Rdd.h[i]=tmpV.b[i];     } } Re=#U;</pre>
<code>Rdd=membh (Rs)</code>	Assembler mapped to: <code>"Rdd=membh""(Rs+#0)"</code>
<code>Rdd=membh (Rs+#s11:2)</code>	<pre>apply_extension(#s); EA=Rs+#s; {     tmpV = *EA;     for (i=0;i&lt;4;i++) {         Rdd.h[i]=tmpV.b[i];     } }</pre>
<code>Rdd=membh (Rt&lt;&lt;#u2+#U6)</code>	<pre>apply_extension(#U); EA=#U+(Rt&lt;&lt;#u); {     tmpV = *EA;     for (i=0;i&lt;4;i++) {         Rdd.h[i]=tmpV.b[i];     } }</pre>
<code>Rdd=membh (Rx++#s4:2)</code>	<pre>EA=Rx; Rx=Rx+#s; {     tmpV = *EA;     for (i=0;i&lt;4;i++) {         Rdd.h[i]=tmpV.b[i];     } }</pre>
<code>Rdd=membh (Rx++#s4:2:circ (Mu) )</code>	<pre>EA=Rx; Rx=Rx=circ_add(Rx, #s, MuV); {     tmpV = *EA;     for (i=0;i&lt;4;i++) {         Rdd.h[i]=tmpV.b[i];     } }</pre>

Syntax	Behavior
<code>Rdd=membh (Rx++I:circ (Mu) )</code>	<pre>EA=Rx; Rx=Rx=circ_add (Rx, I&lt;&lt;2, MuV); {     tmpV = *EA;     for (i=0; i&lt;4; i++) {         Rdd.h[i]=tmpV.b[i];     } }</pre>
<code>Rdd=membh (Rx++Mu)</code>	<pre>EA=Rx; Rx=Rx+MuV; {     tmpV = *EA;     for (i=0; i&lt;4; i++) {         Rdd.h[i]=tmpV.b[i];     } }</pre>
<code>Rdd=membh (Rx++Mu:brev)</code>	<pre>EA=Rx.h[1]   brev (Rx.h[0]); Rx=Rx+MuV; {     tmpV = *EA;     for (i=0; i&lt;4; i++) {         Rdd.h[i]=tmpV.b[i];     } }</pre>
<code>Rdd=memubh (Re=#U6)</code>	<pre>apply_extension (#U); EA=#U; {     tmpV = *EA;     for (i=0; i&lt;4; i++) {         Rdd.h[i]=tmpV.ub[i];     } } Re=#U;</pre>
<code>Rdd=memubh (Rs+#s11:2)</code>	<pre>apply_extension (#s); EA=Rs+#s; {     tmpV = *EA;     for (i=0; i&lt;4; i++) {         Rdd.h[i]=tmpV.ub[i];     } }</pre>
<code>Rdd=memubh (Rt&lt;&lt;#u2+#U6)</code>	<pre>apply_extension (#U); EA=#U+(Rt&lt;&lt;#u); {     tmpV = *EA;     for (i=0; i&lt;4; i++) {         Rdd.h[i]=tmpV.ub[i];     } }</pre>



Syntax	Behavior
<code>Rdd=memubh (Rx++#s4:2)</code>	<pre>EA=Rx; Rx=Rx+#s; {     tmpV = *EA;     for (i=0;i&lt;4;i++) {         Rdd.h[i]=tmpV.ub[i];     } }</pre>
<code>Rdd=memubh (Rx++#s4:2:circ (Mu))</code>	<pre>EA=Rx; Rx=Rx=circ_add (Rx, #s, MuV); {     tmpV = *EA;     for (i=0;i&lt;4;i++) {         Rdd.h[i]=tmpV.ub[i];     } }</pre>
<code>Rdd=memubh (Rx++I:circ (Mu))</code>	<pre>EA=Rx; Rx=Rx=circ_add (Rx, I&lt;&lt;2, MuV); {     tmpV = *EA;     for (i=0;i&lt;4;i++) {         Rdd.h[i]=tmpV.ub[i];     } }</pre>
<code>Rdd=memubh (Rx++Mu)</code>	<pre>EA=Rx; Rx=Rx+MuV; {     tmpV = *EA;     for (i=0;i&lt;4;i++) {         Rdd.h[i]=tmpV.ub[i];     } }</pre>
<code>Rdd=memubh (Rx++Mu:brev)</code>	<pre>EA=Rx.h[1]   brev (Rx.h[0]); Rx=Rx+MuV; {     tmpV = *EA;     for (i=0;i&lt;4;i++) {         Rdd.h[i]=tmpV.ub[i];     } }</pre>

**Class: LD (slots 0,1)**

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			Amode			Type			U	s5					Parse		d5															
1	0	0	1	0	i	i	0	0	0	1	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=membh(Rs+#s11:1)	
1	0	0	1	0	i	i	0	0	1	1	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=memubh(Rs+#s11:1)	
1	0	0	1	0	i	i	0	1	0	1	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	d	d	d	d	d	Rdd=memubh(Rs+#s11:2)	
1	0	0	1	0	i	i	0	1	1	1	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	d	d	d	d	d	Rdd=membh(Rs+#s11:2)	

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS		Amode				Type				UN	x5					Parse		u1											d5				
1	0	0	1	1	0	0	0	0	0	1	x	x	x	x	x	P	P	u	0	-	-	0	i	i	i	i	d	d	d	d	d	Rd=membh(Rx++#s4:1:circ(Mu))	
1	0	0	1	1	0	0	0	0	0	1	x	x	x	x	x	P	P	u	0	-	-	1	-	0	-	-	d	d	d	d	d	Rd=membh(Rx++l:circ(Mu))	
1	0	0	1	1	0	0	0	0	1	1	x	x	x	x	x	P	P	u	0	-	-	0	i	i	i	i	d	d	d	d	d	Rd=memubh(Rx++#s4:1:circ(Mu))	
1	0	0	1	1	0	0	0	0	1	1	x	x	x	x	x	P	P	u	0	-	-	1	-	0	-	-	d	d	d	d	d	Rd=memubh(Rx++l:circ(Mu))	
1	0	0	1	1	0	0	0	1	0	1	x	x	x	x	x	P	P	u	0	-	-	0	i	i	i	i	d	d	d	d	d	Rdd=memubh(Rx++#s4:2:circ(Mu))	
1	0	0	1	1	0	0	0	1	0	1	x	x	x	x	x	P	P	u	0	-	-	1	-	0	-	-	d	d	d	d	d	Rdd=memubh(Rx++l:circ(Mu))	
1	0	0	1	1	0	0	0	1	1	1	x	x	x	x	x	P	P	u	0	-	-	0	i	i	i	i	d	d	d	d	d	Rdd=membh(Rx++#s4:2:circ(Mu))	
1	0	0	1	1	0	0	0	1	1	1	x	x	x	x	x	P	P	u	0	-	-	1	-	0	-	-	d	d	d	d	d	Rdd=membh(Rx++l:circ(Mu))	
ICLASS		Amode				Type				UN	e5					Parse												d5					
1	0	0	1	1	0	1	0	0	0	1	e	e	e	e	e	P	P	0	1	l	l	l	l	-	l	l	d	d	d	d	d	Rd=membh(Re=#U6)	
ICLASS		Amode				Type				UN	x5					Parse												d5					
1	0	0	1	1	0	1	0	0	0	1	x	x	x	x	x	P	P	0	0	-	-	-	i	i	i	i	d	d	d	d	d	Rd=membh(Rx++#s4:1)	
ICLASS		Amode				Type				UN	e5					Parse												d5					
1	0	0	1	1	0	1	0	0	1	1	e	e	e	e	e	P	P	0	1	l	l	l	l	-	l	l	d	d	d	d	d	Rd=memubh(Re=#U6)	
ICLASS		Amode				Type				UN	x5					Parse												d5					
1	0	0	1	1	0	1	0	0	1	1	x	x	x	x	x	P	P	0	0	-	-	-	i	i	i	i	d	d	d	d	d	Rd=memubh(Rx++#s4:1)	
ICLASS		Amode				Type				UN	e5					Parse												d5					
1	0	0	1	1	0	1	0	1	0	1	e	e	e	e	e	P	P	0	1	l	l	l	l	-	l	l	d	d	d	d	d	Rdd=memubh(Re=#U6)	
ICLASS		Amode				Type				UN	x5					Parse												d5					
1	0	0	1	1	0	1	0	1	0	1	x	x	x	x	x	P	P	0	0	-	-	-	i	i	i	i	d	d	d	d	d	Rdd=memubh(Rx++#s4:2)	
ICLASS		Amode				Type				UN	e5					Parse												d5					
1	0	0	1	1	0	1	0	1	1	1	e	e	e	e	e	P	P	0	1	l	l	l	l	-	l	l	d	d	d	d	d	Rdd=membh(Re=#U6)	
ICLASS		Amode				Type				UN	x5					Parse												d5					
1	0	0	1	1	0	1	0	1	1	1	x	x	x	x	x	P	P	0	0	-	-	-	i	i	i	i	d	d	d	d	d	Rdd=membh(Rx++#s4:2)	
ICLASS		Amode				Type				UN	t5					Parse												d5					
1	0	0	1	1	1	0	0	0	0	1	t	t	t	t	t	P	P	i	1	l	l	l	l	i	l	l	d	d	d	d	d	Rd=membh(Rt<<#u2+#U6)	
ICLASS		Amode				Type				UN	x5					Parse		u1											d5				
1	0	0	1	1	1	0	0	0	0	1	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	d	d	d	d	d	Rd=membh(Rx++Mu)	
ICLASS		Amode				Type				UN	t5					Parse												d5					
1	0	0	1	1	1	0	0	0	1	1	t	t	t	t	t	P	P	i	1	l	l	l	l	i	l	l	d	d	d	d	d	Rd=memubh(Rt<<#u2+#U6)	
ICLASS		Amode				Type				UN	x5					Parse		u1											d5				
1	0	0	1	1	1	0	0	0	1	1	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	d	d	d	d	d	Rd=memubh(Rx++Mu)	
ICLASS		Amode				Type				UN	t5					Parse												d5					
1	0	0	1	1	1	0	0	1	0	1	t	t	t	t	t	P	P	i	1	l	l	l	l	i	l	l	d	d	d	d	d	Rdd=memubh(Rt<<#u2+#U6)	

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			Amode			Type		U	x5					Parse		u1					d5											
1	0	0	1	1	1	0	0	1	0	1	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	d	d	d	d	d	Rdd=memubh(Rx++Mu)
ICLASS			Amode			Type		U	t5					Parse							d5											
1	0	0	1	1	1	0	0	1	1	1	t	t	t	t	t	P	P	i	1	l	l	l	l	i	l	l	d	d	d	d	d	Rdd=membh(Rt<<#u2+#U6)
ICLASS			Amode			Type		U	x5					Parse		u1					d5											
1	0	0	1	1	1	0	0	1	1	1	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	d	d	d	d	d	Rdd=membh(Rx++Mu)
1	0	0	1	1	1	1	0	0	0	1	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	d	d	d	d	d	Rd=membh(Rx++Mu:brev)
1	0	0	1	1	1	1	0	0	1	1	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	d	d	d	d	d	Rd=memubh(Rx++Mu:brev)
1	0	0	1	1	1	1	0	1	0	1	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	d	d	d	d	d	Rdd=memubh(Rx++Mu:brev)
1	0	0	1	1	1	1	0	1	1	1	x	x	x	x	x	P	P	u	0	-	-	-	-	0	-	-	d	d	d	d	d	Rdd=membh(Rx++Mu:brev)

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
Parse	Packet/loop parse bits
d5	Field to encode register d
e5	Field to encode register e
s5	Field to encode register s
t5	Field to encode register t
u1	Field to encode register u
x5	Field to encode register x

## 11.6 MEMOP

The MEMOP instruction class includes simple operations on values in memory.

MEMOP instructions are executable on slot 0.

### Operation on memory byte

Perform ALU or bit operation on the memory byte at the effective address.

Syntax	Behavior
<code>memb(Rs+#u6:0)=clrbit(#U5)</code>	<pre> apply_extension(#u); EA=Rs+#u; tmp = *EA; tmp &amp;= (~(1&lt;&lt;#U)); *EA = tmp; </pre>
<code>memb(Rs+#u6:0)=setbit(#U5)</code>	<pre> apply_extension(#u); EA=Rs+#u; tmp = *EA; tmp  = (1&lt;&lt;#U); *EA = tmp; </pre>
<code>memb(Rs+#u6:0)[+-]=#U5</code>	<pre> apply_extension(#u); EA=Rs[+-]#u; tmp = *EA; tmp [+-]= #U; *EA = tmp; </pre>
<code>memb(Rs+#u6:0)[+&amp;]=Rt</code>	<pre> apply_extension(#u); EA=Rs+#u; tmp = *EA; tmp [+&amp;]= Rt; *EA = tmp; </pre>

### Class: MEMOP (slots 0)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS											s5					Parse					t5											
0	0	1	1	1	1	1	0	-	0	0	s	s	s	s	s	P	P	0	i	i	i	i	i	i	0	0	t	t	t	t	t	memb(Rs+#u6:0)+=Rt
0	0	1	1	1	1	1	0	-	0	0	s	s	s	s	s	P	P	0	i	i	i	i	i	i	0	1	t	t	t	t	t	memb(Rs+#u6:0)-=Rt
0	0	1	1	1	1	1	0	-	0	0	s	s	s	s	s	P	P	0	i	i	i	i	i	i	1	0	t	t	t	t	t	memb(Rs+#u6:0)&=Rt
0	0	1	1	1	1	1	0	-	0	0	s	s	s	s	s	P	P	0	i	i	i	i	i	i	1	1	t	t	t	t	t	memb(Rs+#u6:0) =Rt
ICLASS											s5					Parse																
0	0	1	1	1	1	1	1	-	0	0	s	s	s	s	s	P	P	0	i	i	i	i	i	i	0	0	I	I	I	I	I	memb(Rs+#u6:0)+=#U5
0	0	1	1	1	1	1	1	-	0	0	s	s	s	s	s	P	P	0	i	i	i	i	i	i	0	1	I	I	I	I	I	memb(Rs+#u6:0)-=#U5
0	0	1	1	1	1	1	1	-	0	0	s	s	s	s	s	P	P	0	i	i	i	i	i	i	1	0	I	I	I	I	I	memb(Rs+#u6:0)=clrbit(#U5)
0	0	1	1	1	1	1	1	-	0	0	s	s	s	s	s	P	P	0	i	i	i	i	i	i	1	1	I	I	I	I	I	memb(Rs+#u6:0)=setbit(#U5)

<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s
t5	Field to encode register t

## Operation on memory halfword

Perform ALU or bit operation on the memory halfword at the effective address.

Syntax	Behavior
<code>memh(Rs+#u6:1)=clrbit(#U5)</code>	<pre> apply_extension(#u); EA=Rs+#u; tmp = *EA; tmp &amp;= (~(1&lt;&lt;#U)); *EA = tmp; </pre>
<code>memh(Rs+#u6:1)=setbit(#U5)</code>	<pre> apply_extension(#u); EA=Rs+#u; tmp = *EA; tmp  = (1&lt;&lt;#U); *EA = tmp; </pre>
<code>memh(Rs+#u6:1)[+-]=#U5</code>	<pre> apply_extension(#u); EA=Rs[+-]#u; tmp = *EA; tmp [+-]= #U; *EA = tmp; </pre>
<code>memh(Rs+#u6:1)[+- &amp;]=Rt</code>	<pre> apply_extension(#u); EA=Rs+#u; tmp = *EA; tmp [+- &amp;]= Rt; *EA = tmp; </pre>

### Class: MEMOP (slots 0)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS											s5					Parse				t5												
0	0	1	1	1	1	1	0	-	0	1	s	s	s	s	s	P	P	0	i	i	i	i	i	i	0	0	t	t	t	t	t	memh(Rs+#u6:1)+=Rt
0	0	1	1	1	1	1	0	-	0	1	s	s	s	s	s	P	P	0	i	i	i	i	i	i	0	1	t	t	t	t	t	memh(Rs+#u6:1)-=Rt
0	0	1	1	1	1	1	0	-	0	1	s	s	s	s	s	P	P	0	i	i	i	i	i	i	1	0	t	t	t	t	t	memh(Rs+#u6:1)&=Rt
0	0	1	1	1	1	1	0	-	0	1	s	s	s	s	s	P	P	0	i	i	i	i	i	i	1	1	t	t	t	t	t	memh(Rs+#u6:1) =Rt
ICLASS											s5					Parse																
0	0	1	1	1	1	1	1	-	0	1	s	s	s	s	s	P	P	0	i	i	i	i	i	i	0	0	l	l	l	l	l	memh(Rs+#u6:1)+=#U5
0	0	1	1	1	1	1	1	-	0	1	s	s	s	s	s	P	P	0	i	i	i	i	i	i	0	1	l	l	l	l	l	memh(Rs+#u6:1)-=#U5
0	0	1	1	1	1	1	1	-	0	1	s	s	s	s	s	P	P	0	i	i	i	i	i	i	1	0	l	l	l	l	l	memh(Rs+#u6:1)=clrbit(#U5)
0	0	1	1	1	1	1	1	-	0	1	s	s	s	s	s	P	P	0	i	i	i	i	i	i	1	1	l	l	l	l	l	memh(Rs+#u6:1)=setbit(#U5)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s
t5	Field to encode register t

## Operation on memory word

Perform ALU or bit operation on the memory word at the effective address.

Syntax	Behavior
<code>memw(Rs+#u6:2)=clrbit(#U5)</code>	<pre> apply_extension(#u); EA=Rs+#u; tmp = *EA; tmp &amp;= (~(1&lt;&lt;#U)); *EA = tmp; </pre>
<code>memw(Rs+#u6:2)=setbit(#U5)</code>	<pre> apply_extension(#u); EA=Rs+#u; tmp = *EA; tmp  = (1&lt;&lt;#U); *EA = tmp; </pre>
<code>memw(Rs+#u6:2)[+-]=#U5</code>	<pre> apply_extension(#u); EA=Rs[+-]#u; tmp = *EA; tmp [+-]= #U; *EA = tmp; </pre>
<code>memw(Rs+#u6:2)[+- &amp;]=Rt</code>	<pre> apply_extension(#u); EA=Rs+#u; tmp = *EA; tmp [+- &amp;]= Rt; *EA = tmp; </pre>

### Class: MEMOP (slots 0)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS											s5					Parse		t5														
0	0	1	1	1	1	1	0	-	1	0	s	s	s	s	s	P	P	0	i	i	i	i	i	0	0	t	t	t	t	t		memw(Rs+#u6:2)+=Rt
0	0	1	1	1	1	1	0	-	1	0	s	s	s	s	s	P	P	0	i	i	i	i	i	0	1	t	t	t	t	t		memw(Rs+#u6:2)=Rt
0	0	1	1	1	1	1	0	-	1	0	s	s	s	s	s	P	P	0	i	i	i	i	i	1	0	t	t	t	t	t		memw(Rs+#u6:2)&=Rt
0	0	1	1	1	1	1	0	-	1	0	s	s	s	s	s	P	P	0	i	i	i	i	i	1	1	t	t	t	t	t		memw(Rs+#u6:2) =Rt
ICLASS											s5					Parse		t5														
0	0	1	1	1	1	1	1	-	1	0	s	s	s	s	s	P	P	0	i	i	i	i	i	0	0	l	l	l	l	l		memw(Rs+#u6:2)+=#U5
0	0	1	1	1	1	1	1	-	1	0	s	s	s	s	s	P	P	0	i	i	i	i	i	0	1	l	l	l	l	l		memw(Rs+#u6:2)-=#U5
0	0	1	1	1	1	1	1	-	1	0	s	s	s	s	s	P	P	0	i	i	i	i	i	1	0	l	l	l	l	l		memw(Rs+#u6:2)=clrbit(#U5)
0	0	1	1	1	1	1	1	-	1	0	s	s	s	s	s	P	P	0	i	i	i	i	i	1	1	l	l	l	l	l		memw(Rs+#u6:2)=setbit(#U5)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s
t5	Field to encode register t

## 11.7 NV

The NV instruction class includes instructions that take the register source operand from another instruction in the same packet.

NV instructions are executable on slot 0.

### 11.7.1 NV J

The NV J instruction subclass includes jump instructions that take the register source operand from another instruction in the same packet.

#### Jump to address condition on new register value

Compare a register or constant against the value produced by a slot 1 instruction. If the comparison is true, the program counter is changed to a target address, relative to the current PC.

This instruction is executable only on slot 0.

Syntax	Behavior
<pre>if ([!]cmp.eq(Ns.new,#-1)) jump:&lt;hint&gt; #r9:2</pre>	<pre>if ((Ns.new[!]=(-1))) {   apply_extension(#r);   #r=#r &amp; ~PCALIGN_MASK;   PC=PC+#r; }</pre>
<pre>if ([!]cmp.eq(Ns.new,#U5)) jump:&lt;hint&gt; #r9:2</pre>	<pre>if ((Ns.new[!]=(#U))) {   apply_extension(#r);   #r=#r &amp; ~PCALIGN_MASK;   PC=PC+#r; }</pre>
<pre>if ([!]cmp.eq(Ns.new,Rt)) jump:&lt;hint&gt; #r9:2</pre>	<pre>if ((Ns.new[!]=Rt)) {   apply_extension(#r);   #r=#r &amp; ~PCALIGN_MASK;   PC=PC+#r; }</pre>
<pre>if ([!]cmp.gt(Ns.new,#-1)) jump:&lt;hint&gt; #r9:2</pre>	<pre>if ([!](Ns.new&gt;(-1))) {   apply_extension(#r);   #r=#r &amp; ~PCALIGN_MASK;   PC=PC+#r; }</pre>
<pre>if ([!]cmp.gt(Ns.new,#U5)) jump:&lt;hint&gt; #r9:2</pre>	<pre>if ([!](Ns.new&gt;(#U))) {   apply_extension(#r);   #r=#r &amp; ~PCALIGN_MASK;   PC=PC+#r; }</pre>
<pre>if ([!]cmp.gt(Ns.new,Rt)) jump:&lt;hint&gt; #r9:2</pre>	<pre>if ([!](Ns.new&gt;Rt)) {   apply_extension(#r);   #r=#r &amp; ~PCALIGN_MASK;   PC=PC+#r; }</pre>



Syntax	Behavior
<code>if ([!]cmp.gt(Rt,Ns.new) jump:&lt;hint&gt; #r9:2</code>	<code>if ([!](Rt&gt;Ns.new)) {   apply_extension(#r);   #r=#r &amp; ~PCALIGN_MASK;   PC=PC+#r; }</code>
<code>if ([!]cmp.gtu(Ns.new,#U5) jump:&lt;hint&gt; #r9:2</code>	<code>if ([!](Ns.new.uw[0]&gt;(#U))) {   apply_extension(#r);   #r=#r &amp; ~PCALIGN_MASK;   PC=PC+#r; }</code>
<code>if ([!]cmp.gtu(Ns.new,Rt) jump:&lt;hint&gt; #r9:2</code>	<code>if ([!](Ns.new.uw[0]&gt;Rt.uw[0])) {   apply_extension(#r);   #r=#r &amp; ~PCALIGN_MASK;   PC=PC+#r; }</code>
<code>if ([!]cmp.gtu(Rt,Ns.new) jump:&lt;hint&gt; #r9:2</code>	<code>if ([!](Rt.uw[0]&gt;Ns.new.uw[0])) {   apply_extension(#r);   #r=#r &amp; ~PCALIGN_MASK;   PC=PC+#r; }</code>
<code>if ([!]tstbit(Ns.new,#0) jump:&lt;hint&gt; #r9:2</code>	<code>if ([!](Ns.new &amp; 1)) {   apply_extension(#r);   #r=#r &amp; ~PCALIGN_MASK;   PC=PC+#r; }</code>

**Class: NV (slots 0)**

**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS													s3			Parse		t5															
0	0	1	0	0	0	0	0	0	0	0	i	i	-	s	s	s	P	P	0	t	t	t	t	t	i	i	i	i	i	i	i	-	if (cmp.eq(Ns.new,Rt)) jump:nt #r9:2
0	0	1	0	0	0	0	0	0	0	0	i	i	-	s	s	s	P	P	1	t	t	t	t	t	i	i	i	i	i	i	i	-	if (cmp.eq(Ns.new,Rt)) jump:t #r9:2
0	0	1	0	0	0	0	0	0	1	i	i	-	s	s	s	P	P	0	t	t	t	t	t	i	i	i	i	i	i	i	-	if (!cmp.eq(Ns.new,Rt)) jump:nt #r9:2	
0	0	1	0	0	0	0	0	0	1	i	i	-	s	s	s	P	P	1	t	t	t	t	t	i	i	i	i	i	i	i	-	if (!cmp.eq(Ns.new,Rt)) jump:t #r9:2	
0	0	1	0	0	0	0	0	1	0	i	i	-	s	s	s	P	P	0	t	t	t	t	t	i	i	i	i	i	i	i	-	if (cmp.gt(Ns.new,Rt)) jump:nt #r9:2	
0	0	1	0	0	0	0	0	1	0	i	i	-	s	s	s	P	P	1	t	t	t	t	t	i	i	i	i	i	i	i	-	if (cmp.gt(Ns.new,Rt)) jump:t #r9:2	
0	0	1	0	0	0	0	0	1	1	i	i	-	s	s	s	P	P	0	t	t	t	t	t	i	i	i	i	i	i	i	-	if (!cmp.gt(Ns.new,Rt)) jump:nt #r9:2	
0	0	1	0	0	0	0	0	1	1	i	i	-	s	s	s	P	P	1	t	t	t	t	t	i	i	i	i	i	i	i	-	if (!cmp.gt(Ns.new,Rt)) jump:t #r9:2	
0	0	1	0	0	0	0	1	0	0	i	i	-	s	s	s	P	P	0	t	t	t	t	t	i	i	i	i	i	i	i	-	if (cmp.gtu(Ns.new,Rt)) jump:nt #r9:2	
0	0	1	0	0	0	0	1	0	0	i	i	-	s	s	s	P	P	1	t	t	t	t	t	i	i	i	i	i	i	i	-	if (cmp.gtu(Ns.new,Rt)) jump:t #r9:2	
0	0	1	0	0	0	0	1	0	1	i	i	-	s	s	s	P	P	0	t	t	t	t	t	i	i	i	i	i	i	i	-	if (!cmp.gtu(Ns.new,Rt)) jump:nt #r9:2	

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
0	0	1	0	0	0	0	1	0	1	i	i	-	s	s	s	P	P	1	t	t	t	t	t	t	i	i	i	i	i	i	i	-	if (!cmp.gtu(Ns.new,Rt)) jump:t #r9:2
0	0	1	0	0	0	0	1	1	0	i	i	-	s	s	s	P	P	0	t	t	t	t	t	t	i	i	i	i	i	i	i	-	if (cmp.gt(Rt,Ns.new)) jump:nt #r9:2
0	0	1	0	0	0	0	1	1	0	i	i	-	s	s	s	P	P	1	t	t	t	t	t	t	i	i	i	i	i	i	i	-	if (cmp.gt(Rt,Ns.new)) jump:t #r9:2
0	0	1	0	0	0	0	1	1	1	i	i	-	s	s	s	P	P	0	t	t	t	t	t	t	i	i	i	i	i	i	i	-	if (!cmp.gt(Rt,Ns.new)) jump:nt #r9:2
0	0	1	0	0	0	0	1	1	1	i	i	-	s	s	s	P	P	1	t	t	t	t	t	t	i	i	i	i	i	i	i	-	if (!cmp.gt(Rt,Ns.new)) jump:t #r9:2
0	0	1	0	0	0	1	0	0	0	i	i	-	s	s	s	P	P	0	t	t	t	t	t	t	i	i	i	i	i	i	i	-	if (cmp.gtu(Rt,Ns.new)) jump:nt #r9:2
0	0	1	0	0	0	1	0	0	0	i	i	-	s	s	s	P	P	1	t	t	t	t	t	t	i	i	i	i	i	i	i	-	if (cmp.gtu(Rt,Ns.new)) jump:t #r9:2
0	0	1	0	0	0	1	0	0	1	i	i	-	s	s	s	P	P	0	t	t	t	t	t	t	i	i	i	i	i	i	i	-	if (!cmp.gtu(Rt,Ns.new)) jump:nt #r9:2
0	0	1	0	0	0	1	0	0	1	i	i	-	s	s	s	P	P	1	t	t	t	t	t	t	i	i	i	i	i	i	i	-	if (!cmp.gtu(Rt,Ns.new)) jump:t #r9:2
ICLASS													s3		Parse																		
0	0	1	0	0	1	0	0	0	0	i	i	-	s	s	s	P	P	0	l	l	l	l	l	l	i	i	i	i	i	i	-	if (cmp.eq(Ns.new,#U5)) jump:nt #r9:2	
0	0	1	0	0	1	0	0	0	0	i	i	-	s	s	s	P	P	1	l	l	l	l	l	l	i	i	i	i	i	i	-	if (cmp.eq(Ns.new,#U5)) jump:t #r9:2	
0	0	1	0	0	1	0	0	0	1	i	i	-	s	s	s	P	P	0	l	l	l	l	l	l	i	i	i	i	i	i	-	if (!cmp.eq(Ns.new,#U5)) jump:nt #r9:2	
0	0	1	0	0	1	0	0	0	1	i	i	-	s	s	s	P	P	1	l	l	l	l	l	l	i	i	i	i	i	i	-	if (!cmp.eq(Ns.new,#U5)) jump:t #r9:2	
0	0	1	0	0	1	0	0	1	0	i	i	-	s	s	s	P	P	0	l	l	l	l	l	l	i	i	i	i	i	i	-	if (cmp.gt(Ns.new,#U5)) jump:nt #r9:2	
0	0	1	0	0	1	0	0	1	0	i	i	-	s	s	s	P	P	1	l	l	l	l	l	l	i	i	i	i	i	i	-	if (cmp.gt(Ns.new,#U5)) jump:t #r9:2	
0	0	1	0	0	1	0	0	1	1	i	i	-	s	s	s	P	P	0	l	l	l	l	l	l	i	i	i	i	i	i	-	if (!cmp.gt(Ns.new,#U5)) jump:nt #r9:2	
0	0	1	0	0	1	0	0	1	1	i	i	-	s	s	s	P	P	1	l	l	l	l	l	l	i	i	i	i	i	i	-	if (!cmp.gt(Ns.new,#U5)) jump:t #r9:2	
0	0	1	0	0	1	0	1	0	0	i	i	-	s	s	s	P	P	0	l	l	l	l	l	l	i	i	i	i	i	i	-	if (cmp.gtu(Ns.new,#U5)) jump:nt #r9:2	
0	0	1	0	0	1	0	1	0	0	i	i	-	s	s	s	P	P	1	l	l	l	l	l	l	i	i	i	i	i	i	-	if (cmp.gtu(Ns.new,#U5)) jump:t #r9:2	
0	0	1	0	0	1	0	1	0	1	i	i	-	s	s	s	P	P	0	l	l	l	l	l	l	i	i	i	i	i	i	-	if (!cmp.gtu(Ns.new,#U5)) jump:nt #r9:2	
0	0	1	0	0	1	0	1	0	1	i	i	-	s	s	s	P	P	1	l	l	l	l	l	l	i	i	i	i	i	i	-	if (!cmp.gtu(Ns.new,#U5)) jump:t #r9:2	
0	0	1	0	0	1	0	1	1	0	i	i	-	s	s	s	P	P	0	-	-	-	-	-	-	i	i	i	i	i	i	-	if (tstbit(Ns.new,#0)) jump:nt #r9:2	
0	0	1	0	0	1	0	1	1	0	i	i	-	s	s	s	P	P	1	-	-	-	-	-	-	i	i	i	i	i	i	-	if (tstbit(Ns.new,#0)) jump:t #r9:2	
0	0	1	0	0	1	0	1	1	1	i	i	-	s	s	s	P	P	0	-	-	-	-	-	-	i	i	i	i	i	i	-	if (!tstbit(Ns.new,#0)) jump:nt #r9:2	
0	0	1	0	0	1	0	1	1	1	i	i	-	s	s	s	P	P	1	-	-	-	-	-	-	i	i	i	i	i	i	-	if (!tstbit(Ns.new,#0)) jump:t #r9:2	
0	0	1	0	0	1	1	0	0	0	i	i	-	s	s	s	P	P	0	-	-	-	-	-	-	i	i	i	i	i	i	-	if (cmp.eq(Ns.new,#-1)) jump:nt #r9:2	
0	0	1	0	0	1	1	0	0	0	i	i	-	s	s	s	P	P	1	-	-	-	-	-	-	i	i	i	i	i	i	-	if (cmp.eq(Ns.new,#-1)) jump:t #r9:2	
0	0	1	0	0	1	1	0	0	1	i	i	-	s	s	s	P	P	0	-	-	-	-	-	-	i	i	i	i	i	i	-	if (!cmp.eq(Ns.new,#-1)) jump:nt #r9:2	

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	1	0	0	1	1	0	0	1	i	i	-	s	s	s	P	P	1	-	-	-	-	-	i	i	i	i	i	i	i	-	if (!cmp.eq(Ns.new,#-1)) jump:t #r9:2
0	0	1	0	0	1	1	0	1	0	i	i	-	s	s	s	P	P	0	-	-	-	-	-	i	i	i	i	i	i	i	-	if (cmp.gt(Ns.new,#-1)) jump:nt #r9:2
0	0	1	0	0	1	1	0	1	0	i	i	-	s	s	s	P	P	1	-	-	-	-	-	i	i	i	i	i	i	i	-	if (cmp.gt(Ns.new,#-1)) jump:t #r9:2
0	0	1	0	0	1	1	0	1	1	i	i	-	s	s	s	P	P	0	-	-	-	-	-	i	i	i	i	i	i	i	-	if (!cmp.gt(Ns.new,#-1)) jump:nt #r9:2
0	0	1	0	0	1	1	0	1	1	i	i	-	s	s	s	P	P	1	-	-	-	-	-	i	i	i	i	i	i	i	-	if (!cmp.gt(Ns.new,#-1)) jump:t #r9:2

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
s3	Field to encode register s
t5	Field to encode register t

## 11.7.2 NV ST

The NV ST instruction subclass includes store instructions which take the register source operand from another instruction in the same packet.

### Store new-value byte

Store the least-significant byte in a source register in memory at the effective address.

Syntax	Behavior
<code>memb (Re=#U6) =Nt.new</code>	<code>apply_extension (#U);</code> <code>EA=#U;</code> <code>*EA = Nt.new.b[0];</code> <code>Re=#U;</code>
<code>memb (Rs+#s11:0) =Nt.new</code>	<code>apply_extension (#s);</code> <code>EA=Rs+#s;</code> <code>*EA = Nt.new.b[0];</code>
<code>memb (Rs+Ru&lt;&lt;#u2) =Nt.new</code>	<code>EA=Rs+ (Ru&lt;&lt;#u);</code> <code>*EA = Nt.new.b[0];</code>
<code>memb (Ru&lt;&lt;#u2+#U6) =Nt.new</code>	<code>apply_extension (#U);</code> <code>EA=#U+ (Ru&lt;&lt;#u);</code> <code>*EA = Nt.new.b[0];</code>
<code>memb (Rx++#s4:0) =Nt.new</code>	<code>EA=Rx;</code> <code>Rx=Rx+#s;</code> <code>*EA = Nt.new.b[0];</code>
<code>memb (Rx++#s4:0:circ (Mu)) =Nt.new</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, #s, MuV);</code> <code>*EA = Nt.new.b[0];</code>
<code>memb (Rx++I:circ (Mu)) =Nt.new</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, I&lt;&lt;0, MuV);</code> <code>*EA = Nt.new.b[0];</code>
<code>memb (Rx++Mu) =Nt.new</code>	<code>EA=Rx;</code> <code>Rx=Rx+MuV;</code> <code>*EA = Nt.new.b[0];</code>
<code>memb (Rx++Mu:brev) =Nt.new</code>	<code>EA=Rx.h[1]   brev (Rx.h[0]);</code> <code>Rx=Rx+MuV;</code> <code>*EA = Nt.new.b[0];</code>
<code>memb (gp+#u16:0) =Nt.new</code>	<code>apply_extension (#u);</code> <code>EA=(Constant_extended ? (0) : GP)+#u;</code> <code>*EA = Nt.new.b[0];</code>

### Class: NV (slots 0)

#### Notes

- Forms of this instruction that use a new-value operand produced in the packet must execute on slot 0.
- This instruction can execute only in slot 0, even though it is an ST instruction.

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS										s5					Parse			u5					t3									
0	0	1	1	1	0	1	1	1	0	1	s	s	s	s	s	P	P	i	u	u	u	u	u	i	-	-	0	0	t	t	t	memb(Rs+Ru<<#u2)=Nt.new
ICLASS			Type							Parse					t3																	
0	1	0	0	1	i	i	0	1	0	1	i	i	i	i	i	P	P	i	0	0	t	t	t	i	i	i	i	i	i	i	i	memb(gp+#u16:0)=Nt.new
ICLASS			Amode		Type		UN	s5					Parse			t3																
1	0	1	0	0	i	i	1	1	0	1	s	s	s	s	s	P	P	i	0	0	t	t	t	i	i	i	i	i	i	i	i	memb(Rs+#s11:0)=Nt.new
ICLASS			Amode		Type		UN	x5					Parse			u1	t3															
1	0	1	0	1	0	0	1	1	0	1	x	x	x	x	x	P	P	u	0	0	t	t	t	0	-	-	-	-	1	-	memb(Rx++l:circ(Mu))=Nt.new	
1	0	1	0	1	0	0	1	1	0	1	x	x	x	x	x	P	P	u	0	0	t	t	t	0	i	i	i	i	-	0	-	memb(Rx++#s4:0:circ(Mu))=Nt.new
ICLASS			Amode		Type		UN	e5					Parse			t3																
1	0	1	0	1	0	1	1	1	0	1	e	e	e	e	e	P	P	0	0	0	t	t	t	1	-	l	l	l	l	l	l	memb(Re=#U6)=Nt.new
ICLASS			Amode		Type		UN	x5					Parse			t3																
1	0	1	0	1	0	1	1	1	0	1	x	x	x	x	x	P	P	0	0	0	t	t	t	0	i	i	i	i	-	0	-	memb(Rx++#s4:0)=Nt.new
ICLASS			Amode		Type		UN	u5					Parse			t3																
1	0	1	0	1	1	0	1	1	0	1	u	u	u	u	u	P	P	i	0	0	t	t	t	1	i	l	l	l	l	l	l	memb(Ru<<#u2+#U6)=Nt.new
ICLASS			Amode		Type		UN	x5					Parse			u1	t3															
1	0	1	0	1	1	0	1	1	0	1	x	x	x	x	x	P	P	u	0	0	t	t	t	0	-	-	-	-	-	-	-	memb(Rx++Mu)=Nt.new
1	0	1	0	1	1	1	1	1	0	1	x	x	x	x	x	P	P	u	0	0	t	t	t	0	-	-	-	-	-	-	-	memb(Rx++Mu:brev)=Nt.new

Field name	Description
ICLASS	Instruction class
Type	Type
Parse	Packet/loop parse bits
e5	Field to encode register e
s5	Field to encode register s
t3	Field to encode register t
u1	Field to encode register u
u5	Field to encode register u
x5	Field to encode register x
Amode	Amode
UN	Unsigned

## Store new-value byte conditionally

Store the least-significant byte in a source register in memory at the effective address.

This instruction is conditional based on a predicate value. If the predicate is true, the instruction is performed, otherwise it is treated as a NOP.

Syntax	Behavior
<pre>if ([!]Pv[.new]) memb(#u6)=Nt.new</pre>	<pre>apply_extension(#u); EA=#u; if ([!]Pv[.new][0]) {     *EA = Nt[.new].b[0]; } else {     NOP; }</pre>
<pre>if ([!]Pv[.new]) memb(Rs+#u6:0)=Nt.new</pre>	<pre>apply_extension(#u); EA=Rs+#u; if ([!]Pv[.new][0]) {     *EA = Nt[.new].b[0]; } else {     NOP; }</pre>
<pre>if ([!]Pv[.new]) memb(Rs+Ru&lt;&lt;#u2)=Nt.new</pre>	<pre>EA=Rs+(Ru&lt;&lt;#u); if ([!]Pv[.new][0]) {     *EA = Nt[.new].b[0]; } else {     NOP; }</pre>
<pre>if ([!]Pv[.new]) memb(Rx++#s4:0)=Nt.new</pre>	<pre>EA=Rx; if ([!]Pv[.new][0]){     Rx=Rx+#s;     *EA = Nt[.new].b[0]; } else {     NOP; }</pre>

### Class: NV (slots 0)

#### Notes

- Forms of this instruction which use a new-value operand produced in the packet must execute on slot 0.
- This instruction can execute only in slot 0, even though it is an ST instruction.

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS												s5					Parse		u5					t3								
0	0	1	1	0	1	0	0	1	0	1	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	0	0	t	t	t	if (Pv) memb(Rs+Ru<<#u2)=Nt.new
0	0	1	1	0	1	0	1	1	0	1	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	0	0	t	t	t	if (!Pv) memb(Rs+Ru<<#u2)=Nt.new
0	0	1	1	0	1	1	0	1	0	1	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	0	0	t	t	t	if (Pv.new) memb(Rs+Ru<<#u2)=Nt.new
0	0	1	1	0	1	1	1	1	0	1	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	0	0	t	t	t	if (!Pv.new) memb(Rs+Ru<<#u2)=Nt.new
ICLASS				Sense		PredNew		Type			s5					Parse		t3														
0	1	0	0	0	0	0	0	1	0	1	s	s	s	s	s	P	P	i	0	0	t	t	t	i	i	i	i	i	0	v	v	if (Pv) memb(Rs+#u6:0)=Nt.new
0	1	0	0	0	0	1	0	1	0	1	s	s	s	s	s	P	P	i	0	0	t	t	t	i	i	i	i	i	0	v	v	if (Pv.new) memb(Rs+#u6:0)=Nt.new
0	1	0	0	0	1	0	0	1	0	1	s	s	s	s	s	P	P	i	0	0	t	t	t	i	i	i	i	i	0	v	v	if (!Pv) memb(Rs+#u6:0)=Nt.new
0	1	0	0	0	1	1	0	1	0	1	s	s	s	s	s	P	P	i	0	0	t	t	t	i	i	i	i	i	0	v	v	if (!Pv.new) memb(Rs+#u6:0)=Nt.new
ICLASS				Amode			Type			UN	x5					Parse		t3														
1	0	1	0	1	0	1	1	1	0	1	x	x	x	x	x	P	P	1	0	0	t	t	t	0	i	i	i	i	0	v	v	if (Pv) memb(Rx++#s4:0)=Nt.new
1	0	1	0	1	0	1	1	1	0	1	x	x	x	x	x	P	P	1	0	0	t	t	t	0	i	i	i	i	1	v	v	if (!Pv) memb(Rx++#s4:0)=Nt.new
1	0	1	0	1	0	1	1	1	0	1	x	x	x	x	x	P	P	1	0	0	t	t	t	1	i	i	i	i	0	v	v	if (Pv.new) memb(Rx++#s4:0)=Nt.new
1	0	1	0	1	0	1	1	1	0	1	x	x	x	x	x	P	P	1	0	0	t	t	t	1	i	i	i	i	1	v	v	if (!Pv.new) memb(Rx++#s4:0)=Nt.new
ICLASS				Amode			Type			UN	Parse					t3																
1	0	1	0	1	1	1	1	1	0	1	-	-	-	i	i	P	P	0	0	0	t	t	t	1	i	i	i	i	0	v	v	if (Pv) memb(#u6)=Nt.new
1	0	1	0	1	1	1	1	1	0	1	-	-	-	i	i	P	P	0	0	0	t	t	t	1	i	i	i	i	1	v	v	if (!Pv) memb(#u6)=Nt.new
1	0	1	0	1	1	1	1	1	0	1	-	-	-	i	i	P	P	1	0	0	t	t	t	1	i	i	i	i	0	v	v	if (Pv.new) memb(#u6)=Nt.new
1	0	1	0	1	1	1	1	1	0	1	-	-	-	i	i	P	P	1	0	0	t	t	t	1	i	i	i	i	1	v	v	if (!Pv.new) memb(#u6)=Nt.new

<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
Type	Type
PredNew	PredNew
Sense	Sense
Parse	Packet/loop parse bits
s5	Field to encode register s
t3	Field to encode register t
u5	Field to encode register u
v2	Field to encode register v
x5	Field to encode register x
Amode	Amode
UN	Unsigned

## Store new-value halfword

Store the upper or lower 16-bits of a source register in memory at the effective address.

Syntax	Behavior
<code>memh (Re=#U6)=Nt.new</code>	<code>apply_extension(#U);</code> <code>EA=#U;</code> <code>*EA = Nt.new.h[0];</code> <code>Re=#U;</code>
<code>memh (Rs+#s11:1)=Nt.new</code>	<code>apply_extension(#s);</code> <code>EA=Rs+#s;</code> <code>*EA = Nt.new.h[0];</code>
<code>memh (Rs+Ru&lt;&lt;#u2)=Nt.new</code>	<code>EA=Rs+ (Ru&lt;&lt;#u);</code> <code>*EA = Nt.new.h[0];</code>
<code>memh (Ru&lt;&lt;#u2+#U6)=Nt.new</code>	<code>apply_extension(#U);</code> <code>EA=#U+ (Ru&lt;&lt;#u);</code> <code>*EA = Nt.new.h[0];</code>
<code>memh (Rx++#s4:1)=Nt.new</code>	<code>EA=Rx;</code> <code>Rx=Rx+#s;</code> <code>*EA = Nt.new.h[0];</code>
<code>memh (Rx++#s4:1:circ (Mu))=Nt.new</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, #s, MuV);</code> <code>*EA = Nt.new.h[0];</code>
<code>memh (Rx++I:circ (Mu))=Nt.new</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, I&lt;&lt;1, MuV);</code> <code>*EA = Nt.new.h[0];</code>
<code>memh (Rx++Mu)=Nt.new</code>	<code>EA=Rx;</code> <code>Rx=Rx+MuV;</code> <code>*EA = Nt.new.h[0];</code>
<code>memh (Rx++Mu:brev)=Nt.new</code>	<code>EA=Rx.h[1]   brev (Rx.h[0]);</code> <code>Rx=Rx+MuV;</code> <code>*EA = Nt.new.h[0];</code>
<code>memh (gp+#u16:1)=Nt.new</code>	<code>apply_extension(#u);</code> <code>EA=(Constant_extended ? (0) : GP)+#u;</code> <code>*EA = Nt.new.h[0];</code>

### Class: NV (slots 0)

#### Notes

- Forms of this instruction that use a new-value operand produced in the packet must execute on slot 0.
- This instruction can execute only in slot 0, even though it is an ST instruction.



### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS										s5					Parse		u5					t3										
0	0	1	1	1	0	1	1	1	0	1	s	s	s	s	s	P	P	i	u	u	u	u	u	i	-	-	0	1	t	t	t	memh(Rs+Ru<<#u2)=Nt.new
ICLASS			Type							Parse					t3																	
0	1	0	0	1	i	i	0	1	0	1	i	i	i	i	i	P	P	i	0	1	t	t	t	i	i	i	i	i	i	i	i	memh(gp+#u16:1)=Nt.new
ICLASS			Amode		Type		UN	s5					Parse		t3																	
1	0	1	0	0	i	i	1	1	0	1	s	s	s	s	s	P	P	i	0	1	t	t	t	i	i	i	i	i	i	i	i	memh(Rs+#s11:1)=Nt.new
ICLASS			Amode		Type		UN	x5					Parse		u1	t3																
1	0	1	0	1	0	0	1	1	0	1	x	x	x	x	x	P	P	u	0	1	t	t	t	0	-	-	-	-	1	-	memh(Rx++l:circ(Mu))=Nt.new	
1	0	1	0	1	0	0	1	1	0	1	x	x	x	x	x	P	P	u	0	1	t	t	t	0	i	i	i	i	-	0	-	memh(Rx++#s4:1:circ(Mu))=Nt.new
ICLASS			Amode		Type		UN	e5					Parse		t3																	
1	0	1	0	1	0	1	1	1	0	1	e	e	e	e	e	P	P	0	0	1	t	t	t	1	-	l	l	l	l	l	l	memh(Re=#U6)=Nt.new
ICLASS			Amode		Type		UN	x5					Parse		t3																	
1	0	1	0	1	0	1	1	1	0	1	x	x	x	x	x	P	P	0	0	1	t	t	t	0	i	i	i	i	-	0	-	memh(Rx++#s4:1)=Nt.new
ICLASS			Amode		Type		UN	u5					Parse		t3																	
1	0	1	0	1	1	0	1	1	0	1	u	u	u	u	u	P	P	i	0	1	t	t	t	1	i	l	l	l	l	l	l	memh(Ru<<#u2+#U6)=Nt.new
ICLASS			Amode		Type		UN	x5					Parse		u1	t3																
1	0	1	0	1	1	0	1	1	0	1	x	x	x	x	x	P	P	u	0	1	t	t	t	0	-	-	-	-	-	-	-	memh(Rx++Mu)=Nt.new
1	0	1	0	1	1	1	1	1	0	1	x	x	x	x	x	P	P	u	0	1	t	t	t	0	-	-	-	-	-	-	-	memh(Rx++Mu:brev)=Nt.new

Field name	Description
ICLASS	Instruction class
Type	Type
Parse	Packet/loop parse bits
e5	Field to encode register e
s5	Field to encode register s
t3	Field to encode register t
u1	Field to encode register u
u5	Field to encode register u
x5	Field to encode register x
Amode	Amode
UN	Unsigned

## Store new-value halfword conditionally

Store the upper or lower 16 bits of a source register in memory at the effective address.

This instruction is conditional based on a predicate value. If the predicate is true, the instruction is performed, otherwise it is treated as a NOP.

Syntax	Behavior
<code>if ([!]Pv[.new]) memh (#u6)=Nt.new</code>	<pre> apply_extension(#u); EA=#u; if ([!]Pv[.new][0]) {     *EA = Nt[.new].h[0]; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memh (Rs+#u6:1)=Nt.new</code>	<pre> apply_extension(#u); EA=Rs+#u; if ([!]Pv[.new][0]) {     *EA = Nt[.new].h[0]; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memh (Rs+Ru&lt;&lt;#u2)=Nt.new</code>	<pre> EA=Rs+(Ru&lt;&lt;#u); if ([!]Pv[.new][0]) {     *EA = Nt[.new].h[0]; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memh (Rx++#s4:1)=Nt.new</code>	<pre> EA=Rx; if ([!]Pv[.new][0]) {     Rx=Rx+#s;     *EA = Nt[.new].h[0]; } else {     NOP; } </pre>

### Class: NV (slots 0)

#### Notes

- Forms of this instruction that use a new-value operand produced in the packet must execute on slot 0.
- This instruction can execute only in slot 0, even though it is an ST instruction.

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS											s5					Parse		u5					t3									
0	0	1	1	0	1	0	0	1	0	1	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	0	1	t	t	t	if (Pv) memh(Rs+Ru<<#u2)=Nt.new
0	0	1	1	0	1	0	1	1	0	1	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	0	1	t	t	t	if (!Pv) memh(Rs+Ru<<#u2)=Nt.new
0	0	1	1	0	1	1	0	1	0	1	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	0	1	t	t	t	if (Pv.new) memh(Rs+Ru<<#u2)=Nt.new
0	0	1	1	0	1	1	1	1	0	1	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	0	1	t	t	t	if (!Pv.new) memh(Rs+Ru<<#u2)=Nt.new
ICLASS			Sense		PredNew		Type			s5					Parse		t3															
0	1	0	0	0	0	0	0	1	0	1	s	s	s	s	s	P	P	i	0	1	t	t	t	i	i	i	i	i	0	v	v	if (Pv) memh(Rs+#u6:1)=Nt.new
0	1	0	0	0	0	1	0	1	0	1	s	s	s	s	s	P	P	i	0	1	t	t	t	i	i	i	i	i	0	v	v	if (Pv.new) memh(Rs+#u6:1)=Nt.new
0	1	0	0	0	1	0	0	1	0	1	s	s	s	s	s	P	P	i	0	1	t	t	t	i	i	i	i	i	0	v	v	if (!Pv) memh(Rs+#u6:1)=Nt.new
0	1	0	0	0	1	1	0	1	0	1	s	s	s	s	s	P	P	i	0	1	t	t	t	i	i	i	i	i	0	v	v	if (!Pv.new) memh(Rs+#u6:1)=Nt.new
ICLASS			Amode			Type			UN	x5					Parse		t3															
1	0	1	0	1	0	1	1	1	0	1	x	x	x	x	x	P	P	1	0	1	t	t	t	0	i	i	i	i	0	v	v	if (Pv) memh(Rx++#s4:1)=Nt.new
1	0	1	0	1	0	1	1	1	0	1	x	x	x	x	x	P	P	1	0	1	t	t	t	0	i	i	i	i	1	v	v	if (!Pv) memh(Rx++#s4:1)=Nt.new
1	0	1	0	1	0	1	1	1	0	1	x	x	x	x	x	P	P	1	0	1	t	t	t	1	i	i	i	i	0	v	v	if (Pv.new) memh(Rx++#s4:1)=Nt.new
1	0	1	0	1	0	1	1	1	0	1	x	x	x	x	x	P	P	1	0	1	t	t	t	1	i	i	i	i	1	v	v	if (!Pv.new) memh(Rx++#s4:1)=Nt.new
ICLASS			Amode			Type			UN						Parse		t3															
1	0	1	0	1	1	1	1	1	0	1	-	-	-	i	i	P	P	0	0	1	t	t	t	1	i	i	i	i	0	v	v	if (Pv) memh(#u6)=Nt.new
1	0	1	0	1	1	1	1	1	0	1	-	-	-	i	i	P	P	0	0	1	t	t	t	1	i	i	i	i	1	v	v	if (!Pv) memh(#u6)=Nt.new
1	0	1	0	1	1	1	1	1	0	1	-	-	-	i	i	P	P	1	0	1	t	t	t	1	i	i	i	i	0	v	v	if (Pv.new) memh(#u6)=Nt.new
1	0	1	0	1	1	1	1	1	0	1	-	-	-	i	i	P	P	1	0	1	t	t	t	1	i	i	i	i	1	v	v	if (!Pv.new) memh(#u6)=Nt.new

<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
Type	Type
PredNew	PredNew
Sense	Sense
Parse	Packet/loop parse bits
s5	Field to encode register s
t3	Field to encode register t
u5	Field to encode register u
v2	Field to encode register v
x5	Field to encode register x
Amode	Amode
UN	Unsigned

## Store new-value word

Store a 32-bit register in memory at the effective address.

Syntax	Behavior
<code>memw (Re=#U6)=Nt.new</code>	<code>apply_extension (#U);</code> <code>EA=#U;</code> <code>*EA = Nt.new;</code> <code>Re=#U;</code>
<code>memw (Rs+#s11:2)=Nt.new</code>	<code>apply_extension (#s);</code> <code>EA=Rs+#s;</code> <code>*EA = Nt.new;</code>
<code>memw (Rs+Ru&lt;&lt;#u2)=Nt.new</code>	<code>EA=Rs+ (Ru&lt;&lt;#u);</code> <code>*EA = Nt.new;</code>
<code>memw (Ru&lt;&lt;#u2+#U6)=Nt.new</code>	<code>apply_extension (#U);</code> <code>EA=#U+ (Ru&lt;&lt;#u);</code> <code>*EA = Nt.new;</code>
<code>memw (Rx++#s4:2)=Nt.new</code>	<code>EA=Rx;</code> <code>Rx=Rx+#s;</code> <code>*EA = Nt.new;</code>
<code>memw (Rx++#s4:2:circ (Mu))=Nt.new</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, #s, MuV);</code> <code>*EA = Nt.new;</code>
<code>memw (Rx++I:circ (Mu))=Nt.new</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, I&lt;&lt;2, MuV);</code> <code>*EA = Nt.new;</code>
<code>memw (Rx++Mu)=Nt.new</code>	<code>EA=Rx;</code> <code>Rx=Rx+MuV;</code> <code>*EA = Nt.new;</code>
<code>memw (Rx++Mu:brev)=Nt.new</code>	<code>EA=Rx.h [1]   brev (Rx.h [0]);</code> <code>Rx=Rx+MuV;</code> <code>*EA = Nt.new;</code>
<code>memw (gp+#u16:2)=Nt.new</code>	<code>apply_extension (#u);</code> <code>EA=(Constant_extended ? (0) : GP)+#u;</code> <code>*EA = Nt.new;</code>

### Class: NV (slots 0)

#### Notes

- Forms of this instruction that use a new-value operand produced in the packet must execute on slot 0.
- This instruction can execute only in slot 0, even though it is an ST instruction.

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS											s5					Parse			u5					t3								
0	0	1	1	1	0	1	1	1	0	1	s	s	s	s	s	P	P	i	u	u	u	u	u	i	-	-	1	0	t	t	t	memw(Rs+Ru<<#u2)=Nt.new
ICLASS			Type								Parse					t3																
0	1	0	0	1	i	i	0	1	0	1	i	i	i	i	i	P	P	i	1	0	t	t	t	i	i	i	i	i	i	i	i	memw(gp+#u16:2)=Nt.new
ICLASS			Amode		Type		UN	s5					Parse			t3																
1	0	1	0	0	i	i	1	1	0	1	s	s	s	s	s	P	P	i	1	0	t	t	t	i	i	i	i	i	i	i	i	memw(Rs+#s11:2)=Nt.new
ICLASS			Amode		Type		UN	x5					Parse			u1	t3															
1	0	1	0	1	0	0	1	1	0	1	x	x	x	x	x	P	P	u	1	0	t	t	t	0	-	-	-	-	1	-	memw(Rx++l:circ(Mu))=Nt.new	
1	0	1	0	1	0	0	1	1	0	1	x	x	x	x	x	P	P	u	1	0	t	t	t	0	i	i	i	i	-	0	-	memw(Rx++#s4:2:circ(Mu))=Nt.new
ICLASS			Amode		Type		UN	e5					Parse			t3																
1	0	1	0	1	0	1	1	1	0	1	e	e	e	e	e	P	P	0	1	0	t	t	t	1	-	l	l	l	l	l	l	memw(Re=#U6)=Nt.new
ICLASS			Amode		Type		UN	x5					Parse			t3																
1	0	1	0	1	0	1	1	1	0	1	x	x	x	x	x	P	P	0	1	0	t	t	t	0	i	i	i	i	-	0	-	memw(Rx++#s4:2)=Nt.new
ICLASS			Amode		Type		UN	u5					Parse			t3																
1	0	1	0	1	1	0	1	1	0	1	u	u	u	u	u	P	P	i	1	0	t	t	t	1	i	l	l	l	l	l	l	memw(Ru<<#u2+#U6)=Nt.new
ICLASS			Amode		Type		UN	x5					Parse			u1	t3															
1	0	1	0	1	1	0	1	1	0	1	x	x	x	x	x	P	P	u	1	0	t	t	t	0	-	-	-	-	-	-	-	memw(Rx++Mu)=Nt.new
1	0	1	0	1	1	1	1	1	0	1	x	x	x	x	x	P	P	u	1	0	t	t	t	0	-	-	-	-	-	-	-	memw(Rx++Mu:brev)=Nt.new

Field name	Description
ICLASS	Instruction class
Type	Type
Parse	Packet/loop parse bits
e5	Field to encode register e
s5	Field to encode register s
t3	Field to encode register t
u1	Field to encode register u
u5	Field to encode register u
x5	Field to encode register x
Amode	Amode
UN	Unsigned

## Store new-value word conditionally

Store a 32-bit register in memory at the effective address.

This instruction is conditional based on a predicate value. If the predicate is true, the instruction is performed, otherwise it is treated as a NOP.

Syntax	Behavior
<code>if ([!]Pv[.new]) memw(#u6)=Nt.new</code>	<pre> apply_extension(#u); EA=#u; if ([!]Pv[.new][0]) {     *EA = Nt[.new]; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memw(Rs+#u6:2)=Nt.new</code>	<pre> apply_extension(#u); EA=Rs+#u; if ([!]Pv[.new][0]) {     *EA = Nt[.new]; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memw(Rs+Ru&lt;&lt;#u2)=Nt.new</code>	<pre> EA=Rs+(Ru&lt;&lt;#u); if ([!]Pv[.new][0]) {     *EA = Nt[.new]; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memw(Rx++#s4:2)=Nt.new</code>	<pre> EA=Rx; if ([!]Pv[.new][0]) {     Rx=Rx+#s;     *EA = Nt[.new]; } else {     NOP; } </pre>

### Class: NV (slots 0)

#### Notes

- Forms of this instruction that use a new-value operand produced in the packet must execute on slot 0.
- This instruction can execute only in slot 0, even though it is an ST instruction.

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS												s5					Parse		u5					t3								
0	0	1	1	0	1	0	0	1	0	1	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	1	0	t	t	t	if (Pv) memw(Rs+Ru<<#u2)=Nt.new
0	0	1	1	0	1	0	1	1	0	1	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	1	0	t	t	t	if (!Pv) memw(Rs+Ru<<#u2)=Nt.new
0	0	1	1	0	1	1	0	1	0	1	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	1	0	t	t	t	if (Pv.new) memw(Rs+Ru<<#u2)=Nt.new
0	0	1	1	0	1	1	1	1	0	1	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	1	0	t	t	t	if (!Pv.new) memw(Rs+Ru<<#u2)=Nt.new
ICLASS				Sense		PredNew		Type			s5					Parse		t3														
0	1	0	0	0	0	0	0	1	0	1	s	s	s	s	s	P	P	i	1	0	t	t	t	i	i	i	i	i	0	v	v	if (Pv) memw(Rs+#u6:2)=Nt.new
0	1	0	0	0	0	1	0	1	0	1	s	s	s	s	s	P	P	i	1	0	t	t	t	i	i	i	i	i	0	v	v	if (Pv.new) memw(Rs+#u6:2)=Nt.new
0	1	0	0	0	1	0	0	1	0	1	s	s	s	s	s	P	P	i	1	0	t	t	t	i	i	i	i	i	0	v	v	if (!Pv) memw(Rs+#u6:2)=Nt.new
0	1	0	0	0	1	1	0	1	0	1	s	s	s	s	s	P	P	i	1	0	t	t	t	i	i	i	i	i	0	v	v	if (!Pv.new) memw(Rs+#u6:2)=Nt.new
ICLASS				Amode			Type			UN	x5					Parse		t3														
1	0	1	0	1	0	1	1	1	0	1	x	x	x	x	x	P	P	1	1	0	t	t	t	0	i	i	i	i	0	v	v	if (Pv) memw(Rx++#s4:2)=Nt.new
1	0	1	0	1	0	1	1	1	0	1	x	x	x	x	x	P	P	1	1	0	t	t	t	0	i	i	i	i	1	v	v	if (!Pv) memw(Rx++#s4:2)=Nt.new
1	0	1	0	1	0	1	1	1	0	1	x	x	x	x	x	P	P	1	1	0	t	t	t	1	i	i	i	i	0	v	v	if (Pv.new) memw(Rx++#s4:2)=Nt.new
1	0	1	0	1	0	1	1	1	0	1	x	x	x	x	x	P	P	1	1	0	t	t	t	1	i	i	i	i	1	v	v	if (!Pv.new) memw(Rx++#s4:2)=Nt.new
ICLASS				Amode			Type			UN	Parse					t3																
1	0	1	0	1	1	1	1	1	0	1	-	-	-	i	i	P	P	0	1	0	t	t	t	1	i	i	i	i	0	v	v	if (Pv) memw(#u6)=Nt.new
1	0	1	0	1	1	1	1	1	0	1	-	-	-	i	i	P	P	0	1	0	t	t	t	1	i	i	i	i	1	v	v	if (!Pv) memw(#u6)=Nt.new
1	0	1	0	1	1	1	1	1	0	1	-	-	-	i	i	P	P	1	1	0	t	t	t	1	i	i	i	i	0	v	v	if (Pv.new) memw(#u6)=Nt.new
1	0	1	0	1	1	1	1	1	0	1	-	-	-	i	i	P	P	1	1	0	t	t	t	1	i	i	i	i	1	v	v	if (!Pv.new) memw(#u6)=Nt.new

<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
Type	Type
PredNew	PredNew
Sense	Sense
Parse	Packet/loop parse bits
s5	Field to encode register s
t3	Field to encode register t
u5	Field to encode register u
v2	Field to encode register v
x5	Field to encode register x
Amode	Amode
UN	Unsigned

## 11.8 ST

The ST instruction class includes store instructions, used to store values in memory.

ST instructions are executable on slot 0 and slot 1.

### Store doubleword

Store a 64-bit register pair in memory at the effective address.

Syntax	Behavior
<code>memd (Re=#U6) =Rtt</code>	<code>apply_extension (#U) ; EA=#U; *EA = Rtt; Re=#U;</code>
<code>memd (Rs+#s11:3) =Rtt</code>	<code>apply_extension (#s) ; EA=Rs+#s; *EA = Rtt;</code>
<code>memd (Rs+Ru&lt;&lt;#u2) =Rtt</code>	<code>EA=Rs+ (Ru&lt;&lt;#u) ; *EA = Rtt;</code>
<code>memd (Ru&lt;&lt;#u2+#U6) =Rtt</code>	<code>apply_extension (#U) ; EA=#U+ (Ru&lt;&lt;#u) ; *EA = Rtt;</code>
<code>memd (Rx++#s4:3) =Rtt</code>	<code>EA=Rx; Rx=Rx+#s; *EA = Rtt;</code>
<code>memd (Rx++#s4:3:circ (Mu) ) =Rtt</code>	<code>EA=Rx; Rx=Rx=circ_add (Rx, #s, MuV) ; *EA = Rtt;</code>
<code>memd (Rx++I:circ (Mu) ) =Rtt</code>	<code>EA=Rx; Rx=Rx=circ_add (Rx, I&lt;&lt;3, MuV) ; *EA = Rtt;</code>
<code>memd (Rx++Mu) =Rtt</code>	<code>EA=Rx; Rx=Rx+MuV; *EA = Rtt;</code>
<code>memd (Rx++Mu:brev) =Rtt</code>	<code>EA=Rx.h[1]   brev (Rx.h[0]) ; Rx=Rx+MuV; *EA = Rtt;</code>
<code>memd (gp+#u16:3) =Rtt</code>	<code>apply_extension (#u) ; EA= (Constant_extended ? (0) : GP)+#u; *EA = Rtt;</code>



### Class: ST (slots 0,1)

#### Intrinsics

memd(Rx++#s4:3:circ(Mu))=Rtt      void Q6\_memd\_IMP\_circ(void\*\* StartAddress, Word32 Is4\_3, Word32 Mu, Word64 Rtt, void\* BaseAddress)

memd(Rx++I:circ(Mu))=Rtt      void Q6\_memd\_MP\_circ(void\*\* StartAddress, Word32 Mu, Word64 Rtt, void\* BaseAddress)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS											s5					Parse		u5					t5										
0	0	1	1	1	0	1	1	1	1	0	s	s	s	s	s	P	P	i	u	u	u	u	u	i	-	-	t	t	t	t	t	memd(Rs+Ru<<#u2)=Rtt	
ICLASS				Type							s5					Parse		t5															
0	1	0	0	1	i	i	0	1	1	0	i	i	i	i	i	P	P	i	t	t	t	t	t	i	i	i	i	i	i	i	i	memd(gp+#u16:3)=Rtt	
ICLASS		Amode		Type			UN	s5					Parse		t5																		
1	0	1	0	0	i	i	1	1	1	0	s	s	s	s	s	P	P	i	t	t	t	t	t	i	i	i	i	i	i	i	i	memd(Rs+#s11:3)=Rtt	
ICLASS		Amode		Type			UN	x5					Parse		u1	t5																	
1	0	1	0	1	0	0	1	1	1	0	x	x	x	x	x	P	P	u	t	t	t	t	t	0	-	-	-	-	-	1	-	memd(Rx++I:circ(Mu))=Rtt	
ICLASS		Amode		Type			UN	x5					Parse		u1	t5																	
1	0	1	0	1	0	0	1	1	1	0	x	x	x	x	x	P	P	u	t	t	t	t	t	0	i	i	i	i	-	0	-	memd(Rx++#s4:3:circ(Mu))=Rtt	
ICLASS		Amode		Type			UN	e5					Parse		t5																		
1	0	1	0	1	0	1	1	1	1	0	e	e	e	e	e	P	P	0	t	t	t	t	t	1	-	l	l	l	l	l	l	memd(Re=#U6)=Rtt	
ICLASS		Amode		Type			UN	x5					Parse		t5																		
1	0	1	0	1	0	1	1	1	1	0	x	x	x	x	x	P	P	0	t	t	t	t	t	0	i	i	i	i	-	0	-	memd(Rx++#s4:3)=Rtt	
ICLASS		Amode		Type			UN	u5					Parse		t5																		
1	0	1	0	1	1	0	1	1	1	0	u	u	u	u	u	P	P	i	t	t	t	t	t	1	i	l	l	l	l	l	l	l	memd(Ru<<#u2+#U6)=Rtt
ICLASS		Amode		Type			UN	x5					Parse		u1	t5																	
1	0	1	0	1	1	0	1	1	1	0	x	x	x	x	x	P	P	u	t	t	t	t	t	0	-	-	-	-	-	-	-	-	memd(Rx++Mu)=Rtt
ICLASS		Amode		Type			UN	x5					Parse		u1	t5																	
1	0	1	0	1	1	1	1	1	1	0	x	x	x	x	x	P	P	u	t	t	t	t	t	0	-	-	-	-	-	-	-	-	memd(Rx++Mu:brev)=Rtt

Field name	Description
ICLASS	Instruction class
Type	Type
Parse	Packet/loop parse bits
e5	Field to encode register e
s5	Field to encode register s
t5	Field to encode register t
u1	Field to encode register u
u5	Field to encode register u
x5	Field to encode register x
Amode	Amode
UN	Unsigned

## Store-release doubleword

Store a 64-bit register pair in memory at the effective address. The store-release memory operation is observed after all preceding memory operations have been observed at the local point of serialization. A different order may be observed at the global point of serialization. (see Ordering and Synchronization).

When the :st (same domain) option is specified, the preceding memory operations are those that were committed on any thread with the same consistency domain before this instruction was committed.

When the :at (all threads) option is specified, the preceding memory operations are those that were committed on any thread before this instruction was committed.

The store release address is limited to certain memory regions. The following are excluded memory regions: AHB memory space, AXI M2 memory space, Hexagon memory cut-out is excluded with the exception of addressable TCM and VTCM memory, and memory with the CCCC types 2, 3, or 4 are excluded. The :st option does not apply to cache operation by index or global cache operation. The :st option does not apply a consistency domain to vector operations, but instead uses a per hardware thread ordering scope.

Syntax	Behavior
memd_rl(Rs):at=Rtt	EA=Rs; *EA = Rtt
memd_rl(Rs):st=Rtt	EA=Rs; *EA = Rtt

### Class: ST (slots 0)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				Amode				Type		UN	s5					Parse		t5					d2									
1	0	1	0	0	0	0	0	1	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	-	-	0	0	1	0	d	d	memd_rl(Rs):at=Rtt
1	0	1	0	0	0	0	0	1	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	-	-	1	0	1	0	d	d	memd_rl(Rs):st=Rtt

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d2	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
Amode	Amode
Type	Type
UN	Unsigned

## Store doubleword conditionally

Store a 64-bit register pair in memory at the effective address.

This instruction is conditional based on a predicate value. If the predicate is true, the instruction is performed, otherwise it is treated as a NOP.

Syntax	Behavior
<code>if ([!]Pv[.new]) memd(#u6)=Rtt</code>	<pre> apply_extension(#u); EA=#u; if ([!]Pv[.new][0]) {     *EA = Rtt; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memd(Rs+#u6:3)=Rtt</code>	<pre> apply_extension(#u); EA=Rs+#u; if ([!]Pv[.new][0]) {     *EA = Rtt; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memd(Rs+Ru&lt;&lt;#u2)=Rtt</code>	<pre> EA=Rs+(Ru&lt;&lt;#u); if ([!]Pv[.new][0]) {     *EA = Rtt; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memd(Rx++#s4:3)=Rtt</code>	<pre> EA=Rx; if ([!]Pv[.new][0]) {     Rx=Rx+#s;     *EA = Rtt; } else {     NOP; } </pre>

**Class: ST (slots 0,1)**

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS											s5					Parse		u5					t5									
0	0	1	1	0	1	0	0	1	1	0	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	t	t	t	t	t	if (Pv) memd(Rs+Ru<<#u2)=Rtt
0	0	1	1	0	1	0	1	1	1	0	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	t	t	t	t	t	if (!Pv) memd(Rs+Ru<<#u2)=Rtt
0	0	1	1	0	1	1	0	1	1	0	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	t	t	t	t	t	if (Pv.new) memd(Rs+Ru<<#u2)=Rtt
0	0	1	1	0	1	1	1	1	1	0	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	t	t	t	t	t	if (!Pv.new) memd(Rs+Ru<<#u2)=Rtt
ICLASS				Se	Pr	Type				s5					Parse		t5															
0	1	0	0	0	0	0	0	1	1	0	s	s	s	s	s	P	P	i	t	t	t	t	t	i	i	i	i	i	0	v	v	if (Pv) memd(Rs+#u6:3)=Rtt

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
0	1	0	0	0	0	1	0	1	1	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	0	v	v	if (Pv.new) memd(Rs+#u6:3)=Rtt
0	1	0	0	0	1	0	0	1	1	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	0	v	v	if (!Pv) memd(Rs+#u6:3)=Rtt
0	1	0	0	0	1	1	0	1	1	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	0	v	v	if (!Pv.new) memd(Rs+#u6:3)=Rtt
ICLASS			Amode			Type			UN	x5					Parse		t5																
1	0	1	0	1	0	1	1	1	1	0	x	x	x	x	x	P	P	1	t	t	t	t	t	0	i	i	i	i	0	v	v	if (Pv) memd(Rx++#s4:3)=Rtt	
1	0	1	0	1	0	1	1	1	1	0	x	x	x	x	x	P	P	1	t	t	t	t	t	0	i	i	i	i	1	v	v	if (!Pv) memd(Rx++#s4:3)=Rtt	
1	0	1	0	1	0	1	1	1	1	0	x	x	x	x	x	P	P	1	t	t	t	t	t	1	i	i	i	i	0	v	v	if (Pv.new) memd(Rx++#s4:3)=Rtt	
1	0	1	0	1	0	1	1	1	1	0	x	x	x	x	x	P	P	1	t	t	t	t	t	1	i	i	i	i	1	v	v	if (!Pv.new) memd(Rx++#s4:3)=Rtt	
ICLASS			Amode			Type			UN						Parse		t5																
1	0	1	0	1	1	1	1	1	1	0	-	-	-	i	i	P	P	0	t	t	t	t	t	1	i	i	i	i	0	v	v	if (Pv) memd(#u6)=Rtt	
1	0	1	0	1	1	1	1	1	1	0	-	-	-	i	i	P	P	0	t	t	t	t	t	1	i	i	i	i	1	v	v	if (!Pv) memd(#u6)=Rtt	
1	0	1	0	1	1	1	1	1	1	0	-	-	-	i	i	P	P	1	t	t	t	t	t	1	i	i	i	i	0	v	v	if (Pv.new) memd(#u6)=Rtt	
1	0	1	0	1	1	1	1	1	1	0	-	-	-	i	i	P	P	1	t	t	t	t	t	1	i	i	i	i	1	v	v	if (!Pv.new) memd(#u6)=Rtt	

<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
Type	Type
PredNew	PredNew
Sense	Sense
Parse	Packet/loop parse bits
s5	Field to encode register s
t5	Field to encode register t
u5	Field to encode register u
v2	Field to encode register v
x5	Field to encode register x
Amode	Amode
UN	Unsigned

## Store byte

Store the least-significant byte in a source register at the effective address.

Syntax	Behavior
<code>memb (Re=#U6) =Rt</code>	<code>apply_extension (#U);</code> <code>EA=#U;</code> <code>*EA = Rt.b[0];</code> <code>Re=#U;</code>
<code>memb (Rs+#s11:0) =Rt</code>	<code>apply_extension (#s);</code> <code>EA=Rs+#s;</code> <code>*EA = Rt.b[0];</code>
<code>memb (Rs+#u6:0) =#S8</code>	<code>EA=Rs+#u;</code> <code>apply_extension (#S);</code> <code>*EA = #S;</code>
<code>memb (Rs+Ru&lt;&lt;#u2) =Rt</code>	<code>EA=Rs+ (Ru&lt;&lt;#u);</code> <code>*EA = Rt.b[0];</code>
<code>memb (Ru&lt;&lt;#u2+#U6) =Rt</code>	<code>apply_extension (#U);</code> <code>EA=#U+ (Ru&lt;&lt;#u);</code> <code>*EA = Rt.b[0];</code>
<code>memb (Rx++#s4:0) =Rt</code>	<code>EA=Rx;</code> <code>Rx=Rx+#s;</code> <code>*EA = Rt.b[0];</code>
<code>memb (Rx++#s4:0:circ (Mu) ) =Rt</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, #s, MuV);</code> <code>*EA = Rt.b[0];</code>
<code>memb (Rx++I:circ (Mu) ) =Rt</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, I&lt;&lt;0, MuV);</code> <code>*EA = Rt.b[0];</code>
<code>memb (Rx++Mu) =Rt</code>	<code>EA=Rx;</code> <code>Rx=Rx+MuV;</code> <code>*EA = Rt.b[0];</code>
<code>memb (Rx++Mu:brev) =Rt</code>	<code>EA=Rx.h[1]   brev (Rx.h[0]);</code> <code>Rx=Rx+MuV;</code> <code>*EA = Rt.b[0];</code>
<code>memb (gp+#u16:0) =Rt</code>	<code>apply_extension (#u);</code> <code>EA=(Constant_extended ? (0) : GP)+#u;</code> <code>*EA = Rt.b[0];</code>

### Class: ST (slots 0,1)

#### Intrinsics

<code>memb (Rx++#s4:0:circ (Mu) ) = Rt</code>	<code>void Q6_memb_IMR_circ (void** StartAddress, Word32 Is4_0, Word32 Mu, Word32 Rt, void* BaseAddress)</code>
<code>memb (Rx++I:circ (Mu) ) =Rt</code>	<code>void Q6_memb_MR_circ (void** StartAddress, Word32 Mu, Word32 Rt, void* BaseAddress)</code>

## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0			
ICLASS											s5					Parse		u5					t5											
0	0	1	1	1	0	1	1	0	0	0	s	s	s	s	s	P	P	i	u	u	u	u	u	i	-	-	t	t	t	t	t	memb(Rs+Ru<<#u2)=Rt		
ICLASS											s5					Parse																		
0	0	1	1	1	1	0	-	-	0	0	s	s	s	s	s	P	P	l	i	i	i	i	i	i	l	l	l	l	l	l	l	memb(Rs+#u6:0)=#S8		
ICLASS			Type													Parse		t5																
0	1	0	0	1	i	i	0	0	0	0	i	i	i	i	i	P	P	i	t	t	t	t	t	t	i	i	i	i	i	i	i	i	memb(gp+#u16:0)=Rt	
ICLASS			Amode		Type		UN	s5					Parse		t5																			
1	0	1	0	0	i	i	1	0	0	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	i	i	i	memb(Rs+#s11:0)=Rt	
ICLASS			Amode		Type		UN	x5					Parse		u1	t5																		
1	0	1	0	1	0	0	1	0	0	0	x	x	x	x	x	P	P	u	t	t	t	t	t	t	0	-	-	-	-	-	1	-	memb(Rx++l:circ(Mu))=Rt	
1	0	1	0	1	0	0	1	0	0	0	x	x	x	x	x	P	P	u	t	t	t	t	t	t	0	i	i	i	i	i	-	0	-	memb(Rx++#s4:0:circ(Mu))=Rt
ICLASS			Amode		Type		UN	e5					Parse		t5																			
1	0	1	0	1	0	1	1	0	0	0	e	e	e	e	e	P	P	0	t	t	t	t	t	t	1	-	l	l	l	l	l	l	l	memb(Re=#U6)=Rt
ICLASS			Amode		Type		UN	x5					Parse		t5																			
1	0	1	0	1	0	1	1	0	0	0	x	x	x	x	x	P	P	0	t	t	t	t	t	t	0	i	i	i	i	i	-	0	-	memb(Rx++#s4:0)=Rt
ICLASS			Amode		Type		UN	u5					Parse		t5																			
1	0	1	0	1	1	0	1	0	0	0	u	u	u	u	u	P	P	i	t	t	t	t	t	t	1	i	l	l	l	l	l	l	l	memb(Ru<<#u2+#U6)=Rt
ICLASS			Amode		Type		UN	x5					Parse		u1	t5																		
1	0	1	0	1	1	0	1	0	0	0	x	x	x	x	x	P	P	u	t	t	t	t	t	t	0	-	-	-	-	-	-	-	-	memb(Rx++Mu)=Rt
1	0	1	0	1	1	1	1	0	0	0	x	x	x	x	x	P	P	u	t	t	t	t	t	t	0	-	-	-	-	-	-	-	-	memb(Rx++Mu:brev)=Rt

Field name	Description
ICLASS	Instruction class
Type	Type
Parse	Packet/loop parse bits
e5	Field to encode register e
s5	Field to encode register s
t5	Field to encode register t
u1	Field to encode register u
u5	Field to encode register u
x5	Field to encode register x
Amode	Amode
UN	Unsigned

## Store byte conditionally

Store the least-significant byte in a source register at the effective address.

This instruction is conditional based on a predicate value. If the predicate is true, the instruction is performed, otherwise it is treated as a NOP.

Syntax	Behavior
<code>if ([!]Pv[.new]) memb(#u6)=Rt</code>	<pre> apply_extension(#u); EA=#u; if ([!]Pv[.new][0]) {     *EA = Rt.b[0]; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memb(Rs+#u6:0)=#S6</code>	<pre> EA=Rs+#u; if ([!]Pv[.new][0]){     apply_extension(#S);     *EA = #S; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memb(Rs+#u6:0)=Rt</code>	<pre> apply_extension(#u); EA=Rs+#u; if ([!]Pv[.new][0]) {     *EA = Rt.b[0]; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memb(Rs+Ru&lt;&lt;#u2)=Rt</code>	<pre> EA=Rs+(Ru&lt;&lt;#u); if ([!]Pv[.new][0]) {     *EA = Rt.b[0]; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memb(Rx++#s4:0)=Rt</code>	<pre> EA=Rx; if ([!]Pv[.new][0]){     Rx=Rx+#s;     *EA = Rt.b[0]; } else {     NOP; } </pre>

**Class: ST (slots 0,1)**

**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS											s5					Parse		u5					t5										
0	0	1	1	0	1	0	0	0	0	0	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	t	t	t	t	t	if (Pv) memb(Rs+Ru<<#u2)=Rt	
0	0	1	1	0	1	0	1	0	0	0	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	t	t	t	t	t	if (!Pv) memb(Rs+Ru<<#u2)=Rt	
0	0	1	1	0	1	1	0	0	0	0	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	t	t	t	t	t	if (Pv.new) memb(Rs+Ru<<#u2)=Rt	
0	0	1	1	0	1	1	1	0	0	0	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	t	t	t	t	t	if (!Pv.new) memb(Rs+Ru<<#u2)=Rt	
ICLASS											s5					Parse																	
0	0	1	1	1	0	0	0	0	0	0	s	s	s	s	s	P	P	i	i	i	i	i	i	i	v	v	i	i	i	i	i	if (Pv) memb(Rs+#u6:0)=#S6	
0	0	1	1	1	0	0	0	1	0	0	s	s	s	s	s	P	P	i	i	i	i	i	i	i	v	v	i	i	i	i	i	if (!Pv) memb(Rs+#u6:0)=#S6	
0	0	1	1	1	0	0	1	0	0	0	s	s	s	s	s	P	P	i	i	i	i	i	i	i	v	v	i	i	i	i	i	if (Pv.new) memb(Rs+#u6:0)=#S6	
0	0	1	1	1	0	0	1	1	0	0	s	s	s	s	s	P	P	i	i	i	i	i	i	i	v	v	i	i	i	i	i	if (!Pv.new) memb(Rs+#u6:0)=#S6	
ICLASS			Sense	PredNew	Type			s5					Parse		t5																		
0	1	0	0	0	0	0	0	0	0	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	0	v	v	if (Pv) memb(Rs+#u6:0)=Rt
0	1	0	0	0	0	1	0	0	0	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	0	v	v	if (Pv.new) memb(Rs+#u6:0)=Rt
0	1	0	0	0	1	0	0	0	0	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	0	v	v	if (!Pv) memb(Rs+#u6:0)=Rt
0	1	0	0	0	1	1	0	0	0	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	0	v	v	if (!Pv.new) memb(Rs+#u6:0)=Rt
ICLASS			Amode		Type		UN	x5					Parse		t5																		
1	0	1	0	1	0	1	1	0	0	0	x	x	x	x	x	P	P	1	t	t	t	t	t	0	i	i	i	i	0	v	v	if (Pv) memb(Rx++#s4:0)=Rt	
1	0	1	0	1	0	1	1	0	0	0	x	x	x	x	x	P	P	1	t	t	t	t	t	0	i	i	i	i	1	v	v	if (!Pv) memb(Rx++#s4:0)=Rt	
1	0	1	0	1	0	1	1	0	0	0	x	x	x	x	x	P	P	1	t	t	t	t	t	1	i	i	i	i	0	v	v	if (Pv.new) memb(Rx++#s4:0)=Rt	
1	0	1	0	1	0	1	1	0	0	0	x	x	x	x	x	P	P	1	t	t	t	t	t	1	i	i	i	i	1	v	v	if (!Pv.new) memb(Rx++#s4:0)=Rt	
ICLASS			Amode		Type		UN						Parse		t5																		
1	0	1	0	1	1	1	1	0	0	0	-	-	-	i	i	P	P	0	t	t	t	t	t	1	i	i	i	i	0	v	v	if (Pv) memb(#u6)=Rt	
1	0	1	0	1	1	1	1	0	0	0	-	-	-	i	i	P	P	0	t	t	t	t	t	1	i	i	i	i	1	v	v	if (!Pv) memb(#u6)=Rt	
1	0	1	0	1	1	1	1	0	0	0	-	-	-	i	i	P	P	1	t	t	t	t	t	1	i	i	i	i	0	v	v	if (Pv.new) memb(#u6)=Rt	
1	0	1	0	1	1	1	1	0	0	0	-	-	-	i	i	P	P	1	t	t	t	t	t	1	i	i	i	i	1	v	v	if (!Pv.new) memb(#u6)=Rt	

<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
Type	Type
PredNew	PredNew
Sense	Sense
Parse	Packet/loop parse bits
s5	Field to encode register s



<b>Field name</b>	<b>Description</b>
t5	Field to encode register t
u5	Field to encode register u
v2	Field to encode register v
x5	Field to encode register x
Amode	Amode
UN	Unsigned

## Store halfword

Store the upper or lower 16-bits of a source register at the effective address.

Syntax	Behavior
<code>memh (Re=#U6) =Rt.H</code>	<code>apply_extension (#U);</code> <code>EA=#U;</code> <code>*EA = Rt.h[1];</code> <code>Re=#U;</code>
<code>memh (Re=#U6) =Rt</code>	<code>apply_extension (#U);</code> <code>EA=#U;</code> <code>*EA = Rt.h[0];</code> <code>Re=#U;</code>
<code>memh (Rs+#s11:1) =Rt.H</code>	<code>apply_extension (#s);</code> <code>EA=Rs+#s;</code> <code>*EA = Rt.h[1];</code>
<code>memh (Rs+#s11:1) =Rt</code>	<code>apply_extension (#s);</code> <code>EA=Rs+#s;</code> <code>*EA = Rt.h[0];</code>
<code>memh (Rs+#u6:1) =#S8</code>	<code>EA=Rs+#u;</code> <code>apply_extension (#S);</code> <code>*EA = #S;</code>
<code>memh (Rs+Ru&lt;&lt;#u2) =Rt.H</code>	<code>EA=Rs+ (Ru&lt;&lt;#u);</code> <code>*EA = Rt.h[1];</code>
<code>memh (Rs+Ru&lt;&lt;#u2) =Rt</code>	<code>EA=Rs+ (Ru&lt;&lt;#u);</code> <code>*EA = Rt.h[0];</code>
<code>memh (Ru&lt;&lt;#u2+#U6) =Rt.H</code>	<code>apply_extension (#U);</code> <code>EA=#U+ (Ru&lt;&lt;#u);</code> <code>*EA = Rt.h[1];</code>
<code>memh (Ru&lt;&lt;#u2+#U6) =Rt</code>	<code>apply_extension (#U);</code> <code>EA=#U+ (Ru&lt;&lt;#u);</code> <code>*EA = Rt.h[0];</code>
<code>memh (Rx++#s4:1) =Rt.H</code>	<code>EA=Rx;</code> <code>Rx=Rx+#s;</code> <code>*EA = Rt.h[1];</code>
<code>memh (Rx++#s4:1) =Rt</code>	<code>EA=Rx;</code> <code>Rx=Rx+#s;</code> <code>*EA = Rt.h[0];</code>
<code>memh (Rx++#s4:1:circ (Mu)) =Rt.H</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, #s, MuV);</code> <code>*EA = Rt.h[1];</code>
<code>memh (Rx++#s4:1:circ (Mu)) =Rt</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, #s, MuV);</code> <code>*EA = Rt.h[0];</code>
<code>memh (Rx++I:circ (Mu)) =Rt.H</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, I&lt;&lt;1, MuV);</code> <code>*EA = Rt.h[1];</code>
<code>memh (Rx++I:circ (Mu)) =Rt</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, I&lt;&lt;1, MuV);</code> <code>*EA = Rt.h[0];</code>

Syntax	Behavior
memh (Rx++Mu) =Rt .H	EA=Rx; Rx=Rx+MuV; *EA = Rt.h[1];
memh (Rx++Mu) =Rt	EA=Rx; Rx=Rx+MuV; *EA = Rt.h[0];
memh (Rx++Mu:brev) =Rt .H	EA=Rx.h[1]   brev (Rx.h[0]); Rx=Rx+MuV; *EA = Rt.h[1];
memh (Rx++Mu:brev) =Rt	EA=Rx.h[1]   brev (Rx.h[0]); Rx=Rx+MuV; *EA = Rt.h[0];
memh (gp+#u16:1) =Rt .H	apply_extension (#u); EA= (Constant_extended ? (0) : GP)+#u; *EA = Rt.h[1];
memh (gp+#u16:1) =Rt	apply_extension (#u); EA= (Constant_extended ? (0) : GP)+#u; *EA = Rt.h[0];

### Class: ST (slots 0,1)

#### Intrinsics

memh (Rx++#s4:1:circ (Mu) ) =Rt .H	void Q6_memh_IMRh_circ (void** StartAddress, Word32 Is4_1, Word32 Mu, Word32 Rt, void* BaseAddress)
memh (Rx++#s4:1:circ (Mu) ) =Rt	void Q6_memh_IMR_circ (void** StartAddress, Word32 Is4_1, Word32 Mu, Word32 Rt, void* BaseAddress)
memh (Rx++I:circ (Mu) ) =Rt .H	void Q6_memh_MRh_circ (void** StartAddress, Word32 Mu, Word32 Rt, void* BaseAddress)
memh (Rx++I:circ (Mu) ) =Rt	void Q6_memh_MR_circ (void** StartAddress, Word32 Mu, Word32 Rt, void* BaseAddress)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS											s5					Parse		u5					t5										
0	0	1	1	1	0	1	1	0	1	0	s	s	s	s	s	P	P	i	u	u	u	u	u	i	-	-	t	t	t	t	t	memh (Rs+Ru<<#u2)=Rt	
0	0	1	1	1	0	1	1	0	1	1	s	s	s	s	s	P	P	i	u	u	u	u	u	i	-	-	t	t	t	t	t	memh (Rs+Ru<<#u2)=Rt.H	
ICLASS											s5					Parse																	
0	0	1	1	1	1	0	-	-	0	1	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	i	i	i	i	i	memh (Rs+#u6:1)=#S8	
ICLASS											Type		Parse		t5																		
0	1	0	0	1	i	i	0	0	1	0	i	i	i	i	i	P	P	i	t	t	t	t	t	t	i	i	i	i	i	i	i	i	memh (gp+#u16:1)=Rt
0	1	0	0	1	i	i	0	0	1	1	i	i	i	i	i	P	P	i	t	t	t	t	t	t	i	i	i	i	i	i	i	i	memh (gp+#u16:1)=Rt.H
ICLASS			Amode		Type		UN	s5					Parse		t5																		
1	0	1	0	0	i	i	1	0	1	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	i	i	i	memh (Rs+#s11:1)=Rt
1	0	1	0	0	i	i	1	0	1	1	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	i	i	i	memh (Rs+#s11:1)=Rt.H

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS		Amode				Type				UN	x5					Parse		u1	t5														
1	0	1	0	1	0	0	1	0	1	0	x	x	x	x	x	P	P	u	t	t	t	t	t	0	-	-	-	-	-	1	-	memh(Rx++!:circ(Mu))=Rt	
1	0	1	0	1	0	0	1	0	1	0	x	x	x	x	x	P	P	u	t	t	t	t	t	0	i	i	i	i	-	0	-	memh(Rx++#s4:1:circ(Mu))=Rt	
1	0	1	0	1	0	0	1	0	1	1	x	x	x	x	x	P	P	u	t	t	t	t	t	0	-	-	-	-	-	1	-	memh(Rx++!:circ(Mu))=Rt.H	
1	0	1	0	1	0	0	1	0	1	1	x	x	x	x	x	P	P	u	t	t	t	t	t	0	i	i	i	i	-	0	-	memh(Rx++#s4:1:circ(Mu))=Rt.H	
ICLASS		Amode				Type				UN	e5					Parse			t5														
1	0	1	0	1	0	1	1	0	1	0	e	e	e	e	e	P	P	0	t	t	t	t	t	1	-	l	l	l	l	l	l	l	memh(Re=#U6)=Rt
ICLASS		Amode				Type				UN	x5					Parse			t5														
1	0	1	0	1	0	1	1	0	1	0	x	x	x	x	x	P	P	0	t	t	t	t	t	0	i	i	i	i	-	0	-	memh(Rx++#s4:1)=Rt	
ICLASS		Amode				Type				UN	e5					Parse			t5														
1	0	1	0	1	0	1	1	0	1	1	e	e	e	e	e	P	P	0	t	t	t	t	t	1	-	l	l	l	l	l	l	l	memh(Re=#U6)=Rt.H
ICLASS		Amode				Type				UN	x5					Parse			t5														
1	0	1	0	1	0	1	1	0	1	1	x	x	x	x	x	P	P	0	t	t	t	t	t	0	i	i	i	i	-	0	-	memh(Rx++#s4:1)=Rt.H	
ICLASS		Amode				Type				UN	u5					Parse			t5														
1	0	1	0	1	1	0	1	0	1	0	u	u	u	u	u	P	P	i	t	t	t	t	t	1	i	l	l	l	l	l	l	l	memh(Ru<<#u2+#U6)=Rt
ICLASS		Amode				Type				UN	x5					Parse		u1	t5														
1	0	1	0	1	1	0	1	0	1	0	x	x	x	x	x	P	P	u	t	t	t	t	t	0	-	-	-	-	-	-	-	-	memh(Rx++Mu)=Rt
ICLASS		Amode				Type				UN	u5					Parse			t5														
1	0	1	0	1	1	0	1	0	1	1	u	u	u	u	u	P	P	i	t	t	t	t	t	1	i	l	l	l	l	l	l	l	memh(Ru<<#u2+#U6)=Rt.H
ICLASS		Amode				Type				UN	x5					Parse		u1	t5														
1	0	1	0	1	1	0	1	0	1	1	x	x	x	x	x	P	P	u	t	t	t	t	t	0	-	-	-	-	-	-	-	-	memh(Rx++Mu)=Rt.H
1	0	1	0	1	1	1	1	0	1	0	x	x	x	x	x	P	P	u	t	t	t	t	t	0	-	-	-	-	-	-	-	-	memh(Rx++Mu:brev)=Rt
1	0	1	0	1	1	1	1	0	1	1	x	x	x	x	x	P	P	u	t	t	t	t	t	0	-	-	-	-	-	-	-	-	memh(Rx++Mu:brev)=Rt.H

Field name	Description
ICLASS	Instruction class
Type	Type
Parse	Packet/loop parse bits
e5	Field to encode register e
s5	Field to encode register s
t5	Field to encode register t
u1	Field to encode register u
u5	Field to encode register u
x5	Field to encode register x
Amode	Amode
UN	Unsigned

## Store halfword conditionally

Store the upper or lower 16-bits of a source register in memory at the effective address.

This instruction is conditional based on a predicate value. If the predicate is true, the instruction is performed, otherwise it is treated as a NOP.

Syntax	Behavior
<code>if ([!]Pv[.new]) memh(#u6)=Rt.H</code>	<pre> apply_extension(#u); EA=#u; if ([!]Pv[.new][0]) {     *EA = Rt.h[1]; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memh(#u6)=Rt</code>	<pre> apply_extension(#u); EA=#u; if ([!]Pv[.new][0]) {     *EA = Rt.h[0]; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memh(Rs+#u6:1)=#S6</code>	<pre> EA=Rs+#u; if ([!]Pv[.new][0]){     apply_extension(#S);     *EA = #S; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memh(Rs+#u6:1)=Rt.H</code>	<pre> apply_extension(#u); EA=Rs+#u; if ([!]Pv[.new][0]) {     *EA = Rt.h[1]; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memh(Rs+#u6:1)=Rt</code>	<pre> apply_extension(#u); EA=Rs+#u; if ([!]Pv[.new][0]) {     *EA = Rt.h[0]; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memh(Rs+Ru&lt;&lt;#u2)=Rt.H</code>	<pre> EA=Rs+(Ru&lt;&lt;#u); if ([!]Pv[.new][0]) {     *EA = Rt.h[1]; } else {     NOP; } </pre>

Syntax	Behavior
<code>if ([!]Pv[.new]) memh(Rs+Ru&lt;&lt;#u2)=Rt</code>	EA=Rs+(Ru<<#u); if ([!]Pv[.new][0]) { *EA = Rt.h[0]; } else { NOP; }
<code>if ([!]Pv[.new]) memh(Rx++#s4:1)=Rt.H</code>	EA=Rx; if ([!]Pv[.new][0]){ Rx=Rx+#s; *EA = Rt.h[1]; } else { NOP; }
<code>if ([!]Pv[.new]) memh(Rx++#s4:1)=Rt</code>	EA=Rx; if ([!]Pv[.new][0]){ Rx=Rx+#s; *EA = Rt.h[0]; } else { NOP; }

**Class: ST (slots 0,1)**

**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS											s5					Parse		u5					t5										
0	0	1	1	0	1	0	0	0	1	0	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	t	t	t	t	t	if (Pv) memh(Rs+Ru<<#u2)=Rt	
0	0	1	1	0	1	0	0	0	1	1	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	t	t	t	t	t	if (Pv) memh(Rs+Ru<<#u2)=Rt.H	
0	0	1	1	0	1	0	1	0	1	0	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	t	t	t	t	t	if (!Pv) memh(Rs+Ru<<#u2)=Rt	
0	0	1	1	0	1	0	1	0	1	1	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	t	t	t	t	t	if (!Pv) memh(Rs+Ru<<#u2)=Rt.H	
0	0	1	1	0	1	1	0	0	1	0	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	t	t	t	t	t	if (Pv.new) memh(Rs+Ru<<#u2)=Rt	
0	0	1	1	0	1	1	0	0	1	1	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	t	t	t	t	t	if (Pv.new) memh(Rs+Ru<<#u2)=Rt.H	
0	0	1	1	0	1	1	1	0	1	0	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	t	t	t	t	t	if (!Pv.new) memh(Rs+Ru<<#u2)=Rt	
0	0	1	1	0	1	1	1	0	1	1	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	t	t	t	t	t	if (!Pv.new) memh(Rs+Ru<<#u2)=Rt.H	
ICLASS											s5					Parse																	
0	0	1	1	1	0	0	0	0	0	1	s	s	s	s	s	P	P	i	i	i	i	i	i	i	v	v	i	i	i	i	if (Pv) memh(Rs+#u6:1)=#S6		
0	0	1	1	1	0	0	0	1	0	1	s	s	s	s	s	P	P	i	i	i	i	i	i	i	v	v	i	i	i	i	if (!Pv) memh(Rs+#u6:1)=#S6		
0	0	1	1	1	0	0	1	0	0	1	s	s	s	s	s	P	P	i	i	i	i	i	i	i	v	v	i	i	i	i	if (Pv.new) memh(Rs+#u6:1)=#S6		
0	0	1	1	1	0	0	1	1	0	1	s	s	s	s	s	P	P	i	i	i	i	i	i	i	v	v	i	i	i	i	if (!Pv.new) memh(Rs+#u6:1)=#S6		
ICLASS		Sense		Pred New		Type		s5					Parse		t5																		
0	1	0	0	0	0	0	0	0	1	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	0	v	v	if (Pv) memh(Rs+#u6:1)=Rt

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
0	1	0	0	0	0	0	0	0	1	1	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	0	v	v	if (Pv) memh(Rs+#u6:1)=Rt.H
0	1	0	0	0	0	1	0	0	1	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	0	v	v	if (Pv.new) memh(Rs+#u6:1)=Rt
0	1	0	0	0	0	1	0	0	1	1	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	0	v	v	if (Pv.new) memh(Rs+#u6:1)=Rt.H
0	1	0	0	0	1	0	0	0	1	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	0	v	v	if (!Pv) memh(Rs+#u6:1)=Rt
0	1	0	0	0	1	0	0	0	1	1	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	0	v	v	if (!Pv) memh(Rs+#u6:1)=Rt.H
0	1	0	0	0	1	1	0	0	1	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	0	v	v	if (!Pv.new) memh(Rs+#u6:1)=Rt
0	1	0	0	0	1	1	0	0	1	1	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	0	v	v	if (!Pv.new) memh(Rs+#u6:1)=Rt.H
ICLASS		Amode		Type		U	x5					Parse		t5																			
1	0	1	0	1	0	1	1	0	1	0	x	x	x	x	x	P	P	1	t	t	t	t	t	0	i	i	i	i	0	v	v	if (Pv) memh(Rx++#s4:1)=Rt	
1	0	1	0	1	0	1	1	0	1	0	x	x	x	x	x	P	P	1	t	t	t	t	t	0	i	i	i	i	1	v	v	if (!Pv) memh(Rx++#s4:1)=Rt	
1	0	1	0	1	0	1	1	0	1	0	x	x	x	x	x	P	P	1	t	t	t	t	t	1	i	i	i	i	0	v	v	if (Pv.new) memh(Rx++#s4:1)=Rt	
1	0	1	0	1	0	1	1	0	1	0	x	x	x	x	x	P	P	1	t	t	t	t	t	1	i	i	i	i	1	v	v	if (!Pv.new) memh(Rx++#s4:1)=Rt	
1	0	1	0	1	0	1	1	0	1	1	x	x	x	x	x	P	P	1	t	t	t	t	t	0	i	i	i	i	0	v	v	if (Pv) memh(Rx++#s4:1)=Rt.H	
1	0	1	0	1	0	1	1	0	1	1	x	x	x	x	x	P	P	1	t	t	t	t	t	0	i	i	i	i	1	v	v	if (!Pv) memh(Rx++#s4:1)=Rt.H	
1	0	1	0	1	0	1	1	0	1	1	x	x	x	x	x	P	P	1	t	t	t	t	t	1	i	i	i	i	0	v	v	if (Pv.new) memh(Rx++#s4:1)=Rt.H	
1	0	1	0	1	0	1	1	0	1	1	x	x	x	x	x	P	P	1	t	t	t	t	t	1	i	i	i	i	1	v	v	if (!Pv.new) memh(Rx++#s4:1)=Rt.H	
ICLASS		Amode		Type		U						Parse		t5																			
1	0	1	0	1	1	1	1	0	1	0	-	-	-	i	i	P	P	0	t	t	t	t	t	1	i	i	i	i	0	v	v	if (Pv) memh(#u6)=Rt	
1	0	1	0	1	1	1	1	0	1	0	-	-	-	i	i	P	P	0	t	t	t	t	t	1	i	i	i	i	1	v	v	if (!Pv) memh(#u6)=Rt	
1	0	1	0	1	1	1	1	0	1	0	-	-	-	i	i	P	P	1	t	t	t	t	t	1	i	i	i	i	0	v	v	if (Pv.new) memh(#u6)=Rt	
1	0	1	0	1	1	1	1	0	1	0	-	-	-	i	i	P	P	1	t	t	t	t	t	1	i	i	i	i	1	v	v	if (!Pv.new) memh(#u6)=Rt	
1	0	1	0	1	1	1	1	0	1	1	-	-	-	i	i	P	P	0	t	t	t	t	t	1	i	i	i	i	0	v	v	if (Pv) memh(#u6)=Rt.H	
1	0	1	0	1	1	1	1	0	1	1	-	-	-	i	i	P	P	0	t	t	t	t	t	1	i	i	i	i	1	v	v	if (!Pv) memh(#u6)=Rt.H	
1	0	1	0	1	1	1	1	0	1	1	-	-	-	i	i	P	P	1	t	t	t	t	t	1	i	i	i	i	0	v	v	if (Pv.new) memh(#u6)=Rt.H	
1	0	1	0	1	1	1	1	0	1	1	-	-	-	i	i	P	P	1	t	t	t	t	t	1	i	i	i	i	1	v	v	if (!Pv.new) memh(#u6)=Rt.H	

<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
Type	Type
PredNew	PredNew
Sense	Sense
Parse	Packet/loop parse bits
s5	Field to encode register s
t5	Field to encode register t
u5	Field to encode register u
v2	Field to encode register v

<b>Field name</b>	<b>Description</b>
x5	Field to encode register x
Amode	Amode
UN	Unsigned



## Release

The release memory operation is observed after all preceding memory operations have been observed at the local point of serialization. A different order can be observed at the global point of serialization (see Ordering and Synchronization). No data is modified by this instruction.

When the `:st` (same domain) option is specified, the preceding memory operations are those that were committed on any thread with the same consistency domain before this instruction was committed.

When the `:at` (all threads) option is specified, the preceding memory operations are those that were committed on any thread before this instruction was committed.

The store release address is limited to certain memory regions. The following memory regions are excluded:

- AHB memory space
- AXI M2 memory space
- Hexagon memory cut-out is excluded with the exception of addressable TCM and VTCM memory
- Memory with the CCCC types 2, 3, or 4

The `:st` option does not apply to cache operation by index or global cache operation. The `:st` option does not apply a consistency domain to vector operations, but instead uses a per hardware thread ordering scope.

Syntax	Behavior
<code>release(Rs):at</code>	EA=Rs; *EA = Rs
<code>release(Rs):st</code>	EA=Rs; *EA = Rs

### Class: ST (slots 0)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				Amode				Type	U N	s5					Parse		t5					d2										
1	0	1	0	0	0	0	0	1	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	-	-	0	0	1	1	d	d	release(Rs):at
1	0	1	0	0	0	0	0	1	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	-	-	1	0	1	1	d	d	release(Rs):st

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d2	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
Amode	Amode
Type	Type
UN	Unsigned

## Store word

Store a 32-bit register in memory at the effective address.

Syntax	Behavior
<code>memw (Re=#U6) =Rt</code>	<code>apply_extension (#U);</code> <code>EA=#U;</code> <code>*EA = Rt;</code> <code>Re=#U;</code>
<code>memw (Rs+#s11:2) =Rt</code>	<code>apply_extension (#s);</code> <code>EA=Rs+#s;</code> <code>*EA = Rt;</code>
<code>memw (Rs+#u6:2) =#S8</code>	<code>EA=Rs+#u;</code> <code>apply_extension (#S);</code> <code>*EA = #S;</code>
<code>memw (Rs+Ru&lt;&lt;#u2) =Rt</code>	<code>EA=Rs+ (Ru&lt;&lt;#u);</code> <code>*EA = Rt;</code>
<code>memw (Ru&lt;&lt;#u2+#U6) =Rt</code>	<code>apply_extension (#U);</code> <code>EA=#U+ (Ru&lt;&lt;#u);</code> <code>*EA = Rt;</code>
<code>memw (Rx++#s4:2) =Rt</code>	<code>EA=Rx;</code> <code>Rx=Rx+#s;</code> <code>*EA = Rt;</code>
<code>memw (Rx++#s4:2:circ (Mu)) =Rt</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, #s, MuV);</code> <code>*EA = Rt;</code>
<code>memw (Rx++I:circ (Mu)) =Rt</code>	<code>EA=Rx;</code> <code>Rx=Rx=circ_add (Rx, I&lt;&lt;2, MuV);</code> <code>*EA = Rt;</code>
<code>memw (Rx++Mu) =Rt</code>	<code>EA=Rx;</code> <code>Rx=Rx+MuV;</code> <code>*EA = Rt;</code>
<code>memw (Rx++Mu:brev) =Rt</code>	<code>EA=Rx.h[1]   brev (Rx.h[0]);</code> <code>Rx=Rx+MuV;</code> <code>*EA = Rt;</code>
<code>memw (gp+#u16:2) =Rt</code>	<code>apply_extension (#u);</code> <code>EA= (Constant_extended ? (0) : GP)+#u;</code> <code>*EA = Rt;</code>

### Class: ST (slots 0,1)

#### Intrinsics

```
memw (Rx++#s4:2:circ (Mu)) = Rt  void Q6_memw_IMR_circ (void** StartAddress, Word32
                                Is4_2, Word32 Mu, Word32 Rt, void* BaseAddress)
```

```
memw (Rx++I:circ (Mu)) =Rt      void Q6_memw_MR_circ (void** StartAddress, Word32 Mu,
                                Word32 Rt, void* BaseAddress)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS											s5					Parse		u5					t5									
0	0	1	1	1	0	1	1	1	0	0	s	s	s	s	s	P	P	i	u	u	u	u	u	i	-	-	t	t	t	t	t	memw(Rs+Ru<<#u2)=Rt
ICLASS											s5					Parse																
0	0	1	1	1	1	0	-	-	1	0	s	s	s	s	s	P	P	l	i	i	i	i	i	i	l	l	l	l	l	l	l	memw(Rs+#u6:2)=#S8
ICLASS			Type								s5					Parse		t5														
0	1	0	0	1	i	i	0	1	0	0	i	i	i	i	i	P	P	i	t	t	t	t	t	i	i	i	i	i	i	i	i	memw(gp+#u16:2)=Rt
ICLASS			Amode		Type		UN	s5					Parse		t5																	
1	0	1	0	0	i	i	1	1	0	0	s	s	s	s	s	P	P	i	t	t	t	t	t	i	i	i	i	i	i	i	i	memw(Rs+#s11:2)=Rt
ICLASS			Amode		Type		UN	x5					Parse		u1	t5																
1	0	1	0	1	0	0	1	1	0	0	x	x	x	x	x	P	P	u	t	t	t	t	t	0	-	-	-	-	-	1	-	memw(Rx++l:circ(Mu))=Rt
1	0	1	0	1	0	0	1	1	0	0	x	x	x	x	x	P	P	u	t	t	t	t	t	0	i	i	i	i	-	0	-	memw(Rx++#s4:2:circ(Mu))=Rt
ICLASS			Amode		Type		UN	e5					Parse		t5																	
1	0	1	0	1	0	1	1	1	0	0	e	e	e	e	e	P	P	0	t	t	t	t	t	1	-	l	l	l	l	l	l	memw(Re=#U6)=Rt
ICLASS			Amode		Type		UN	x5					Parse		t5																	
1	0	1	0	1	0	1	1	1	0	0	x	x	x	x	x	P	P	0	t	t	t	t	t	0	i	i	i	i	-	0	-	memw(Rx++#s4:2)=Rt
ICLASS			Amode		Type		UN	u5					Parse		t5																	
1	0	1	0	1	1	0	1	1	0	0	u	u	u	u	u	P	P	i	t	t	t	t	t	1	i	l	l	l	l	l	l	memw(Ru<<#u2+#U6)=Rt
ICLASS			Amode		Type		UN	x5					Parse		u1	t5																
1	0	1	0	1	1	0	1	1	0	0	x	x	x	x	x	P	P	u	t	t	t	t	t	0	-	-	-	-	-	-	-	memw(Rx++Mu)=Rt
1	0	1	0	1	1	1	1	1	0	0	x	x	x	x	x	P	P	u	t	t	t	t	t	0	-	-	-	-	-	-	-	memw(Rx++Mu.brev)=Rt

Field name	Description
ICLASS	Instruction class
Type	Type
Parse	Packet/loop parse bits
e5	Field to encode register e
s5	Field to encode register s
t5	Field to encode register t
u1	Field to encode register u
u5	Field to encode register u
x5	Field to encode register x
Amode	Amode
UN	Unsigned

## Store-release word

Store a 32-bit register in memory at the effective address. The store-release memory operation is observed after all preceding memory operations have been observed at the local point of serialization. A different order can be observed at the global point of serialization (see Ordering and Synchronization).

When the `:st` (same domain) option is specified, the preceding memory operations are those that were committed on any thread with the same consistency domain before this instruction was committed.

When the `:at` (all threads) option is specified, the preceding memory operations are those that were committed on any thread before this instruction was committed.

The store release address is limited to certain memory regions. The following are excluded memory regions: AHB memory space, AXI M2 memory space, Hexagon memory cut-out is excluded with the exception of addressable TCM and VTCM memory, and memory with the CCCC types 2, 3, or 4 are excluded. The `:st` option does not apply to cache operation by index or global cache operation. The `:st` option does not apply a consistency domain to vector operations, but instead uses a per hardware thread ordering scope.

Syntax	Behavior
<code>memw_rl(Rs):at=Rt</code>	<code>EA=Rs;</code> <code>*EA = Rt</code>
<code>memw_rl(Rs):st=Rt</code>	<code>EA=Rs;</code> <code>*EA = Rt</code>

### Class: ST (slots 0)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				Amode				Type		UN	s5					Parse		t5					d2									
1	0	1	0	0	0	0	0	1	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	0	0	1	0	d	d	memw_rl(Rs):at=Rt
1	0	1	0	0	0	0	0	1	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	1	0	1	0	d	d	memw_rl(Rs):st=Rt

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d2	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
Amode	Amode
Type	Type
UN	Unsigned

## Store word conditionally

Store a 32-bit register in memory at the effective address.

This instruction is conditional based on a predicate value. If the predicate is true, the instruction is performed, otherwise it is treated as a NOP.

Syntax	Behavior
<code>if ([!]Pv[.new]) memw(#u6)=Rt</code>	<pre> apply_extension(#u); EA=#u; if ([!]Pv[.new][0]) {     *EA = Rt; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memw(Rs+#u6:2)=#S6</code>	<pre> EA=Rs+#u; if ([!]Pv[.new][0]){     apply_extension(#S);     *EA = #S; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memw(Rs+#u6:2)=Rt</code>	<pre> apply_extension(#u); EA=Rs+#u; if ([!]Pv[.new][0]) {     *EA = Rt; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memw(Rs+Ru&lt;&lt;#u2)=Rt</code>	<pre> EA=Rs+(Ru&lt;&lt;#u); if ([!]Pv[.new][0]) {     *EA = Rt; } else {     NOP; } </pre>
<code>if ([!]Pv[.new]) memw(Rx++#s4:2)=Rt</code>	<pre> EA=Rx; if ([!]Pv[.new][0]){     Rx=Rx+#s;     *EA = Rt; } else {     NOP; } </pre>

**Class: ST (slots 0,1)**

**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS											s5					Parse		u5					t5										
0	0	1	1	0	1	0	0	1	0	0	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	t	t	t	t	t	if (Pv) memw(Rs+Ru<<#u2)=Rt	
0	0	1	1	0	1	0	1	1	0	0	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	t	t	t	t	t	if (!Pv) memw(Rs+Ru<<#u2)=Rt	
0	0	1	1	0	1	1	0	1	0	0	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	t	t	t	t	t	if (Pv.new) memw(Rs+Ru<<#u2)=Rt	
0	0	1	1	0	1	1	1	1	0	0	s	s	s	s	s	P	P	i	u	u	u	u	u	i	v	v	t	t	t	t	t	if (!Pv.new) memw(Rs+Ru<<#u2)=Rt	
ICLASS											s5					Parse																	
0	0	1	1	1	0	0	0	0	1	0	s	s	s	s	s	P	P	i	i	i	i	i	i	i	v	v	i	i	i	i	i	if (Pv) memw(Rs+#u6:2)=#S6	
0	0	1	1	1	0	0	0	1	1	0	s	s	s	s	s	P	P	i	i	i	i	i	i	i	v	v	i	i	i	i	i	if (!Pv) memw(Rs+#u6:2)=#S6	
0	0	1	1	1	0	0	1	0	1	0	s	s	s	s	s	P	P	i	i	i	i	i	i	i	v	v	i	i	i	i	i	if (Pv.new) memw(Rs+#u6:2)=#S6	
0	0	1	1	1	0	0	1	1	1	0	s	s	s	s	s	P	P	i	i	i	i	i	i	i	v	v	i	i	i	i	i	if (!Pv.new) memw(Rs+#u6:2)=#S6	
ICLASS			Sense		PredNew		Type			s5					Parse		t5																
0	1	0	0	0	0	0	0	1	0	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	0	v	v	if (Pv) memw(Rs+#u6:2)=Rt
0	1	0	0	0	0	1	0	1	0	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	0	v	v	if (Pv.new) memw(Rs+#u6:2)=Rt
0	1	0	0	0	1	0	0	1	0	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	0	v	v	if (!Pv) memw(Rs+#u6:2)=Rt
0	1	0	0	0	1	1	0	1	0	0	s	s	s	s	s	P	P	i	t	t	t	t	t	t	i	i	i	i	i	0	v	v	if (!Pv.new) memw(Rs+#u6:2)=Rt
ICLASS			Amode			Type			UN	x5					Parse		t5																
1	0	1	0	1	0	1	1	1	0	0	x	x	x	x	x	P	P	1	t	t	t	t	t	0	i	i	i	i	0	v	v	if (Pv) memw(Rx++#s4:2)=Rt	
1	0	1	0	1	0	1	1	1	0	0	x	x	x	x	x	P	P	1	t	t	t	t	t	0	i	i	i	i	1	v	v	if (!Pv) memw(Rx++#s4:2)=Rt	
1	0	1	0	1	0	1	1	1	0	0	x	x	x	x	x	P	P	1	t	t	t	t	t	1	i	i	i	i	0	v	v	if (Pv.new) memw(Rx++#s4:2)=Rt	
1	0	1	0	1	0	1	1	1	0	0	x	x	x	x	x	P	P	1	t	t	t	t	t	1	i	i	i	i	1	v	v	if (!Pv.new) memw(Rx++#s4:2)=Rt	
ICLASS			Amode			Type			UN						Parse		t5																
1	0	1	0	1	1	1	1	1	0	0	-	-	-	i	i	P	P	0	t	t	t	t	t	1	i	i	i	i	0	v	v	if (Pv) memw(#u6)=Rt	
1	0	1	0	1	1	1	1	1	0	0	-	-	-	i	i	P	P	0	t	t	t	t	t	1	i	i	i	i	1	v	v	if (!Pv) memw(#u6)=Rt	
1	0	1	0	1	1	1	1	1	0	0	-	-	-	i	i	P	P	1	t	t	t	t	t	1	i	i	i	i	0	v	v	if (Pv.new) memw(#u6)=Rt	
1	0	1	0	1	1	1	1	1	0	0	-	-	-	i	i	P	P	1	t	t	t	t	t	1	i	i	i	i	1	v	v	if (!Pv.new) memw(#u6)=Rt	

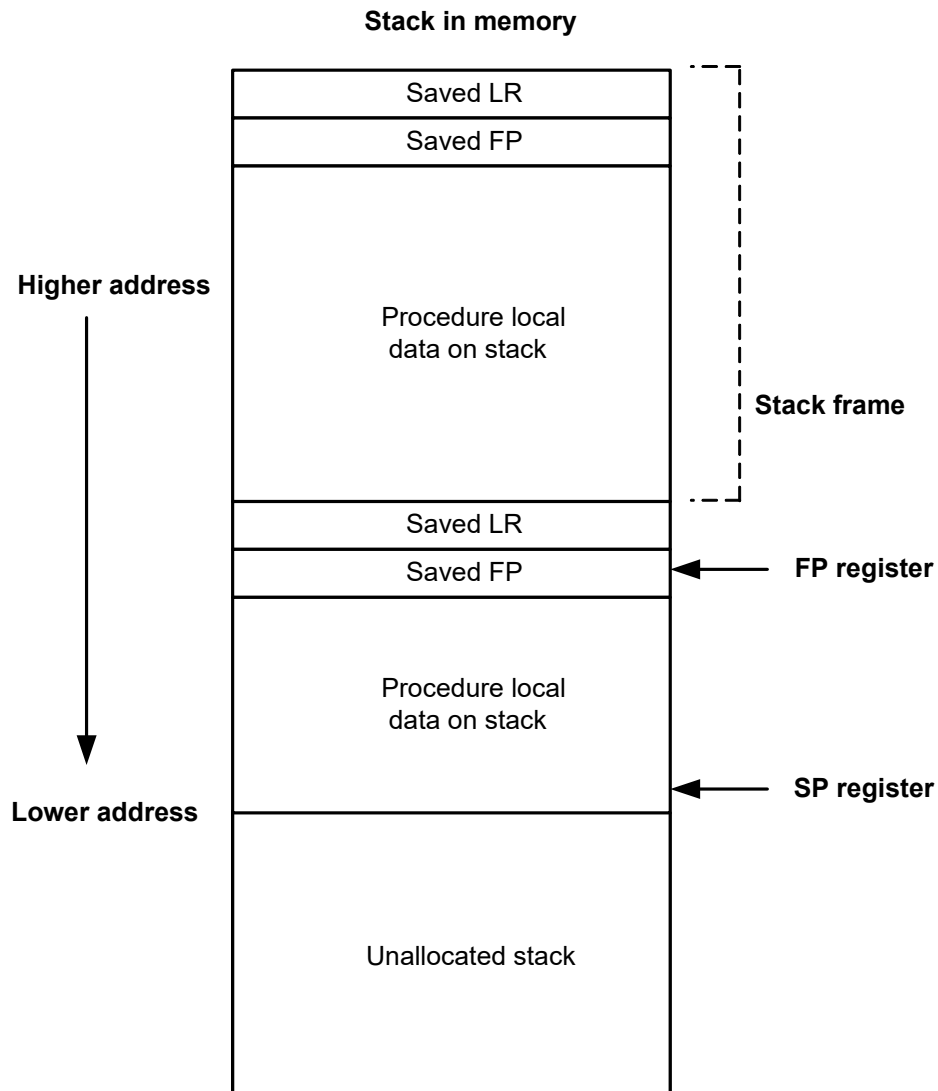
<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
Type	Type
PredNew	PredNew
Sense	Sense
Parse	Packet/loop parse bits
s5	Field to encode register s

<b>Field name</b>	<b>Description</b>
t5	Field to encode register t
u5	Field to encode register u
v2	Field to encode register v
x5	Field to encode register x
Amode	Amode
UN	Unsigned

## Allocate stack frame

Allocate a stack frame on the call stack. This instruction first pushes LR and FP to the top of stack. It then subtracts an unsigned immediate from SP to allocate room for local variables. FP is set to the address of the old frame pointer on the stack.

The following figure shows the stack layout.



### Syntax

```
allocframe (#u11:3)
```

```
allocframe (Rx, #u11:3) :raw
```

### Behavior

```
Assembler mapped to:  
"allocframe (r29, #u11:3) :raw"
```

```
EA=Rx+8;  
*EA = frame_scramble((LR << 32) | FP);  
FP=EA;  
frame_check_limit(EA-#u);  
Rx = EA-#u;
```



**Class: ST (slots 0)****Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS				Amode			Type		U N	x5					Parse																		
1	0	1	0	0	0	0	0	1	0	0	x	x	x	x	x	P	P	0	0	0	i	i	i	i	i	i	i	i	i	i	i	i	allocframe(Rx,#u11:3):raw

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
x5	Field to encode register x
Amode	Amode
Type	Type
UN	Unsigned

## 11.9 SYSTEM

The SYSTEM instruction class includes instructions for managing system resources.

### 11.9.1 SYSTEM USER

The SYSTEM USER instruction subclass includes instructions which allow user access to system resources.

#### Load locked

This memory lock instruction performs a word or double-word locked load.

This instruction returns the contents of the memory at address Rs and also reserves a lock reservation at that address. For more information, see [Atomic operations](#).

Syntax	Behavior
Rd=memw_locked(Rs)	EA=Rs; Rd = *EA;
Rdd=memd_locked(Rs)	EA=Rs; Rdd = *EA;

#### Class: SYSTEM (slots 0)

#### Notes

- This instruction is only grouped with ALU32 or nonfloating-point XTYPE instructions.

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			Amode			Type			U	s5					Parse		d5															
1	0	0	1	0	0	1	0	0	0	0	s	s	s	s	s	P	P	0	0	0	-	-	-	0	0	0	d	d	d	d	d	Rd=memw_locked(Rs)
1	0	0	1	0	0	1	0	0	0	0	s	s	s	s	s	P	P	0	1	0	-	-	-	0	0	0	d	d	d	d	d	Rdd=memd_locked(Rs)

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s

## Store conditional

This memory lock instruction performs a word or double-word conditional store operation.

If the address reservation is held by this thread and there have been no intervening accesses to the memory location, the store is performed and the predicate is set to true. Otherwise, the store is not performed and the predicate returns false. For more information, see [Atomic operations](#).

Syntax	Behavior
<code>memd_locked(Rs, Pd)=Rtt</code>	<pre>EA=Rs; if (lock_valid) {     *EA = Rtt;     Pd = 0xff;     lock_valid = 0; } else {     Pd = 0; }</pre>
<code>memw_locked(Rs, Pd)=Rt</code>	<pre>EA=Rs; if (lock_valid) {     *EA = Rt;     Pd = 0xff;     lock_valid = 0; } else {     Pd = 0; }</pre>

**Class: SYSTEM (slots 0)**

### Notes

- This instruction may only be grouped with ALU32 or non-floating-point XTYPE instructions.
- The predicate generated by this instruction can not be used as a .new predicate, nor can it be automatically ANDed with another predicate.

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		Amode			Type	UN	s5					Parse		t5					d2													
1	0	1	0	0	0	0	0	1	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	-	0	0	d	d	memw_locked(Rs,Pd)=Rt
1	0	1	0	0	0	0	0	1	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	-	-	-	-	0	0	d	d	memd_locked(Rs,Pd)=Rtt

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d2	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
Amode	Amode
Type	Type
UN	Unsigned

## Zero a cache line

The `dczeroa` instruction clears 32 bytes of memory.

If the memory is marked write-back cacheable, a cache line is allocated in the data cache and 32 bytes are cleared.

If the memory is write-through or write-back, 32 bytes of zeros are sent to memory.

This instruction is useful for efficiently handling write-only data by pre-allocating lines in the cache.

The address must be 32-byte aligned. If not, an unaligned error exception is raised.

If this instruction appears in a packet, slot 1 must be A-type or empty.

### Syntax

```
dczeroa(Rs)
```

### Behavior

```
EA=Rs;
dcache_zero_addr(EA);
```

### Class: SYSTEM (slots 0)

### Notes

- A packet containing this instruction must have slot 1 either empty or executing an ALU32 instruction.

### Intrinsics

```
dczeroa(Rs)
```

```
void Q6_dczeroa_A(Address a)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS				Amode				Type	UN	s5					Parse																		
1	0	1	0	0	0	0	0	1	1	0	s	s	s	s	s	P	P	0	-	-	-	-	-	-	-	-	-	-	-	-	-	-	dczeroa(Rs)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s
Amode	Amode
Type	Type
UN	Unsigned

## Memory barrier

The barrier instruction establishes a memory barrier to ensure proper ordering between load/store accesses within a consistency domain before the barrier instruction and accesses after the barrier instruction.

All scalar loads, stores, and cache operation by address within a consistency domain before the barrier are globally observable before any access after the barrier can be observed.

The use of this instruction is system-dependent.

### Syntax

```
barrier
```

### Behavior

```
memory_barrier;
```

### Class: SYSTEM (slots 0)

### Notes

- This is a solo instruction. It must not be grouped with other instructions in a packet.

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS				Amode			Type			UN						Parse																	
1	0	1	0	1	0	0	0	0	0	0	-	-	-	-	-	P	P	-	-	-	-	-	-	0	0	0	-	-	-	-	-	-	barrier

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
Amode	Amode
Type	Type
UN	Unsigned

## Breakpoint

The brkpt instruction causes the program to enter Debug mode if enabled by ISDB.

Execution control is handed to ISDB and the program does not proceed until directed by the debugger.

If ISDB is disabled, this instruction is treated as a NOP.

### Syntax

```
brkpt
```

### Behavior

```
Enter Debug mode;
```

### Class: SYSTEM (slot 3)

### Notes

- This is a solo instruction. It must not be grouped with other instructions in a packet.

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS			sm														Parse																
0	1	1	0	1	1	0	0	0	0	0	1	-	-	-	-	-	P	P	-	-	-	-	-	-	0	0	0	-	-	-	-	-	brkpt

Field name	Description
sm	Supervisor mode only
ICLASS	Instruction class
Parse	Packet/loop parse bits

## Data cache prefetch

The dcfetch instruction prefetches the data at address Rs + unsigned immediate.

This instruction is a hint to the memory system, and is handled in an implementation-dependent manner.

Syntax	Behavior
dcfetch(Rs)	Assembler mapped to: "dcfetch(Rs+#0)"
dcfetch(Rs+#u11:3)	EA=Rs+#u; dcache_fetch(EA);

**Class: SYSTEM (slots 0)**

### Intrinsics

```
dcfetch(Rs) void Q6_dcfetch_A(Address a)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			Amode			Type			UN	s5					Parse																	
1	0	0	1	0	1	0	0	0	0	0	s	s	s	s	s	P	P	0	-	-	i	i	i	i	i	i	i	i	i	i	i	dcfetch(Rs+#u11:3)

Field name	Description
ICLASS	Instruction class
Amode	Amode
Type	Type
UN	Unsigned
Parse	Packet/loop parse bits
s5	Field to encode register s





<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s
Amode	Amode
Type	Type
UN	Unsigned

## Send value to DIAG trace

These instructions send the sources to the external DIAG trace.

Syntax	Behavior
diag (Rs)	
diag0 (Rss, Rtt)	
diag1 (Rss, Rtt)	

### Class: SYSTEM (slot 3)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				sm							s5					Parse																
0	1	1	0	0	0	1	0	0	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	1	-	-	-	-	-	diag(Rs)
ICLASS				sm							s5					Parse		t5														
0	1	1	0	0	0	1	0	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	-	-	-	-	-	diag0(Rss,Rtt)
0	1	1	0	0	0	1	0	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	-	-	-	-	-	diag1(Rss,Rtt)

Field name	Description
sm	Supervisor mode only
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s
t5	Field to encode register t

## Instruction cache maintenance user operations

The `icinva` instruction looks up the address given in `Rs` in the instruction cache. If a translation for `Rs` cannot be found, a TLB-miss-X exception with cause code 0x62 is indicated.

If a translation is found but the user does not have proper permissions to the page to invalidate, the instruction converts to a NOP.

If a translation is found and the user has proper permissions, all ways of all 32 byte segments matching `Rs[11:5]` in any instruction cache are invalidated.

Syntax	Behavior
<code>icinva(Rs)</code>	<code>EA=Rs;</code> <code>icache_inv_addr(EA);</code>

### Class: SYSTEM (slot 2)

#### Notes

- This is a solo instruction. It must not be grouped with other instructions in a packet.

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS			s5											Parse																			
0	1	0	1	0	1	1	0	1	1	0	s	s	s	s	s	P	P	0	0	0	-	-	-	-	-	-	-	-	-	-	-	-	icinva(Rs)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s

## Instruction synchronization

The isync instruction ensures that all previous instructions have committed before continuing to the next instruction.

This instruction should execute after the following events (when subsequent instructions must observe the results of the event):

- After modifying the TLB with a TLBW instruction
- After modifying the SSR register
- After modifying the SYSCFG register
- After any instruction cache maintenance operation
- After modifying the TID register

Syntax	Behavior
<code>isync</code>	<code>instruction_sync;</code>

### Class: SYSTEM (slot 2)

#### Notes

- This is a solo instruction. It must not be grouped with other instructions in a packet.

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS																Parse																
0	1	0	1	0	1	1	1	1	1	0	0	0	0	0	0	P	P	0	-	-	-	0	0	0	0	0	0	0	0	1	0	isync

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits

## L2 cache prefetch

The L2fetch instruction initiates background prefetching into the L2 cache.

R<sub>s</sub> specifies the 32-bit virtual start address. There are two forms of this instruction.

In the first form, the dimensions of the area to prefetch are encoded in source register R<sub>t</sub> as follows:

R<sub>t</sub>[15:8] = Width of a fetch block in bytes.

R<sub>t</sub>[7:0] = Height: the number of Width-sized blocks to fetch.

R<sub>t</sub>[31:16] = Stride: an unsigned byte offset which is used to increment the pointer after each Width-sized block is fetched.

In the second form, the operands are encoded in register pair R<sub>tt</sub> as follows:

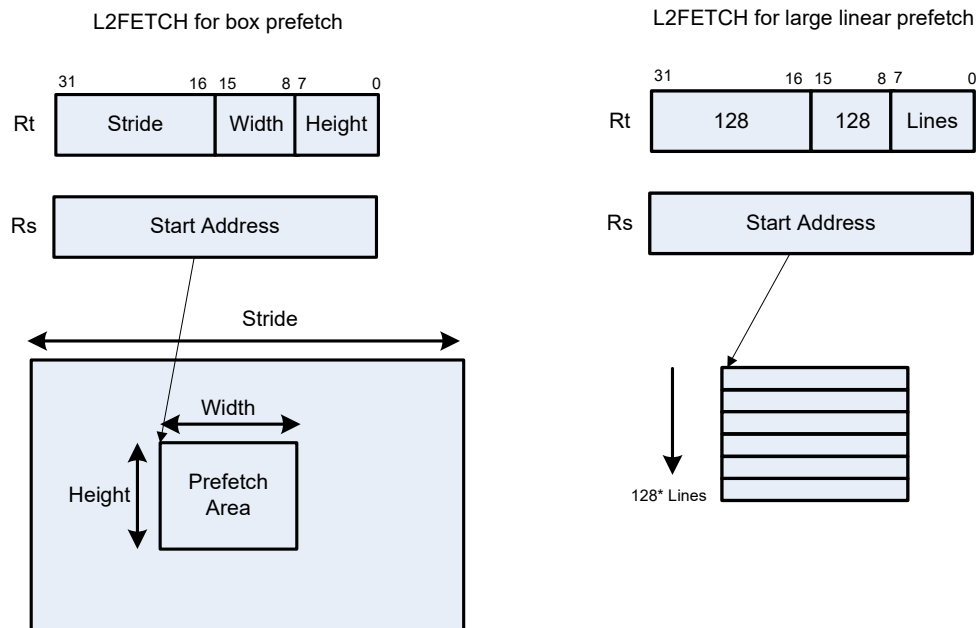
R<sub>tt</sub>[31:16] = Width of a fetch block in bytes.

R<sub>tt</sub>[15:0] = Height: the number of Width-sized blocks to fetch.

R<sub>tt</sub>[47:32] = Stride: an unsigned byte offset that is used to increment the pointer after each Width-sized block is fetched.

R<sub>tt</sub>[48] = Direction. If clear, perform the prefetches in row major form, meaning fetch cache lines in a row before proceeding to the next row. If the bit is set, prefetch in column major form, meaning fetch all cache lines in a column before proceeding to the next column.

The following figure shows two examples of using the L2FETCH instruction.



In the box prefetch, a 2D range of memory is defined within a larger frame. The second example shows prefetch for a large linear area of memory, which has size Lines \* 128.

L2FETCH is nonblocking. After the instruction is initiated, the program continues on to the next instruction while the prefetching is performed in the background. L2fetch can bring in either code or data to the L2 cache. If the lines of interest are already in the L2, no action is performed. If the lines are missing from the L2\$, the hardware attempts to fetch them from the system memory.

The hardware prefetch engine continues to request all lines in the programmed memory range. The prefetching hardware makes a best-effort to prefetch the requested data, and attempts to perform prefetching at a lower priority than demand fetches. This prevents prefetch from adding traffic while the system is under heavy load.

If a program initiates a new L2FETCH while an older L2FETCH operation is still pending, the new request is queued, up to three deep. If three L2FETCHes are already pending, the oldest request is dropped. During the time a L2 prefetch is active for a thread, the USR:PFA status bit is set to indicate that prefetches are in progress. The programmer can use this bit to decide whether to start a new L2FETCH before the previous one completes.

Executing an L2fetch with any subfield programmed as zero cancels all pending prefetches by the calling thread.

The implementation is free to drop prefetches when needed.

Syntax	Behavior
<code>l2fetch(Rs, Rt)</code>	<code>l2fetch(Rs, INFO);</code>
<code>l2fetch(Rs, Rtt)</code>	<code>l2fetch(Rs, INFO);</code>

## Class: SYSTEM (slots 0)

### Notes

- This instruction can only be grouped with ALU32 or non-floating-point XTYPE instructions.

### Intrinsics

<code>l2fetch(Rs, Rt)</code>	<code>void Q6_l2fetch_AR(Address a, Word32 Rt)</code>
<code>l2fetch(Rs, Rtt)</code>	<code>void Q6_l2fetch_AP(Address a, Word64 Rtt)</code>

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			Amode			Type			U	s5					Parse		t5															
1	0	1	0	0	1	1	0	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	-	-	-	-	-	l2fetch(Rs,Rt)
1	0	1	0	0	1	1	0	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	-	-	-	-	-	l2fetch(Rs,Rtt)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s
t5	Field to encode register t
Amode	Amode
Type	Type

<b>Field name</b>	<b>Description</b>
UN	Unsigned

## Pause

The PAUSE instruction pauses execution for a specified period of time.

During the pause duration, the program enters a low-power state and does not fetch and execute instructions. The instruction provides a short immediate that indicates the pause duration. The program will pause for at most the number of cycles specified in the immediate plus 8. The minimum pause is 0 cycles, and the maximum pause is implementation-defined.

An interrupt to the program exits the paused state.

System events, such as hardware or DMA completion, can trigger exits from Pause mode.

An implementation is free to pause for durations shorter than (immediate+8), but not longer.

This instruction is useful for implementing user-level low-power synchronization operations, such as spin locks or wait-for-event signaling.

### Syntax

```
pause(#u10)
```

### Behavior

```
Pause for #u cycles;
```

### Class: SYSTEM (slot 2)

### Notes

- This is a solo instruction. It must not be grouped with other instructions in a packet.

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS																Parse																	
0	1	0	1	0	1	0	0	0	1	-	-	-	-	i	i	P	P	-	i	i	i	i	i	i	-	-	-	i	i	i	-	-	pause(#u10)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits



## Memory thread synchronization

The syncht instruction synchronizes memory.

Outstanding memory operations, including cached and uncached loads and stores, are completed before the processor continues to the next instruction. This ensures that certain memory operations are performed in the desired order (for example, when accessing I/O devices).

After performing a syncht operation, the processor ceases fetching and executing instructions from the program until all outstanding memory operations of that program complete.

In multithreaded or multi-core environments, SYNCHT is not concerned with other execution contexts.

The use of this instruction is system-dependent.

Syntax	Behavior
Rd=dmsyncht	Rd = DM0;
syncht	memory_synch;

### Class: SYSTEM (slots 0)

#### Notes

- This is a solo instruction. It must not be grouped with other instructions in a packet.
- This is a monitor-level feature. If performed in User or Guest mode, a privilege error exception occurs.

#### Intrinsics

Rd=dmsyncht `Word32 Q6_R_dmsyncht()`

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				Amode			Type		UN	Parse								d5														
1	0	1	0	1	0	0	0	0	0	0	-	-	-	-	-	P	P	-	-	-	-	0	1	1	1	d	d	d	d	d	Rd=dmsyncht	
ICLASS				Amode			Type		UN	Parse																						
1	0	1	0	1	0	0	0	0	1	0	-	-	-	-	-	P	P	-	-	-	-	-	-	-	-	-	-	-	-	-	-	syncht

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
Amode	Amode
Type	Type
UN	Unsigned

## Send value to ETM trace

The trace instruction takes the value of register Rs and emits it to the ETM trace.

The ETM block must be enabled, and the thread must have permissions to perform tracing. The contents of Rs are user-defined.

### Syntax

```
trace(Rs)
```

### Behavior

```
Send value to ETM trace;
```

### Class: SYSTEM (slot 3)

### Notes

- This instruction may only be grouped with ALU32 or non-floating-point XTYPE instructions.

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			sm			s5					Parse																					
0	1	1	0	0	0	1	0	0	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	0	-	-	-	-	-	trace(Rs)

Field name	Description
sm	Supervisor mode only
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s

## Trap

The trap instruction causes a precise exception.

Executing a trap instruction sets the EX bit in SSR to 1, which disables interrupts and enables Supervisor mode. The program then jumps to the vector location (either TRAP0 or TRAP1). The instruction specifies a n 8-bit immediate field. This field is copied into the system status register cause field.

Upon returning from the service routine with a RTE, execution resumes at the packet after the TRAP instruction.

These instructions are generally intended for user code to request services from the operating system. Two TRAP instructions are provided so the OS can optimize for fast service routines and slower service routines.

Syntax	Behavior
trap0(#u8)	SSR.CAUSE = #u; TRAP "0";
trap1(#u8)	Assembler mapped to: "trap1(R0, #u8)"
trap1(Rx, #u8)	if (!can_handle_trap1_virtinsn(#u)) { SSR.CAUSE = #u; TRAP "1"; } else if (#u == 1) { VMRTE; } else if (#u == 3) { VMSETIE; } else if (#u == 4) { VMGETIE; } else if (#u == 6) { VMSPSWAP;

### Class: SYSTEM (slot 2)

#### Notes

- This is a solo instruction. It must not be grouped with other instructions in a packet.

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS																Parse																	
0	1	0	1	0	1	0	0	0	0	-	-	-	-	-	-	P	P	-	i	i	i	i	i	-	-	-	i	i	i	-	-	trap0(#u8)	
ICLASS																x5					Parse												
0	1	0	1	0	1	0	0	1	0	-	x	x	x	x	x	P	P	-	i	i	i	i	i	-	-	-	i	i	i	-	-	trap1(Rx,#u8)	

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
x5	Field to encode register x

## Unpause

The unpause instruction resumes threads whose execution has stalled with a pause instruction.

### Syntax

```
unpause
```

### Behavior

```
Unpause threads currently in pause state;
```

### Class: SYSTEM (slot 2)

### Notes

- This is a solo instruction. It must not be grouped with other instructions in a packet.

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS																Parse																	
0	1	0	1	0	1	1	1	1	1	1	1	-	-	-	-	-	P	P	0	1	-	-	-	-	0	0	0	-	-	-	-	-	unpause

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits

## 11.10 XTYPE

The XTYPE instruction class includes instructions that perform most of the data processing done by the Hexagon processor.

XTYPE instructions are executable on slot 2 or slot 3.

### 11.10.1 XTYPE ALU

The XTYPE ALU instruction subclass includes instructions that perform arithmetic and logical operations.

#### Absolute value doubleword

Take the absolute value of the 64-bit source register and place it in the destination register.

#### Syntax

```
Rdd=abs(Rss)
```

#### Behavior

```
Rdd = ABS(Rss);
```

#### Class: XTYPE (slots 2,3)

#### Intrinsics

```
Rdd=abs(Rss)
```

```
Word64 Q6_P_abs_P(Word64 Rss)
```

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp		s5					Parse		MinOp			d5												
1	0	0	0	0	0	0	0	1	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	1	0	d	d	d	d	d	Rdd=abs(Rss)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Absolute value word

Take the absolute value of the source register and place it in the destination register.

The 32-bit absolute value is available with optional saturation. The single case of saturation is when the source register is equal to 0x8000\_0000, the destination saturates to 0x7fff\_ffff.

### Syntax

```
Rd=abs(Rs) [:sat]
```

### Behavior

```
Rd = [sat32] (ABS (sxt32->64 (Rs))) ;
```

### Class: XTYPE (slots 2,3)

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

### Intrinsics

```
Rd=abs(Rs)
```

```
Word32 Q6_R_abs_R(Word32 Rs)
```

```
Rd=abs(Rs) :sat
```

```
Word32 Q6_R_abs_R_sat(Word32 Rs)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp				s5					Parse		MinOp					d5								
1	0	0	0	1	1	0	0	1	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	0	d	d	d	d	d	Rd=abs(Rs)
1	0	0	0	1	1	0	0	1	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	1	d	d	d	d	d	Rd=abs(Rs):sat

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Add and accumulate

Add Rs and Rt or a signed immediate, then add or subtract the resulting value. The result is saved in Rx.

Syntax	Behavior
Rd=add(Rs,add(Ru,#s6))	Rd = Rs + Ru + apply_extension(#s);
Rd=add(Rs,sub(#s6,Ru))	Rd = Rs - Ru + apply_extension(#s);
Rx+=add(Rs,#s8)	apply_extension(#s); Rx=Rx + Rs + #s;
Rx+=add(Rs,Rt)	Rx=Rx + Rs + Rt;
Rx-=add(Rs,#s8)	apply_extension(#s); Rx=Rx - (Rs + #s);
Rx-=add(Rs,Rt)	Rx=Rx - (Rs + Rt);

### Class: XTYPE (slots 2,3)

#### Intrinsics

Rd=add(Rs,add(Ru,#s6))	Word32 Q6_R_add_add_RRI(Word32 Rs, Word32 Ru, Word32 Is6)
Rd=add(Rs,sub(#s6,Ru))	Word32 Q6_R_add_sub_RIR(Word32 Rs, Word32 Is6, Word32 Ru)
Rx+=add(Rs,#s8)	Word32 Q6_R_addacc_RI(Word32 Rx, Word32 Rs, Word32 Is8)
Rx+=add(Rs,Rt)	Word32 Q6_R_addacc_RR(Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=add(Rs,#s8)	Word32 Q6_R_addnac_RI(Word32 Rx, Word32 Rs, Word32 Is8)
Rx-=add(Rs,Rt)	Word32 Q6_R_addnac_RR(Word32 Rx, Word32 Rs, Word32 Rt)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS		RegType				s5					Parse		d5					u5															
1	1	0	1	1	0	1	1	0	i	i	s	s	s	s	s	P	P	i	d	d	d	d	d	i	i	i	u	u	u	u	u	Rd=add(Rs,add(Ru,#s6))	
1	1	0	1	1	0	1	1	1	i	i	s	s	s	s	s	P	P	i	d	d	d	d	d	i	i	i	u	u	u	u	u	Rd=add(Rs,sub(#s6,Ru))	
ICLASS		RegType				MajOp		s5					Parse		MinOp					x5													
1	1	1	0	0	0	1	0	0	-	-	s	s	s	s	s	P	P	0	i	i	i	i	i	i	i	i	x	x	x	x	x	Rx+=add(Rs,#s8)	
1	1	1	0	0	0	1	0	1	-	-	s	s	s	s	s	P	P	0	i	i	i	i	i	i	i	i	x	x	x	x	x	Rx-=add(Rs,#s8)	
ICLASS		RegType				MajOp		s5					Parse		t5					MinOp		x5											
1	1	1	0	1	1	1	1	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	t	0	0	1	x	x	x	x	x	Rx+=add(Rs,Rt)
1	1	1	0	1	1	1	1	1	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	t	0	0	1	x	x	x	x	x	Rx-=add(Rs,Rt)

<b>Field name</b>	<b>Description</b>
RegType	Register type
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
u5	Field to encode register u
x5	Field to encode register x



## Add doublewords

The first form of this instruction adds two 32-bit registers. If the result overflows 32 bits, the result is saturated to 0x7FFF\_FFFF for a positive result, or 0x8000\_0000 for a negative result. A 32-bit nonsaturating register add is a ALU32-class instruction and can execute on any slot.

The second instruction form sign-extends a 32-bit register Rt to 64-bits and performs a 64-bit add with Rss. The result is stored in Rdd.

The third instruction form adds 64-bit registers Rss and Rtt and places the result in Rdd.

The final instruction form adds two 64-bit registers Rss and Rtt. If the result overflows 64 bits, it is saturated to 0x7fff\_ffff\_ffff\_ffff for a positive result, or 0x8000\_0000\_0000\_0000 for a negative result.

Syntax	Behavior
<code>Rd=add(Rs,Rt):sat:deprecated</code>	<code>Rd=sat<sub>32</sub>(Rs+Rt);</code>
<code>Rdd=add(Rs,Rtt)</code>	<pre>if ("Rs &amp; 1") {     Assembler mapped to: "Rdd=add(Rss,Rtt):raw:hi"; } else {     Assembler mapped to: "Rdd=add(Rss,Rtt):raw:lo"; }</pre>
<code>Rdd=add(Rss,Rtt)</code>	<code>Rdd=Rss+Rtt;</code>
<code>Rdd=add(Rss,Rtt):raw:hi</code>	<code>Rdd=Rtt+sxt<sub>32-&gt;64</sub>(Rss.w[1]);</code>
<code>Rdd=add(Rss,Rtt):raw:lo</code>	<code>Rdd=Rtt+sxt<sub>32-&gt;64</sub>(Rss.w[0]);</code>
<code>Rdd=add(Rss,Rtt):sat</code>	<code>Rdd=sat<sub>64</sub>(Rss+Rtt);</code>

### Class: XTYPE (slots 2,3)

#### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the Status Register is set. OVF remains set until explicitly cleared by a transfer to the status register.

## Intrinsics

Rdd=add(Rs,Rtt)	Word64 Q6_P_add_RP(Word32 Rs, Word64 Rtt)
Rdd=add(Rss,Rtt)	Word64 Q6_P_add_PP(Word64 Rss, Word64 Rtt)
Rdd=add(Rss,Rtt):sat	Word64 Q6_P_add_PP_sat(Word64 Rss, Word64 Rtt)

## Encoding

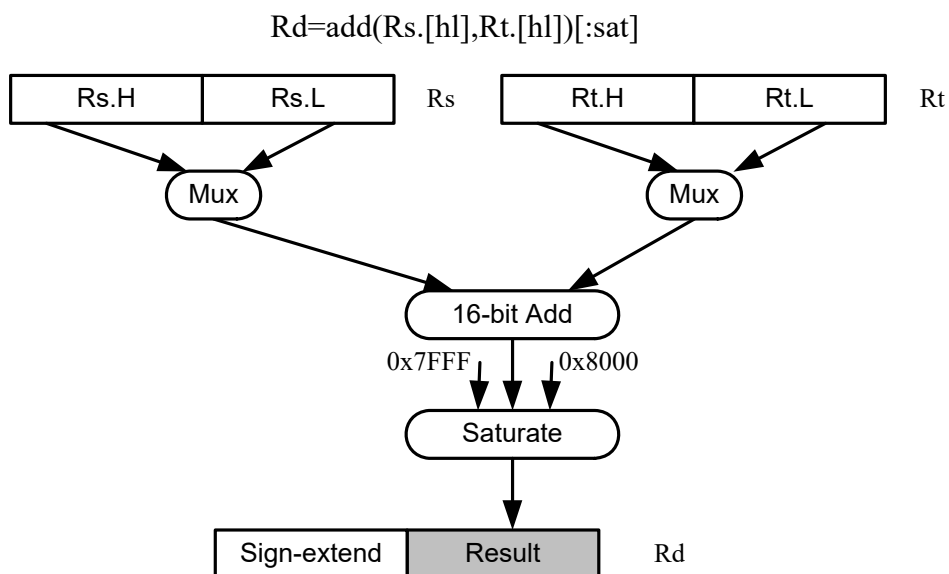
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5					MinOp			d5										
1	1	0	1	0	0	1	1	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	1	d	d	d	d	d	Rdd=add(Rss,Rtt)
1	1	0	1	0	0	1	1	0	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	1	d	d	d	d	d	Rdd=add(Rss,Rtt):sat
1	1	0	1	0	0	1	1	0	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	0	d	d	d	d	d	Rdd=add(Rss,Rtt):raw:lo
1	1	0	1	0	0	1	1	0	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	1	d	d	d	d	d	Rdd=add(Rss,Rtt):raw:hi
1	1	0	1	0	1	0	1	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	-	-	d	d	d	d	d	Rd=add(Rs,Rt):sat:depreca ted

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Add halfword

Perform a 16-bit add with optional saturation, and place the result in either the upper or lower half of a register. If the result goes in the upper half, the sources are any high or low halfword of Rs and Rt. The lower 16 bits of the result are zeroed.

If the result is placed in the lower 16 bits of Rd, the Rs source can be either high or low, but the other source must be the low halfword of Rt. In this case, the upper halfword of Rd is the sign-extension of the low halfword.



### Syntax

```
Rd=add(Rt.L, Rs.[HL])[:sat]
```

```
Rd=add(Rt.[HL], Rs.[HL])[:sat]:  
<<16
```

### Behavior

```
Rd=[sat16](Rt.h[0]+Rs.h[01]);
```

```
Rd=( [sat16](Rt.h[01]+Rs.h[01]) ) <<16;
```

### Class: XTYPE (slots 2,3)

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

## Intrinsics

Rd=add(Rt.H,Rs.H):<<16	Word32 Q6_R_add_RhRh_s16(Word32 Rt, Word32 Rs)
Rd=add(Rt.H,Rs.H):sat:<<16	Word32 Q6_R_add_RhRh_sat_s16(Word32 Rt, Word32 Rs)
Rd=add(Rt.H,Rs.L):<<16	Word32 Q6_R_add_RhRl_s16(Word32 Rt, Word32 Rs)
Rd=add(Rt.H,Rs.L):sat:<<16	Word32 Q6_R_add_RhRl_sat_s16(Word32 Rt, Word32 Rs)
Rd=add(Rt.L,Rs.H)	Word32 Q6_R_add_RlRh(Word32 Rt, Word32 Rs)
Rd=add(Rt.L,Rs.H):<<16	Word32 Q6_R_add_RlRh_s16(Word32 Rt, Word32 Rs)
Rd=add(Rt.L,Rs.H):sat	Word32 Q6_R_add_RlRh_sat(Word32 Rt, Word32 Rs)
Rd=add(Rt.L,Rs.H):sat:<<16	Word32 Q6_R_add_RlRh_sat_s16(Word32 Rt, Word32 Rs)
Rd=add(Rt.L,Rs.L)	Word32 Q6_R_add_RlRl(Word32 Rt, Word32 Rs)
Rd=add(Rt.L,Rs.L):<<16	Word32 Q6_R_add_RlRl_s16(Word32 Rt, Word32 Rs)
Rd=add(Rt.L,Rs.L):sat	Word32 Q6_R_add_RlRl_sat(Word32 Rt, Word32 Rs)
Rd=add(Rt.L,Rs.L):sat:<<16	Word32 Q6_R_add_RlRl_sat_s16(Word32 Rt, Word32 Rs)

## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				s5					Parse		t5					MinOp		d5										
1	1	0	1	0	1	0	1	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	d	d	d	d	Rd=add(Rt.L,Rs.L)	
1	1	0	1	0	1	0	1	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	-	d	d	d	d	Rd=add(Rt.L,Rs.H)	
1	1	0	1	0	1	0	1	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	-	d	d	d	d	Rd=add(Rt.L,Rs.L):sat	
1	1	0	1	0	1	0	1	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	-	d	d	d	d	Rd=add(Rt.L,Rs.H):sat	
1	1	0	1	0	1	0	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	d	d	d	d	Rd=add(Rt.L,Rs.L):<<16	
1	1	0	1	0	1	0	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	d	d	d	d	Rd=add(Rt.L,Rs.H):<<16	
1	1	0	1	0	1	0	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	d	d	d	d	Rd=add(Rt.H,Rs.L):<<16	
1	1	0	1	0	1	0	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	d	d	d	d	Rd=add(Rt.H,Rs.H):<<16	
1	1	0	1	0	1	0	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	0	d	d	d	d	Rd=add(Rt.L,Rs.L):sat:<<16	
1	1	0	1	0	1	0	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	1	d	d	d	d	Rd=add(Rt.L,Rs.H):sat:<<16	
1	1	0	1	0	1	0	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	0	d	d	d	d	Rd=add(Rt.H,Rs.L):sat:<<16	
1	1	0	1	0	1	0	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	1	d	d	d	d	Rd=add(Rt.H,Rs.H):sat:<<16	

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Add or subtract doublewords with carry

Add or subtract with carry. Predicate register Px is used as an extra input and output.

For adds, add the LSB of the predicate to the sum of the two input pairs.

For subtracts, the predicate is considered a not-borrow. The LSB of the predicate is added to the first source register and the logical complement of the second argument.

The carry-out from the sum is saved in predicate Px.

These instructions allow efficient addition or subtraction of numbers larger than 64 bits.

Syntax	Behavior
<code>Rdd=add(Rss,Rtt,Px):carry</code>	PREDUSE_TIMING; $Rdd = Rss + Rtt + Px[0];$ $Px = \text{carry\_from\_add}(Rss,Rtt,Px[0]) ? 0xff : 0x00;$
<code>Rdd=sub(Rss,Rtt,Px):carry</code>	PREDUSE_TIMING; $Rdd = Rss + \sim Rtt + Px[0];$ $Px = \text{carry\_from\_add}(Rss,\sim Rtt,Px[0]) ? 0xff : 0x00;$

### Class: XTYPE (slots 2,3)

#### Notes

- The predicate generated by this instruction can not be used as a .new predicate, nor can it be automatically ANDed with another predicate.

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				Maj		s5					Parse		t5					x2		d5									
1	1	0	0	0	0	1	0	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	x	x	d	d	d	d	d	Rdd=add(Rss,Rtt,Px):carry
1	1	0	0	0	0	1	0	1	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	x	x	d	d	d	d	d	Rdd=sub(Rss,Rtt,Px):carry

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x2	Field to encode register x
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Clip to unsigned

Clip input to unsigned integer.

### Syntax

```
Rd=clip(Rs, #u5)
```

### Behavior

```
Rd=MIN((1<<#u)-1, MAX(Rs, -(1<<#u)));
```

**Class: XTYPE (slots 2,3)**

### Notes

- This instruction can only execute on a core with the Hexagon audio extensions

### Intrinsics

```
Rd=clip(Rs, #u5)
```

```
Word32 Q6_R_clip_RI(Word32 Rs, Word32 Iu5)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse		MinOp					d5											
1	0	0	0	1	0	0	0	1	1	0	s	s	s	s	s	P	P	0	i	i	i	i	i	1	0	1	d	d	d	d	d	Rd=clip(Rs,#u5)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Logical doublewords

Perform bitwise logical AND, OR, XOR, and NOT operations.

The source and destination registers are 64-bit.

For 32-bit logical operations, see the ALU32 logical instructions.

Syntax	Behavior
Rdd=and(Rss,Rtt)	Rdd=Rss&Rtt;
Rdd=and(Rtt,~Rss)	Rdd = (Rtt & ~Rss);
Rdd=not(Rss)	Rdd=~Rss;
Rdd=or(Rss,Rtt)	Rdd=Rss Rtt;
Rdd=or(Rtt,~Rss)	Rdd = (Rtt   ~Rss);
Rdd=xor(Rss,Rtt)	Rdd=Rss^Rtt;

### Class: XTYPE (slots 2,3)

#### Intrinsics

Rdd=and(Rss,Rtt)	Word64 Q6_P_and_PP(Word64 Rss, Word64 Rtt)
Rdd=and(Rtt,~Rss)	Word64 Q6_P_and_PnP(Word64 Rtt, Word64 Rss)
Rdd=not(Rss)	Word64 Q6_P_not_P(Word64 Rss)
Rdd=or(Rss,Rtt)	Word64 Q6_P_or_PP(Word64 Rss, Word64 Rtt)
Rdd=or(Rtt,~Rss)	Word64 Q6_P_or_PnP(Word64 Rtt, Word64 Rss)
Rdd=xor(Rss,Rtt)	Word64 Q6_P_xor_PP(Word64 Rss, Word64 Rtt)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse		MinOp			d5													
1	0	0	0	0	0	0	0	1	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	0	d	d	d	d	d	Rdd=not(Rss)
ICLASS			RegType				s5					Parse		t5			MinOp			d5												
1	1	0	1	0	0	1	1	1	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	d	d	d	d	d	Rdd=and(Rss,Rtt)
1	1	0	1	0	0	1	1	1	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	d	d	d	d	d	Rdd=and(Rtt,~Rss)
1	1	0	1	0	0	1	1	1	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	d	d	d	d	d	Rdd=or(Rss,Rtt)
1	1	0	1	0	0	1	1	1	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	d	d	d	d	d	Rdd=or(Rtt,~Rss)
1	1	0	1	0	0	1	1	1	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	0	d	d	d	d	d	Rdd=xor(Rss,Rtt)

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

---

<b>Field name</b>	<b>Description</b>
MajOp	Major opcode



## Logical-logical doublewords

Perform a logical operation of the two source operands, then perform a second logical operation of the result with the destination register Rxx.

The source and destination registers are 64-bit.

### Syntax

```
Rxx ^= xor (Rss, Rtt)
```

### Behavior

```
Rxx ^= Rss ^ Rtt;
```

### Class: XTYPE (slots 2,3)

### Intrinsics

```
Rxx ^= xor (Rss, Rtt)
```

```
Word64 Q6_P_xorxacc_PP (Word64 Rxx, Word64  
Rss, Word64 Rtt)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS				RegType				Maj		s5					Parse		t5					Min		x5									
1	1	0	0	1	0	1	0	1	0	-	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	x	x	x	x	x		Rxx ^= xor(Rss,Rtt)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Logical-logical words

Perform a logical operation of the two source operands, then perform a second logical operation of the result with the destination register Rx.

The source and destination registers are 32-bit.

Syntax	Behavior
<code>Rx=or (Ru, and (Rx, #s10))</code>	<code>Rx = Ru   (Rx &amp; apply_extension(#s));</code>
<code>Rx[&amp; ^]=and (Rs, Rt)</code>	<code>Rx [ &amp;^]= (Rs [ &amp;^] Rt);</code>
<code>Rx[&amp; ^]=and (Rs, ~Rt)</code>	<code>Rx [ &amp;^]= (Rs [ &amp;^] ~Rt);</code>
<code>Rx[&amp; ^]=or (Rs, Rt)</code>	<code>Rx [ &amp;^]= (Rs [ &amp;^] Rt);</code>
<code>Rx[&amp; ^]=xor (Rs, Rt)</code>	<code>Rx [ &amp;^]=Rs [ &amp;^]Rt;</code>
<code>Rx =and (Rs, #s10)</code>	<code>Rx = Rx   (Rs &amp; apply_extension(#s));</code>
<code>Rx =or (Rs, #s10)</code>	<code>Rx = Rx   (Rs   apply_extension(#s));</code>

**Class: XTYPE (slots 2,3)****Intrinsics**

Rx&=and(Rs,Rt)	Word32 Q6_R_andand_RR(Word32 Rx, Word32 Rs, Word32 Rt)
Rx&=and(Rs,~Rt)	Word32 Q6_R_andand_RnR(Word32 Rx, Word32 Rs, Word32 Rt)
Rx&=or(Rs,Rt)	Word32 Q6_R_orand_RR(Word32 Rx, Word32 Rs, Word32 Rt)
Rx&=xor(Rs,Rt)	Word32 Q6_R_xorand_RR(Word32 Rx, Word32 Rs, Word32 Rt)
Rx=or(Ru, and(Rx, #s10))	Word32 Q6_R_or_and_RRI(Word32 Ru, Word32 Rx, Word32 Is10)
Rx^=and(Rs,Rt)	Word32 Q6_R_andxacc_RR(Word32 Rx, Word32 Rs, Word32 Rt)
Rx^=and(Rs,~Rt)	Word32 Q6_R_andxacc_RnR(Word32 Rx, Word32 Rs, Word32 Rt)
Rx^=or(Rs,Rt)	Word32 Q6_R_orxacc_RR(Word32 Rx, Word32 Rs, Word32 Rt)
Rx^=xor(Rs,Rt)	Word32 Q6_R_xorxacc_RR(Word32 Rx, Word32 Rs, Word32 Rt)
Rx =and(Rs,#s10)	Word32 Q6_R_andor_RI(Word32 Rx, Word32 Rs, Word32 Is10)
Rx =and(Rs,Rt)	Word32 Q6_R_andor_RR(Word32 Rx, Word32 Rs, Word32 Rt)
Rx =and(Rs,~Rt)	Word32 Q6_R_andor_RnR(Word32 Rx, Word32 Rs, Word32 Rt)
Rx =or(Rs,#s10)	Word32 Q6_R_oror_RI(Word32 Rx, Word32 Rs, Word32 Is10)
Rx =or(Rs,Rt)	Word32 Q6_R_oror_RR(Word32 Rx, Word32 Rs, Word32 Rt)
Rx =xor(Rs,Rt)	Word32 Q6_R_xoror_RR(Word32 Rx, Word32 Rs, Word32 Rt)

**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS		RegType				s5					Parse					x5																	
1	1	0	1	1	0	1	0	0	0	i	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	x	x	x	x	x	Rx =and(Rs,#s10)	
ICLASS		RegType				x5					Parse					u5																	
1	1	0	1	1	0	1	0	0	1	i	x	x	x	x	x	P	P	i	i	i	i	i	i	i	i	i	u	u	u	u	u	Rx=or(Ru,and(Rx,#s10))	
ICLASS		RegType				s5					Parse					x5																	
1	1	0	1	1	0	1	0	1	0	i	s	s	s	s	s	P	P	i	i	i	i	i	i	i	i	i	x	x	x	x	x	Rx =or(Rs,#s10)	
ICLASS		RegType				MajOp		s5					Parse					t5			MinOp		x5										
1	1	1	0	1	1	1	1	0	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	0	x	x	x	x	x	Rx =and(Rs,~Rt)
1	1	1	0	1	1	1	1	0	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	x	x	x	x	x	Rx&=and(Rs,~Rt)	
1	1	1	0	1	1	1	1	0	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	0	x	x	x	x	x	Rx^=and(Rs,~Rt)	
1	1	1	0	1	1	1	1	0	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	x	x	x	x	x	Rx&=and(Rs,Rt)	
1	1	1	0	1	1	1	1	0	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	x	x	x	x	x	Rx&=or(Rs,Rt)	
1	1	1	0	1	1	1	1	0	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	0	x	x	x	x	x	Rx&=xor(Rs,Rt)	

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
1	1	1	0	1	1	1	1	0	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	1	x	x	x	x	x	Rx =and(Rs,Rt)
1	1	1	0	1	1	1	1	1	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	1	x	x	x	x	x	Rx^=xor(Rs,Rt)
1	1	1	0	1	1	1	1	1	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	x	x	x	x	x	Rx =or(Rs,Rt)
1	1	1	0	1	1	1	1	1	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	x	x	x	x	x	Rx =xor(Rs,Rt)
1	1	1	0	1	1	1	1	1	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	0	x	x	x	x	x	Rx^=and(Rs,Rt)
1	1	1	0	1	1	1	1	1	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	1	x	x	x	x	x	Rx^=or(Rs,Rt)

Field name	Description
RegType	Register type
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
Parse	Packet/loop parse bits
s5	Field to encode register s
t5	Field to encode register t
u5	Field to encode register u
x5	Field to encode register x

## Maximum words

Select either the signed or unsigned maximum of two source registers and place in a destination register Rdd.

Syntax	Behavior
Rd=max(Rs,Rt)	Rd = max(Rs,Rt);
Rd=maxu(Rs,Rt)	Rd = max(Rs.uw[0],Rt.uw[0]);

### Class: XTYPE (slots 2,3)

#### Intrinsics

Rd=max(Rs,Rt)	Word32 Q6_R_max_RR(Word32 Rs, Word32 Rt)
Rd=maxu(Rs,Rt)	UWord32 Q6_R_maxu_RR(Word32 Rs, Word32 Rt)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5					MinOp			d5										
1	1	0	1	0	1	0	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	-	-	d	d	d	d	d	Rd=max(Rs,Rt)
1	1	0	1	0	1	0	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	-	-	d	d	d	d	d	Rd=maxu(Rs,Rt)

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Maximum doublewords

Select either the signed or unsigned maximum of two 64-bit source registers and place in a destination register.

Syntax	Behavior
<code>Rdd=max(Rss,Rtt)</code>	<code>Rdd = max(Rss,Rtt);</code>
<code>Rdd=maxu(Rss,Rtt)</code>	<code>Rdd = max(Rss.u64,Rtt.u64);</code>

### Class: XTYPE (slots 2,3)

#### Intrinsics

<code>Rdd=max(Rss,Rtt)</code>	<code>Word64 Q6_P_max_PP(Word64 Rss, Word64 Rtt)</code>
<code>Rdd=maxu(Rss,Rtt)</code>	<code>UWord64 Q6_P_maxu_PP(Word64 Rss, Word64 Rtt)</code>

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5					MinOp			d5										
1	1	0	1	0	0	1	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	0	d	d	d	d	d	Rdd=max(Rss,Rtt)
1	1	0	1	0	0	1	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	1	d	d	d	d	d	Rdd=maxu(Rss,Rtt)

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Minimum words

Select either the signed or unsigned minimum of two source registers and place in destination register Rd.

Syntax	Behavior
Rd=min(Rt,Rs)	Rd = min(Rt,Rs);
Rd=minu(Rt,Rs)	Rd = min(Rt.uw[0],Rs.uw[0]);

### Class: XTYPE (slots 2,3)

#### Intrinsics

Rd=min(Rt,Rs)

Word32 Q6\_R\_min\_RR(Word32 Rt, Word32 Rs)

Rd=minu(Rt,Rs)

UWord32 Q6\_R\_minu\_RR(Word32 Rt, Word32 Rs)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5					MinOp			d5										
1	1	0	1	0	1	0	1	1	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	-	-	d	d	d	d	d	Rd=min(Rt,Rs)
1	1	0	1	0	1	0	1	1	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	-	-	d	d	d	d	d	Rd=minu(Rt,Rs)

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Minimum doublewords

Select either the signed or unsigned minimum of two 64-bit source registers and place in the destination register Rdd.

Syntax	Behavior
<code>Rdd=min(Rtt,Rss)</code>	<code>Rdd = min(Rtt,Rss);</code>
<code>Rdd=minu(Rtt,Rss)</code>	<code>Rdd = min(Rtt.u64,Rss.u64);</code>

### Class: XTYPE (slots 2,3)

#### Intrinsics

<code>Rdd=min(Rtt,Rss)</code>	<code>Word64 Q6_P_min_PP(Word64 Rtt, Word64 Rss)</code>
<code>Rdd=minu(Rtt,Rss)</code>	<code>UWord64 Q6_P_minu_PP(Word64 Rtt, Word64 Rss)</code>

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5					MinOp			d5										
1	1	0	1	0	0	1	1	1	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	0	d	d	d	d	d	Rdd=min(Rtt,Rss)
1	1	0	1	0	0	1	1	1	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	1	d	d	d	d	d	Rdd=minu(Rtt,Rss)

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t



## Modulo wrap

Wrap the Rs value into the modulo range from 0 to Rt.

If Rs is greater than or equal to Rt, wrap it to the bottom of the range by subtracting Rt.

If Rs is less than zero, wrap it to the top of the range by adding Rt.

Otherwise, when Rs fits within the range, no adjustment is necessary. The result is returned in register Rd.

### Syntax

```
Rd=modwrap(Rs,Rt)
```

### Behavior

```
if (Rs < 0) {
    Rd = Rs + Rt.uw[0];
} else if (Rs.uw[0] >= Rt.uw[0]) {
    Rd = Rs - Rt.uw[0];
} else {
    Rd = Rs;
}
```

### Class: XTYPE (slots 2,3)

### Intrinsics

```
Rd=modwrap(Rs,Rt)
```

```
Word32 Q6_R_modwrap_RR(Word32 Rs, Word32
Rt)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType		s5					Parse		t5					MinOp		d5														
1	1	0	1	0	0	1	1	1	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	1	d	d	d	d	d	Rd=modwrap(Rs,Rt)

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Negate

The first form of this instruction performs a negate on a 32-bit register with saturation. If the input is 0x80000000, the result is saturated to 0x7fffffff. The nonsaturating 32-bit register negate is a ALU32-class instruction and can execute on any slot.

The second form of this instruction negates a 64-bit source register and places the result in destination Rdd.

Syntax	Behavior
Rd=neg(Rs):sat	Rd = sat <sub>32</sub> (-Rs.s64);
Rdd=neg(Rss)	Rdd = -Rss;

### Class: XTYPE (slots 2,3)

#### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

#### Intrinsics

Rd=neg(Rs):sat	Word32 Q6_R_neg_R_sat(Word32 Rs)
Rdd=neg(Rss)	Word64 Q6_P_neg_P(Word64 Rss)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp		s5					Parse		MinOp			d5												
1	0	0	0	0	0	0	0	1	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	1	d	d	d	d	d	Rdd=neg(Rss)
1	0	0	0	1	1	0	0	1	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	1	0	d	d	d	d	d	Rd=neg(Rs):sat

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Round

Perform either arithmetic (.5 is rounded up) or convergent (.5 is rounded towards even) rounding to any bit location.

Arithmetic rounding has optional saturation. In this version, the result is saturated to a 32-bit number after adding the rounding constant. After the rounding and saturation have been performed, the final result is right shifted using a sign-extending shift.

Syntax	Behavior
<code>Rd=cround(Rs,#u5)</code>	<code>Rd = (#u==0)?Rs:convround(Rs,2**(#u-1))&gt;&gt;#u;</code>
<code>Rd=cround(Rs,Rt)</code>	<code>Rd = (zxt<sub>5-&gt;32</sub>(Rt)==0)?Rs:convround(Rs,2** (zxt<sub>5-&gt;32</sub>(Rt) - 1))&gt;&gt;zxt<sub>5-&gt;32</sub>(Rt);</code>
<code>Rd=round(Rs,#u5)[:sat]</code>	<code>Rd = ([sat<sub>32</sub>]((#u==0)?(Rs):round(Rs,2**(#u-1))))&gt;&gt;#u;</code>
<code>Rd=round(Rs,Rt)[:sat]</code>	<code>Rd = ([sat<sub>32</sub>]((zxt<sub>5-&gt;32</sub>(Rt)==0)?(Rs):round(Rs,2** (zxt<sub>5-&gt;32</sub>(Rt) - 1))))&gt;&gt;zxt<sub>5-&gt;32</sub>(Rt);</code>
<code>Rd=round(Rss):sat</code>	<code>tmp=sat<sub>64</sub>(Rss+0x080000000ULL); Rd = tmp.w[1];</code>
<code>Rdd=cround(Rss,#u6)</code>	<pre>if (#u == 0) {     Rdd = Rss; } else if ((Rss &amp; (size8s_t)((1LL &lt;&lt; (#u - 1)) - 1LL)) == 0) {     src_128 = sxt<sub>64-&gt;128</sub>(Rss);     rndbit_128 = sxt<sub>64-&gt;128</sub>(1LL);     rndbit_128 = (rndbit_128 &lt;&lt; #u);     rndbit_128 = (rndbit_128 &amp; src_128);     rndbit_128 = (size8s_t) (rndbit_128 &gt;&gt; 1);     tmp128 = src_128+rndbit_128;     tmp128 = (size8s_t) (tmp128 &gt;&gt; #u);     Rdd = sxt<sub>128-&gt;64</sub>(tmp128); } else {     size16s_t rndbit_128 = sxt<sub>64-&gt;128</sub>((1LL &lt;&lt; (#u - 1)));     size16s_t src_128 = sxt<sub>64-&gt;128</sub>(Rss);     size16s_t tmp128 = src_128+rndbit_128;     tmp128 = (size8s_t) (tmp128 &gt;&gt; #u);     Rdd = sxt<sub>128-&gt;64</sub>(tmp128); }</pre>

Syntax	Behavior
Rdd=cround(Rss,Rt)	<pre> if (zxt<sub>6-&gt;32</sub>(Rt) == 0) {     Rdd = Rss; } else if ((Rss &amp; (size8s_t)((1LL &lt;&lt; (zxt<sub>6-&gt;32</sub>(Rt) - 1)) - 1LL)) == 0) {     src_128 = sxt<sub>64-&gt;128</sub>(Rss);     rndbit_128 = sxt<sub>64-&gt;128</sub>(1LL);     rndbit_128 = (rndbit_128 &lt;&lt; zxt<sub>6-&gt;32</sub>(Rt));     rndbit_128 = (rndbit_128 &amp; src_128);     rndbit_128 = (size8s_t) (rndbit_128 &gt;&gt; 1);     tmp128 = src_128+rndbit_128;     tmp128 = (size8s_t) (tmp128 &gt;&gt; zxt<sub>6-&gt;32</sub>(Rt));     Rdd = sxt<sub>128-&gt;64</sub>(tmp128); } else {     size16s_t rndbit_128 = sxt<sub>64-&gt;128</sub>((1LL &lt;&lt; (zxt<sub>6-&gt;32</sub>(Rt) - 1));     size16s_t src_128 = sxt<sub>64-&gt;128</sub>(Rss);     size16s_t tmp128 = src_128+rndbit_128;     tmp128 = (size8s_t) (tmp128 &gt;&gt; zxt<sub>6-&gt;32</sub>(Rt));     Rdd = sxt<sub>128-&gt;64</sub>(tmp128); } </pre>

**Class: XTYPE (slots 2,3)****Notes**

- This instruction can only execute on a core with the Hexagon audio extensions
- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

## Intrinsics

Rd=cround(Rs, #u5)	Word32 Q6_R_cround_RI(Word32 Rs, Word32 Iu5)
Rd=cround(Rs, Rt)	Word32 Q6_R_cround_RR(Word32 Rs, Word32 Rt)
Rd=round(Rs, #u5)	Word32 Q6_R_round_RI(Word32 Rs, Word32 Iu5)
Rd=round(Rs, #u5) :sat	Word32 Q6_R_round_RI_sat(Word32 Rs, Word32 Iu5)
Rd=round(Rs, Rt)	Word32 Q6_R_round_RR(Word32 Rs, Word32 Rt)
Rd=round(Rs, Rt) :sat	Word32 Q6_R_round_RR_sat(Word32 Rs, Word32 Rt)
Rd=round(Rss) :sat	Word32 Q6_R_round_P_sat(Word64 Rss)
Rdd=cround(Rss, #u6)	Word64 Q6_P_cround_PI(Word64 Rss, Word32 Iu6)
Rdd=cround(Rss, Rt)	Word64 Q6_P_cround_PR(Word64 Rss, Word32 Rt)

## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse				MinOp			d5											
1	0	0	0	1	0	0	0	1	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	1	d	d	d	d	d	Rd=round(Rss):sat
1	0	0	0	1	1	0	0	1	1	1	s	s	s	s	s	P	P	0	i	i	i	i	i	0	0	-	d	d	d	d	d	Rd=cround(Rs,#u5)
1	0	0	0	1	1	0	0	1	1	1	s	s	s	s	s	P	P	0	i	i	i	i	i	1	0	-	d	d	d	d	d	Rd=round(Rs,#u5)
1	0	0	0	1	1	0	0	1	1	1	s	s	s	s	s	P	P	0	i	i	i	i	i	1	1	-	d	d	d	d	d	Rd=round(Rs,#u5):sat
1	0	0	0	1	1	0	0	1	1	1	s	s	s	s	s	P	P	i	i	i	i	i	i	0	1	-	d	d	d	d	d	Rdd=cround(Rss,#u6)
ICLASS			RegType				Maj		s5					Parse				t5			Min		d5									
1	1	0	0	0	1	1	0	1	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	d	d	d	d	d	Rd=cround(Rs,Rt)
1	1	0	0	0	1	1	0	1	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	-	d	d	d	d	d	Rdd=cround(Rss,Rt)
1	1	0	0	0	1	1	0	1	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	-	d	d	d	d	d	Rd=round(Rs,Rt)
1	1	0	0	0	1	1	0	1	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	-	d	d	d	d	d	Rd=round(Rs,Rt):sat

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
MajOp	Major opcode
MinOp	Minor opcode
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Subtract doublewords

Subtract the 64-bit register *Rss* from register *Rtt*.

Syntax	Behavior
$Rd = \text{sub}(Rt, Rs) : \text{sat} : \text{deprecated}$	$Rd = \text{sat}_{32}(Rt - Rs);$
$Rdd = \text{sub}(Rtt, Rss)$	$Rdd = Rtt - Rss;$

**Class: XTYPE (slots 2,3)**

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

### Intrinsics

$Rdd = \text{sub}(Rtt, Rss)$  `Word64 Q6_P_sub_PP(Word64 Rtt, Word64 Rss)`

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5					MinOp			d5										
1	1	0	1	0	0	1	1	0	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	1	d	d	d	d	d	Rdd=sub(Rtt,Rss)
1	1	0	1	0	1	0	1	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	-	-	d	d	d	d	d	Rd=sub(Rt,Rs):sat:depreca ted

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Subtract and accumulate words

Subtract Rs from Rt, then add the resulting value with Rx. The result is saved in Rx.

### Syntax

$$Rx += \text{sub}(Rt, Rs)$$

### Behavior

$$Rx = Rx + Rt - Rs;$$

**Class: XTYPE (slots 2,3)**

### Intrinsics

$$Rx += \text{sub}(Rt, Rs)$$

$$\text{Word32 } Q6\_R\_subacc\_RR(\text{Word32 } Rx, \text{Word32 } Rt, \text{Word32 } Rs)$$

### Encoding

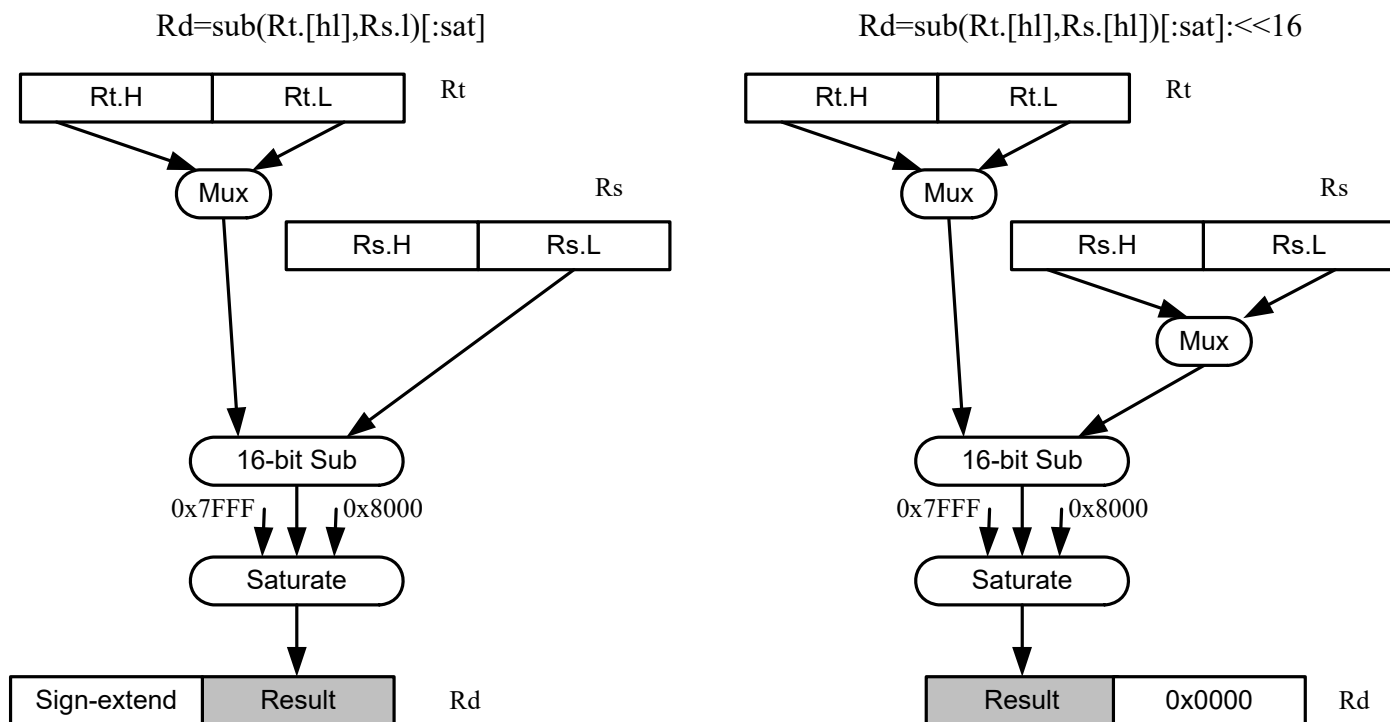
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			x5							
1	1	1	0	1	1	1	1	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	1	x	x	x	x	x	Rx+=sub(Rt,Rs)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Subtract halfword

Perform a 16-bit subtract with optional saturation and place the result in either the upper or lower half of a register. If the result goes in the upper half, the sources can be any high or low halfword of Rs and Rt. The lower 16 bits of the result are zeroed.

If the result is placed in the lower 16 bits of Rd, the Rs source can be either high or low, but the other source must be the low halfword of Rt. In this case, the upper halfword of Rd is the sign-extension of the low halfword.



Syntax	Behavior
<code>Rd=sub(Rt.L,Rs.[HL])[:sat]</code>	<code>Rd=[sat<sub>16</sub>](Rt.h[0]-Rs.h[01]);</code>
<code>Rd=sub(Rt.[HL],Rs.[HL])[:sat]&lt;&lt;16</code>	<code>Rd=( [sat<sub>16</sub>](Rt.h[01]-Rs.h[01]))&lt;&lt;16;</code>

**Class: XTYPE (slots 2,3)**

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.



## Intrinsics

Rd=sub(Rt.H,Rs.H):<<16	Word32 Q6_R_sub_RhRh_s16(Word32 Rt, Word32 Rs)
Rd=sub(Rt.H,Rs.H):sat:<<16	Word32 Q6_R_sub_RhRh_sat_s16(Word32 Rt, Word32 Rs)
Rd=sub(Rt.H,Rs.L):<<16	Word32 Q6_R_sub_RhRl_s16(Word32 Rt, Word32 Rs)
Rd=sub(Rt.H,Rs.L):sat:<<16	Word32 Q6_R_sub_RhRl_sat_s16(Word32 Rt, Word32 Rs)
Rd=sub(Rt.L,Rs.H)	Word32 Q6_R_sub_RlRh(Word32 Rt, Word32 Rs)
Rd=sub(Rt.L,Rs.H):<<16	Word32 Q6_R_sub_RlRh_s16(Word32 Rt, Word32 Rs)
Rd=sub(Rt.L,Rs.H):sat	Word32 Q6_R_sub_RlRh_sat(Word32 Rt, Word32 Rs)
Rd=sub(Rt.L,Rs.H):sat:<<16	Word32 Q6_R_sub_RlRh_sat_s16(Word32 Rt, Word32 Rs)
Rd=sub(Rt.L,Rs.L)	Word32 Q6_R_sub_RlRl(Word32 Rt, Word32 Rs)
Rd=sub(Rt.L,Rs.L):<<16	Word32 Q6_R_sub_RlRl_s16(Word32 Rt, Word32 Rs)
Rd=sub(Rt.L,Rs.L):sat	Word32 Q6_R_sub_RlRl_sat(Word32 Rt, Word32 Rs)
Rd=sub(Rt.L,Rs.L):sat:<<16	Word32 Q6_R_sub_RlRl_sat_s16(Word32 Rt, Word32 Rs)

## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				s5					Parse		t5					MinOp		d5												
1	1	0	1	0	1	0	1	0	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	d	d	d	d	d	Rd=sub(Rt.L,Rs.L)
1	1	0	1	0	1	0	1	0	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	-	d	d	d	d	d	Rd=sub(Rt.L,Rs.H)
1	1	0	1	0	1	0	1	0	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	-	d	d	d	d	d	Rd=sub(Rt.L,Rs.L):sat
1	1	0	1	0	1	0	1	0	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	-	d	d	d	d	d	Rd=sub(Rt.L,Rs.H):sat
1	1	0	1	0	1	0	1	0	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	d	d	d	d	d	Rd=sub(Rt.L,Rs.L):<<16
1	1	0	1	0	1	0	1	0	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	d	d	d	d	d	Rd=sub(Rt.L,Rs.H):<<16
1	1	0	1	0	1	0	1	0	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	d	d	d	d	d	Rd=sub(Rt.H,Rs.L):<<16
1	1	0	1	0	1	0	1	0	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	d	d	d	d	d	Rd=sub(Rt.H,Rs.H):<<16
1	1	0	1	0	1	0	1	0	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	0	d	d	d	d	d	Rd=sub(Rt.L,Rs.L):sat:<<16
1	1	0	1	0	1	0	1	0	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	1	d	d	d	d	d	Rd=sub(Rt.L,Rs.H):sat:<<16
1	1	0	1	0	1	0	1	0	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	0	d	d	d	d	d	Rd=sub(Rt.H,Rs.L):sat:<<16
1	1	0	1	0	1	0	1	0	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	1	d	d	d	d	d	Rd=sub(Rt.H,Rs.H):sat:<<16

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Sign extend word to doubleword

Sign-extend a 32-bit word to a 64-bit doubleword.

### Syntax

`Rdd=sxtw(Rs)`

### Behavior

`Rdd = sxt32->64(Rs);`

**Class: XTYPE (slots 2,3)**

### Intrinsics

`Rdd=sxtw(Rs)`

`Word64 Q6_P_sxtw_R(Word32 Rs)`

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse		MinOp			d5													
1	0	0	0	0	1	0	0	0	1	-	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	-	d	d	d	d	d	Rdd=sxtw(Rs)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Vector absolute value halfwords

Take the absolute value of each of the four halfwords in the 64-bit source vector *Rss*. Place the result in *Rdd*.

Saturation is optionally available.

Syntax	Behavior
<code>Rdd=vabsh(Rss)</code>	<pre>for (i=0;i&lt;4;i++) {   Rdd.h[i]=ABS(Rss.h[i]); }</pre>
<code>Rdd=vabsh(Rss):sat</code>	<pre>for (i=0;i&lt;4;i++) {   Rdd.h[i]=sat<sub>16</sub>(ABS(Rss.h[i])); }</pre>

**Class: XTYPE (slots 2,3)**

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

### Intrinsics

<code>Rdd=vabsh(Rss)</code>	<code>Word64 Q6_P_vabsh_P(Word64 Rss)</code>
<code>Rdd=vabsh(Rss):sat</code>	<code>Word64 Q6_P_vabsh_P_sat(Word64 Rss)</code>

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp		s5					Parse		MinOp			d5												
1	0	0	0	0	0	0	0	0	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	0	d	d	d	d	d	Rdd=vabsh(Rss)
1	0	0	0	0	0	0	0	0	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	1	d	d	d	d	d	Rdd=vabsh(Rss):sat

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Vector absolute value words

Take the absolute value of each of the two words in the 64-bit source vector *Rss*. Place the result in *Rdd*.

Saturation is optionally available.

Syntax	Behavior
<code>Rdd=vabsw(Rss)</code>	<pre>for (i=0;i&lt;2;i++) {   Rdd.w[i]=ABS(Rss.w[i]); }</pre>
<code>Rdd=vabsw(Rss):sat</code>	<pre>for (i=0;i&lt;2;i++) {   Rdd.w[i]=sat<sub>32</sub>(ABS(Rss.w[i])); }</pre>

**Class: XTYPE (slots 2,3)**

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

### Intrinsics

<code>Rdd=vabsw(Rss)</code>	<code>Word64 Q6_P_vabsw_P(Word64 Rss)</code>
<code>Rdd=vabsw(Rss):sat</code>	<code>Word64 Q6_P_vabsw_P_sat(Word64 Rss)</code>

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp		s5					Parse		MinOp			d5												
1	0	0	0	0	0	0	0	0	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	1	0	d	d	d	d	d	Rdd=vabsw(Rss)
1	0	0	0	0	0	0	0	0	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	1	1	d	d	d	d	d	Rdd=vabsw(Rss):sat

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Vector absolute difference bytes

For each element in the source vector *Rss*, subtract the corresponding element in source vector *Rtt*. Take the absolute value of the results, and store into *Rdd*.

Syntax	Behavior
<code>Rdd=vabsdiffb(Rtt,Rss)</code>	<pre>for (i=0;i&lt;8;i++) {   Rdd.b[i]=ABS(Rtt.b[i] - Rss.b[i]); }</pre>
<code>Rdd=vabsdiffub(Rtt,Rss)</code>	<pre>for (i=0;i&lt;8;i++) {   Rdd.b[i]=ABS(Rtt.ub[i] - Rss.ub[i]); }</pre>

### Class: XTYPE (slots 2,3)

#### Intrinsics

`Rdd=vabsdiffb(Rtt,Rss)`    Word64 Q6\_P\_vabsdiffb\_PP(Word64 Rtt, Word64 Rss)

`Rdd=vabsdiffub(Rtt,Rss)`    Word64 Q6\_P\_vabsdiffub\_PP(Word64 Rtt, Word64 Rss)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp				s5					Parse		t5					MinOp			d5					
1	1	1	0	1	0	0	0	1	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	d	d	d	d	d	Rdd=vabsdiffub(Rtt,Rss)
1	1	1	0	1	0	0	0	1	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	d	d	d	d	d	Rdd=vabsdiffb(Rtt,Rss)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector absolute difference halfwords

For each element in the source vector *Rss*, subtract the corresponding element in source vector *Rtt*. Take the absolute value of the results, and store into *Rdd*.

### Syntax

```
Rdd=vabsdiffh(Rtt,Rss)
```

### Behavior

```
for (i=0;i<4;i++) {
    Rdd.h[i]=ABS(Rtt.h[i] - Rss.h[i]);
}
```

**Class: XTYPE (slots 2,3)**

### Intrinsics

```
Rdd=vabsdiffh(Rtt,Rss)
```

```
Word64 Q6_P_vabsdiffh_PP(Word64 Rtt, Word64
Rss)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse		t5					MinOp			d5								
1	1	1	0	1	0	0	0	0	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	d	d	d	d	d	Rdd=vabsdiffh(Rtt,Rss)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector absolute difference words

For each element in the source vector *Rss*, subtract the corresponding element in source vector *Rtt*. Take the absolute value of the results, and store into *Rdd*.

### Syntax

```
Rdd=vabsdiffw(Rtt,Rss)
```

### Behavior

```
for (i=0;i<2;i++) {
    Rdd.w[i]=ABS(Rtt.w[i] - Rss.w[i]);
}
```

**Class: XTYPE (slots 2,3)**

### Intrinsics

```
Rdd=vabsdiffw(Rtt,Rss)
```

```
Word64 Q6_P_vabsdiffw_PP(Word64 Rtt, Word64
Rss)
```

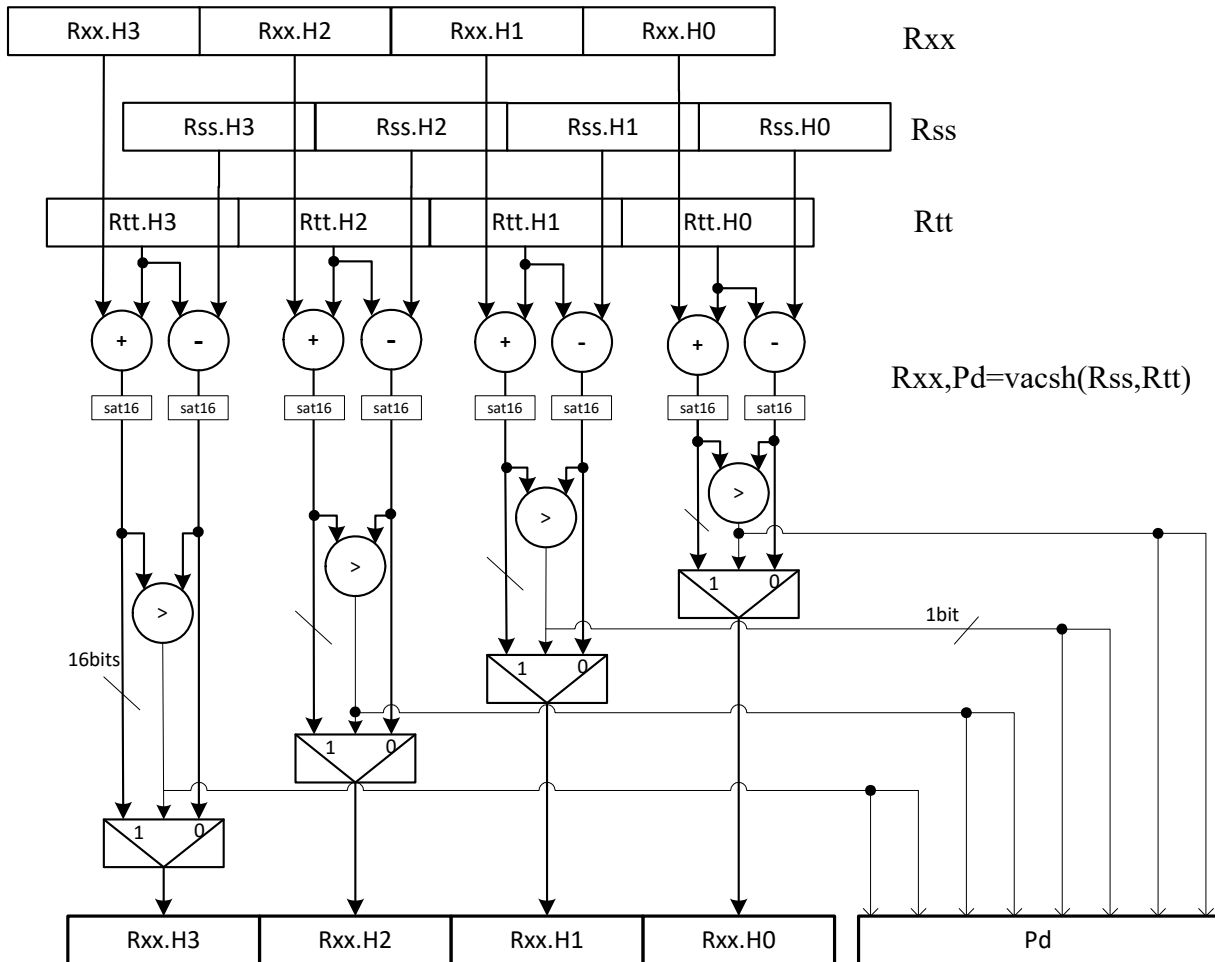
### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse		t5					MinOp			d5								
1	1	1	0	1	0	0	0	0	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	d	d	d	d	d	Rdd=vabsdiffw(Rtt,Rss)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector add compare and select maximum bytes

Add each byte element in Rxx and Rtt, and compare the resulting sums with the corresponding differences between Rss and Rtt. Store the maximum value of each compare in Rxx, and set the corresponding bits in a predicate destination to '1' if the compare result is greater, '0' if not. Each sum and difference is saturated to 8 bits before the compare, and the compare operation is a signed byte compare.



**Syntax**

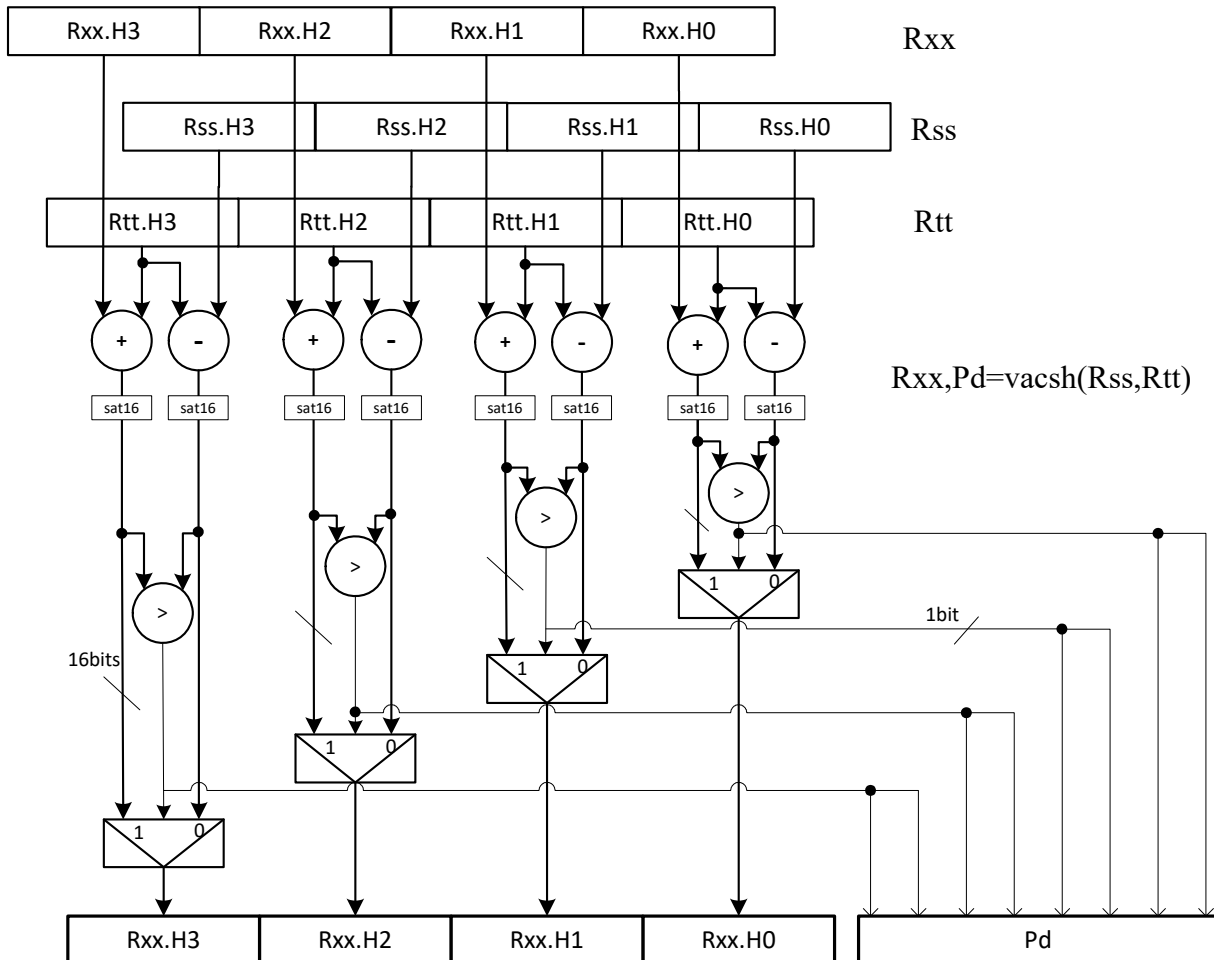
**Behavior**

**Class: N/A**



## Vector add compare and select maximum halfwords

Add each halfword element in Rxx and Rtt, and compare the resulting sums with the corresponding differences between Rss and Rtt. Store the maximum value of each compare in Rxx, and set the corresponding bits in a predicate destination to '11' if the compare result is greater, '00' if not. Each sum and difference is saturated to 16 bits before the compare, and the compare operation is a signed halfword compare.



### Syntax

$Rxx, Pe = \text{vacsh}(Rss, Rtt)$

### Behavior

```
for (i = 0; i < 4; i++) {
  xv = (int) Rxx.h[i];
  sv = (int) Rss.h[i];
  tv = (int) Rtt.h[i];
  xv = xv + tv;
  sv = sv - tv;
  Pe.i*2 = (xv > sv);
  Pe.i*2+1 = (xv > sv);
  Rxx.h[i] = sat16(max(xv, sv));
}
```

**Class: XTYPE (slots 2,3)****Notes**

- The predicate generated by this instruction cannot be used as a .new predicate, nor can it be automatically AND'd with another predicate.
- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp				s5					Parse		t5					e2		x5						
1	1	1	0	1	0	1	0	1	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	e	e	x	x	x	x	x	Rxx,Pe=vacsh(Rss,Rtt)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
e2	Field to encode register e
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Vector add halfwords

Add each of the four halfwords in 64-bit vector *Rss* to the corresponding halfword in vector *Rtt*.

Optionally saturate each 16-bit addition to either a signed or unsigned 16-bit value. Applying saturation to the *vaddh* instruction clamps the result to the signed range 0x8000 to 0x7fff, whereas applying saturation to the *vadduh* instruction ensures that the unsigned result falls within the range 0 to 0xffff. When saturation is not needed, use the *vaddh* form.

For the 32-bit version of this vector operation, see the ALU32 instructions.

Syntax	Behavior
<code>Rdd=vaddh(Rss,Rtt)[:sat]</code>	<pre>for (i=0;i&lt;4;i++) {   Rdd.h[i]=[sat<sub>16</sub>](Rss.h[i]+Rtt.h[i]); }</pre>
<code>Rdd=vadduh(Rss,Rtt):sat</code>	<pre>for (i=0;i&lt;4;i++) {   Rdd.h[i]=usat<sub>16</sub>(Rss.uh[i]+Rtt.uh[i]); }</pre>

### Class: XTYPE (slots 2,3)

#### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

#### Intrinsics

<code>Rdd=vaddh(Rss,Rtt)</code>	<code>Word64 Q6_P_vaddh_PP(Word64 Rss, Word64 Rtt)</code>
<code>Rdd=vaddh(Rss,Rtt):sat</code>	<code>Word64 Q6_P_vaddh_PP_sat(Word64 Rss, Word64 Rtt)</code>
<code>Rdd=vadduh(Rss,Rtt):sat</code>	<code>Word64 Q6_P_vadduh_PP_sat(Word64 Rss, Word64 Rtt)</code>

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5					MinOp		d5											
1	1	0	1	0	0	1	1	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	d	d	d	d	d	Rdd=vaddh(Rss,Rtt)
1	1	0	1	0	0	1	1	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	d	d	d	d	d	Rdd=vaddh(Rss,Rtt):sat
1	1	0	1	0	0	1	1	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	0	d	d	d	d	d	Rdd=vadduh(Rss,Rtt):sat

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits

<b>Field name</b>	<b>Description</b>
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector add halfwords with saturate and pack to unsigned bytes

Add the four 16-bit halfwords of Rss to the four 16-bit halfwords of Rtt. The results are saturated to unsigned 8-bits and packed in destination register Rd.

### Syntax

```
Rd=vaddhub(Rss,Rtt):sat
```

### Behavior

```
for (i=0;i<4;i++) {
    Rd.b[i]=usat8(Rss.h[i]+Rtt.h[i]);
}
```

**Class: XTYPE (slots 2,3)**

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

### Intrinsics

```
Rd=vaddhub(Rss,Rtt):sat Word32 Q6_R_vaddhub_PP_sat(Word64 Rss, Word64 Rtt)
```

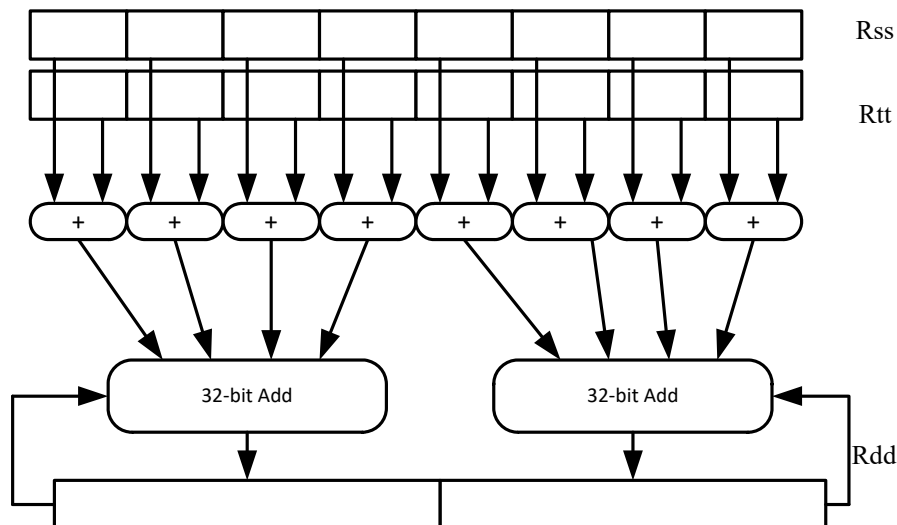
### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				Maj		s5					Parse		t5			Min		d5											
1	1	0	0	0	0	0	1	0	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	d	d	d	d	d	Rd=vaddhub(Rss,Rtt):sat

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Vector reduce add unsigned bytes

For each byte in the source vector *Rss*, add the corresponding byte in the source vector *Rtt*. Add the four upper intermediate results and optionally the upper word of the destination. Add the four lower results and optionally the lower word of the destination.



### Syntax

`Rdd=vraddub(Rss,Rtt)`

```
Rdd = 0;
for (i=0;i<4;i++) {
    Rdd.w[0]=(Rdd.w[0] + (Rss.ub[i]+Rtt.ub[i]));
}
for (i=4;i<8;i++) {
    Rdd.w[1]=(Rdd.w[1] + (Rss.ub[i]+Rtt.ub[i]));
}
```

`Rxx+=vraddub(Rss,Rtt)`

```
for (i = 0; i < 4; i++) {
    Rxx.w[0]=(Rxx.w[0] + (Rss.ub[i]+Rtt.ub[i]));
}
for (i = 4; i < 8; i++) {
    Rxx.w[1]=(Rxx.w[1] + (Rss.ub[i]+Rtt.ub[i]));
}
```

### Behavior

## Class: XTYPE (slots 2,3)

### Intrinsics

`Rdd=vraddub(Rss,Rtt)`

`Word64 Q6_P_vraddub_PP(Word64 Rss, Word64 Rtt)`

`Rxx+=vraddub(Rss,Rtt)`

`Word64 Q6_P_vraddubacc_PP(Word64 Rxx, Word64 Rss, Word64 Rtt)`

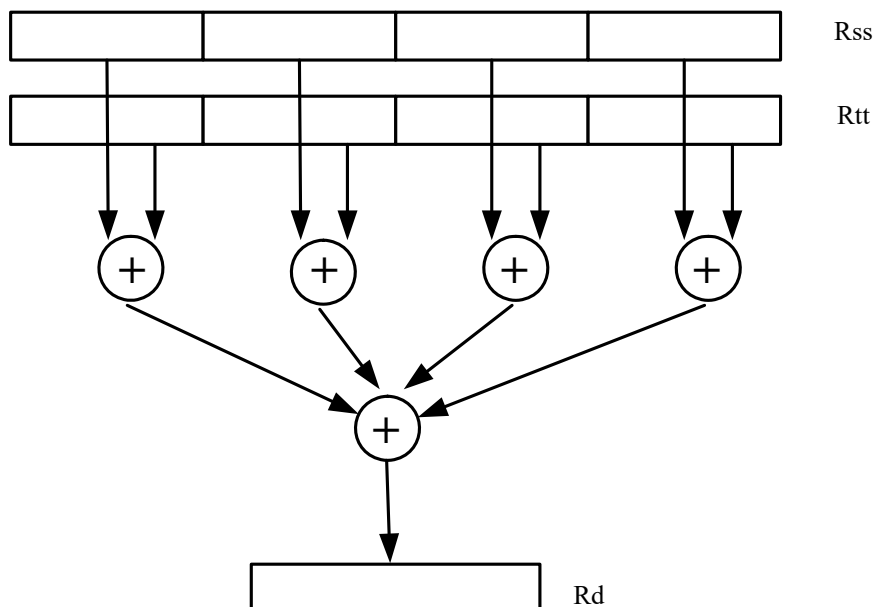
## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp				s5					Parse		t5					MinOp			d5					
1	1	1	0	1	0	0	0	0	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	d	d	d	d	d	Rdd=vraddub(Rss,Rtt)
ICLASS				RegType				MajOp				s5					Parse		t5					MinOp			x5					
1	1	1	0	1	0	1	0	0	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	x	x	x	x	x	Rxx+=vraddub(Rss,Rtt)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Vector reduce add halfwords

For each halfword in the source vector Rss, add the corresponding halfword in the source vector Rtt. Add these intermediate results together, and place the result in Rd.



### Syntax

`Rd=vraddh(Rss,Rtt)`

`Rd=vradduh(Rss,Rtt)`

### Behavior

```
Rd = 0;
for (i=0;i<4;i++) {
    Rd += (Rss.h[i]+Rtt.h[i]);
}
```

```
Rd = 0;
for (i=0;i<4;i++) {
    Rd += (Rss.uh[i]+Rtt.uh[i]);
}
```

### Class: XTYPE (slots 2,3)

### Intrinsics

`Rd=vraddh(Rss,Rtt)`

Word32 Q6\_R\_vraddh\_PP(Word64 Rss, Word64 Rtt)

`Rd=vradduh(Rss,Rtt)`

Word32 Q6\_R\_vradduh\_PP(Word64 Rss, Word64 Rtt)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				MajOp				s5				Parse		t5				MinOp		d5										
1	1	1	0	1	0	0	1	0	-	-	s	s	s	s	s	P	P	0	t	t	t	t	t	-	0	1	d	d	d	d	d	Rd=vradduh(Rss,Rtt)
1	1	1	0	1	0	0	1	0	-	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	1	d	d	d	d	d	Rd=vraddh(Rss,Rtt)



<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector add bytes

Add each of the eight bytes in 64-bit vector *Rss* to the corresponding byte in vector *Rtt*. Optionally, saturate each 8-bit addition to an unsigned value between 0 and 255. The eight results are stored in destination register *Rdd*.

Syntax	Behavior
<code>Rdd=vaddb(Rss,Rtt)</code>	Assembler mapped to: <code>"Rdd=vaddub(Rss,Rtt)"</code>
<code>Rdd=vaddub(Rss,Rtt)[:sat]</code>	<pre>for (i = 0; i &lt; 8; i++) {     Rdd.b[i]=[usat<sub>8</sub>](Rss.ub[i]+Rtt.ub[i]); }</pre>

### Class: XTYPE (slots 2,3)

#### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

#### Intrinsics

<code>Rdd=vaddb(Rss,Rtt)</code>	<code>Word64 Q6_P_vaddb_PP(Word64 Rss, Word64 Rtt)</code>
<code>Rdd=vaddub(Rss,Rtt)</code>	<code>Word64 Q6_P_vaddub_PP(Word64 Rss, Word64 Rtt)</code>
<code>Rdd=vaddub(Rss,Rtt):sat</code> at	<code>Word64 Q6_P_vaddub_PP_sat(Word64 Rss, Word64 Rtt)</code>

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				s5					Parse		t5					MinOp			d5											
1	1	0	1	0	0	1	1	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	d	d	d	d	d	Rdd=vaddub(Rss,Rtt)
1	1	0	1	0	0	1	1	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	d	d	d	d	d	Rdd=vaddub(Rss,Rtt):sat

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector add words

Add each of the two words in 64-bit vector *Rss* to the corresponding word in vector *Rtt*. Optionally, saturate each 32-bit addition to a signed value between 0x80000000 and 0x7fffffff. The two word results are stored in destination register *Rdd*.

### Syntax

```
Rdd=vaddw(Rss,Rtt)[:sat]
```

### Behavior

```
for (i=0;i<2;i++) {
    Rdd.w[i]=[sat32](Rss.w[i]+Rtt.w[i]);
}
```

### Class: XTYPE (slots 2,3)

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

### Intrinsics

```
Rdd=vaddw(Rss,Rtt)
```

```
Word64 Q6_P_vaddw_PP(Word64 Rss, Word64 Rtt)
```

```
Rdd=vaddw(Rss,Rtt):sat
```

```
Word64 Q6_P_vaddw_PP_sat(Word64 Rss, Word64 Rtt)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				s5					Parse		t5					MinOp		d5												
1	1	0	1	0	0	1	1	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	1	d	d	d	d	d	Rdd=vaddw(Rss,Rtt)
1	1	0	1	0	0	1	1	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	0	d	d	d	d	d	Rdd=vaddw(Rss,Rtt):sat

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector average halfwords

Average each of the four halfwords in the 64-bit source vector *Rss* with the corresponding halfword in *Rtt*. The average operation performed on each halfword adds the two halfwords and shifts the result right by one bit. Unsigned average uses a logical right shift (shift in 0), whereas signed average uses an arithmetic right shift (shift in the sign bit). If the round option is used, 0x0001 is also added to each result before shifting. This operation does not overflow. When a summation (before right shift by 1) causes an overflow of 32 bits, the value shifted in is the most-significant carry out.

The signed average and negative average halfwords is available with optional convergent rounding. In convergent rounding, if the two LSBs after the addition/subtraction are 11, a rounding constant of 1 is added, otherwise a 0 is added. This result is then shifted right by one bit. Convergent rounding accumulates less error than arithmetic rounding.

Syntax	Behavior
<code>Rdd=vavgh(Rss,Rtt)</code>	<pre>for (i=0;i&lt;4;i++) {   Rdd.h[i]=(Rss.h[i]+Rtt.h[i])&gt;&gt;1; }</pre>
<code>Rdd=vavgh(Rss,Rtt):crnd</code>	<pre>for (i=0;i&lt;4;i++) {   Rdd.h[i]=convround(Rss.h[i]+Rtt.h[i])&gt;&gt;1; }</pre>
<code>Rdd=vavgh(Rss,Rtt):rnd</code>	<pre>for (i=0;i&lt;4;i++) {   Rdd.h[i]=(Rss.h[i]+Rtt.h[i]+1)&gt;&gt;1; }</pre>
<code>Rdd=vavguh(Rss,Rtt)</code>	<pre>for (i=0;i&lt;4;i++) {   Rdd.h[i]=(Rss.uh[i]+Rtt.uh[i])&gt;&gt;1; }</pre>
<code>Rdd=vavguh(Rss,Rtt):rnd</code>	<pre>for (i=0;i&lt;4;i++) {   Rdd.h[i]=(Rss.uh[i]+Rtt.uh[i]+1)&gt;&gt;1; }</pre>
<code>Rdd=vnavgh(Rtt,Rss)</code>	<pre>for (i=0;i&lt;4;i++) {   Rdd.h[i]=(Rtt.h[i]-Rss.h[i])&gt;&gt;1; }</pre>
<code>Rdd=vnavgh(Rtt,Rss):crnd:sat</code>	<pre>for (i=0;i&lt;4;i++) {   Rdd.h[i]=sat<sub>16</sub>(convround(Rtt.h[i]- Rss.h[i])&gt;&gt;1); }</pre>
<code>Rdd=vnavgh(Rtt,Rss):rnd:sat</code>	<pre>for (i=0;i&lt;4;i++) {   Rdd.h[i]=sat<sub>16</sub>((Rtt.h[i]-Rss.h[i]+1)&gt;&gt;1); }</pre>

### Class: XTYPE (slots 2,3)

#### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

## Intrinsics

Rdd=vavgh(Rss,Rtt)	Word64 Q6_P_vavgh_PP(Word64 Rss, Word64 Rtt)
Rdd=vavgh(Rss,Rtt):crnd	Word64 Q6_P_vavgh_PP_crnd(Word64 Rss, Word64 Rtt)
Rdd=vavgh(Rss,Rtt):rnd	Word64 Q6_P_vavgh_PP_rnd(Word64 Rss, Word64 Rtt)
Rdd=vavguh(Rss,Rtt)	Word64 Q6_P_vavguh_PP(Word64 Rss, Word64 Rtt)
Rdd=vavguh(Rss,Rtt):rnd	Word64 Q6_P_vavguh_PP_rnd(Word64 Rss, Word64 Rtt)
Rdd=vnavgh(Rtt,Rss)	Word64 Q6_P_vnavgh_PP(Word64 Rtt, Word64 Rss)
Rdd=vnavgh(Rtt,Rss):crnd:sat	Word64 Q6_P_vnavgh_PP_crnd_sat(Word64 Rtt, Word64 Rss)
Rdd=vnavgh(Rtt,Rss):rnd:sat	Word64 Q6_P_vnavgh_PP_rnd_sat(Word64 Rtt, Word64 Rss)

## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5					MinOp		d5											
1	1	0	1	0	0	1	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	d	d	d	d	d	Rdd=vavgh(Rss,Rtt)
1	1	0	1	0	0	1	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	d	d	d	d	d	Rdd=vavgh(Rss,Rtt):rnd
1	1	0	1	0	0	1	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	0	d	d	d	d	d	Rdd=vavgh(Rss,Rtt):crnd
1	1	0	1	0	0	1	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	1	d	d	d	d	d	Rdd=vavguh(Rss,Rtt)
1	1	0	1	0	0	1	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	-	d	d	d	d	d	Rdd=vavguh(Rss,Rtt):rnd
1	1	0	1	0	0	1	1	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	d	d	d	d	d	Rdd=vnavgh(Rtt,Rss)
1	1	0	1	0	0	1	1	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	d	d	d	d	d	Rdd=vnavgh(Rtt,Rss):rnd:sat
1	1	0	1	0	0	1	1	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	d	d	d	d	d	Rdd=vnavgh(Rtt,Rss):crnd:sat

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector average unsigned bytes

Average each of the eight unsigned bytes in the 64-bit source vector Rss with the corresponding byte in Rtt. The average operation performed on each byte is the sum of the two bytes shifted right by 1 bit. When the round option is used, 0x01 is also added to each result before shifting. This operation does not overflow. When a summation (before right shift by 1) causes an overflow of 8 bits, the value shifted in is the most-significant carry out.

Syntax	Behavior
<code>Rdd=vavgub(Rss,Rtt)</code>	<pre>for (i = 0; i &lt; 8; i++) {     Rdd.b[i]=((Rss.ub[i] + Rtt.ub[i])&gt;&gt;1); }</pre>
<code>Rdd=vavgub(Rss,Rtt):rnd</code>	<pre>for (i = 0; i &lt; 8; i++) {     Rdd.b[i]=((Rss.ub[i]+Rtt.ub[i]+1)&gt;&gt;1); }</pre>

### Class: XTYPE (slots 2,3)

#### Intrinsics

<code>Rdd=vavgub(Rss,Rtt)</code>	<code>Word64 Q6_P_vavgub_PP(Word64 Rss, Word64 Rtt)</code>
<code>Rdd=vavgub(Rss,Rtt):rnd</code>	<code>Word64 Q6_P_vavgub_PP_rnd(Word64 Rss, Word64 Rtt)</code>

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5					MinOp			d5										
1	1	0	1	0	0	1	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	d	d	d	d	d	Rdd=vavgub(Rss,Rtt)
1	1	0	1	0	0	1	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	d	d	d	d	d	Rdd=vavgub(Rss,Rtt):rnd

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector average words

Average each of the two words in the 64-bit source vector *Rss* with the corresponding word in *Rtt*. The average operation performed on each halfword adds the two words and shifts the result right by 1 bit. Unsigned average uses a logical right shift (shift in 0), whereas signed average uses an arithmetic right shift (shift in the sign bit). When the round option is used, 0x1 is also added to each result before shifting. This operation does not overflow. When a summation (before right shift by 1) causes an overflow of 32 bits, the value shifted in is the most-significant carry out.

The signed average and negative average words are available with optional convergent rounding. In convergent rounding, if the two LSBs after the addition/subtraction are 11, a rounding constant of 1 is added, otherwise a 0 is added. This result is then shifted right by one bit. Convergent rounding accumulates less error than arithmetic rounding.

Syntax	Behavior
<code>Rdd=vavguw(Rss,Rtt)[:rnd]</code>	<pre>for (i=0;i&lt;2;i++) {     Rdd.w[i]=(zxt<sub>32-&gt;33</sub>(Rss.uw[i])+zxt<sub>32-&gt;33</sub>(Rtt.uw[i])+1)&gt;&gt;1; }</pre>
<code>Rdd=vavgw(Rss,Rtt):crnd</code>	<pre>for (i=0;i&lt;2;i++) {     Rdd.w[i]=(convround(sxt<sub>32-&gt;33</sub>(Rss.w[i])+sxt<sub>32-&gt;33</sub>(Rtt.w[i]))&gt;&gt;1); }</pre>
<code>Rdd=vavgw(Rss,Rtt)[:rnd]</code>	<pre>for (i=0;i&lt;2;i++) {     Rdd.w[i]=(sxt<sub>32-&gt;33</sub>(Rss.w[i])+sxt<sub>32-&gt;33</sub>(Rtt.w[i])+1)&gt;&gt;1; }</pre>
<code>Rdd=vnavgw(Rtt,Rss)</code>	<pre>for (i=0;i&lt;2;i++) {     Rdd.w[i]=(sxt<sub>32-&gt;33</sub>(Rtt.w[i])-sxt<sub>32-&gt;33</sub>(Rss.w[i]))&gt;&gt;1; }</pre>
<code>Rdd=vnavgw(Rtt,Rss):crnd:sat</code>	<pre>for (i=0;i&lt;2;i++) {     Rdd.w[i]=sat<sub>32</sub>(convround(sxt<sub>32-&gt;33</sub>(Rtt.w[i])-sxt<sub>32-&gt;33</sub>(Rss.w[i]))&gt;&gt;1); }</pre>
<code>Rdd=vnavgw(Rtt,Rss):rnd:sat</code>	<pre>for (i=0;i&lt;2;i++) {     Rdd.w[i]=sat<sub>32</sub>((sxt<sub>32-&gt;33</sub>(Rtt.w[i])-sxt<sub>32-&gt;33</sub>(Rss.w[i])+1)&gt;&gt;1); }</pre>

### Class: XTYPE (slots 2,3)

#### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

## Intrinsics

Rdd=vavgw(Rss,Rtt)	Word64 Q6_P_vavgw_PP(Word64 Rss, Word64 Rtt)
Rdd=vavgw(Rss,Rtt):rnd	Word64 Q6_P_vavgw_PP_rnd(Word64 Rss, Word64 Rtt)
Rdd=vavgw(Rss,Rtt)	Word64 Q6_P_vavgw_PP(Word64 Rss, Word64 Rtt)
Rdd=vavgw(Rss,Rtt):crnd	Word64 Q6_P_vavgw_PP_crnd(Word64 Rss, Word64 Rtt)
Rdd=vavgw(Rss,Rtt):rnd	Word64 Q6_P_vavgw_PP_rnd(Word64 Rss, Word64 Rtt)
Rdd=vnavgw(Rtt,Rss)	Word64 Q6_P_vnavgw_PP(Word64 Rtt, Word64 Rss)
Rdd=vnavgw(Rtt,Rss):crnd: sat	Word64 Q6_P_vnavgw_PP_crnd_sat(Word64 Rtt, Word64 Rss)
Rdd=vnavgw(Rtt,Rss):rnd: sat	Word64 Q6_P_vnavgw_PP_rnd_sat(Word64 Rtt, Word64 Rss)

## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				s5					Parse		t5					MinOp			d5											
1	1	0	1	0	0	1	1	0	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	d	d	d	d	d	Rdd=vavgw(Rss,Rtt)
1	1	0	1	0	0	1	1	0	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	d	d	d	d	d	Rdd=vavgw(Rss,Rtt):rnd
1	1	0	1	0	0	1	1	0	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	d	d	d	d	d	Rdd=vavgw(Rss,Rtt):crnd
1	1	0	1	0	0	1	1	0	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	d	d	d	d	d	Rdd=vavgw(Rss,Rtt):rnd
1	1	0	1	0	0	1	1	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	d	d	d	d	d	Rdd=vnavgw(Rtt,Rss)
1	1	0	1	0	0	1	1	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	-	d	d	d	d	d	Rdd=vnavgw(Rtt,Rss):rnd: sat
1	1	0	1	0	0	1	1	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	-	d	d	d	d	d	Rdd=vnavgw(Rtt,Rss):crnd: sat

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t



## Vector clip to unsigned

Clip input to an unsigned integer.

### Syntax

```
Rdd=vclip(Rss,#u5)
```

### Behavior

```
tmp=MIN((1<<#u)-1,MAX(Rss.w[0],-(1<<#u)));
Rdd.w[0]=tmp;
tmp=MIN((1<<#u)-1,MAX(Rss.w[1],-(1<<#u)));
Rdd.w[1]=tmp;
```

**Class: XTYPE (slots 2,3)**

### Notes

- This instruction can only execute on a core with the Hexagon audio extensions

### Intrinsics

```
Rdd=vclip(Rss,#u5)
```

```
Word64 Q6_P_vclip_PI(Word64 Rss, Word32
Iu5)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp		s5					Parse		MinOp			d5												
1	0	0	0	1	0	0	0	1	1	0	s	s	s	s	s	P	P	0	i	i	i	i	i	1	1	0	d	d	d	d	d	Rdd=vclip(Rss,#u5)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Vector conditional negate

Based on bits in Rt, conditionally negate halves in Rss.

Syntax	Behavior
<code>Rdd=vcnegh(Rss,Rt)</code>	<pre>for (i = 0; i &lt; 4; i++) {   if (Rt.i) {     Rdd.h[i]=sat<sub>16</sub>(-Rss.h[i]);   } else {     Rdd.h[i]=Rss.h[i];   } }</pre>
<code>Rxx+=vrcnegh(Rss,Rt)</code>	<pre>for (i = 0; i &lt; 4; i++) {   if (Rt.i) {     Rxx += -Rss.h[i];   } else {     Rxx += Rss.h[i];   } }</pre>

**Class: XTYPE (slots 2,3)**

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

### Intrinsics

<code>Rdd=vcnegh(Rss,Rt)</code>	<code>Word64 Q6_P_vcnegh_PR(Word64 Rss, Word32 Rt)</code>
<code>Rxx+=vrcnegh(Rss,Rt)</code>	<code>Word64 Q6_P_vrcneghacc_PR(Word64 Rxx, Word64 Rss, Word32 Rt)</code>

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				Maj		s5					Parse		t5			Min		d5												
1	1	0	0	0	0	1	1	1	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	-	d	d	d	d	d	Rdd=vcnegh(Rss,Rt)
ICLASS		RegType				Maj		s5					Parse		t5			Min		x5												
1	1	0	0	1	0	1	1	0	0	1	s	s	s	s	s	P	P	1	t	t	t	t	t	1	1	1	x	x	x	x	x	Rxx+=vrcnegh(Rss,Rt)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

<b>Field name</b>	<b>Description</b>
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Vector maximum bytes

Compare each of the eight unsigned bytes in the 64-bit source vector Rss to the corresponding byte in Rtt. For each comparison, select the maximum of the two bytes and place that byte in the corresponding location in Rdd.

Syntax	Behavior
Rdd=vmaxb(Rtt,Rss)	<pre>for (i = 0; i &lt; 8; i++) {     Rdd.b[i]=max(Rtt.b[i],Rss.b[i]); }</pre>
Rdd=vmaxub(Rtt,Rss)	<pre>for (i = 0; i &lt; 8; i++) {     Rdd.b[i]=max(Rtt.ub[i],Rss.ub[i]); }</pre>

### Class: XTYPE (slots 2,3)

#### Intrinsics

Rdd=vmaxb(Rtt,Rss)	Word64 Q6_P_vmaxb_PP(Word64 Rtt, Word64 Rss)
Rdd=vmaxub(Rtt,Rss)	Word64 Q6_P_vmaxub_PP(Word64 Rtt, Word64 Rss)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5				MinOp			d5											
1	1	0	1	0	0	1	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	d	d	d	d	d	Rdd=vmaxub(Rtt,Rss)
1	1	0	1	0	0	1	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	0	d	d	d	d	d	Rdd=vmaxb(Rtt,Rss)

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector maximum halfwords

Compare each of the four halfwords in the 64-bit source vector Rss to the corresponding halfword in Rtt. For each comparison, select the maximum of the two halfwords and place that halfword in the corresponding location in Rdd. Comparisons are available in both signed and unsigned form.

Syntax	Behavior
<code>Rdd=vmaxh(Rtt,Rss)</code>	<pre>for (i = 0; i &lt; 4; i++) {     Rdd.h[i]=max(Rtt.h[i],Rss.h[i]); }</pre>
<code>Rdd=vmaxuh(Rtt,Rss)</code>	<pre>for (i = 0; i &lt; 4; i++) {     Rdd.h[i]=max(Rtt.uh[i],Rss.uh[i]); }</pre>

### Class: XTYPE (slots 2,3)

#### Intrinsics

<code>Rdd=vmaxh(Rtt,Rss)</code>	<code>Word64 Q6_P_vmaxh_PP(Word64 Rtt, Word64 Rss)</code>
<code>Rdd=vmaxuh(Rtt,Rss)</code>	<code>Word64 Q6_P_vmaxuh_PP(Word64 Rtt, Word64 Rss)</code>

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5					MinOp			d5										
1	1	0	1	0	0	1	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	d	d	d	d	d	Rdd=vmaxh(Rtt,Rss)
1	1	0	1	0	0	1	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	d	d	d	d	d	Rdd=vmaxuh(Rtt,Rss)

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector reduce maximum halfwords

Register Rxx contains a maximum value in the low word and the address of that maximum value in the high word. Register Rss contains a vector of four halfword values, and register Ru contains the address of this data. The instruction finds the maximum halfword between the previous maximum in Rxx[0] and the four values in Rss. The address of the new maximum is stored in Rxx[1].

Syntax	Behavior
Rxx=vrmaxh(Rss,Ru)	<pre> max = Rxx.h[0]; addr = Rxx.w[1]; for (i = 0; i &lt; 4; i++) {     if (max &lt; Rss.h[i]) {         max = Rss.h[i];         addr = Ru   i&lt;&lt;1;     } } Rxx.w[0]=max; Rxx.w[1]=addr; </pre>
Rxx=vrmaxuh(Rss,Ru)	<pre> max = Rxx.uh[0]; addr = Rxx.w[1]; for (i = 0; i &lt; 4; i++) {     if (max &lt; Rss.uh[i]) {         max = Rss.uh[i];         addr = Ru   i&lt;&lt;1;     } } Rxx.w[0]=max; Rxx.w[1]=addr; </pre>

### Class: XTYPE (slots 2,3)

#### Intrinsics

Rxx=vrmaxh(Rss,Ru)	Word64 Q6_P_vrmaxh_PR(Word64 Rxx, Word64 Rss, Word32 Ru)
Rxx=vrmaxuh(Rss,Ru)	Word64 Q6_P_vrmaxuh_PR(Word64 Rxx, Word64 Rss, Word32 Ru)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType			Maj		s5					Parse		x5				Min		u5												
1	1	0	0	1	0	1	1	0	0	1	s	s	s	s	s	P	P	0	x	x	x	x	x	0	0	1	u	u	u	u	u	Rxx=vrmaxh(Rss,Ru)
1	1	0	0	1	0	1	1	0	0	1	s	s	s	s	s	P	P	1	x	x	x	x	x	0	0	1	u	u	u	u	u	Rxx=vrmaxuh(Rss,Ru)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s
u5	Field to encode register u
x5	Field to encode register x

<b>Field name</b>	<b>Description</b>
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Vector reduce maximum words

Find the maximum word between the previous maximum in Rxx[0] and the two values in Rss. The address of the new maximum is stored in Rxx[1].

Register Rxx contains a maximum value in the low word and the address of that maximum value in the high word. Register Rss contains a vector of two word values, and register Ru contains the address of this data.

Syntax	Behavior
<code>Rxx=vrmaxuw(Rss,Ru)</code>	<pre> max = Rxx.uw[0]; addr = Rxx.w[1]; for (i = 0; i &lt; 2; i++) {     if (max &lt; Rss.uw[i]) {         max = Rss.uw[i];         addr = Ru   i&lt;&lt;2;     } } Rxx.w[0]=max; Rxx.w[1]=addr; </pre>
<code>Rxx=vrmaxw(Rss,Ru)</code>	<pre> max = Rxx.w[0]; addr = Rxx.w[1]; for (i = 0; i &lt; 2; i++) {     if (max &lt; Rss.w[i]) {         max = Rss.w[i];         addr = Ru   i&lt;&lt;2;     } } Rxx.w[0]=max; Rxx.w[1]=addr; </pre>

### Class: XTYPE (slots 2,3)

#### Intrinsics

`Rxx=vrmaxuw(Rss,Ru)` Word64 Q6\_P\_vrmaxuw\_PR(Word64 Rxx, Word64 Rss, Word32 Ru)

`Rxx=vrmaxw(Rss,Ru)` Word64 Q6\_P\_vrmaxw\_PR(Word64 Rxx, Word64 Rss, Word32 Ru)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType			Maj		s5					Parse		x5			Min		u5													
1	1	0	0	1	0	1	1	0	0	1	s	s	s	s	s	P	P	0	x	x	x	x	x	0	1	0	u	u	u	u	u	Rxx=vrmaxw(Rss,Ru)
1	1	0	0	1	0	1	1	0	0	1	s	s	s	s	s	P	P	1	x	x	x	x	x	0	1	0	u	u	u	u	u	Rxx=vrmaxuw(Rss,Ru)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s
u5	Field to encode register u



<b>Field name</b>	<b>Description</b>
x5	Field to encode register x
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Vector maximum words

Compare each of the two words in the 64-bit source vector Rss to the corresponding word in Rtt. For each comparison, select the maximum of the two words and place that word in the corresponding location in Rdd.

Comparisons are available in both signed and unsigned form.

Syntax	Behavior
<code>Rdd=vmaxuw(Rtt,Rss)</code>	<pre>for (i = 0; i &lt; 2; i++) {     Rdd.w[i]=max(Rtt.uw[i],Rss.uw[i]); }</pre>
<code>Rdd=vmaxw(Rtt,Rss)</code>	<pre>for (i = 0; i &lt; 2; i++) {     Rdd.w[i]=max(Rtt.w[i],Rss.w[i]); }</pre>

### Class: XTYPE (slots 2,3)

#### Intrinsics

`Rdd=vmaxuw(Rtt,Rss)`

Word64 Q6\_P\_vmaxuw\_PP(Word64 Rtt, Word64 Rss)

`Rdd=vmaxw(Rtt,Rss)`

Word64 Q6\_P\_vmaxw\_PP(Word64 Rtt, Word64 Rss)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5					MinOp			d5										
1	1	0	1	0	0	1	1	1	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	1	d	d	d	d	d	Rdd=vmaxuw(Rtt,Rss)
1	1	0	1	0	0	1	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	d	d	d	d	d	Rdd=vmaxw(Rtt,Rss)

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector minimum bytes

Compare each of the eight unsigned bytes in the 64-bit source vector *Rss* to the corresponding byte in *Rtt*. For each comparison, select the minimum of the two bytes and place that byte in the corresponding location in *Rdd*.

Syntax	Behavior
<code>Rdd, Pe=vminub(Rtt, Rss)</code>	<pre>for (i = 0; i &lt; 8; i++) {     Pe.i = (Rtt.ub[i] &gt; Rss.ub[i]);     Rdd.b[i]=min(Rtt.ub[i],Rss.ub[i]); }</pre>
<code>Rdd=vminb(Rtt, Rss)</code>	<pre>for (i = 0; i &lt; 8; i++) {     Rdd.b[i]=min(Rtt.b[i],Rss.b[i]); }</pre>
<code>Rdd=vminub(Rtt, Rss)</code>	<pre>for (i = 0; i &lt; 8; i++) {     Rdd.b[i]=min(Rtt.ub[i],Rss.ub[i]); }</pre>

### Class: XTYPE (slots 2,3)

#### Notes

- The predicate generated by this instruction cannot be used as a .new predicate, nor can it be automatically ANDed with another predicate.

#### Intrinsics

`Rdd=vminb(Rtt, Rss)`

Word64 Q6\_P\_vminb\_PP(Word64 Rtt, Word64 Rss)

`Rdd=vminub(Rtt, Rss)`

Word64 Q6\_P\_vminub\_PP(Word64 Rtt, Word64 Rss)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5				MinOp			d5											
1	1	0	1	0	0	1	1	1	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	d	d	d	d	d	Rdd=vminub(Rtt,Rss)
1	1	0	1	0	0	1	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	1	d	d	d	d	d	Rdd=vminb(Rtt,Rss)
ICLASS			RegType				MajOp			s5					Parse		t5				e2		d5									
1	1	1	0	1	0	1	0	1	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	e	e	d	d	d	d	d	Rdd,Pe=vminub(Rtt,Rss)

Field name	Description
RegType	Register type
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
Parse	Packet/loop parse bits
d5	Field to encode register d

<b>Field name</b>	<b>Description</b>
e2	Field to encode register e
s5	Field to encode register s
t5	Field to encode register t

## Vector minimum halfwords

Compare each of the four halfwords in the 64-bit source vector Rss to the corresponding halfword in Rtt. For each comparison, select the minimum of the two halfwords and place that halfword in the corresponding location in Rdd.

Comparisons are available in both signed and unsigned form.

Syntax	Behavior
<code>Rdd=vminh(Rtt,Rss)</code>	<pre>for (i = 0; i &lt; 4; i++) {     Rdd.h[i]=min(Rtt.h[i],Rss.h[i]); }</pre>
<code>Rdd=vminuh(Rtt,Rss)</code>	<pre>for (i = 0; i &lt; 4; i++) {     Rdd.h[i]=min(Rtt.uh[i],Rss.uh[i]); }</pre>

### Class: XTYPE (slots 2,3)

#### Intrinsics

`Rdd=vminh(Rtt,Rss)`

Word64 Q6\_P\_vminh\_PP(Word64 Rtt, Word64 Rss)

`Rdd=vminuh(Rtt,Rss)`

Word64 Q6\_P\_vminuh\_PP(Word64 Rtt, Word64 Rss)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5					MinOp			d5										
1	1	0	1	0	0	1	1	1	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	d	d	d	d	d	Rdd=vminh(Rtt,Rss)
1	1	0	1	0	0	1	1	1	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	d	d	d	d	d	Rdd=vminuh(Rtt,Rss)

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector reduce minimum halfwords

Find the minimum halfword between the previous minimum in Rxx[0] and the four values in Rss. The address of the new minimum is stored in Rxx[1].

Register Rxx contains a minimum value in the low word and the address of that minimum value in the high word. Register Rss contains a vector of four halfword values, and register Ru contains the address of this data.

Syntax	Behavior
<code>Rxx=vrminh(Rss,Ru)</code>	<pre> min = Rxx.h[0]; addr = Rxx.w[1]; for (i = 0; i &lt; 4; i++) {     if (min &gt; Rss.h[i]) {         min = Rss.h[i];         addr = Ru   i&lt;&lt;1;     } } Rxx.w[0]=min; Rxx.w[1]=addr; </pre>
<code>Rxx=vrminuh(Rss,Ru)</code>	<pre> min = Rxx.uh[0]; addr = Rxx.w[1]; for (i = 0; i &lt; 4; i++) {     if (min &gt; Rss.uh[i]) {         min = Rss.uh[i];         addr = Ru   i&lt;&lt;1;     } } Rxx.w[0]=min; Rxx.w[1]=addr; </pre>

### Class: XTYPE (slots 2,3)

#### Intrinsics

<code>Rxx=vrminh(Rss,Ru)</code>	Word64 Q6_P_vrminh_PR(Word64 Rxx, Word64 Rss, Word32 Ru)
<code>Rxx=vrminuh(Rss,Ru)</code>	Word64 Q6_P_vrminuh_PR(Word64 Rxx, Word64 Rss, Word32 Ru)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType		Maj		s5					Parse		x5			Min		u5														
1	1	0	0	1	0	1	1	0	0	1	s	s	s	s	s	P	P	0	x	x	x	x	x	1	0	1	u	u	u	u	u	Rxx=vrminh(Rss,Ru)
1	1	0	0	1	0	1	1	0	0	1	s	s	s	s	s	P	P	1	x	x	x	x	x	1	0	1	u	u	u	u	u	Rxx=vrminuh(Rss,Ru)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s
u5	Field to encode register u

<b>Field name</b>	<b>Description</b>
x5	Field to encode register x
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Vector reduce minimum words

Find the minimum word between the previous minimum in Rxx[0] and the two values in Rss. The address of the new minimum is stored in Rxx[1].

Register Rxx contains a minimum value in the low word and the address of that minimum value in the high word. Register Rss contains a vector of two word values, and register Ru contains the address of this data.

Syntax	Behavior
<code>Rxx=vrminuw(Rss,Ru)</code>	<pre> min = Rxx.uw[0]; addr = Rxx.w[1]; for (i = 0; i &lt; 2; i++) {     if (min &gt; Rss.uw[i]) {         min = Rss.uw[i];         addr = Ru   i&lt;&lt;2;     } } Rxx.w[0]=min; Rxx.w[1]=addr; </pre>
<code>Rxx=vrminw(Rss,Ru)</code>	<pre> min = Rxx.w[0]; addr = Rxx.w[1]; for (i = 0; i &lt; 2; i++) {     if (min &gt; Rss.w[i]) {         min = Rss.w[i];         addr = Ru   i&lt;&lt;2;     } } Rxx.w[0]=min; Rxx.w[1]=addr; </pre>

### Class: XTYPE (slots 2,3)

#### Intrinsics

`Rxx=vrminuw(Rss,Ru)` Word64 Q6\_P\_vrminuw\_PR(Word64 Rxx, Word64 Rss, Word32 Ru)

`Rxx=vrminw(Rss,Ru)` Word64 Q6\_P\_vrminw\_PR(Word64 Rxx, Word64 Rss, Word32 Ru)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType			Maj		s5					Parse		x5			Min		u5													
1	1	0	0	1	0	1	1	0	0	1	s	s	s	s	s	P	P	0	x	x	x	x	x	1	1	0	u	u	u	u	u	Rxx=vrminw(Rss,Ru)
1	1	0	0	1	0	1	1	0	0	1	s	s	s	s	s	P	P	1	x	x	x	x	x	1	1	0	u	u	u	u	u	Rxx=vrminuw(Rss,Ru)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s
u5	Field to encode register u



<b>Field name</b>	<b>Description</b>
x5	Field to encode register x
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Vector minimum words

Compare each of the two words in the 64-bit source vector Rss to the corresponding word in Rtt. For each comparison, select the minimum of the two words and place that word in the corresponding location in Rdd.

Comparisons are available in both signed and unsigned form.

Syntax	Behavior
<code>Rdd=vminuw(Rtt,Rss)</code>	<pre>for (i = 0; i &lt; 2; i++) {     Rdd.w[i]=min(Rtt.uw[i],Rss.uw[i]); }</pre>
<code>Rdd=vminw(Rtt,Rss)</code>	<pre>for (i = 0; i &lt; 2; i++) {     Rdd.w[i]=min(Rtt.w[i],Rss.w[i]); }</pre>

### Class: XTYPE (slots 2,3)

#### Intrinsics

`Rdd=vminuw(Rtt,Rss)`

Word64 Q6\_P\_vminuw\_PP(Word64 Rtt, Word64 Rss)

`Rdd=vminw(Rtt,Rss)`

Word64 Q6\_P\_vminw\_PP(Word64 Rtt, Word64 Rss)

#### Encoding

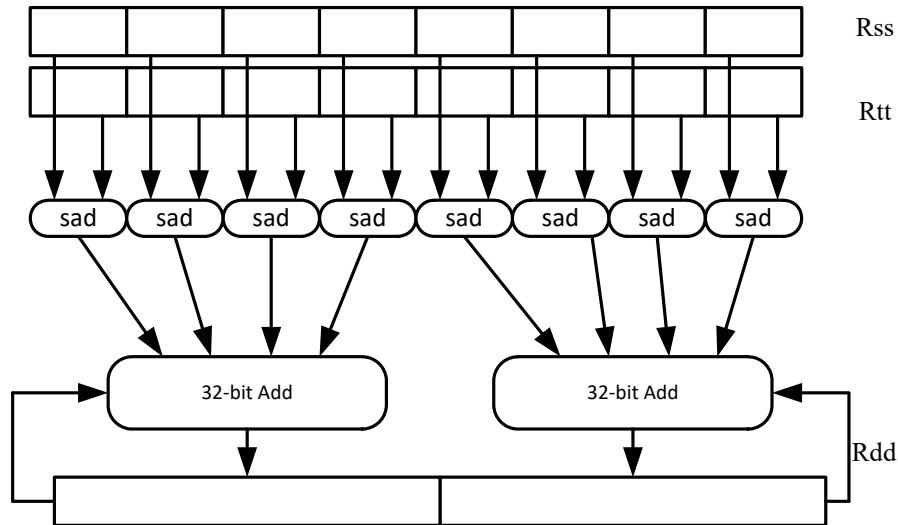
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5					MinOp		d5											
1	1	0	1	0	0	1	1	1	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	d	d	d	d	d	Rdd=vminw(Rtt,Rss)
1	1	0	1	0	0	1	1	1	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	0	d	d	d	d	d	Rdd=vminuw(Rtt,Rss)

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector sum of absolute differences unsigned bytes

For each byte in the source vector  $R_{ss}$ , subtract the corresponding byte in source vector  $R_{tt}$ . Take the absolute value of the intermediate results, and the upper four together and add the lower four together. Optionally, add the destination upper and lower words to these results.

This instruction is useful in determining distance between two vectors, in applications such as motion estimation.



### Syntax

```
Rdd=vrsadub(Rss,Rtt)
```

```
Rxx+=vrsadub(Rss,Rtt)
```

### Behavior

```
Rdd = 0;
for (i = 0; i < 4; i++) {
    Rdd.w[0] = (Rdd.w[0] + ABS((Rss.ub[i] -
    Rtt.ub[i])));
}
for (i = 4; i < 8; i++) {
    Rdd.w[1] = (Rdd.w[1] + ABS((Rss.ub[i] -
    Rtt.ub[i])));
}
```

```
for (i = 0; i < 4; i++) {
    Rxx.w[0] = (Rxx.w[0] + ABS((Rss.ub[i] -
    Rtt.ub[i])));
}
for (i = 4; i < 8; i++) {
    Rxx.w[1] = (Rxx.w[1] + ABS((Rss.ub[i] -
    Rtt.ub[i])));
}
```

**Class: XTYPE (slots 2,3)****Intrinsics**

Rdd=vrsadub(Rss,Rtt) Word64 Q6\_P\_vrsadub\_PP(Word64 Rss, Word64 Rtt)

Rxx+=vrsadub(Rss,Rtt) Word64 Q6\_P\_vrsadubacc\_PP(Word64 Rxx, Word64 Rss, Word64 Rtt)

**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse		t5					MinOp			d5								
1	1	1	0	1	0	0	0	0	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	0	d	d	d	d	d	Rdd=vrsadub(Rss,Rtt)
ICLASS			RegType				MajOp		s5					Parse		t5					MinOp			x5								
1	1	1	0	1	0	1	0	0	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	0	x	x	x	x	x	Rxx+=vrsadub(Rss,Rtt)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Vector subtract halfwords

Subtract each of the four halfwords in 64-bit vector *Rss* from the corresponding halfword in vector *Rtt*.

Optionally, saturate each 16-bit addition to either a signed or unsigned 16-bit value. Applying saturation to the *vsubh* instruction clamps the result to the signed range 0x8000 to 0x7fff, whereas applying saturation to the *vsubuh* instruction ensures that the unsigned result falls within the range 0 to 0xffff.

When saturation is not needed, use the *vsubh* instruction.

Syntax	Behavior
<code>Rdd=vsubh(Rtt,Rss)[:sat]</code>	<pre>for (i=0;i&lt;4;i++) {   Rdd.h[i]=[sat<sub>16</sub>](Rtt.h[i]-Rss.h[i]); }</pre>
<code>Rdd=vsubuh(Rtt,Rss):sat</code>	<pre>for (i=0;i&lt;4;i++) {   Rdd.h[i]=usat<sub>16</sub>(Rtt.uh[i]-Rss.uh[i]); }</pre>

### Class: XTYPE (slots 2,3)

#### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

#### Intrinsics

<code>Rdd=vsubh(Rtt,Rss)</code>	<code>Word64 Q6_P_vsubh_PP(Word64 Rtt, Word64 Rss)</code>
<code>Rdd=vsubh(Rtt,Rss):sat</code>	<code>Word64 Q6_P_vsubh_PP_sat(Word64 Rtt, Word64 Rss)</code>
<code>Rdd=vsubuh(Rtt,Rss):sat</code>	<code>Word64 Q6_P_vsubuh_PP_sat(Word64 Rtt, Word64 Rss)</code>

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5					MinOp		d5											
1	1	0	1	0	0	1	1	0	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	d	d	d	d	d	Rdd=vsubh(Rtt,Rss)
1	1	0	1	0	0	1	1	0	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	d	d	d	d	d	Rdd=vsubh(Rtt,Rss):sat
1	1	0	1	0	0	1	1	0	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	0	d	d	d	d	d	Rdd=vsubuh(Rtt,Rss):sat

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits

<b>Field name</b>	<b>Description</b>
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector subtract bytes

Subtract each of the eight bytes in 64-bit vector *Rss* from the corresponding byte in vector *Rtt*.

Optionally, saturate each 8-bit subtraction to an unsigned value between 0 and 255. The eight results are stored in destination register *Rdd*.

Syntax	Behavior
<code>Rdd=vsubb(Rss,Rtt)</code>	Assembler mapped to: <code>"Rdd=vsubub(Rss,Rtt)"</code>
<code>Rdd=vsubub(Rtt,Rss)[:sat]</code>	<pre>for (i = 0; i &lt; 8; i++) {     Rdd.b[i]=[usat<sub>8</sub>] (Rtt.ub[i]-Rss.ub[i]); }</pre>

### Class: XTYPE (slots 2,3)

#### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

#### Intrinsics

<code>Rdd=vsubb(Rss,Rtt)</code>	<code>Word64 Q6_P_vsubb_PP(Word64 Rss, Word64 Rtt)</code>
<code>Rdd=vsubub(Rtt,Rss)</code>	<code>Word64 Q6_P_vsubub_PP(Word64 Rtt, Word64 Rss)</code>
<code>Rdd=vsubub(Rtt,Rss):sat</code>	<code>Word64 Q6_P_vsubub_PP_sat(Word64 Rtt, Word64 Rss)</code>

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5					MinOp			d5										
1	1	0	1	0	0	1	1	0	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	d	d	d	d	d	Rdd=vsubub(Rtt,Rss)
1	1	0	1	0	0	1	1	0	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	d	d	d	d	d	Rdd=vsubub(Rtt,Rss):sat

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector subtract words

Subtract each of the two words in 64-bit vector *Rss* from the corresponding word in vector *Rtt*.

Optionally, saturate each 32-bit subtraction to a signed value between 0x8000\_0000 and 0x7fff\_ffff. The two word results are stored in destination register *Rdd*.

### Syntax

```
Rdd=vsubw(Rtt,Rss)[:sat]
```

### Behavior

```
for (i=0;i<2;i++) {
    Rdd.w[i]=[sat32](Rtt.w[i]-Rss.w[i]);
}
```

**Class: XTYPE (slots 2,3)**

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

### Intrinsics

```
Rdd=vsubw(Rtt,Rss)
```

```
Word64 Q6_P_vsubw_PP(Word64 Rtt, Word64 Rss)
```

```
Rdd=vsubw(Rtt,Rss):sat
```

```
Word64 Q6_P_vsubw_PP_sat(Word64 Rtt, Word64 Rss)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType									s5					Parse		t5					MinOp			d5						
1	1	0	1	0	0	1	1	0	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	1	d	d	d	d	d	Rdd=vsubw(Rtt,Rss)
1	1	0	1	0	0	1	1	0	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	0	d	d	d	d	d	Rdd=vsubw(Rtt,Rss):sat

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t



## 11.10.2 XTYPE BIT

The XTYPE BIT instruction subclass includes instructions for bit manipulation.

### Count leading

Count leading zeros (cl0) counts the number of consecutive zeros starting with the most significant bit.

Count leading ones (cl1) counts the number of consecutive ones starting with the most significant bit.

Count leading bits (clb) counts both leading ones and leading zeros and then selects the maximum.

The normamt instruction returns the number of leading bits minus one.

For a two's-complement number, the number of leading zeros is zero for negative numbers. The number of leading ones is zero for positive numbers.

The number of leading bits can be used to judge the magnitude of the value.

Syntax	Behavior
Rd=add(clb(Rs), #s6)	Rd = (max(count_leading_ones(Rs), count_leading_ones(~Rs)))+#s;
Rd=add(clb(Rss), #s6)	Rd = (max(count_leading_ones(Rss), count_leading_ones(~Rss)))+#s;
Rd=cl0(Rs)	Rd = count_leading_ones(~Rs);
Rd=cl0(Rss)	Rd = count_leading_ones(~Rss);
Rd=cl1(Rs)	Rd = count_leading_ones(Rs);
Rd=cl1(Rss)	Rd = count_leading_ones(Rss);
Rd=clb(Rs)	Rd = max(count_leading_ones(Rs), count_leading_ones(~Rs));
Rd=clb(Rss)	Rd = max(count_leading_ones(Rss), count_leading_ones(~Rss));
Rd=normamt(Rs)	if (Rs == 0) { Rd = 0; } else { Rd = (max(count_leading_ones(Rs), count_leading_ones(~Rs)))-1; }
Rd=normamt(Rss)	if (Rss == 0) { Rd = 0; } else { Rd = (max(count_leading_ones(Rss), count_leading_ones(~Rss)))-1; }

**Class: XTYPE (slots 2,3)****Intrinsics**

Rd=add(clb(Rs),#s6)	Word32 Q6_R_add_clb_RI(Word32 Rs, Word32 Is6)
Rd=add(clb(Rss),#s6)	Word32 Q6_R_add_clb_PI(Word64 Rss, Word32 Is6)
Rd=cl0(Rs)	Word32 Q6_R_cl0_R(Word32 Rs)
Rd=cl0(Rss)	Word32 Q6_R_cl0_P(Word64 Rss)
Rd=cl1(Rs)	Word32 Q6_R_cl1_R(Word32 Rs)
Rd=cl1(Rss)	Word32 Q6_R_cl1_P(Word64 Rss)
Rd=clb(Rs)	Word32 Q6_R_clb_R(Word32 Rs)
Rd=clb(Rss)	Word32 Q6_R_clb_P(Word64 Rss)
Rd=normamt(Rs)	Word32 Q6_R_normamt_R(Word32 Rs)
Rd=normamt(Rss)	Word32 Q6_R_normamt_P(Word64 Rss)

**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse				MinOp			d5											
1	0	0	0	1	0	0	0	0	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	0	d	d	d	d	d	Rd=clb(Rss)
1	0	0	0	1	0	0	0	0	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	0	1	0	d	d	d	d	d	Rd=cl0(Rss)
1	0	0	0	1	0	0	0	0	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	0	d	d	d	d	d	Rd=cl1(Rss)
1	0	0	0	1	0	0	0	0	1	1	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	0	d	d	d	d	d	Rd=normamt(Rss)
1	0	0	0	1	0	0	0	0	1	1	s	s	s	s	s	P	P	i	i	i	i	i	i	0	1	0	d	d	d	d	d	Rd=add(clb(Rss),#s6)
1	0	0	0	1	1	0	0	0	0	1	s	s	s	s	s	P	P	i	i	i	i	i	i	0	0	0	d	d	d	d	d	Rd=add(clb(Rs),#s6)
1	0	0	0	1	1	0	0	0	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	0	d	d	d	d	d	Rd=clb(Rs)
1	0	0	0	1	1	0	0	0	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	1	d	d	d	d	d	Rd=cl0(Rs)
1	0	0	0	1	1	0	0	0	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	1	0	d	d	d	d	d	Rd=cl1(Rs)
1	0	0	0	1	1	0	0	0	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	1	1	d	d	d	d	d	Rd=normamt(Rs)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Count population

The population count (popcount) instruction counts the number set bits in Rss.

### Syntax

```
Rd=popcount (Rss)
```

### Behavior

```
Rd = count_ones (Rss);
```

**Class: XTYPE (slots 2,3)**

### Intrinsics

```
Rd=popcount (Rss)
```

```
Word32 Q6_R_popcount_P (Word64 Rss)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp				s5					Parse		MinOp				d5									
1	0	0	0	1	0	0	0	0	1	1	s	s	s	s	s	P	P	-	-	-	-	-	-	0	1	1	d	d	d	d	d	Rd=popcount(Rss)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Count trailing

Count trailing zeros (ct0) counts the number of consecutive zeros starting with the least significant bit.

Count trailing ones (ct1) counts the number of consecutive ones starting with the least significant bit.

Syntax	Behavior
Rd=ct0(Rs)	Rd = count_leading_ones(~reverse_bits(Rs));
Rd=ct0(Rss)	Rd = count_leading_ones(~reverse_bits(Rss));
Rd=ct1(Rs)	Rd = count_leading_ones(reverse_bits(Rs));
Rd=ct1(Rss)	Rd = count_leading_ones(reverse_bits(Rss));

### Class: XTYPE (slots 2,3)

#### Intrinsics

Rd=ct0(Rs)	Word32 Q6_R_ct0_R(Word32 Rs)
Rd=ct0(Rss)	Word32 Q6_R_ct0_P(Word64 Rss)
Rd=ct1(Rs)	Word32 Q6_R_ct1_R(Word32 Rs)
Rd=ct1(Rss)	Word32 Q6_R_ct1_P(Word64 Rss)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse				MinOp			d5										
1	0	0	0	1	0	0	0	1	1	1	s	s	s	s	s	P	P	-	-	-	-	-	-	0	1	0	d	d	d	d	d	Rd=ct0(Rss)
1	0	0	0	1	0	0	0	1	1	1	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	0	d	d	d	d	d	Rd=ct1(Rss)
1	0	0	0	1	1	0	0	0	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	0	d	d	d	d	d	Rd=ct0(Rs)
1	0	0	0	1	1	0	0	0	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	1	d	d	d	d	d	Rd=ct1(Rs)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

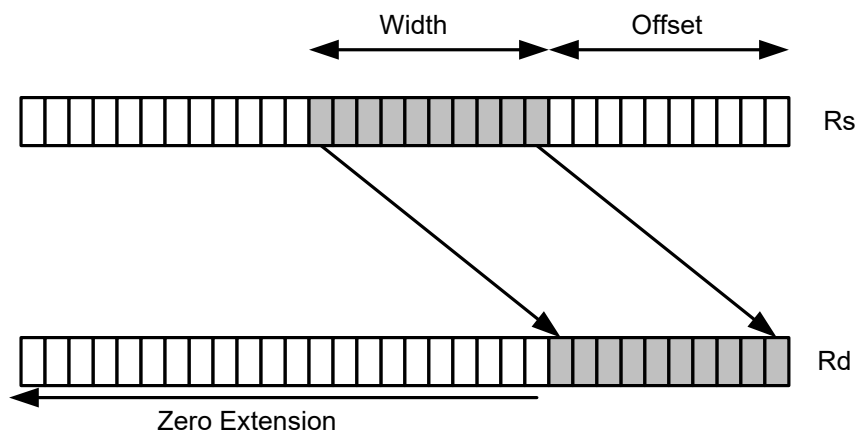
## Extract bit field

Extract a bit field from the source register (or register pair) and deposit into the least significant bits of the destination register (or register pair). The other, more significant bits in the destination are either cleared or sign-extended, depending on the instruction.

The width of the extracted field is obtained from the first immediate or from the most-significant word of Rtt. The field offset is obtained from either the second immediate or from the least-significant word of Rtt.

For register-based extract, where Rtt supplies the offset and width, the offset value is treated as a signed 7-bit number. If this value is negative, the source register Rss is shifted left (the reverse direction). Width number of bits are then taken from the least-significant portion of this result.

When the shift amount and/or offset captures data beyond the most significant end of the input, these bits are taken as zero.



Syntax	Behavior
<code>Rd=extract (Rs, #u5, #U5)</code>	<pre>width=#u; offset=#U; Rd = sxt_width-&gt;32 ((Rs &gt;&gt; offset));</pre>
<code>Rd=extract (Rs, Rtt)</code>	<pre>width=zxt_6-&gt;32 ((Rtt.w[1])); offset=sxt_7-&gt;32 ((Rtt.w[0])); Rd = sxt_width-&gt;64 ((offset&gt;0) ? (zxt_32-&gt;64 (zxt_32-&gt;64 (Rs)))&gt;&gt;&gt; offset) : (zxt_32-&gt;64 (zxt_32-&gt;64 (Rs))&lt;&lt;offset));</pre>
<code>Rd=extractu (Rs, #u5, #U5)</code>	<pre>width=#u; offset=#U; Rd = zxt_width-&gt;32 ((Rs &gt;&gt; offset));</pre>
<code>Rd=extractu (Rs, Rtt)</code>	<pre>width=zxt_6-&gt;32 ((Rtt.w[1])); offset=sxt_7-&gt;32 ((Rtt.w[0])); Rd = zxt_width-&gt;64 ((offset&gt;0) ? (zxt_32-&gt;64 (zxt_32-&gt;64 (Rs)))&gt;&gt;&gt; offset) : (zxt_32-&gt;64 (zxt_32-&gt;64 (Rs))&lt;&lt;offset));</pre>
<code>Rdd=extract (Rss, #u6, #U6)</code>	<pre>width=#u; offset=#U; Rdd = sxt_width-&gt;64 ((Rss &gt;&gt; offset));</pre>

Syntax	Behavior
Rdd=extract (Rss, Rtt)	width=zxt <sub>6-&gt;32</sub> ((Rtt.w[1])); offset=sxt <sub>7-&gt;32</sub> ((Rtt.w[0])); Rdd = sxt <sub>width-&gt;64</sub> ((offset>0)?(Rss>>>offset):(Rss<<<offset));
Rdd=extractu (Rss, #u6, #U6)	width=#u; offset=#U; Rdd = zxt <sub>width-&gt;64</sub> ((Rss >> offset));
Rdd=extractu (Rss, Rtt)	width=zxt <sub>6-&gt;32</sub> ((Rtt.w[1])); offset=sxt <sub>7-&gt;32</sub> ((Rtt.w[0])); Rdd = zxt <sub>width-&gt;64</sub> ((offset>0)?(Rss>>>offset):(Rss<<<offset));

**Class: XTYPE (slots 2,3)**

**Intrinsics**

Rd=extract (Rs, #u5, #U5)	Word32 Q6_R_extract_RII(Word32 Rs, Word32 Iu5, Word32 IU5)
Rd=extract (Rs, Rtt)	Word32 Q6_R_extract_RP(Word32 Rs, Word64 Rtt)
Rd=extractu (Rs, #u5, #U5)	Word32 Q6_R_extractu_RII(Word32 Rs, Word32 Iu5, Word32 IU5)
Rd=extractu (Rs, Rtt)	Word32 Q6_R_extractu_RP(Word32 Rs, Word64 Rtt)
Rdd=extract (Rss, #u6, #U6)	Word64 Q6_P_extract_PII(Word64 Rss, Word32 Iu6, Word32 IU6)
Rdd=extract (Rss, Rtt)	Word64 Q6_P_extract_PP(Word64 Rss, Word64 Rtt)
Rdd=extractu (Rss, #u6, #U6)	Word64 Q6_P_extractu_PII(Word64 Rss, Word32 Iu6, Word32 IU6)
Rdd=extractu (Rss, Rtt)	Word64 Q6_P_extractu_PP(Word64 Rss, Word64 Rtt)

**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse			MinOp			d5											
1	0	0	0	0	0	0	1	I	I	I	s	s	s	s	s	P	P	i	i	i	i	i	i	I	I	I	d	d	d	d	d	Rdd=extractu(Rss,#u6,#U6)
1	0	0	0	1	0	1	0	I	I	I	s	s	s	s	s	P	P	i	i	i	i	i	i	I	I	I	d	d	d	d	d	Rdd=extract(Rss,#u6,#U6)
1	0	0	0	1	1	0	1	0	I	I	s	s	s	s	s	P	P	0	i	i	i	i	i	I	I	I	d	d	d	d	d	Rd=extractu(Rs,#u5,#U5)
1	0	0	0	1	1	0	1	1	I	I	s	s	s	s	s	P	P	0	i	i	i	i	i	I	I	I	d	d	d	d	d	Rd=extract(Rs,#u5,#U5)
ICLASS			RegType				Maj			s5					Parse			t5			Min			d5								
1	1	0	0	0	0	0	1	0	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	d	d	d	d	d	Rdd=extractu(Rss,Rtt)
1	1	0	0	0	0	0	1	1	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	-	d	d	d	d	d	Rdd=extract(Rss,Rtt)
1	1	0	0	1	0	0	1	0	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	d	d	d	d	d	Rd=extractu(Rs,Rtt)
1	1	0	0	1	0	0	1	0	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	-	d	d	d	d	d	Rd=extract(Rs,Rtt)

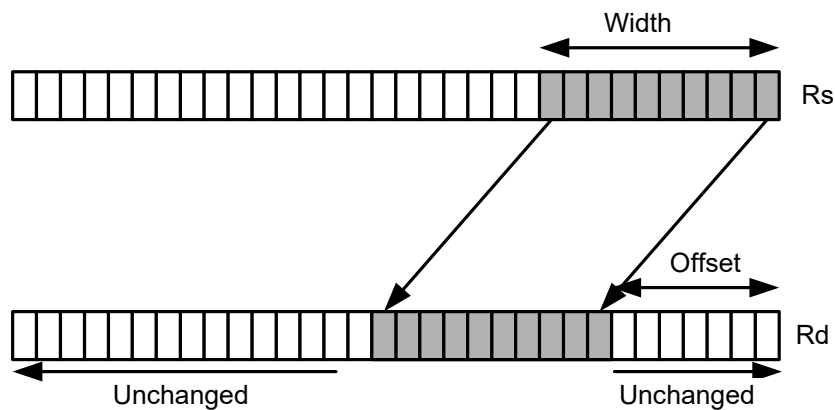
<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
MajOp	Major opcode
MinOp	Minor opcode
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Insert bit field

Replace a bit field in the destination register (or register pair) with bits from the least significant portion of Rs/Rss. The number of bits is obtained from the first immediate or the most-significant word of Rtt. The bits are shifted by the second immediate or the least significant word of Rtt.

When register Rtt specifies the offset, the low 7-bits of Rtt are treated as a signed 7-bit value. If this value is negative, the result is zero.

Shift amounts and offsets that are too large may push bits beyond the end of the destination register., and the bits do not appear in the destination register.



Syntax	Behavior
<code>Rx=insert (Rs, #u5, #U5)</code>	<pre>width=#u; offset=#U; Rx &amp;= ~((1&lt;&lt;width)-1)&lt;&lt;offset); Rx  = ((Rs &amp; ((1&lt;&lt;width)-1)) &lt;&lt; offset);</pre>
<code>Rx=insert (Rs, Rtt)</code>	<pre>width=zxt<sub>6-&gt;32</sub>((Rtt.w[1])); offset=sxt<sub>7-&gt;32</sub>((Rtt.w[0])); mask = ((1&lt;&lt;width)-1); if (offset &lt; 0) {     Rx = 0; } else {     Rx &amp;= ~(mask&lt;&lt;offset);     Rx  = ((Rs &amp; mask) &lt;&lt; offset); }</pre>
<code>Rxx=insert (Rss, #u6, #U6 )</code>	<pre>width=#u; offset=#U; Rxx &amp;= ~((1&lt;&lt;width)-1)&lt;&lt;offset); Rxx  = ((Rss &amp; ((1&lt;&lt;width)-1)) &lt;&lt; offset);</pre>



**Syntax**

```
Rxx=insert(Rss,Rtt)
```

**Behavior**

```
width=zxt6->32((Rtt.w[1]));
offset=sxt7->32((Rtt.w[0]));
mask = ((1<<width)-1);
if (offset < 0) {
    Rxx = 0;
} else {
    Rxx &= ~(mask<<offset);
    Rxx |= ((Rss & mask) << offset);
}
```

**Class: XTYPE (slots 2,3)****Intrinsics**

Rx=insert(Rs,#u5,#U5)	Word32 Q6_R_insert_RII(Word32 Rx, Word32 Rs, Word32 Iu5, Word32 IU5)
Rx=insert(Rs,Rtt)	Word32 Q6_R_insert_RP(Word32 Rx, Word32 Rs, Word64 Rtt)
Rxx=insert(Rss,#u6,#U6)	Word64 Q6_P_insert_PII(Word64 Rxx, Word64 Rss, Word32 Iu6, Word32 IU6)
Rxx=insert(Rss,Rtt)	Word64 Q6_P_insert_PP(Word64 Rxx, Word64 Rss, Word64 Rtt)

**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				MajOp			s5					Parse				MinOp				x5										
1	0	0	0	0	0	1	1	I	I	I	s	s	s	s	s	P	P	i	i	i	i	i	i	I	I	I	x	x	x	x	x	Rxx=insert(Rss,#u6,#U6)
1	0	0	0	1	1	1	1	0	I	I	s	s	s	s	s	P	P	0	i	i	i	i	i	I	I	I	x	x	x	x	x	Rx=insert(Rs,#u5,#U5)
ICLASS		RegType							s5					Parse				t5				x5										
1	1	0	0	1	0	0	0	-	-	-	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	x	x	x	x	x	Rx=insert(Rs,Rtt)
ICLASS		RegType				Maj			s5					Parse				t5				x5										
1	1	0	0	1	0	1	0	0	-	-	s	s	s	s	s	P	P	0	t	t	t	t	t	-	-	-	x	x	x	x	x	Rxx=insert(Rss,Rtt)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x
MajOp	Major opcode
MinOp	Minor opcode
Maj	Major opcode
RegType	Register type

## Interleave/deinterleave

For interleave, bits I+32 of Rss (which are the bits from the upper source word) are placed in the odd bits (I\*2)+1 of Rdd, while bits I of Rss (which are the bits from the lower source word) are placed in the even bits (I\*2) of Rdd.

For deinterleave, the even bits of the source register are placed in the even register of the result pair, and the odd bits of the source register are placed in the odd register of the result pair.

"r1:0 = deinterleave(r1:0)" is the inverse of "r1:0 = interleave(r1:0)".

Syntax	Behavior
Rdd=deinterleave(Rss)	Rdd = deinterleave(ODD, EVEN);
Rdd=interleave(Rss)	Rdd = interleave(Rss.w[1], Rss.w[0]);

### Class: XTYPE (slots 2,3)

#### Intrinsics

Rdd=deinterleave(Rss)      Word64 Q6\_P\_deinterleave\_P(Word64 Rss)

Rdd=interleave(Rss)      Word64 Q6\_P\_interleave\_P(Word64 Rss)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp		s5					Parse		MinOp			d5												
1	0	0	0	0	0	0	0	1	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	0	d	d	d	d	d	Rdd=deinterleave(Rss)
1	0	0	0	0	0	0	0	1	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	1	d	d	d	d	d	Rdd=interleave(Rss)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Linear feedback-shift iteration

Count the number of ones of the logical AND of the two source input values, and take the least significant value of that sum. The first source value is shifted right by one bit, and the parity is placed in the MSB.

### Syntax

```
Rdd=lfs(Rss,Rtt)
```

### Behavior

```
Rdd = (Rss.u64 >> 1) | ((1&count_ones(Rss &
Rtt)).u64<<63);
```

**Class: XTYPE (slots 2,3)**

### Intrinsics

```
Rdd=lfs(Rss,Rtt)
```

```
Word64 Q6_P_lfs_PP(Word64 Rss, Word64 Rtt)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				Maj		s5					Parse		t5					Min		d5									
1	1	0	0	0	0	0	1	1	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	0	d	d	d	d	d	Rdd=lfs(Rss,Rtt)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Masked parity

Count the number of ones of the logical AND of the two source input values, and take the least significant bit of that sum.

Syntax	Behavior
<code>Rd=parity(Rs,Rt)</code>	<code>Rd = 1&amp;count_ones(Rs &amp; Rt);</code>
<code>Rd=parity(Rss,Rtt)</code>	<code>Rd = 1&amp;count_ones(Rss &amp; Rtt);</code>

### Class: XTYPE (slots 2,3)

#### Intrinsics

<code>Rd=parity(Rs,Rt)</code>	<code>Word32 Q6_R_parity_RR(Word32 Rs, Word32 Rt)</code>
<code>Rd=parity(Rss,Rtt)</code>	<code>Word32 Q6_R_parity_PP(Word64 Rss, Word64 Rtt)</code>

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				s5					Parse		t5					d5												
1	1	0	1	0	0	0	0	-	-	-	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=parity(Rss,Rtt)
1	1	0	1	0	1	0	1	1	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rd=parity(Rs,Rt)

Field name	Description
RegType	Register type
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Bit reverse

Reverse the order of bits. The most significant swap with the least significant, bit 30 swaps with bit 1, and so on.

Syntax	Behavior
Rd=brev(Rs)	Rd = reverse_bits(Rs);
Rdd=brev(Rss)	Rdd = reverse_bits(Rss);

### Class: XTYPE (slots 2,3)

#### Intrinsics

Rd=brev(Rs)	Word32 Q6_R_brev_R(Word32 Rs)
Rdd=brev(Rss)	Word64 Q6_P_brev_P(Word64 Rss)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse		MinOp			d5													
1	0	0	0	0	0	0	0	1	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	1	0	d	d	d	d	d	Rdd=brev(Rss)
1	0	0	0	1	1	0	0	0	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	1	0	d	d	d	d	d	Rd=brev(Rs)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Set/clear/toggle bit

Set (to 1), clear (to 0), or toggle a single bit in the source, and place the resulting value in the destination. Indicate the bit to manipulate using an immediate or register value.

If a register is used to indicate the bit position, and the value of the least-significant 7 bits of Rt is out of range, the destination register is unchanged.

Syntax	Behavior
Rd=clrbit(Rs,#u5)	$Rd = (Rs \& (\sim(1 \ll \#u)))$ ;
Rd=clrbit(Rs,Rt)	$Rd = (Rs \& (\sim((sxt_{7 \rightarrow 32}(Rt) > 0) ? (zxt_{32 \rightarrow 64}(1) \ll sxt_{7 \rightarrow 32}(Rt)) : (zxt_{32 \rightarrow 64}(1) \gg sxt_{7 \rightarrow 32}(Rt))))$ );
Rd=setbit(Rs,#u5)	$Rd = (Rs   (1 \ll \#u))$ ;
Rd=setbit(Rs,Rt)	$Rd = (Rs   (sxt_{7 \rightarrow 32}(Rt) > 0) ? (zxt_{32 \rightarrow 64}(1) \ll sxt_{7 \rightarrow 32}(Rt)) : (zxt_{32 \rightarrow 64}(1) \gg sxt_{7 \rightarrow 32}(Rt)))$ ;
Rd=togglebit(Rs,#u5)	$Rd = (Rs \wedge (1 \ll \#u))$ ;
Rd=togglebit(Rs,Rt)	$Rd = (Rs \wedge (sxt_{7 \rightarrow 32}(Rt) > 0) ? (zxt_{32 \rightarrow 64}(1) \ll sxt_{7 \rightarrow 32}(Rt)) : (zxt_{32 \rightarrow 64}(1) \gg sxt_{7 \rightarrow 32}(Rt)))$ ;

### Class: XTYPE (slots 2,3)

#### Intrinsics

Rd=clrbit(Rs,#u5)	Word32 Q6_R_clrbit_RI(Word32 Rs, Word32 Iu5)
Rd=clrbit(Rs,Rt)	Word32 Q6_R_clrbit_RR(Word32 Rs, Word32 Rt)
Rd=setbit(Rs,#u5)	Word32 Q6_R_setbit_RI(Word32 Rs, Word32 Iu5)
Rd=setbit(Rs,Rt)	Word32 Q6_R_setbit_RR(Word32 Rs, Word32 Rt)
Rd=togglebit(Rs,#u5)	Word32 Q6_R_togglebit_RI(Word32 Rs, Word32 Iu5)
Rd=togglebit(Rs,Rt)	Word32 Q6_R_togglebit_RR(Word32 Rs, Word32 Rt)

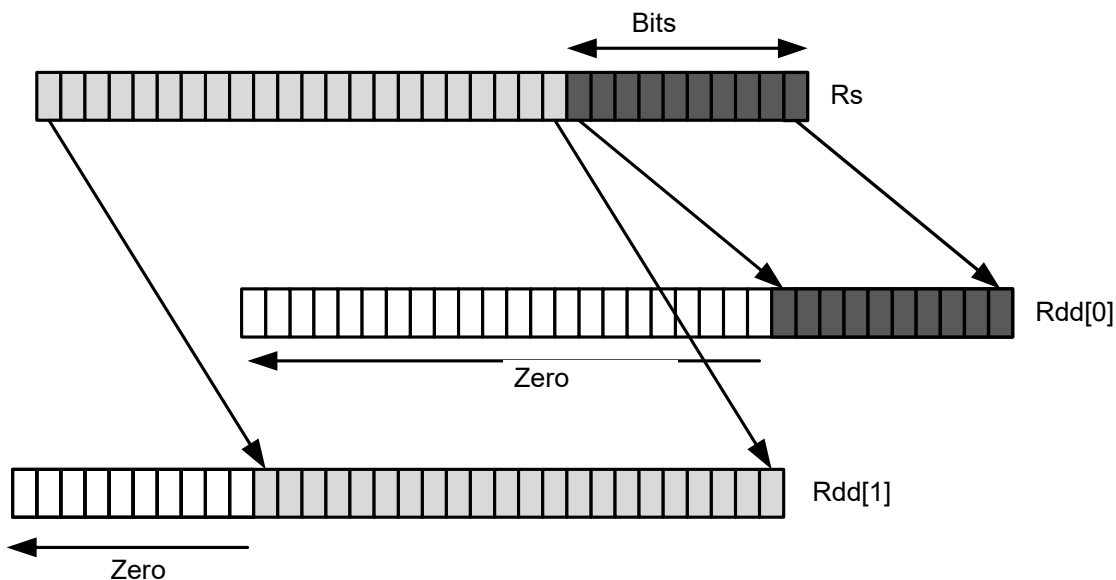
#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse			MinOp			d5												
1	0	0	0	1	1	0	0	1	1	0	s	s	s	s	s	P	P	0	i	i	i	i	i	0	0	0	d	d	d	d	d	Rd=setbit(Rs,#u5)
1	0	0	0	1	1	0	0	1	1	0	s	s	s	s	s	P	P	0	i	i	i	i	i	0	0	1	d	d	d	d	d	Rd=clrbit(Rs,#u5)
1	0	0	0	1	1	0	0	1	1	0	s	s	s	s	s	P	P	0	i	i	i	i	i	0	1	0	d	d	d	d	d	Rd=togglebit(Rs,#u5)
ICLASS			RegType				Maj		s5					Parse			t5			Min		d5										
1	1	0	0	0	1	1	0	1	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	d	d	d	d	d	Rd=setbit(Rs,Rt)
1	1	0	0	0	1	1	0	1	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	-	d	d	d	d	d	Rd=clrbit(Rs,Rt)
1	1	0	0	0	1	1	0	1	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	-	d	d	d	d	d	Rd=togglebit(Rs,Rt)

<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
MajOp	Major opcode
MinOp	Minor opcode
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Split bit field

Split the bit field in a register into upper and lower parts of variable size. The lower part is placed in the lower word of a destination register pair, and the upper part is placed in the upper word of the destination. An immediate value or register Rt is used to determine the bit position of the split.



### Syntax

```
Rdd=bitsplit (Rs, #u5)
```

```
Rdd=bitsplit (Rs, Rt)
```

### Behavior

```
Rdd.w[1]=(Rs>>#u);  
Rdd.w[0]=zxt#u->32(Rs);
```

```
shamt = zxt5->32(Rt);  
Rdd.w[1]=(Rs>>shamt);  
Rdd.w[0]=zxtshamt->32(Rs);
```

## Class: XTYPE (slots 2,3)

### Intrinsics

```
Rdd=bitsplit (Rs, #u5)
```

```
Word64 Q6_P_bitsplit_RI(Word32 Rs, Word32  
Iu5)
```

```
Rdd=bitsplit (Rs, Rt)
```

```
Word64 Q6_P_bitsplit_RR(Word32 Rs, Word32  
Rt)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType			MajOp			s5					Parse			MinOp			d5												
1	0	0	0	1	0	0	0	1	1	0	s	s	s	s	s	P	P	0	i	i	i	i	i	1	0	0	d	d	d	d	d	Rdd=bitsplit(Rs,#u5)
ICLASS			RegType			s5					Parse			t5					d5													
1	1	0	1	0	1	0	0	-	-	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	d	d	d	d	d	Rdd=bitsplit(Rs,Rt)



<b>Field name</b>	<b>Description</b>
RegType	Register type
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
MajOp	Major opcode
MinOp	Minor opcode

## Table index

The table index instruction supports fast lookup tables where the index into the table is stored in a bit-field. The instruction forms the address of a table element by extracting the bit field and inserting it into the appropriate bits of a pointer to the table element.

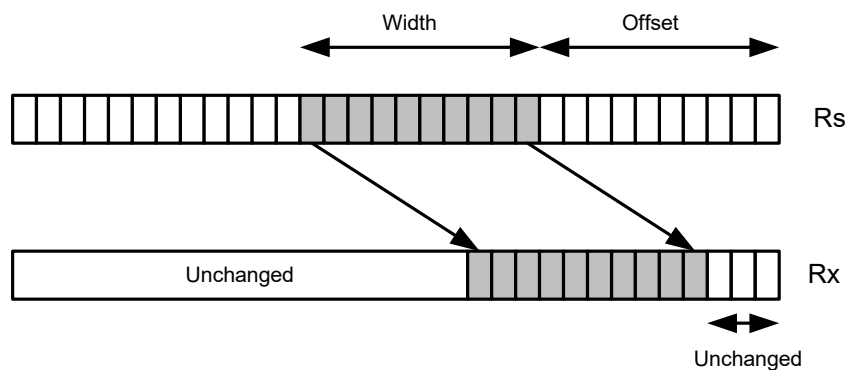
Tables are defined to contain entries of bytes, halfwords, words, or doublewords. The table must align to a power-of-two size greater than or equal to the table size. For example, a 4 K byte table should align to a 4 K byte boundary. This instruction supports tables with a maximum of 32 K table entries.

Register Rx contains a pointer to within the table. Register Rs contains a field to extract and use as a table index. This instruction first extracts the field from register Rs and then inserts it into register Rx. The insertion point is bit 0 for tables of bytes, bit 1 for tables of halfwords, bit 2 for tables of words, and bit 3 for tables of doublewords.

In the assembly syntax, the width and offset values represent the field in Rs to extract. Use unsigned constants to specify the width and offsets in assembly. In the encoded instruction, however, the assembler adjusts these values as follows.

- For `tableidxb`, no adjustment is necessary.
- For `tableidxh`, the assembler encodes `offset-1` in the signed immediate field.
- For `tableidxw`, the assembler encodes `offset-2` in the signed immediate field.
- For `tableidxd`, the assembler encodes `offset-3` in the signed immediate field.

`Rx=TABLEIDX(D)(Rs,#width,#offset)`



Syntax	Behavior
<code>Rx=tableidxb(Rs, #u4, #S6) :raw</code>	<pre>width=#u; offset=#S; field = Rs[(width+offset-1):offset]; Rx[(width-1+0):0]=field;</pre>
<code>Rx=tableidxb(Rs, #u4, #U5)</code>	Assembler mapped to: <code>"Rx=tableidxb(Rs, #u4, #U5) :raw"</code>
<code>Rx=tableidxd(Rs, #u4, #S6) :raw</code>	<pre>width=#u; offset=#S+3; field = Rs[(width+offset-1):offset]; Rx[(width-1+3):3]=field;</pre>

Syntax	Behavior
Rx=tableidxd(Rs, #u4, #U5)	Assembler mapped to: "Rx=tableidxd(Rs, #u4, #U5-3):raw"
Rx=tableidxh(Rs, #u4, #S6):raw	width=#u; offset=#S+1; field = Rs[(width+offset-1):offset]; Rx[(width-1+1):1]=field;
Rx=tableidxh(Rs, #u4, #U5)	Assembler mapped to: "Rx=tableidxh(Rs, #u4, #U5-1):raw"
Rx=tableidxw(Rs, #u4, #S6):raw	width=#u; offset=#S+2; field = Rs[(width+offset-1):offset]; Rx[(width-1+2):2]=field;
Rx=tableidxw(Rs, #u4, #U5)	Assembler mapped to: "Rx=tableidxw(Rs, #u4, #U5-2):raw"

### Class: XTYPE (slots 2,3)

#### Intrinsics

Rx=tableidxb(Rs, #u4, #U5)	Word32 Q6_R_tableidxb_RII(Word32 Rx, Word32 Rs, Word32 Iu4, Word32 IU5)
Rx=tableidxd(Rs, #u4, #U5)	Word32 Q6_R_tableidxd_RII(Word32 Rx, Word32 Rs, Word32 Iu4, Word32 IU5)
Rx=tableidxh(Rs, #u4, #U5)	Word32 Q6_R_tableidxh_RII(Word32 Rx, Word32 Rs, Word32 Iu4, Word32 IU5)
Rx=tableidxw(Rs, #u4, #U5)	Word32 Q6_R_tableidxw_RII(Word32 Rx, Word32 Rs, Word32 Iu4, Word32 IU5)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5								Parse				MinOp				x5							
1	0	0	0	0	1	1	1	0	0	i	s	s	s	s	s	P	P	I	I	I	I	I	I	i	i	i	x	x	x	x	x	Rx=tableidxb(Rs,#u4,#S6):raw
1	0	0	0	0	1	1	1	0	1	i	s	s	s	s	s	P	P	I	I	I	I	I	I	i	i	i	x	x	x	x	x	Rx=tableidxh(Rs,#u4,#S6):raw
1	0	0	0	0	1	1	1	1	0	i	s	s	s	s	s	P	P	I	I	I	I	I	I	i	i	i	x	x	x	x	x	Rx=tableidxw(Rs,#u4,#S6):raw
1	0	0	0	0	1	1	1	1	1	i	s	s	s	s	s	P	P	I	I	I	I	I	I	i	i	i	x	x	x	x	x	Rx=tableidxd(Rs,#u4,#S6):raw

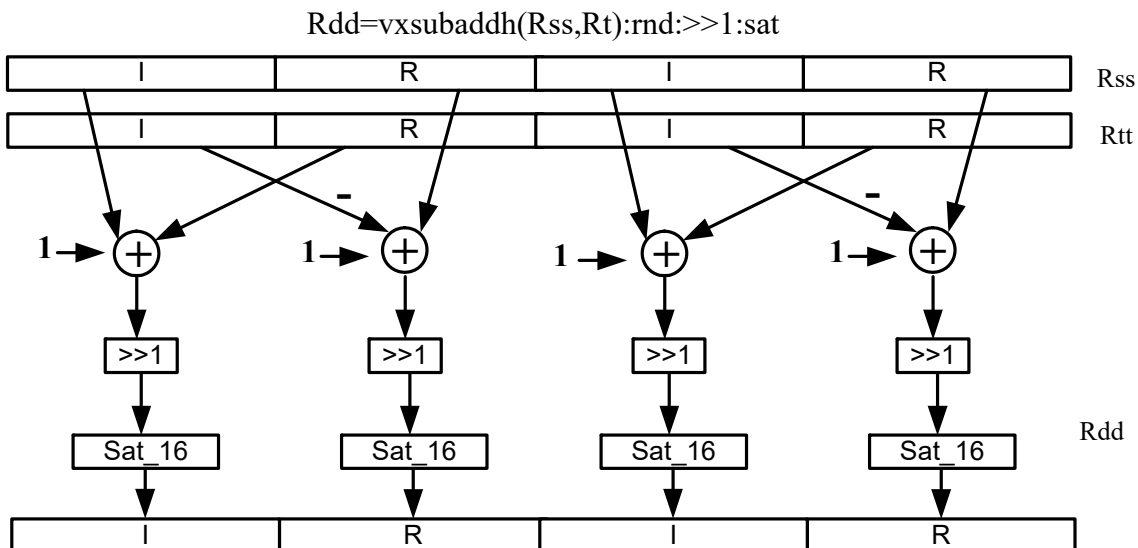
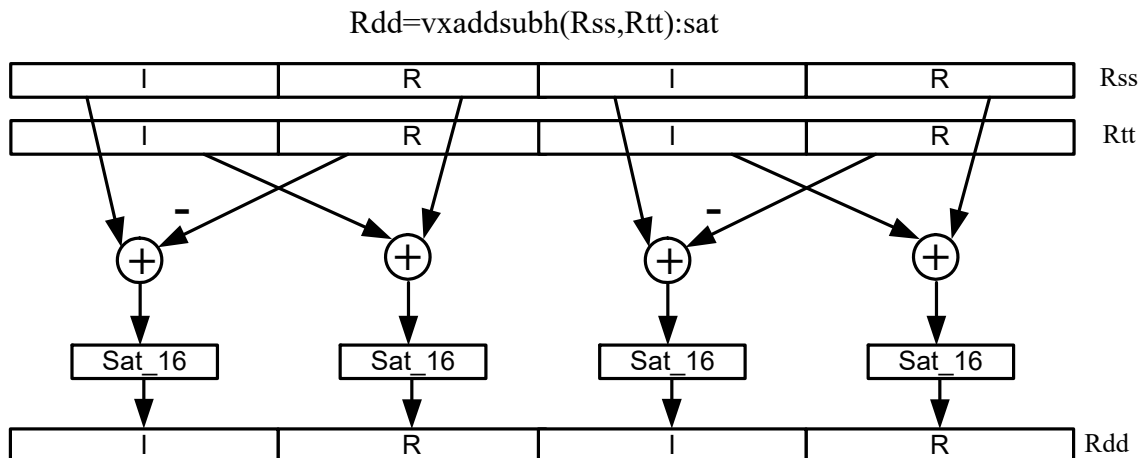
Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s
x5	Field to encode register x
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

### 11.10.3 XTYPE COMPLEX

The XTYPE COMPLEX instruction subclass includes instructions that are for complex math, using imaginary values.

#### Complex add/sub halfwords

Cross vector add-sub or sub-add instructions perform  $X + jY$  and  $X - jY$  complex operations. Each 16-bit result is saturated to 16 bits.



Syntax	Behavior
<code>Rdd=vxaddsubh (Rss, Rtt) :rnd:&gt;&gt;1 :sat</code>	<pre> Rdd.h[0]=sat<sub>16</sub>((Rss.h[0]+Rtt.h[1]+1)&gt;&gt;1) ; Rdd.h[1]=sat<sub>16</sub>((Rss.h[1]-Rtt.h[0]+1)&gt;&gt;1); Rdd.h[2]=sat<sub>16</sub>((Rss.h[2]+Rtt.h[3]+1)&gt;&gt;1) ; Rdd.h[3]=sat<sub>16</sub>((Rss.h[3]-Rtt.h[2]+1)&gt;&gt;1); </pre>
<code>Rdd=vxaddsubh (Rss, Rtt) :sat</code>	<pre> Rdd.h[0]=sat<sub>16</sub>(Rss.h[0]+Rtt.h[1]); Rdd.h[1]=sat<sub>16</sub>(Rss.h[1]-Rtt.h[0]); Rdd.h[2]=sat<sub>16</sub>(Rss.h[2]+Rtt.h[3]); Rdd.h[3]=sat<sub>16</sub>(Rss.h[3]-Rtt.h[2]); </pre>
<code>Rdd=vxsubaddh (Rss, Rtt) :rnd:&gt;&gt;1 :sat</code>	<pre> Rdd.h[0]=sat<sub>16</sub>((Rss.h[0]-Rtt.h[1]+1)&gt;&gt;1); Rdd.h[1]=sat<sub>16</sub>((Rss.h[1]+Rtt.h[0]+1)&gt;&gt;1) ; Rdd.h[2]=sat<sub>16</sub>((Rss.h[2]-Rtt.h[3]+1)&gt;&gt;1); Rdd.h[3]=sat<sub>16</sub>((Rss.h[3]+Rtt.h[2]+1)&gt;&gt;1) ; </pre>
<code>Rdd=vxsubaddh (Rss, Rtt) :sat</code>	<pre> Rdd.h[0]=sat<sub>16</sub>(Rss.h[0]-Rtt.h[1]); Rdd.h[1]=sat<sub>16</sub>(Rss.h[1]+Rtt.h[0]); Rdd.h[2]=sat<sub>16</sub>(Rss.h[2]-Rtt.h[3]); Rdd.h[3]=sat<sub>16</sub>(Rss.h[3]+Rtt.h[2]); </pre>

**Class: XTYPE (slots 2,3)****Notes**

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

## Intrinsics

Rdd=vxaddsubh(Rss,Rtt):rnd:>>1:sat	Word64 Q6_P_vxaddsubh_PP_rnd_rsl_sat(Word64 Rss, Word64 Rtt)
Rdd=vxaddsubh(Rss,Rtt):sat	Word64 Q6_P_vxaddsubh_PP_sat(Word64 Rss, Word64 Rtt)
Rdd=vxsubaddh(Rss,Rtt):rnd:>>1:sat	Word64 Q6_P_vxsubaddh_PP_rnd_rsl_sat(Word64 Rss, Word64 Rtt)
Rdd=vxsubaddh(Rss,Rtt):sat	Word64 Q6_P_vxsubaddh_PP_sat(Word64 Rss, Word64 Rtt)

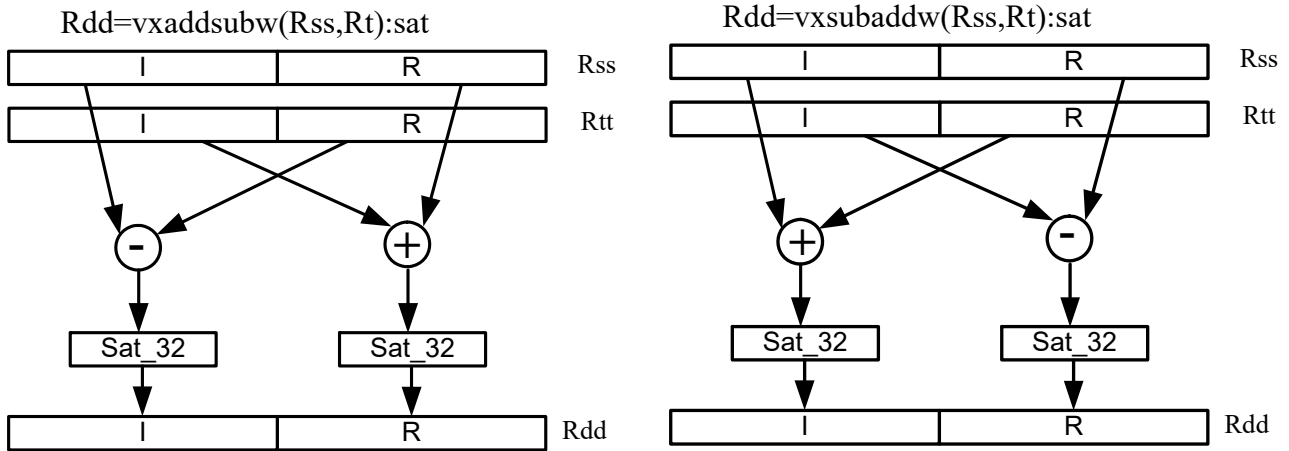
## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				Maj		s5					Parse		t5					Min		d5									
1	1	0	0	0	0	0	1	0	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	0	d	d	d	d	d	Rdd=vxaddsubh(Rss,Rtt):sat
1	1	0	0	0	0	0	1	0	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	0	d	d	d	d	d	Rdd=vxsubaddh(Rss,Rtt):sat
1	1	0	0	0	0	0	1	1	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	d	d	d	d	d	Rdd=vxaddsubh(Rss,Rtt):rnd:>>1:sat
1	1	0	0	0	0	0	1	1	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	-	d	d	d	d	d	Rdd=vxsubaddh(Rss,Rtt):rnd:>>1:sat

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Complex add/sub words

Cross vector add-sub or sub-add instructions perform  $X+jY$  and  $X-jY$  complex operations. Each 32-bit result is saturated to 32 bits.



### Syntax

`Rdd=vxaddsubw(Rss,Rtt):sat`

`Rdd=vxsubaddw(Rss,Rtt):sat`

### Behavior

`Rdd.w[0]=sat32(Rss.w[0]+Rtt.w[1]);`  
`Rdd.w[1]=sat32(Rss.w[1]-Rtt.w[0]);`

`Rdd.w[0]=sat32(Rss.w[0]-Rtt.w[1]);`  
`Rdd.w[1]=sat32(Rss.w[1]+Rtt.w[0]);`

## Class: XTYPE (slots 2,3)

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

### Intrinsics

`Rdd=vxaddsubw(Rss,Rtt):sat`

`Word64 Q6_P_vxaddsubw_PP_sat(Word64 Rss, Word64 Rtt)`

`Rdd=vxsubaddw(Rss,Rtt):sat`

`Word64 Q6_P_vxsubaddw_PP_sat(Word64 Rss, Word64 Rtt)`

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType					Maj		s5					Parse		t5				Min		d5										
1	1	0	0	0	0	0	1	0	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	d	d	d	d	d	
1	1	0	0	0	0	0	1	0	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	d	d	d	d	d	

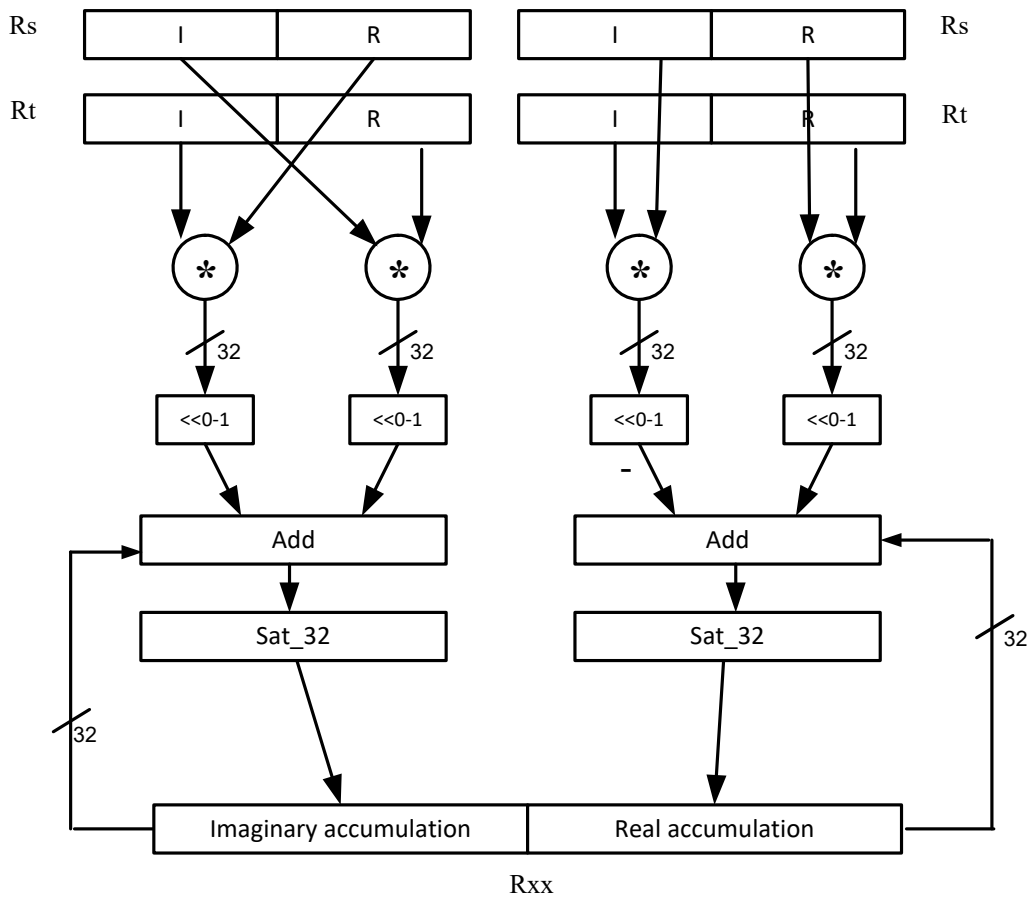
<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
Maj	Major opcode
Min	Minor opcode
RegType	Register type



## Complex multiply

Multiply complex values Rs and Rt. The inputs have a real 16-bit value in the low halfword and an imaginary 16-bit value in the high halfword. Optionally, scale the result by 0-1 bits. Optionally, add a complex accumulator. Saturate the real and imaginary portions to 32-bits. The output has a real 32-bit value in the low word and an imaginary 32-bit value in the high word. The Rt input can be optionally conjugated. Another option is to subtracted the result from the destination rather than accumulate it.

$R_{xx} += \text{cmpy}(R_s, R_t) : \text{sat}$



**Syntax**

$R_{dd} = \text{cmpy}(R_s, R_t) [ : \ll 1 ] : \text{sat}$

$R_{dd}.w[1] = \text{sat}_{32}((R_s.h[1] * R_t.h[0]) [\ll 1] + (R_s.h[0] * R_t.h[1]) [\ll 1]);$   
 $R_{dd}.w[0] = \text{sat}_{32}((R_s.h[0] * R_t.h[0]) [\ll 1] - (R_s.h[1] * R_t.h[1]) [\ll 1]);$

$R_{dd} = \text{cmpy}(R_s, R_t^*) [ : \ll 1 ] : \text{sat}$

$R_{dd}.w[1] = \text{sat}_{32}((R_s.h[1] * R_t.h[0]) [\ll 1] - (R_s.h[0] * R_t.h[1]) [\ll 1]);$   
 $R_{dd}.w[0] = \text{sat}_{32}((R_s.h[0] * R_t.h[0]) [\ll 1] + (R_s.h[1] * R_t.h[1]) [\ll 1]);$

**Behavior**

Syntax	Behavior
$R_{xx} += \text{cmpy}(R_s, R_t) [ : \ll 1 ] : \text{sat}$	$R_{xx}.w[1] = \text{sat}_{32}(R_{xx}.w[1] + (R_s.h[1] * R_t.h[0]) [ \ll 1 ] + (R_s.h[0] * R_t.h[1]) [ \ll 1 ] );$ $R_{xx}.w[0] = \text{sat}_{32}(R_{xx}.w[0] + (R_s.h[0] * R_t.h[0]) [ \ll 1 ] - (R_s.h[1] * R_t.h[1]) [ \ll 1 ] );$
$R_{xx} += \text{cmpy}(R_s, R_t^*) [ : \ll 1 ] : \text{sat}$	$R_{xx}.w[1] = \text{sat}_{32}(R_{xx}.w[1] + (R_s.h[1] * R_t.h[0]) [ \ll 1 ] - (R_s.h[0] * R_t.h[1]) [ \ll 1 ] );$ $R_{xx}.w[0] = \text{sat}_{32}(R_{xx}.w[0] + (R_s.h[0] * R_t.h[0]) [ \ll 1 ] + (R_s.h[1] * R_t.h[1]) [ \ll 1 ] );$
$R_{xx} - = \text{cmpy}(R_s, R_t) [ : \ll 1 ] : \text{sat}$	$R_{xx}.w[1] = \text{sat}_{32}(R_{xx}.w[1] - ((R_s.h[1] * R_t.h[0]) [ \ll 1 ] + (R_s.h[0] * R_t.h[1]) [ \ll 1 ] ));$ $R_{xx}.w[0] = \text{sat}_{32}(R_{xx}.w[0] - ((R_s.h[0] * R_t.h[0]) [ \ll 1 ] - (R_s.h[1] * R_t.h[1]) [ \ll 1 ] ));$
$R_{xx} - = \text{cmpy}(R_s, R_t^*) [ : \ll 1 ] : \text{sat}$	$R_{xx}.w[1] = \text{sat}_{32}(R_{xx}.w[1] - ((R_s.h[1] * R_t.h[0]) [ \ll 1 ] - (R_s.h[0] * R_t.h[1]) [ \ll 1 ] ));$ $R_{xx}.w[0] = \text{sat}_{32}(R_{xx}.w[0] - ((R_s.h[0] * R_t.h[0]) [ \ll 1 ] + (R_s.h[1] * R_t.h[1]) [ \ll 1 ] ));$

**Class: XTYPE (slots 2,3)****Notes**

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

## Intrinsics

Rdd=cmpy (Rs,Rt) :<<1:sat	Word64 Q6_P_cmpy_RR_s1_sat (Word32 Rs, Word32 Rt)
Rdd=cmpy (Rs,Rt) :sat	Word64 Q6_P_cmpy_RR_sat (Word32 Rs, Word32 Rt)
Rdd=cmpy (Rs,Rt*) :<<1:sat	Word64 Q6_P_cmpy_RR_conj_s1_sat (Word32 Rs, Word32 Rt)
Rdd=cmpy (Rs,Rt*) :sat	Word64 Q6_P_cmpy_RR_conj_sat (Word32 Rs, Word32 Rt)
Rxx+=cmpy (Rs,Rt) :<<1:sat	Word64 Q6_P_cmpyacc_RR_s1_sat (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx+=cmpy (Rs,Rt) :sat	Word64 Q6_P_cmpyacc_RR_sat (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx+=cmpy (Rs,Rt*) :<<1:sat	Word64 Q6_P_cmpyacc_RR_conj_s1_sat (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx+=cmpy (Rs,Rt*) :sat	Word64 Q6_P_cmpyacc_RR_conj_sat (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx-=cmpy (Rs,Rt) :<<1:sat	Word64 Q6_P_cmpynac_RR_s1_sat (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx-=cmpy (Rs,Rt) :sat	Word64 Q6_P_cmpynac_RR_sat (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx-=cmpy (Rs,Rt*) :<<1:sat	Word64 Q6_P_cmpynac_RR_conj_s1_sat (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx-=cmpy (Rs,Rt*) :sat	Word64 Q6_P_cmpynac_RR_conj_sat (Word64 Rxx, Word32 Rs, Word32 Rt)

## Encoding

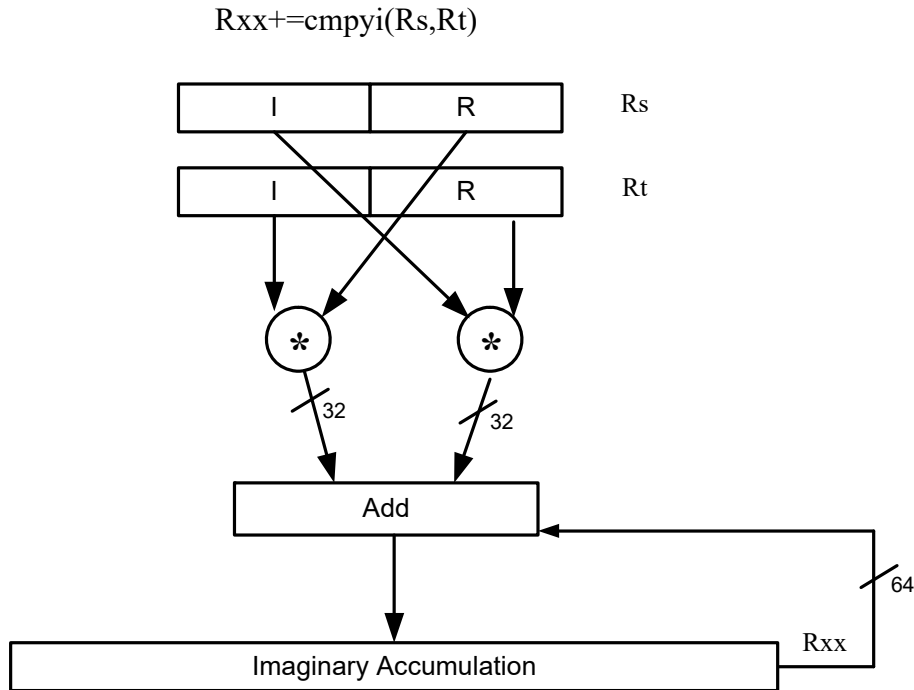
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				MajOp			s5					Parse		t5					MinOp			d5								
1	1	1	0	0	1	0	1	N	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	0	d	d	d	d	d	Rdd=cmpy(Rs,Rt)[:<<N]:sat
1	1	1	0	0	1	0	1	N	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	0	d	d	d	d	d	Rdd=cmpy(Rs,Rt*)[:<<N]:sat
ICLASS		RegType				MajOp			s5					Parse		t5					MinOp			x5								
1	1	1	0	0	1	1	1	N	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	0	x	x	x	x	x	Rxx+=cmpy(Rs,Rt)[:<<N]:sat
1	1	1	0	0	1	1	1	N	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	1	x	x	x	x	x	Rxx-=cmpy(Rs,Rt)[:<<N]:sat
1	1	1	0	0	1	1	1	N	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	0	x	x	x	x	x	Rxx+=cmpy(Rs,Rt*)[:<<N]:sat
1	1	1	0	0	1	1	1	N	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	1	x	x	x	x	x	Rxx-=cmpy(Rs,Rt*)[:<<N]:sat

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d

<b>Field name</b>	<b>Description</b>
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Complex multiply real or imaginary

Multiply complex values  $R_s$  and  $R_t$ . The inputs have a real 16-bit value in the low halfword and an imaginary 16-bit value in the high halfword. Take either the real or imaginary result and optionally accumulate with a 64-bit destination.



Syntax	Behavior
$R_{dd} = \text{cmpyi}(R_s, R_t)$	$R_{dd} = (R_s.h[1] * R_t.h[0]) + (R_s.h[0] * R_t.h[1]);$
$R_{dd} = \text{cmpyr}(R_s, R_t)$	$R_{dd} = (R_s.h[0] * R_t.h[0]) - (R_s.h[1] * R_t.h[1]);$
$R_{xx} += \text{cmpyi}(R_s, R_t)$	$R_{xx} = R_{xx} + (R_s.h[1] * R_t.h[0]) + (R_s.h[0] * R_t.h[1]);$
$R_{xx} += \text{cmpyr}(R_s, R_t)$	$R_{xx} = R_{xx} + (R_s.h[0] * R_t.h[0]) - (R_s.h[1] * R_t.h[1]);$

**Class: XTYPE (slots 2,3)****Intrinsics**

Rdd=cmpyi (Rs,Rt)	Word64 Q6_P_cmpyi_RR(Word32 Rs, Word32 Rt)
Rdd=cmpyr (Rs,Rt)	Word64 Q6_P_cmpyr_RR(Word32 Rs, Word32 Rt)
Rxx+=cmpyi (Rs,Rt)	Word64 Q6_P_cmpyiacc_RR(Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx+=cmpyr (Rs,Rt)	Word64 Q6_P_cmpyracc_RR(Word64 Rxx, Word32 Rs, Word32 Rt)

**Encoding**

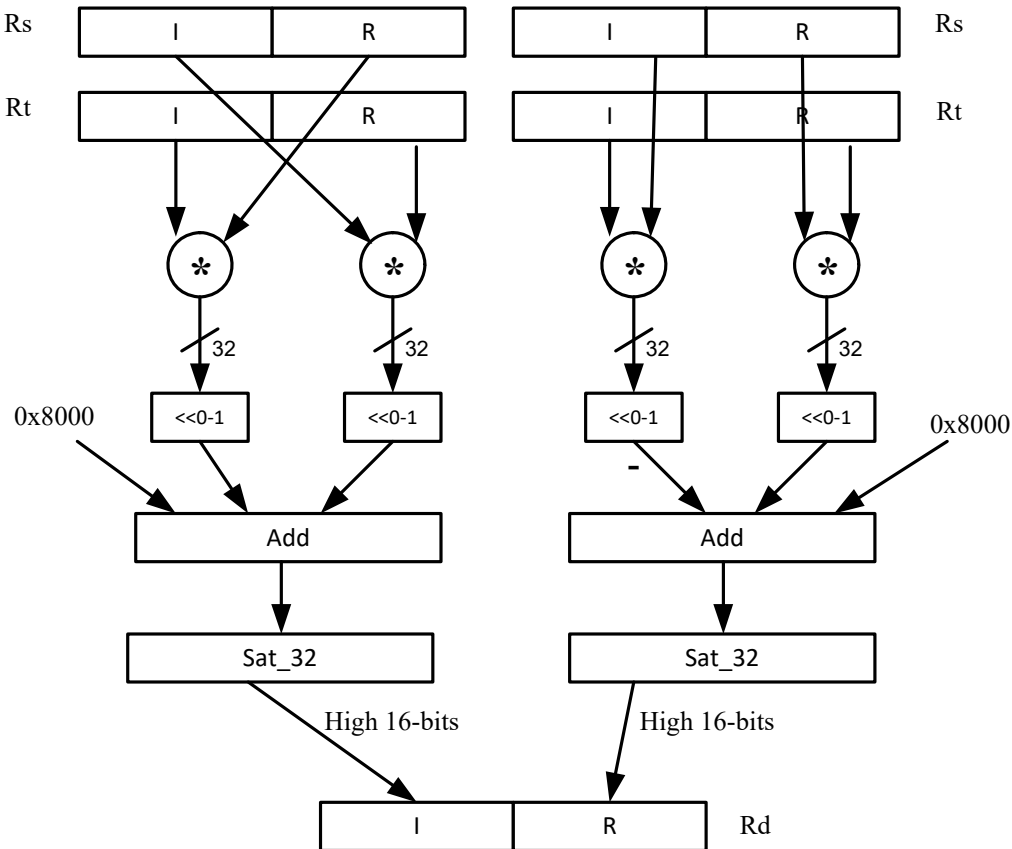
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp		d5								
1	1	1	0	0	1	0	1	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	d	d	d	d	d	Rdd=cmpyi(Rs,Rt)
1	1	1	0	0	1	0	1	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	0	d	d	d	d	d	Rdd=cmpyr(Rs,Rt)
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp		x5								
1	1	1	0	0	1	1	1	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	x	x	x	x	x	Rxx+=cmpyi(Rs,Rt)
1	1	1	0	0	1	1	1	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	0	x	x	x	x	x	Rxx+=cmpyr(Rs,Rt)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Complex multiply with round and pack

Multiply complex values  $R_s$  and  $R_t$ . The inputs have a real 16-bit value in the low halfword and an imaginary 16-bit value in the high halfword. The  $R_t$  input is optionally conjugated. The multiplier results are optionally scaled by 0 to 1 bits. A rounding constant is added to each real and imaginary sum. The real and imaginary parts are individually saturated to 32 bits. The upper 16 bits of each 32-bit results are packed in a 32-bit destination register.

$Rd = \text{cmpy}(Rs, Rt) : \text{rnd} : \text{sat}$



### Syntax

$Rd = \text{cmpy}(Rs, Rt) [ : \ll 1 ] : \text{rnd} : \text{sat}$

$Rd = \text{cmpy}(Rs, Rt^*) [ : \ll 1 ] : \text{rnd} : \text{sat}$

### Behavior

$Rd.h[1] = (\text{sat}_{32}((Rs.h[1] * Rt.h[0]) [\ll 1] + (Rs.h[0] * Rt.h[1]) [\ll 1] + 0x8000)).h[1];$   
 $Rd.h[0] = (\text{sat}_{32}((Rs.h[0] * Rt.h[0]) [\ll 1] - (Rs.h[1] * Rt.h[1]) [\ll 1] + 0x8000)).h[1];$

$Rd.h[1] = (\text{sat}_{32}((Rs.h[1] * Rt.h[0]) [\ll 1] - (Rs.h[0] * Rt.h[1]) [\ll 1] + 0x8000)).h[1];$   
 $Rd.h[0] = (\text{sat}_{32}((Rs.h[0] * Rt.h[0]) [\ll 1] + (Rs.h[1] * Rt.h[1]) [\ll 1] + 0x8000)).h[1];$

**Class: XTYPE (slots 2,3)****Notes**

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

**Intrinsics**

Rd=cmpy(Rs,Rt):<<1:rnd:sat	Word32 Q6_R_cmpy_RR_s1_rnd_sat(Word32 Rs, Word32 Rt)
Rd=cmpy(Rs,Rt):rnd:sat	Word32 Q6_R_cmpy_RR_rnd_sat(Word32 Rs, Word32 Rt)
Rd=cmpy(Rs,Rt*):<<1:rnd:sat	Word32 Q6_R_cmpy_RR_conj_s1_rnd_sat(Word32 Rs, Word32 Rt)
Rd=cmpy(Rs,Rt*):rnd:sat	Word32 Q6_R_cmpy_RR_conj_rnd_sat(Word32 Rs, Word32 Rt)

**Encoding**

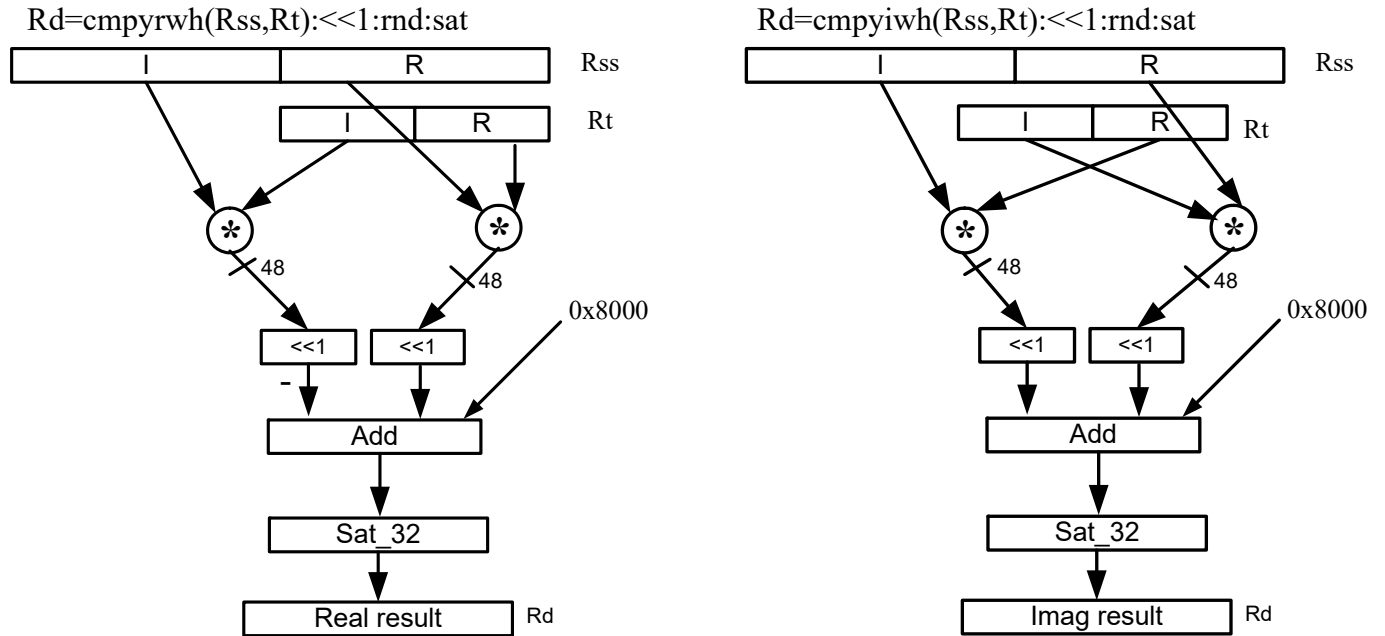
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp		s5					Parse		t5					MinOp		d5								
1	1	1	0	1	1	0	1	N	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	0	d	d	d	d	d	Rd=cmpy(Rs,Rt)[:<<N]:rnd:sat
1	1	1	0	1	1	0	1	N	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	0	d	d	d	d	d	Rd=cmpy(Rs,Rt*)[:<<N]:rnd:sat

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t



## Complex multiply $32 \times 16$

Multiply 32 by 16 bit complex values Rss and Rt. The inputs have a real value in the low part of a register and the imaginary value in the upper part. The multiplier results are scaled by 1 bit and accumulated with a rounding constant. The result is saturated to 32 bits.



### Syntax

`Rd=cmpyiw h(Rss,Rt) : <<1 : rnd : s  
at`

`Rd=cmpyiw h(Rss,Rt*) : <<1 : rnd :  
sat`

`Rd=cmpyrw h(Rss,Rt) : <<1 : rnd : s  
at`

`Rd=cmpyrw h(Rss,Rt*) : <<1 : rnd :  
sat`

### Behavior

$Rd = \text{sat}_{32}((Rss.w[0] * Rt.h[1]) + (Rss.w[1] * Rt.h[0]) + 0x4000) \gg 15;$

$Rd = \text{sat}_{32}((Rss.w[1] * Rt.h[0]) - (Rss.w[0] * Rt.h[1]) + 0x4000) \gg 15;$

$Rd = \text{sat}_{32}((Rss.w[0] * Rt.h[0]) - (Rss.w[1] * Rt.h[1]) + 0x4000) \gg 15;$

$Rd = \text{sat}_{32}((Rss.w[0] * Rt.h[0]) + (Rss.w[1] * Rt.h[1]) + 0x4000) \gg 15;$

### Class: XTYPE (slots 2,3)

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

## Intrinsics

Rd=cmpyiw(Rss,Rt):<<1:rnd:s at Word32 Q6\_R\_cmpyiw\_PR\_s1\_rnd\_sat(Word64 Rss, Word32 Rt)

Rd=cmpyiw(Rss,Rt\*):<<1:rnd:sat Word32 Q6\_R\_cmpyiw\_PR\_conj\_s1\_rnd\_sat(Word64 Rss, Word32 Rt)

Rd=cmpyrw(Rss,Rt):<<1:rnd:s at Word32 Q6\_R\_cmpyrw\_PR\_s1\_rnd\_sat(Word64 Rss, Word32 Rt)

Rd=cmpyrw(Rss,Rt\*):<<1:rnd:sat Word32 Q6\_R\_cmpyrw\_PR\_conj\_s1\_rnd\_sat(Word64 Rss, Word32 Rt)

## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5					Min		d5											
1	1	0	0	0	1	0	1	-	-	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	0	d	d	d	d	d	Rd=cmpyiw(Rss,Rt):<<1:rnd:sat
1	1	0	0	0	1	0	1	-	-	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	1	d	d	d	d	d	Rd=cmpyiw(Rss,Rt*):<<1:rnd:sat
1	1	0	0	0	1	0	1	-	-	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	0	d	d	d	d	d	Rd=cmpyrw(Rss,Rt):<<1:rnd:sat
1	1	0	0	0	1	0	1	-	-	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	1	d	d	d	d	d	Rd=cmpyrw(Rss,Rt*):<<1:rnd:sat

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
Min	Minor opcode
RegType	Register type

## Complex multiply real or imaginary 32-bit

Multiply complex values Rss and Rtt. The inputs have a real 32-bit value in the low word and an imaginary 32-bit value in the high word. Take either the real or imaginary result and optionally accumulate with a 64-bit destination.

Syntax	Behavior
<code>Rd=cmpyiw(Rss,Rtt):&lt;&lt;1:rnd:sat</code>	<pre>tmp128 = sxt<sub>64-&gt;128</sub>((Rss.w[0] * Rtt.w[1])); acc128 = sxt<sub>64-&gt;128</sub>((Rss.w[1] * Rtt.w[0])); const128 = sxt<sub>64-&gt;128</sub>(0x40000000); acc128 = tmp128+acc128; acc128 = acc128+const128; acc128 = (size8s_t) (acc128 &gt;&gt; 31); acc64 = sxt<sub>128-&gt;64</sub>(acc128); Rd = sat<sub>32</sub>(acc64);</pre>
<code>Rd=cmpyiw(Rss,Rtt):&lt;&lt;1:sat</code>	<pre>tmp128 = sxt<sub>64-&gt;128</sub>((Rss.w[0] * Rtt.w[1])); acc128 = sxt<sub>64-&gt;128</sub>((Rss.w[1] * Rtt.w[0])); acc128 = tmp128+acc128; acc128 = (size8s_t) (acc128 &gt;&gt; 31); acc64 = sxt<sub>128-&gt;64</sub>(acc128); Rd = sat<sub>32</sub>(acc64);</pre>
<code>Rd=cmpyiw(Rss,Rtt*):&lt;&lt;1:rnd:sat</code>	<pre>tmp128 = sxt<sub>64-&gt;128</sub>((Rss.w[1] * Rtt.w[0])); acc128 = sxt<sub>64-&gt;128</sub>((Rss.w[0] * Rtt.w[1])); const128 = sxt<sub>64-&gt;128</sub>(0x40000000); acc128 = tmp128-acc128; acc128 = acc128+const128; acc128 = (size8s_t) (acc128 &gt;&gt; 31); acc64 = sxt<sub>128-&gt;64</sub>(acc128); Rd = sat<sub>32</sub>(acc64);</pre>
<code>Rd=cmpyiw(Rss,Rtt*):&lt;&lt;1:sat</code>	<pre>tmp128 = sxt<sub>64-&gt;128</sub>((Rss.w[1] * Rtt.w[0])); acc128 = sxt<sub>64-&gt;128</sub>((Rss.w[0] * Rtt.w[1])); acc128 = tmp128-acc128; acc128 = (size8s_t) (acc128 &gt;&gt; 31); acc64 = sxt<sub>128-&gt;64</sub>(acc128); Rd = sat<sub>32</sub>(acc64);</pre>
<code>Rd=cmpyrw(Rss,Rtt):&lt;&lt;1:rnd:sat</code>	<pre>tmp128 = sxt<sub>64-&gt;128</sub>((Rss.w[0] * Rtt.w[0])); acc128 = sxt<sub>64-&gt;128</sub>((Rss.w[1] * Rtt.w[1])); const128 = sxt<sub>64-&gt;128</sub>(0x40000000); acc128 = tmp128-acc128; acc128 = acc128+const128; acc128 = (size8s_t) (acc128 &gt;&gt; 31); acc64 = sxt<sub>128-&gt;64</sub>(acc128); Rd = sat<sub>32</sub>(acc64);</pre>
<code>Rd=cmpyrw(Rss,Rtt):&lt;&lt;1:sat</code>	<pre>tmp128 = sxt<sub>64-&gt;128</sub>((Rss.w[0] * Rtt.w[0])); acc128 = sxt<sub>64-&gt;128</sub>((Rss.w[1] * Rtt.w[1])); acc128 = tmp128-acc128; acc128 = (size8s_t) (acc128 &gt;&gt; 31); acc64 = sxt<sub>128-&gt;64</sub>(acc128); Rd = sat<sub>32</sub>(acc64);</pre>

Syntax	Behavior
<code>Rd=cmprw (Rss, Rtt*) :&lt;&lt;1:rnd:sat</code>	<pre> tmp128 = sxt<sub>64-&gt;128</sub>((Rss.w[0] * Rtt.w[0])); acc128 = sxt<sub>64-&gt;128</sub>((Rss.w[1] * Rtt.w[1])); const128 = sxt<sub>64-&gt;128</sub>(0x40000000); acc128 = tmp128+acc128; acc128 = acc128+const128; acc128 = (size8s_t) (acc128 &gt;&gt; 31); acc64 = sxt<sub>128-&gt;64</sub>(acc128); Rd = sat<sub>32</sub>(acc64); </pre>
<code>Rd=cmprw (Rss, Rtt*) :&lt;&lt;1:sat</code>	<pre> tmp128 = sxt<sub>64-&gt;128</sub>((Rss.w[0] * Rtt.w[0])); acc128 = sxt<sub>64-&gt;128</sub>((Rss.w[1] * Rtt.w[1])); acc128 = tmp128+acc128; acc128 = (size8s_t) (acc128 &gt;&gt; 31); acc64 = sxt<sub>128-&gt;64</sub>(acc128); Rd = sat<sub>32</sub>(acc64); </pre>
<code>Rdd=cmprw (Rss, Rtt)</code>	<code>Rdd = ((Rss.w[0] * Rtt.w[1]) + (Rss.w[1] * Rtt.w[0]));</code>
<code>Rdd=cmprw (Rss, Rtt*)</code>	<code>Rdd = ((Rss.w[1] * Rtt.w[0]) - (Rss.w[0] * Rtt.w[1]));</code>
<code>Rdd=cmprw (Rss, Rtt)</code>	<code>Rdd = ((Rss.w[0] * Rtt.w[0]) - (Rss.w[1] * Rtt.w[1]));</code>
<code>Rdd=cmprw (Rss, Rtt*)</code>	<code>Rdd = ((Rss.w[0] * Rtt.w[0]) + (Rss.w[1] * Rtt.w[1]));</code>
<code>Rxx+=cmprw (Rss, Rtt)</code>	<code>Rxx += ((Rss.w[0] * Rtt.w[1]) + (Rss.w[1] * Rtt.w[0]));</code>
<code>Rxx+=cmprw (Rss, Rtt*)</code>	<code>Rxx += ((Rss.w[1] * Rtt.w[0]) - (Rss.w[0] * Rtt.w[1]));</code>
<code>Rxx+=cmprw (Rss, Rtt)</code>	<code>Rxx += ((Rss.w[0] * Rtt.w[0]) - (Rss.w[1] * Rtt.w[1]));</code>
<code>Rxx+=cmprw (Rss, Rtt*)</code>	<code>Rxx += ((Rss.w[0] * Rtt.w[0]) + (Rss.w[1] * Rtt.w[1]));</code>

**Class: XTYPE (slots 3)****Notes**

- This instruction can only execute on a core with the Hexagon audio extensions
- A packet with this instruction cannot have a slot 2 multiply instruction.
- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

## Intrinsics

Rd=cmpyiw(Rss,Rtt):<<1:rnd:sat	Word32 Q6_R_cmpyiw_PP_s1_rnd_sat(Word64 Rss, Word64 Rtt)
Rd=cmpyiw(Rss,Rtt):<<1:sat	Word32 Q6_R_cmpyiw_PP_s1_sat(Word64 Rss, Word64 Rtt)
Rd=cmpyiw(Rss,Rtt*):<<1:rnd:sat	Word32 Q6_R_cmpyiw_PP_conj_s1_rnd_sat(Word64 Rss, Word64 Rtt)
Rd=cmpyiw(Rss,Rtt*):<<1:sat	Word32 Q6_R_cmpyiw_PP_conj_s1_sat(Word64 Rss, Word64 Rtt)
Rd=cmpyrw(Rss,Rtt):<<1:rnd:sat	Word32 Q6_R_cmpyrw_PP_s1_rnd_sat(Word64 Rss, Word64 Rtt)
Rd=cmpyrw(Rss,Rtt):<<1:sat	Word32 Q6_R_cmpyrw_PP_s1_sat(Word64 Rss, Word64 Rtt)
Rd=cmpyrw(Rss,Rtt*):<<1:rnd:sat	Word32 Q6_R_cmpyrw_PP_conj_s1_rnd_sat(Word64 Rss, Word64 Rtt)
Rd=cmpyrw(Rss,Rtt*):<<1:sat	Word32 Q6_R_cmpyrw_PP_conj_s1_sat(Word64 Rss, Word64 Rtt)
Rdd=cmpyiw(Rss,Rtt)	Word64 Q6_P_cmpyiw_PP(Word64 Rss, Word64 Rtt)
Rdd=cmpyiw(Rss,Rtt*)	Word64 Q6_P_cmpyiw_PP_conj(Word64 Rss, Word64 Rtt)
Rdd=cmpyrw(Rss,Rtt)	Word64 Q6_P_cmpyrw_PP(Word64 Rss, Word64 Rtt)
Rdd=cmpyrw(Rss,Rtt*)	Word64 Q6_P_cmpyrw_PP_conj(Word64 Rss, Word64 Rtt)
Rxx+=cmpyiw(Rss,Rtt)	Word64 Q6_P_cmpyiwacc_PP(Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=cmpyiw(Rss,Rtt*)	Word64 Q6_P_cmpyiwacc_PP_conj(Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=cmpyrw(Rss,Rtt)	Word64 Q6_P_cmpyrwacc_PP(Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=cmpyrw(Rss,Rtt*)	Word64 Q6_P_cmpyrwacc_PP_conj(Word64 Rxx, Word64 Rss, Word64 Rtt)

## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp				s5				Parse		t5				MinOp		d5									
1	1	1	0	1	0	0	0	0	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	0	d	d	d	d	d	Rdd=cmpyiw(Rss,Rtt)
1	1	1	0	1	0	0	0	1	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	0	d	d	d	d	d	Rdd=cmpyrw(Rss,Rtt)
1	1	1	0	1	0	0	0	1	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	0	d	d	d	d	d	Rdd=cmpyrw(Rss,Rtt*)
1	1	1	0	1	0	0	0	1	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	0	d	d	d	d	d	Rdd=cmpyiw(Rss,Rtt*)
1	1	1	0	1	0	0	1	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	0	d	d	d	d	d	Rd=cmpyiw(Rss,Rtt*):<<1:sat
1	1	1	0	1	0	0	1	0	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	d	d	d	d	d	Rd=cmpyiw(Rss,Rtt):<<1:sat
1	1	1	0	1	0	0	1	0	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	d	d	d	d	d	Rd=cmpyrw(Rss,Rtt):<<1:sat
1	1	1	0	1	0	0	1	0	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	d	d	d	d	d	Rd=cmpyrw(Rss,Rtt*):<<1:sat
1	1	1	0	1	0	0	1	1	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	0	d	d	d	d	d	Rd=cmpyiw(Rss,Rtt*):<<1:rnd:sat

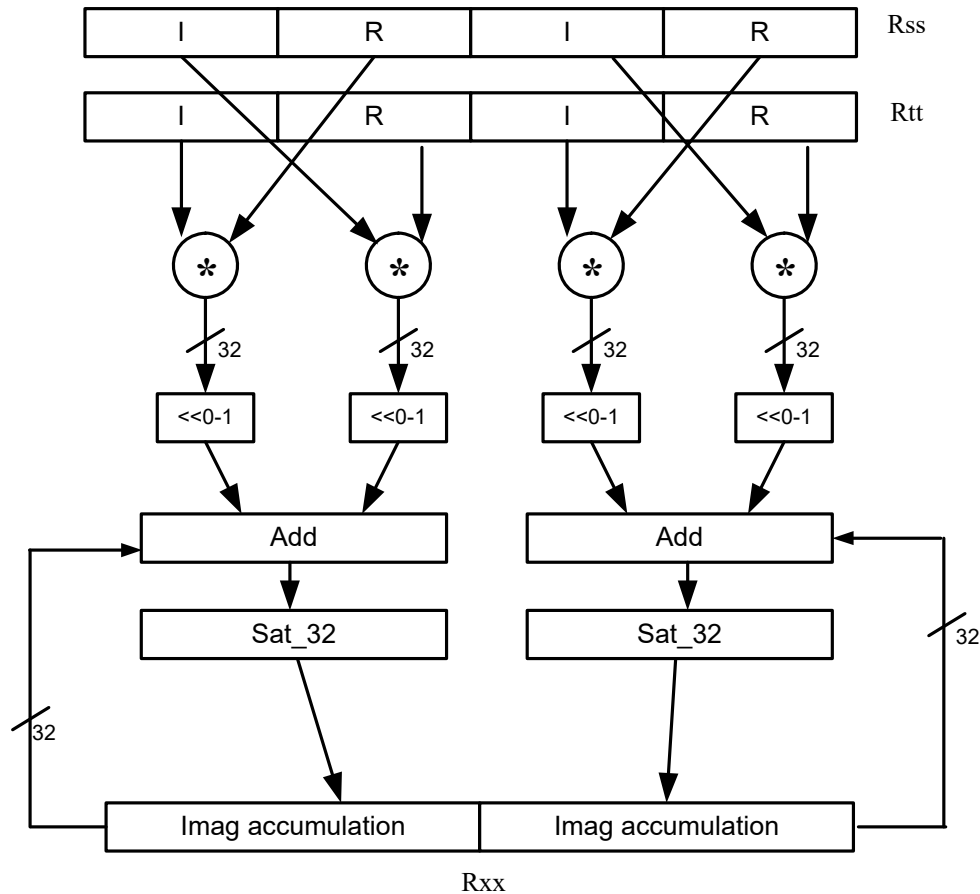
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
1	1	1	0	1	0	0	1	1	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	d	d	d	d	d	Rd=cmpyiw(Rss,Rtt):<<1:rnd:sat
1	1	1	0	1	0	0	1	1	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	d	d	d	d	d	Rd=cmpyrw(Rss,Rtt):<<1:rnd:sat
1	1	1	0	1	0	0	1	1	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	d	d	d	d	d	Rd=cmpyrw(Rss,Rtt*):<<1:rnd:sat
ICLASS			RegType			MajOp			s5					Parse		t5					MinOp		x5									
1	1	1	0	1	0	1	0	0	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	0	x	x	x	x	x	Rxx+=cmpyiw(Rss,Rtt*)
1	1	1	0	1	0	1	0	0	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	0	x	x	x	x	x	Rxx+=cmpyiw(Rss,Rtt)
1	1	1	0	1	0	1	0	1	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	0	x	x	x	x	x	Rxx+=cmpyrw(Rss,Rtt)
1	1	1	0	1	0	1	0	1	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	0	x	x	x	x	x	Rxx+=cmpyrw(Rss,Rtt*)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Vector complex multiply real or imaginary

The inputs *Rss* and *Rtt* are a vector of two complex values. Each complex value is composed of a 16-bit imaginary portion in the upper halfword and a 16-bit real portion in the lower halfword. Generate two complex results, either the real result or the imaginary result. These results are optionally shifted left by 0 to 1 bits, and optionally accumulated with the destination register.

### $Rxx += \text{vcmpyi}(Rss, Rtt) : \text{sat}$



#### Syntax

```
Rdd=vcmpyi (Rss,Rtt) [:<<1]:sat
```

```
Rdd=vcmpyr (Rss,Rtt) [:<<1]:sat
```

#### Behavior

```
Rdd.w[0]=sat32((Rss.h[1] * Rtt.h[0]) + (Rss.h[0] * Rtt.h[1])) [<<1];
Rdd.w[1]=sat32((Rss.h[3] * Rtt.h[2]) + (Rss.h[2] * Rtt.h[3])) [<<1];
```

```
Rdd.w[0]=sat32((Rss.h[0] * Rtt.h[0]) - (Rss.h[1] * Rtt.h[1])) [<<1];
Rdd.w[1]=sat32((Rss.h[2] * Rtt.h[2]) - (Rss.h[3] * Rtt.h[3])) [<<1];
```

Syntax	Behavior
$Rxx += vcmpyi(Rss, Rtt) : sat$	$Rxx.w[0] = sat_{32}(Rxx.w[0] + (Rss.h[1] * Rtt.h[0]) + (Rss.h[0] * Rtt.h[1]) << 0);$ $Rxx.w[1] = sat_{32}(Rxx.w[1] + (Rss.h[3] * Rtt.h[2]) + (Rss.h[2] * Rtt.h[3]) << 0);$
$Rxx += vcmpyr(Rss, Rtt) : sat$	$Rxx.w[0] = sat_{32}(Rxx.w[0] + (Rss.h[0] * Rtt.h[0]) - (Rss.h[1] * Rtt.h[1]) << 0);$ $Rxx.w[1] = sat_{32}(Rxx.w[1] + (Rss.h[2] * Rtt.h[2]) - (Rss.h[3] * Rtt.h[3]) << 0);$

### Class: XTYPE (slots 2,3)

#### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

#### Intrinsics

$Rdd = vcmpyi(Rss, Rtt) : <<1 : sat$	Word64 Q6_P_vcmpyi_PP_s1_sat(Word64 Rss, Word64 Rtt)
$Rdd = vcmpyi(Rss, Rtt) : sat$	Word64 Q6_P_vcmpyi_PP_sat(Word64 Rss, Word64 Rtt)
$Rdd = vcmpyr(Rss, Rtt) : <<1 : sat$	Word64 Q6_P_vcmpyr_PP_s1_sat(Word64 Rss, Word64 Rtt)
$Rdd = vcmpyr(Rss, Rtt) : sat$	Word64 Q6_P_vcmpyr_PP_sat(Word64 Rss, Word64 Rtt)
$Rxx += vcmpyi(Rss, Rtt) : sat$	Word64 Q6_P_vcmpyiacc_PP_sat(Word64 Rxx, Word64 Rss, Word64 Rtt)
$Rxx += vcmpyr(Rss, Rtt) : sat$	Word64 Q6_P_vcmpyracc_PP_sat(Word64 Rxx, Word64 Rss, Word64 Rtt)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse		t5					MinOp			d5								
1	1	1	0	1	0	0	0	N	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	0	d	d	d	d	d	Rdd=vcmpyr(Rss,Rtt)[:<<N]:sat
1	1	1	0	1	0	0	0	N	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	0	d	d	d	d	d	Rdd=vcmpyi(Rss,Rtt)[:<<N]:sat
ICLASS			RegType				MajOp		s5					Parse		t5					MinOp			x5								
1	1	1	0	1	0	1	0	0	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	0	x	x	x	x	x	Rxx+=vcmpyr(Rss,Rtt):sat
1	1	1	0	1	0	1	0	0	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	0	x	x	x	x	x	Rxx+=vcmpyi(Rss,Rtt):sat

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode



<b>Field name</b>	<b>Description</b>
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Vector complex conjugate

Perform a vector complex conjugate of both complex values in vector Rss by negating the imaginary halfwords, and placing the result in destination Rdd.

### Syntax

```
Rdd=vconj(Rss):sat
```

### Behavior

```
Rdd.h[1]=sat16(-Rss.h[1]);
Rdd.h[0]=Rss.h[0];
Rdd.h[3]=sat16(-Rss.h[3]);
Rdd.h[2]=Rss.h[2];
```

### Class: XTYPE (slots 2,3)

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

### Intrinsics

```
Rdd=vconj(Rss):sat
```

```
Word64 Q6_P_vconj_P_sat(Word64 Rss)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp		s5					Parse		MinOp			d5												
1	0	0	0	0	0	0	0	1	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	1	1	d	d	d	d	d	Rdd=vconj(Rss):sat

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Vector complex rotate

Take the least significant bits of Rt, and use these bits to rotate each of the two complex values in the source vector a multiple of 90 degrees. Bits 0 and 1 control the rotation factor for word 0, and bits 2 and 3 control the rotation factor for word 1.

If the rotation control bits are 0, the rotation is 0: the real and imaginary halves of the source appear unchanged and unmoved in the destination.

If the rotation control bits are 1, the rotation is  $-\pi/2$ : the real half of the destination gets the imaginary half of the source, and the imaginary half of the destination gets the negative real half of the source.

If the rotation control bits are 2, the rotation is  $\pi/2$ : the real half of the destination gets the negative imaginary half of the source, and the imaginary half of the destination gets the real half of the source.

If the rotation control bits are 3, the rotation is  $\pi$ : the real half of the destination gets the negative real half of the source, and the imaginary half of the destination gets the negative imaginary half of the source.

Syntax	Behavior
<pre>Rdd=vcrotate (Rss,Rt)</pre>	<pre>tmp = Rt[1:0]; if (tmp == 0) {     Rdd.h[0]=Rss.h[0];     Rdd.h[1]=Rss.h[1]; } else if (tmp == 1) {     Rdd.h[0]=Rss.h[1];     Rdd.h[1]=sat<sub>16</sub>(-Rss.h[0]); } else if (tmp == 2) {     Rdd.h[0]=sat<sub>16</sub>(-Rss.h[1]);     Rdd.h[1]=Rss.h[0]; } else {     Rdd.h[0]=sat<sub>16</sub>(-Rss.h[0]);     Rdd.h[1]=sat<sub>16</sub>(-Rss.h[1]); } tmp = Rt[3:2]; if (tmp == 0) {     Rdd.h[2]=Rss.h[2];     Rdd.h[3]=Rss.h[3]; } else if (tmp == 1) {     Rdd.h[2]=Rss.h[3];     Rdd.h[3]=sat<sub>16</sub>(-Rss.h[2]); } else if (tmp == 2) {     Rdd.h[2]=sat<sub>16</sub>(-Rss.h[3]);     Rdd.h[3]=Rss.h[2]; } else {     Rdd.h[2]=sat<sub>16</sub>(-Rss.h[2]);     Rdd.h[3]=sat<sub>16</sub>(-Rss.h[3]); }</pre>

**Class: XTYPE (slots 2,3)****Notes**

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

**Intrinsics**

Rdd=vcrotate (Rss,Rt)

Word64 Q6\_P\_vcrotate\_PR (Word64 Rss, Word32 Rt)

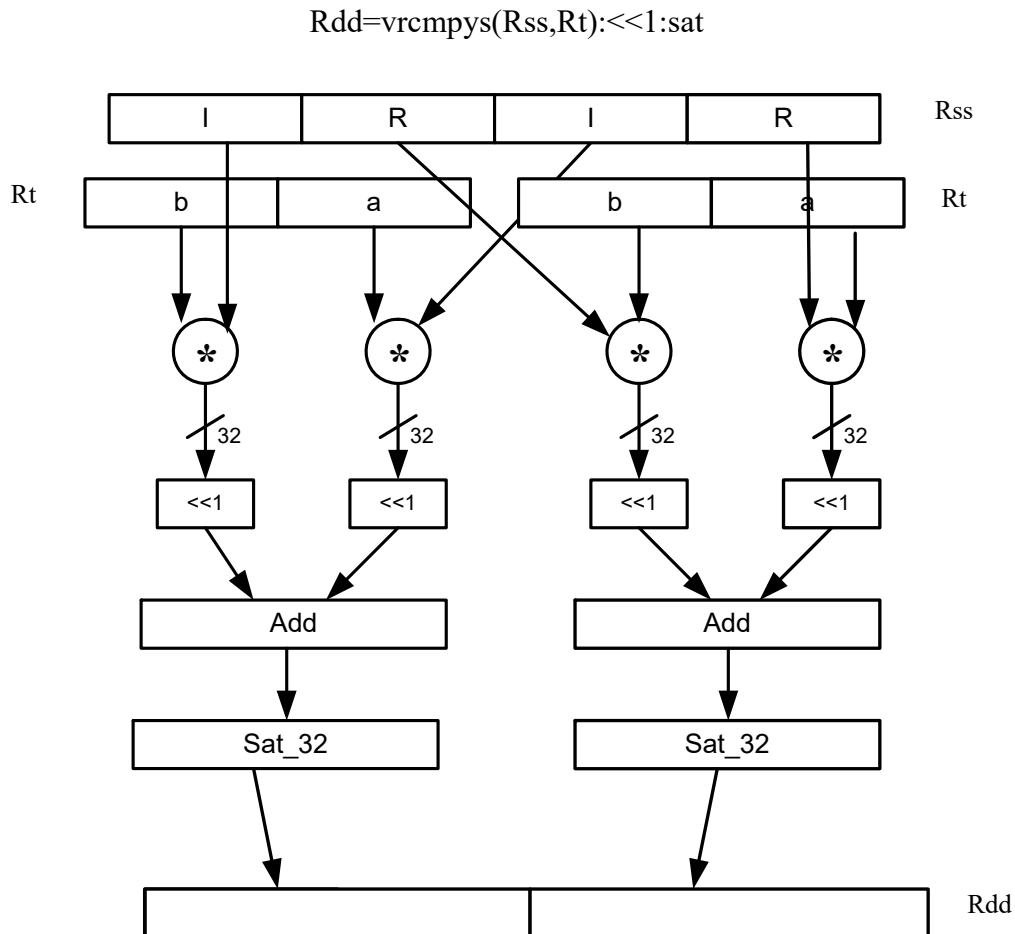
**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				Maj		s5					Parse		t5					Min		d5									
1	1	0	0	0	0	1	1	1	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	d	d	d	d	d	Rdd=vcrotate(Rss,Rt)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Vector reduce complex multiply by scalar

Multiply a complex number by a scalar. Rss contains two complex numbers. The real portions are each multiplied by two scalars contained in register Rt, scaled, summed, optionally accumulated, saturated, and stored in the lower word of Rdd. A similar operation is done on the two imaginary portions of Rss.



### Syntax

```
Rdd=vrcmpys(Rss,Rt):<<1:sat
```

```
Rdd=vrcmpys(Rss,Rtt):<<1:sat:raw:hi:hi
```

### Behavior

```
if ("Rt & 1") {
    Assembler mapped to:
    "Rdd=vrcmpys(Rss,Rtt):<<1:sat:raw:hi";
} else {
    Assembler mapped to:
    "Rdd=vrcmpys(Rss,Rtt):<<1:sat:raw:lo";
}
```

```
Rdd.w[1]=sat32((Rss.h[1] * Rtt.w[1].h[0])<<1 +
(Rss.h[3] * Rtt.w[1].h[1])<<1);
Rdd.w[0]=sat32((Rss.h[0] * Rtt.w[1].h[0])<<1 +
(Rss.h[2] * Rtt.w[1].h[1])<<1);
```

Syntax	Behavior
<code>Rdd=vrcmpys (Rss,Rtt) :&lt;&lt;1:sat:raw:lo</code>	<pre>Rdd.w[1]=sat<sub>32</sub>((Rss.h[1] * Rtt.w[0].h[0])&lt;&lt;1 + (Rss.h[3] * Rtt.w[0].h[1])&lt;&lt;1); Rdd.w[0]=sat<sub>32</sub>((Rss.h[0] * Rtt.w[0].h[0])&lt;&lt;1 + (Rss.h[2] * Rtt.w[0].h[1])&lt;&lt;1);</pre>
<code>Rxx+=vrcmpys (Rss,Rt) :&lt;&lt;1:sat</code>	<pre>if ("Rt &amp; 1") {   Assembler mapped to:   "Rxx+=vrcmpys (Rss,Rtt) :&lt;&lt;1:sat:raw:hi"; } else {   Assembler mapped to:   "Rxx+=vrcmpys (Rss,Rtt) :&lt;&lt;1:sat:raw:lo"; }</pre>
<code>Rxx+=vrcmpys (Rss,Rtt) :&lt;&lt;1:sat:raw:hi</code>	<pre>Rxx.w[1]=sat<sub>32</sub>(Rxx.w[1] + (Rss.h[1] * Rtt.w[1].h[0])&lt;&lt;1 + (Rss.h[3] * Rtt.w[1].h[1])&lt;&lt;1); Rxx.w[0]=sat<sub>32</sub>(Rxx.w[0] + (Rss.h[0] * Rtt.w[1].h[0])&lt;&lt;1 + (Rss.h[2] * Rtt.w[1].h[1])&lt;&lt;1);</pre>
<code>Rxx+=vrcmpys (Rss,Rtt) :&lt;&lt;1:sat:raw:lo</code>	<pre>Rxx.w[1]=sat<sub>32</sub>(Rxx.w[1] + (Rss.h[1] * Rtt.w[0].h[0])&lt;&lt;1 + (Rss.h[3] * Rtt.w[0].h[1])&lt;&lt;1); Rxx.w[0]=sat<sub>32</sub>(Rxx.w[0] + (Rss.h[0] * Rtt.w[0].h[0])&lt;&lt;1 + (Rss.h[2] * Rtt.w[0].h[1])&lt;&lt;1);</pre>

### Class: XTYPE (slots 2,3)

#### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

#### Intrinsics

`Rdd=vrcmpys (Rss,Rt) :<<1:sat` Word64 Q6\_P\_vrcmpys\_PR\_s1\_sat (Word64 Rss, Word32 Rt)  
at

`Rxx+=vrcmpys (Rss,Rt) :<<1:sat` Word64 Q6\_P\_vrcmpysacc\_PR\_s1\_sat (Word64 Rxx, Word64 Rss, Word32 Rt)

## Encoding

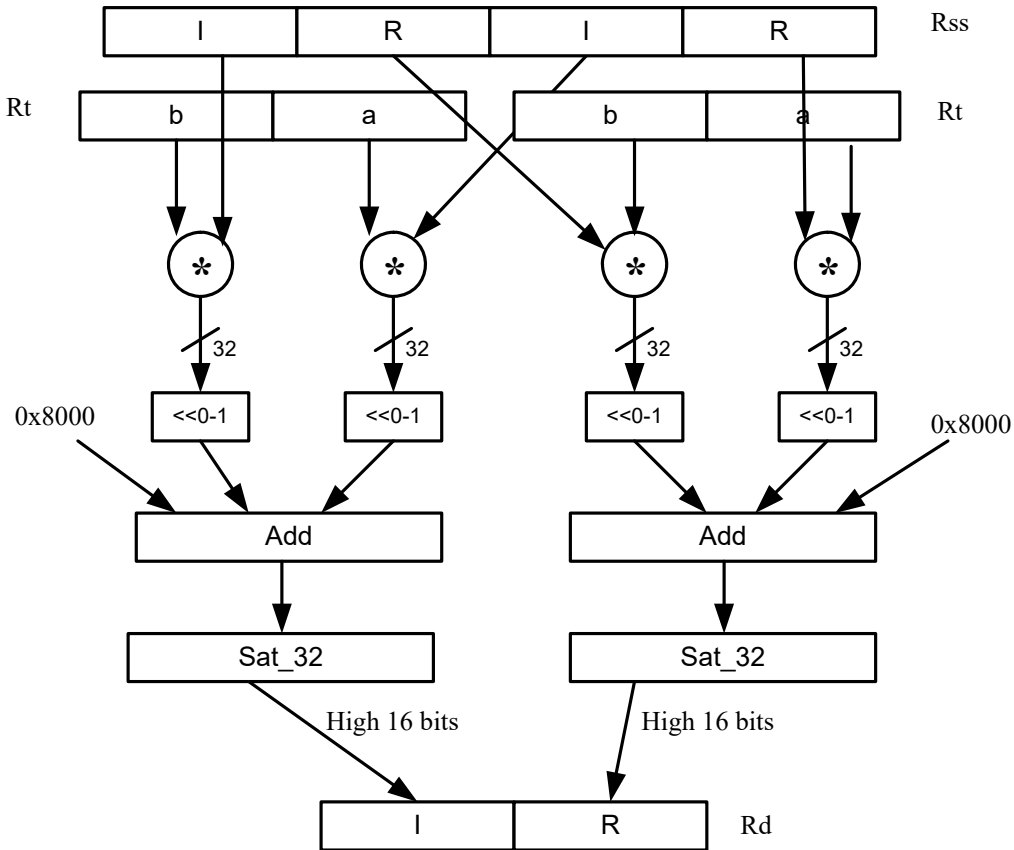
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			d5							
1	1	1	0	1	0	0	0	1	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	0	d	d	d	d	d	Rdd=vrcmpys(Rss,Rtt):<<1:sat:raw:hi
1	1	1	0	1	0	0	0	1	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	0	d	d	d	d	d	Rdd=vrcmpys(Rss,Rtt):<<1:sat:raw:lo
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			x5							
1	1	1	0	1	0	1	0	1	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	0	x	x	x	x	x	Rxx+=vrcmpys(Rss,Rtt):<<1:sat:raw:hi
1	1	1	0	1	0	1	0	1	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	0	x	x	x	x	x	Rxx+=vrcmpys(Rss,Rtt):<<1:sat:raw:lo

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Vector reduce complex multiply by scalar with round and pack

Multiply a complex number by scalar. Rss contains two complex numbers. The real portions are each multiplied by two scalars contained in register Rt, scaled, summed, rounded, and saturated. The upper 16bits of this result are packed in the lower halfword of Rd. A similar operation is done on the two imaginary portions of Rss.

$$Rd = \text{vrcmpys}(Rss, Rt) : \ll 1 : \text{rnd} : \text{sat}$$



### Syntax

```
Rd=vrcmpys(Rss,Rt):<<1:rnd:sat
```

### Behavior

```
if ("Rt & 1") {
    Assembler mapped to:
    "Rd=vrcmpys(Rss,Rtt):<<1:rnd:sat:raw:hi";
} else {
    Assembler mapped to:
    "Rd=vrcmpys(Rss,Rtt):<<1:rnd:sat:raw:lo";
}
```

```
Rd=vrcmpys(Rss,Rtt):<<1:rnd:sat:raw:hi
```

```
Rd.h[1]=sat32((Rss.h[1] * Rtt.w[1].h[0])<<1 +
(Rss.h[3] * Rtt.w[1].h[1])<<1 + 0x8000).h[1];
Rd.h[0]=sat32((Rss.h[0] * Rtt.w[1].h[0])<<1 +
(Rss.h[2] * Rtt.w[1].h[1])<<1 + 0x8000).h[1];
```



**Syntax**

```
Rd=vrcmpys(Rss,Rtt):<<1:rnd:sat:raw:lo
```

**Behavior**

```
Rd.h[1]=sat32((Rss.h[1] * Rtt.w[0].h[0])<<1 +
(Rss.h[3] * Rtt.w[0].h[1])<<1 + 0x8000).h[1];
Rd.h[0]=sat32((Rss.h[0] * Rtt.w[0].h[0])<<1 +
(Rss.h[2] * Rtt.w[0].h[1])<<1 + 0x8000).h[1];
```

**Class: XTYPE (slots 2,3)****Notes**

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

**Intrinsics**

```
Rd=vrcmpys(Rss,Rt):<<1:rnd:sat Word32 Q6_R_vrcmpys_PR_s1_rnd_sat(Word64
at Rss, Word32 Rt)
```

**Encoding**

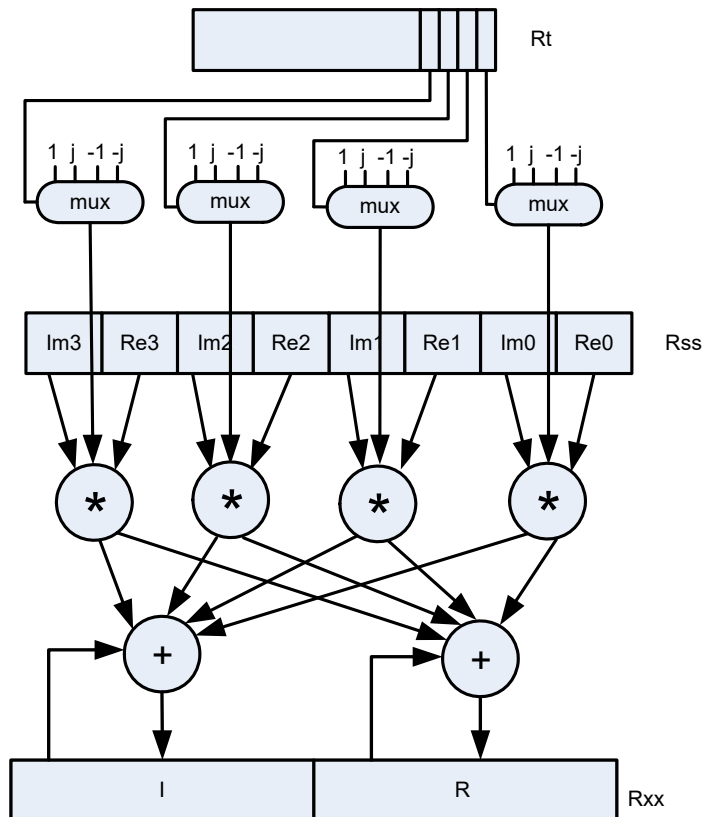
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp				s5					Parse		t5					MinOp			d5						
1	1	1	0	1	0	0	1	1	-	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	0	d	d	d	d	d	Rd=vrcmpys(Rss,Rtt):<<1:rnd:sat:raw:hi
1	1	1	0	1	0	0	1	1	-	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	1	d	d	d	d	d	Rd=vrcmpys(Rss,Rtt):<<1:rnd:sat:raw:lo

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector reduce complex rotate

This instruction is useful for CDMA despreding. An unsigned 2-bit immediate specifies a byte to use in  $R_t$ . Four 2-bit fields in the specified byte each select a rotation amount for one of the four complex numbers in  $R_{ss}$ . The real and imaginary products accumulate and store as a 32-bit complex number in  $R_d$ . Optionally, the destination register can also accumulate.

$R_{xx} += \text{vrcrotate}(R_{ss}, R_t, \#0)$



Syntax	Behavior
<pre>Rdd=vrcrotate (Rss,Rt,#u2)</pre>	<pre>sumr = 0; sumi = 0; control = Rt.ub[#u]; for (i = 0; i &lt; 8; i += 2) {     tmpr = Rss.b[i];     tmpi = Rss.b[i+1];     switch (control &amp; 3) {         case 0: sumr += tmpr;             sumi += tmpi;             break;         case 1: sumr += tmpi;             sumi -= tmpr;             break;         case 2: sumr -= tmpr;             sumi += tmpi;             break;         case 3: sumr -= tmpi;             sumi -= tmpr;             break;     }     control = control &gt;&gt; 2; } Rdd.w[0]=sumr; Rdd.w[1]=sumi;</pre>
<pre>Rxx+=vrcrotate (Rss,Rt,#u2)</pre>	<pre>sumr = 0; sumi = 0; control = Rt.ub[#u]; for (i = 0; i &lt; 8; i += 2) {     tmpr = Rss.b[i];     tmpi = Rss.b[i+1];     switch (control &amp; 3) {         case 0: sumr += tmpr;             sumi += tmpi;             break;         case 1: sumr += tmpi;             sumi -= tmpr;             break;         case 2: sumr -= tmpr;             sumi += tmpi;             break;         case 3: sumr -= tmpi;             sumi -= tmpr;             break;     }     control = control &gt;&gt; 2; } Rxx.w[0]=Rxx.w[0] + sumr; Rxx.w[1]=Rxx.w[1] + sumi;</pre>

**Class: XTYPE (slots 2,3)****Intrinsics**

```
Rdd=vrcrotate(Rss,Rt,#u2) Word64 Q6_P_vrcrotate_PRI(Word64 Rss, Word32
)                               Rt, Word32 Iu2)
```

```
Rxx+=vrcrotate(Rss,Rt,#u Word64 Q6_P_vrcrotateacc_PRI(Word64 Rxx, Word64
2)                               Rss, Word32 Rt, Word32 Iu2)
```

**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				Maj		s5					Parse		t5					Min		d5									
1	1	0	0	0	0	1	1	1	1	-	s	s	s	s	s	P	P	i	t	t	t	t	t	1	1	i	d	d	d	d	d	Rdd=vrcrotate(Rss,Rt,#u2)
ICLASS			RegType				Maj		s5					Parse		t5					x5											
1	1	0	0	1	0	1	1	1	0	1	s	s	s	s	s	P	P	i	t	t	t	t	t	-	-	i	x	x	x	x	x	Rxx+=vrcrotate(Rss,Rt,#u2)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## 11.10.4 XTYPE FP

The XTYPE FP instruction subclass includes instructions for floating point math.

### Floating point addition

Add two floating point values.

Syntax	Behavior
$Rd = sfadd(Rs, Rt)$	$Rd = Rs + Rt;$
$Rdd = dfadd(Rss, Rtt)$	$Rdd = Rss + Rtt;$

### Class: XTYPE (slots 2,3)

#### Intrinsics

$Rd = sfadd(Rs, Rt)$	Word32 Q6_R_sfadd_RR(Word32 Rs, Word32 Rt)
$Rdd = dfadd(Rss, Rtt)$	Word64 Q6_P_dfadd_PP(Word64 Rss, Word64 Rtt)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0				
ICLASS				RegType				MajOp				s5					Parse		t5					MinOp			d5								
1	1	1	0	1	0	0	0	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	1	d	d	d	d	d	Rdd=dfadd(Rss,Rtt)			
1	1	1	0	1	0	1	1	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	d	d	d	d	d	Rd=sfadd(Rs,Rt)			

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Classify floating point value

Classify floating point values. Classes are normal, subnormal, zero, NaN, or infinity. When the number is one of the specified classes, return true.

Syntax	Behavior
<code>Pd=dfclass(Rss,#u5)</code>	<pre>Pd = 0; class = fpclassify(Rss); if (#u.0 &amp;&amp; (class == FP_ZERO)) Pd = 0xff; if (#u.1 &amp;&amp; (class == FP_NORMAL)) Pd = 0xff; if (#u.2 &amp;&amp; (class == FP_SUBNORMAL)) Pd = 0xff; if (#u.3 &amp;&amp; (class == FP_INFINITE)) Pd = 0xff; if (#u.4 &amp;&amp; (class == FP_NAN)) Pd = 0xff; cancel_flags();</pre>
<code>Pd=sfclass(Rs,#u5)</code>	<pre>Pd = 0; class = fpclassify(Rs); if (#u.0 &amp;&amp; (class == FP_ZERO)) Pd = 0xff; if (#u.1 &amp;&amp; (class == FP_NORMAL)) Pd = 0xff; if (#u.2 &amp;&amp; (class == FP_SUBNORMAL)) Pd = 0xff; if (#u.3 &amp;&amp; (class == FP_INFINITE)) Pd = 0xff; if (#u.4 &amp;&amp; (class == FP_NAN)) Pd = 0xff; cancel_flags();</pre>

### Class: XTYPE (slots 2,3)

#### Intrinsics

<code>Pd=dfclass(Rss,#u5)</code>	Byte Q6_p_dfclass_PI(Word64 Rss, Word32 Iu5)
<code>Pd=sfclass(Rs,#u5)</code>	Byte Q6_p_sfclass_RI(Word32 Rs, Word32 Iu5)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse					d2												
1	0	0	0	0	1	0	1	1	1	1	s	s	s	s	s	P	P	0	i	i	i	i	i	-	-	-	-	-	-	d	d	Pd=sfclass(Rs,#u5)
ICLASS			RegType				s5					Parse					d2															
1	1	0	1	1	1	0	0	1	0	0	s	s	s	s	s	P	P	-	0	0	0	i	i	i	i	i	1	0	-	d	d	Pd=dfclass(Rss,#u5)

Field name	Description
RegType	Register type
MajOp	Major opcode
ICLASS	Instruction class

<b>Field name</b>	<b>Description</b>
Parse	Packet/loop parse bits
d2	Field to encode register d
s5	Field to encode register s

## Compare floating point value

Compare floating point values. p0 returns true when at least one value is a NaN, otherwise p0 returns zero.

Syntax	Behavior
Pd=dfcmp.eq(Rss,Rtt)	Pd=Rss==Rtt ? 0xff : 0x00;
Pd=dfcmp.ge(Rss,Rtt)	Pd=Rss>=Rtt ? 0xff : 0x00;
Pd=dfcmp.gt(Rss,Rtt)	Pd=Rss>Rtt ? 0xff : 0x00;
Pd=dfcmp.uo(Rss,Rtt)	Pd=isunordered(Rss,Rtt) ? 0xff : 0x00;
Pd=sfcmp.eq(Rs,Rt)	Pd=Rs==Rt ? 0xff : 0x00;
Pd=sfcmp.ge(Rs,Rt)	Pd=Rs>=Rt ? 0xff : 0x00;
Pd=sfcmp.gt(Rs,Rt)	Pd=Rs>Rt ? 0xff : 0x00;
Pd=sfcmp.uo(Rs,Rt)	Pd=isunordered(Rs,Rt) ? 0xff : 0x00;

### Class: XTYPE (slots 2,3)

#### Intrinsics

Pd=dfcmp.eq(Rss,Rtt)	Byte Q6_p_dfcmp_eq_PP(Word64 Rss, Word64 Rtt)
Pd=dfcmp.ge(Rss,Rtt)	Byte Q6_p_dfcmp_ge_PP(Word64 Rss, Word64 Rtt)
Pd=dfcmp.gt(Rss,Rtt)	Byte Q6_p_dfcmp_gt_PP(Word64 Rss, Word64 Rtt)
Pd=dfcmp.uo(Rss,Rtt)	Byte Q6_p_dfcmp_uo_PP(Word64 Rss, Word64 Rtt)
Pd=sfcmp.eq(Rs,Rt)	Byte Q6_p_sfcmp_eq_RR(Word32 Rs, Word32 Rt)
Pd=sfcmp.ge(Rs,Rt)	Byte Q6_p_sfcmp_ge_RR(Word32 Rs, Word32 Rt)
Pd=sfcmp.gt(Rs,Rt)	Byte Q6_p_sfcmp_gt_RR(Word32 Rs, Word32 Rt)
Pd=sfcmp.uo(Rs,Rt)	Byte Q6_p_sfcmp_uo_RR(Word32 Rs, Word32 Rt)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				Maj				s5				Parse		t5				Min		d2								
1	1	0	0	0	1	1	1	1	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	-	-	-	d	d	Pd=sfcmp.ge(Rs,Rt)
1	1	0	0	0	1	1	1	1	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	-	-	-	d	d	Pd=sfcmp.uo(Rs,Rt)
1	1	0	0	0	1	1	1	1	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	-	-	-	d	d	Pd=sfcmp.eq(Rs,Rt)
1	1	0	0	0	1	1	1	1	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	0	-	-	-	d	d	Pd=sfcmp.gt(Rs,Rt)



31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType					s5					Parse		t5					MinOp			d2									
1	1	0	1	0	0	1	0	1	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	-	-	-	d	d	Pd=dfcmp.eq(Rss,Rtt)
1	1	0	1	0	0	1	0	1	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	-	-	-	d	d	Pd=dfcmp.gt(Rss,Rtt)
1	1	0	1	0	0	1	0	1	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	-	-	-	d	d	Pd=dfcmp.ge(Rss,Rtt)
1	1	0	1	0	0	1	0	1	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	-	-	-	d	d	Pd=dfcmp.uo(Rss,Rtt)

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d2	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
Maj	Major opcode
Min	Minor opcode

## Convert floating point value to other format

Convert floating point values. When rounding is required, it occurs according to the rounding mode.

Syntax	Behavior
<code>Rd=convert_df2sf(Rss)</code>	<code>Rd = conv_df_to_sf(Rss);</code>
<code>Rdd=convert_sf2df(Rs)</code>	<code>Rdd = conv_sf_to_df(Rs);</code>

### Class: XTYPE (slots 2,3)

#### Intrinsics

`Rd=convert_df2sf(Rss)`      `Word32 Q6_R_convert_df2sf_P(Word64 Rss)`

`Rdd=convert_sf2df(Rs)`      `Word64 Q6_P_convert_sf2df_R(Word32 Rs)`

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse		MinOp			d5													
1	0	0	0	0	1	0	0	1	-	-	s	s	s	s	s	P	P	-	-	-	-	-	0	0	0	d	d	d	d	d		Rdd=convert_sf2df(Rs)
1	0	0	0	1	0	0	0	0	0	0	s	s	s	s	s	P	P	-	-	-	-	-	0	0	1	d	d	d	d	d		Rd=convert_df2sf(Rss)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Convert integer to floating point value

Convert floating point values. When rounding is required, it occurs according to the rounding mode unless the :chop option is specified.

Syntax	Behavior
Rd=convert_d2sf(Rss)	Rd = conv_8s_to_sf(Rss.s64);
Rd=convert_ud2sf(Rss)	Rd = conv_8u_to_sf(Rss.u64);
Rd=convert_uw2sf(Rs)	Rd = conv_4u_to_sf(Rs.uw[0]);
Rd=convert_w2sf(Rs)	Rd = conv_4s_to_sf(Rs.s32);
Rdd=convert_d2df(Rss)	Rdd = conv_8s_to_df(Rss.s64);
Rdd=convert_ud2df(Rss)	Rdd = conv_8u_to_df(Rss.u64);
Rdd=convert_uw2df(Rs)	Rdd = conv_4u_to_df(Rs.uw[0]);
Rdd=convert_w2df(Rs)	Rdd = conv_4s_to_df(Rs.s32);

### Class: XTYPE (slots 2,3)

#### Intrinsics

Rd=convert_d2sf(Rss)	Word32 Q6_R_convert_d2sf_P(Word64 Rss)
Rd=convert_ud2sf(Rss)	Word32 Q6_R_convert_ud2sf_P(Word64 Rss)
Rd=convert_uw2sf(Rs)	Word32 Q6_R_convert_uw2sf_R(Word32 Rs)
Rd=convert_w2sf(Rs)	Word32 Q6_R_convert_w2sf_R(Word32 Rs)
Rdd=convert_d2df(Rss)	Word64 Q6_P_convert_d2df_P(Word64 Rss)
Rdd=convert_ud2df(Rss)	Word64 Q6_P_convert_ud2df_P(Word64 Rss)
Rdd=convert_uw2df(Rs)	Word64 Q6_P_convert_uw2df_R(Word32 Rs)
Rdd=convert_w2df(Rs)	Word64 Q6_P_convert_w2df_R(Word32 Rs)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp				s5					Parse				MinOp				d5							
1	0	0	0	0	0	0	0	1	1	1	s	s	s	s	s	P	P	0	-	-	-	-	-	0	1	0	d	d	d	d	d	Rdd=convert_ud2df(Rss)
1	0	0	0	0	0	0	0	1	1	1	s	s	s	s	s	P	P	0	-	-	-	-	-	0	1	1	d	d	d	d	d	Rdd=convert_d2df(Rss)
1	0	0	0	0	1	0	0	1	-	-	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	1	d	d	d	d	d	Rdd=convert_uw2df(Rs)
1	0	0	0	0	1	0	0	1	-	-	s	s	s	s	s	P	P	-	-	-	-	-	-	0	1	0	d	d	d	d	d	Rdd=convert_w2df(Rs)
1	0	0	0	1	0	0	0	0	0	1	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	1	d	d	d	d	d	Rd=convert_ud2sf(Rss)
1	0	0	0	1	0	0	0	0	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	1	d	d	d	d	d	Rd=convert_d2sf(Rss)
1	0	0	0	1	0	1	1	0	0	1	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	0	d	d	d	d	d	Rd=convert_uw2sf(Rs)
1	0	0	0	1	0	1	1	0	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	0	d	d	d	d	d	Rd=convert_w2sf(Rs)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d

<b>Field name</b>	<b>Description</b>
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Convert floating point value to integer

Convert floating point values. When rounding is required, it occurs according to the rounding mode unless the `:chop` option is specified.

If the value is out of range of the destination integer type, the invalid flag is raised and closest integer is chosen, including for infinite inputs. For NaN inputs, the invalid flag is also raised, and the output value is implementation defined.

Syntax	Behavior
<code>Rd=convert_df2uw(Rss)</code>	<code>Rd = conv_df_to_4u(Rss).uw[0];</code>
<code>Rd=convert_df2uw(Rss):chop</code>	<code>round_to_zero();</code> <code>Rd = conv_df_to_4u(Rss).uw[0];</code>
<code>Rd=convert_df2w(Rss)</code>	<code>Rd = conv_df_to_4s(Rss).s32;</code>
<code>Rd=convert_df2w(Rss):chop</code>	<code>round_to_zero();</code> <code>Rd = conv_df_to_4s(Rss).s32;</code>
<code>Rd=convert_sf2uw(Rs)</code>	<code>Rd = conv_sf_to_4u(Rs).uw[0];</code>
<code>Rd=convert_sf2uw(Rs):chop</code>	<code>round_to_zero();</code> <code>Rd = conv_sf_to_4u(Rs).uw[0];</code>
<code>Rd=convert_sf2w(Rs)</code>	<code>Rd = conv_sf_to_4s(Rs).s32;</code>
<code>Rd=convert_sf2w(Rs):chop</code>	<code>round_to_zero();</code> <code>Rd = conv_sf_to_4s(Rs).s32;</code>
<code>Rdd=convert_df2d(Rss)</code>	<code>Rdd = conv_df_to_8s(Rss).s64;</code>
<code>Rdd=convert_df2d(Rss):chop</code>	<code>round_to_zero();</code> <code>Rdd = conv_df_to_8s(Rss).s64;</code>
<code>Rdd=convert_df2ud(Rss)</code>	<code>Rdd = conv_df_to_8u(Rss).u64;</code>
<code>Rdd=convert_df2ud(Rss):chop</code>	<code>round_to_zero();</code> <code>Rdd = conv_df_to_8u(Rss).u64;</code>
<code>Rdd=convert_sf2d(Rs)</code>	<code>Rdd = conv_sf_to_8s(Rs).s64;</code>
<code>Rdd=convert_sf2d(Rs):chop</code>	<code>round_to_zero();</code> <code>Rdd = conv_sf_to_8s(Rs).s64;</code>
<code>Rdd=convert_sf2ud(Rs)</code>	<code>Rdd = conv_sf_to_8u(Rs).u64;</code>
<code>Rdd=convert_sf2ud(Rs):chop</code>	<code>round_to_zero();</code> <code>Rdd = conv_sf_to_8u(Rs).u64;</code>

**Class: XTYPE (slots 2,3)****Intrinsics**

Rd=convert_df2uw(Rss)	Word32 Q6_R_convert_df2uw_P(Word64 Rss)
Rd=convert_df2uw(Rss):chop	Word32 Q6_R_convert_df2uw_P_chop(Word64 Rss)
Rd=convert_df2w(Rss)	Word32 Q6_R_convert_df2w_P(Word64 Rss)
Rd=convert_df2w(Rss):chop	Word32 Q6_R_convert_df2w_P_chop(Word64 Rss)
Rd=convert_sf2uw(Rs)	Word32 Q6_R_convert_sf2uw_R(Word32 Rs)
Rd=convert_sf2uw(Rs):chop	Word32 Q6_R_convert_sf2uw_R_chop(Word32 Rs)
Rd=convert_sf2w(Rs)	Word32 Q6_R_convert_sf2w_R(Word32 Rs)
Rd=convert_sf2w(Rs):chop	Word32 Q6_R_convert_sf2w_R_chop(Word32 Rs)
Rdd=convert_df2d(Rss)	Word64 Q6_P_convert_df2d_P(Word64 Rss)
Rdd=convert_df2d(Rss):chop	Word64 Q6_P_convert_df2d_P_chop(Word64 Rss)
Rdd=convert_df2ud(Rss)	Word64 Q6_P_convert_df2ud_P(Word64 Rss)
Rdd=convert_df2ud(Rss):chop	Word64 Q6_P_convert_df2ud_P_chop(Word64 Rss)
Rdd=convert_sf2d(Rs)	Word64 Q6_P_convert_sf2d_R(Word32 Rs)
Rdd=convert_sf2d(Rs):chop	Word64 Q6_P_convert_sf2d_R_chop(Word32 Rs)
Rdd=convert_sf2ud(Rs)	Word64 Q6_P_convert_sf2ud_R(Word32 Rs)
Rdd=convert_sf2ud(Rs):chop	Word64 Q6_P_convert_sf2ud_R_chop(Word32 Rs)

**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				MajOp		s5					Parse		MinOp			d5														
1	0	0	0	0	0	0	0	1	1	1	s	s	s	s	s	P	P	0	-	-	-	-	0	0	0	d	d	d	d	d	Rdd=convert_df2d(Rss)	
1	0	0	0	0	0	0	0	1	1	1	s	s	s	s	s	P	P	0	-	-	-	-	0	0	1	d	d	d	d	d	Rdd=convert_df2ud(Rss)	
1	0	0	0	0	0	0	0	1	1	1	s	s	s	s	s	P	P	0	-	-	-	-	1	1	0	d	d	d	d	d	Rdd=convert_df2d(Rss):chop	
1	0	0	0	0	0	0	0	1	1	1	s	s	s	s	s	P	P	0	-	-	-	-	1	1	1	d	d	d	d	d	Rdd=convert_df2ud(Rss):chop	
1	0	0	0	0	1	0	0	1	-	-	s	s	s	s	s	P	P	-	-	-	-	-	0	1	1	d	d	d	d	d	Rdd=convert_sf2ud(Rs)	
1	0	0	0	0	1	0	0	1	-	-	s	s	s	s	s	P	P	-	-	-	-	-	1	0	0	d	d	d	d	d	Rdd=convert_sf2d(Rs)	
1	0	0	0	0	1	0	0	1	-	-	s	s	s	s	s	P	P	-	-	-	-	-	1	0	1	d	d	d	d	d	Rdd=convert_sf2ud(Rs):chop	
1	0	0	0	0	1	0	0	1	-	-	s	s	s	s	s	P	P	-	-	-	-	-	0	1	1	d	d	d	d	d	Rdd=convert_sf2d(Rs):chop	
1	0	0	0	1	0	0	0	0	1	1	s	s	s	s	s	P	P	-	-	-	-	-	0	0	1	d	d	d	d	d	Rd=convert_df2uw(Rss)	
1	0	0	0	1	0	0	0	1	0	0	s	s	s	s	s	P	P	-	-	-	-	-	0	0	1	d	d	d	d	d	Rd=convert_df2w(Rss)	
1	0	0	0	1	0	0	0	1	0	1	s	s	s	s	s	P	P	-	-	-	-	-	0	0	1	d	d	d	d	d	Rd=convert_df2uw(Rss):chop	

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
1	0	0	0	1	0	0	0	1	1	1	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	1	d	d	d	d	d	Rd=convert_df2w(Rs):chop
1	0	0	0	1	0	1	1	0	1	1	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	0	0	d	d	d	d	Rd=convert_sf2uw(Rs)
1	0	0	0	1	0	1	1	0	1	1	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	1	d	d	d	d	d	Rd=convert_sf2uw(Rs):chop
1	0	0	0	1	0	1	1	1	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	0	d	d	d	d	d	Rd=convert_sf2w(Rs)
1	0	0	0	1	0	1	1	1	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	1	d	d	d	d	d	Rd=convert_sf2w(Rs):chop

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Floating point extreme value assistance

For divide and square root routines, certain values are problematic for the default routine. These instructions appropriately fix up the numerator (fixupn), denominator (fixupd), or radicand (fixupr) for proper calculations when combined with the divide or square root approximation instructions.

Syntax	Behavior
Rd=sffixupd(Rs,Rt)	(Rs,Rt,Rd,adjust)=recip_common(Rs,Rt); Rd = Rt;
Rd=sffixupn(Rs,Rt)	(Rs,Rt,Rd,adjust)=recip_common(Rs,Rt); Rd = Rs;
Rd=sffixupr(Rs)	(Rs,Rd,adjust)=invsqrt_common(Rs); Rd = Rs;

### Class: XTYPE (slots 2,3)

#### Intrinsics

Rd=sffixupd(Rs,Rt)	Word32 Q6_R_sffixupd_RR(Word32 Rs, Word32 Rt)
Rd=sffixupn(Rs,Rt)	Word32 Q6_R_sffixupn_RR(Word32 Rs, Word32 Rt)
Rd=sffixupr(Rs)	Word32 Q6_R_sffixupr_R(Word32 Rs)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType			MajOp			s5					Parse		MinOp			d5													
1	0	0	0	1	0	1	1	1	0	1	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	0	d	d	d	d	d	Rd=sffixupr(Rs)
ICLASS			RegType			MajOp			s5					Parse		t5			MinOp			d5										
1	1	1	0	1	0	1	1	1	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	d	d	d	d	d	Rd=sffixupn(Rs,Rt)
1	1	1	0	1	0	1	1	1	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	d	d	d	d	d	Rd=sffixupd(Rs,Rt)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t



## Floating point fused multiply-add

Multiply two values, and add to (or subtract from) the accumulator. Full intermediate precision is kept.

Syntax	Behavior
$Rx += sfmpy(Rs, Rt)$	$Rx = fmaf(Rs, Rt, Rx);$
$Rx -= sfmpy(Rs, Rt)$	$Rx = fmaf(-Rs, Rt, Rx);$
$Rxx += dfmpyhh(Rss, Rtt)$	$Rxx = Rss * Rtt$ with partial product $Rxx$ ;
$Rxx += dfmpylh(Rss, Rtt)$	$Rxx += (Rss.uw[0] * (0x00100000   zxt_{20 \rightarrow 64}(Rtt.uw[1]))) \ll 1;$

### Class: XTYPE (slots 2,3)

#### Intrinsics

$Rx += sfmpy(Rs, Rt)$	Word32 Q6_R_sfmpyacc_RR(Word32 Rx, Word32 Rs, Word32 Rt)
$Rx -= sfmpy(Rs, Rt)$	Word32 Q6_R_sfmpynac_RR(Word32 Rx, Word32 Rs, Word32 Rt)
$Rxx += dfmpyhh(Rss, Rtt)$	Word64 Q6_P_dfmpyhacc_PP(Word64 Rxx, Word64 Rss, Word64 Rtt)
$Rxx += dfmpylh(Rss, Rtt)$	Word64 Q6_P_dfmpylhacc_PP(Word64 Rxx, Word64 Rss, Word64 Rtt)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp				s5					Parse		t5					MinOp			x5						
1	1	1	0	1	0	1	0	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	1	x	x	x	x	x	Rxx+=dfmpylh(Rss,Rtt)
1	1	1	0	1	0	1	0	1	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	1	x	x	x	x	x	Rxx+=dfmpyhh(Rss,Rtt)
1	1	1	0	1	1	1	1	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	0	x	x	x	x	x	Rx+=sfmpy(Rs,Rt)
1	1	1	0	1	1	1	1	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	1	x	x	x	x	x	Rx-=sfmpy(Rs,Rt)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Floating point fused multiply-add with scaling

Multiply two values, and add to (or subtract from) the accumulator. Full intermediate precision is kept. Additionally, scale the output.

This instruction has special handling of corner cases. If a multiplicand source is zero and a NaN is not produced, the accumulator is unchanged; thus the sign of a zero accumulator does not change when the product is a true zero.

For single precision, the scaling factor is the predicate taken as a two's compliment number.

For double precision, the scaling factor is twice the predicate taken as a two's compliment number. The implementation can change denormal accumulator values to zero for positive scale factors.

### Syntax

```
Rx+=sfmpy(Rs,Rt,Pu):scale
```

### Behavior

```
PREDUSE_TIMING;
if (isnan(Rx) || isnan(Rs) || isnan(Rt)) Rx =
NaN;
tmp=fmaf(Rs,Rt,Rx) * 2**(Pu);
if (!(Rx == 0.0) && is_true_zero(Rs*Rt)) Rx =
tmp;
```

**Class: XTYPE (slots 2,3)**

### Intrinsics

```
Rx+=sfmpy(Rs,Rt,Pu):scale
```

```
Word32 Q6_R_sfmpyacc_RRp_scale(Word32 Rx,
Word32 Rs, Word32 Rt, Byte Pu)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse		t5					u2		x5									
1	1	1	0	1	1	1	1	0	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	u	u	x	x	x	x	x	Rx+=sfmpy(Rs,Rt,Pu):scale

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
s5	Field to encode register s
t5	Field to encode register t
u2	Field to encode register u
x5	Field to encode register x

## Floating point reciprocal square root approximation

Provides an approximation of the reciprocal square root of the radicand (Rs), if combined with the appropriate fixup instruction. Certain values (such as infinities or zeros) in the numerator or denominator can yield values that are not reciprocal approximations, but yield the correct answer when combined with fixup instructions and the appropriate routines.

For compatibility, exact results of these instructions cannot be relied on. The precision of the approximation for this architecture and later is at least 6.6 bits.

Syntax	Behavior
<code>Rd,Pe=sfinvsqrta(Rs)</code>	<pre> if ((Rs,Rd,adjust)=invsqrt_common(Rs)) {     Pe = adjust;     idx = (Rs &gt;&gt; 17) &amp; 0x7f;     mant = (invsqrt_lut[idx] &lt;&lt; 15);     exp = 127 - ((exponent(Rs) - 127) &gt;&gt; 1)     - 1;     Rd = -1**Rs.31 * 1.MANT * 2**(exp- BIAS); } </pre>

### Class: XTYPE (slots 2,3)

#### Notes

- This instruction provides a certain amount of accuracy. In future versions the accuracy may increase. For future compatibility, avoid dependence on exact values.
- The predicate generated by this instruction cannot be used as a .new predicate, nor can it be automatically ANDed with another predicate.

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0			
ICLASS				RegType			MajOp			s5					Parse		e2			d5														
1	0	0	0	1	0	1	1	1	1	1	s	s	s	s	s	P	P	-	-	-	-	-	-	-	0	e	e	d	d	d	d	d	d	Rd,Pe=sfinvsqrta(Rs)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
e2	Field to encode register e
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Floating point fused multiply-add for library routines

Multiply two values, and add to (or subtract from) the accumulator. Full intermediate precision is kept. This instruction has special handling of corner cases. Addition of infinities with opposite signs, or subtraction of infinities with like signs, is defined as (positive) zero. Rounding is always nearest-even, except for overflows to infinity, which round to maximal finite values. If a multiplicand source is zero and a NaN is not produced, the accumulator is unchanged; thus the sign of a zero accumulator does not change if the product is a true zero. Flags and exceptions do not generate.

Syntax	Behavior
<code>Rx+=sfmpy(Rs,Rt):lib</code>	<pre>round_to_nearest(); infminusinf = ((isinf(Rx)) &amp;&amp; (isinf(Rs*Rt)) &amp;&amp; (Rs ^ Rx ^ Rt.31 != 0)); infinp = (isinf(Rx)    (isinf(Rt))    (isinf(Rs))); if (isnan(Rx)    isnan(Rs)    isnan(Rt)) Rx = NaN; tmp=fmaf(Rs,Rt,Rx); if (!(Rx == 0.0) &amp;&amp; is_true_zero(Rs*Rt)) Rx = tmp; cancel_flags(); if (isinf(Rx) &amp;&amp; !infinp) Rx = Rx - 1; if (infminusinf) Rx = 0;</pre>
<code>Rx- =sfmpy(Rs,Rt):lib</code>	<pre>round_to_nearest(); infminusinf = ((isinf(Rx)) &amp;&amp; (isinf(Rs*Rt)) &amp;&amp; (Rs ^ Rx ^ Rt.31 == 0)); infinp = (isinf(Rx)    (isinf(Rt))    (isinf(Rs))); if (isnan(Rx)    isnan(Rs)    isnan(Rt)) Rx = NaN; tmp=fmaf(-Rs,Rt,Rx); if (!(Rx == 0.0) &amp;&amp; is_true_zero(Rs*Rt)) Rx = tmp; cancel_flags(); if (isinf(Rx) &amp;&amp; !infinp) Rx = Rx - 1; if (infminusinf) Rx = 0;</pre>

### Class: XTYPE (slots 2,3)

#### Intrinsics

`Rx+=sfmpy(Rs,Rt):lib` Word32 Q6\_R\_sfmpyacc\_RR\_lib(Word32 Rx, Word32 Rs, Word32 Rt)

`Rx-=sfmpy(Rs,Rt):lib` Word32 Q6\_R\_sfmpynac\_RR\_lib(Word32 Rx, Word32 Rs, Word32 Rt)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				MajOp				s5				Parse		t5				MinOp			x5									
1	1	1	0	1	1	1	1	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	0	x	x	x	x	x	Rx+=sfmpy(Rs,Rt):lib
1	1	1	0	1	1	1	1	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	1	x	x	x	x	x	Rx-=sfmpy(Rs,Rt):lib

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode

<b>Field name</b>	<b>Description</b>
RegType	Register type
Parse	Packet/loop parse bits
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Create floating-point constant

Using ten bits of immediate, form a floating-point constant.

Syntax	Behavior
<code>Rd=sfmake(#u10):neg</code>	<code>Rd = (127 - 6) &lt;&lt; 23;</code> <code>Rd += (#u &lt;&lt; 17);</code> <code>Rd  = (1 &lt;&lt; 31);</code>
<code>Rd=sfmake(#u10):pos</code>	<code>Rd = (127 - 6) &lt;&lt; 23;</code> <code>Rd += #u &lt;&lt; 17;</code>
<code>Rdd=dfmake(#u10):neg</code>	<code>Rdd = (1023ULL - 6) &lt;&lt; 52;</code> <code>Rdd += (#u) &lt;&lt; 46;</code> <code>Rdd  = ((1ULL) &lt;&lt; 63);</code>
<code>Rdd=dfmake(#u10):pos</code>	<code>Rdd = (1023ULL - 6) &lt;&lt; 52;</code> <code>Rdd += (#u) &lt;&lt; 46;</code>

**Class: XTYPE (slots 2,3)**

### Intrinsics

<code>Rd=sfmake(#u10):neg</code>	<code>Word32 Q6_R_sfmake_I_neg(Word32 Iu10)</code>
<code>Rd=sfmake(#u10):pos</code>	<code>Word32 Q6_R_sfmake_I_pos(Word32 Iu10)</code>
<code>Rdd=dfmake(#u10):neg</code>	<code>Word64 Q6_P_dfmake_I_neg(Word32 Iu10)</code>
<code>Rdd=dfmake(#u10):pos</code>	<code>Word64 Q6_P_dfmake_I_pos(Word32 Iu10)</code>

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType														Parse				MinOp				d5								
1	1	0	1	0	1	1	0	0	0	i	-	-	-	-	-	P	P	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=sfmake(#u10):pos
1	1	0	1	0	1	1	0	0	1	i	-	-	-	-	-	P	P	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd=sfmake(#u10):neg
ICLASS		RegType														Parse								d5								
1	1	0	1	1	0	0	1	0	0	i	-	-	-	-	-	P	P	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rdd=dfmake(#u10):pos
1	1	0	1	1	0	0	1	0	1	i	-	-	-	-	-	P	P	i	i	i	i	i	i	i	i	i	d	d	d	d	d	Rdd=dfmake(#u10):neg

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d

## Floating point maximum

Maximum of two floating point values. If one value is a NaN, the other is chosen.

Syntax	Behavior
<code>Rd=sfmax(Rs,Rt)</code>	<code>Rd = fmaxf(Rs,Rt);</code>
<code>Rdd=dfmax(Rss,Rtt)</code>	<code>Rdd = fmax(Rss,Rtt);</code>

### Class: XTYPE (slots 2,3)

#### Intrinsics

<code>Rd=sfmax(Rs,Rt)</code>	<code>Word32 Q6_R_sfmax_RR(Word32 Rs, Word32 Rt)</code>
<code>Rdd=dfmax(Rss,Rtt)</code>	<code>Word64 Q6_P_dfmax_PP(Word64 Rss, Word64 Rtt)</code>

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			d5							
1	1	1	0	1	0	0	0	0	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	1	d	d	d	d	d	Rdd=dfmax(Rss,Rtt)
1	1	1	0	1	0	1	1	1	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	d	d	d	d	d	Rd=sfmax(Rs,Rt)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Floating point minimum

Minimum of two floating point values. If one value is a NaN, the other is chosen.

Syntax	Behavior
<code>Rd=sfmin(Rs,Rt)</code>	<code>Rd = fmin(Rs,Rt);</code>
<code>Rdd=dfmin(Rss,Rtt)</code>	<code>Rdd = fmin(Rss,Rtt);</code>

**Class: XTYPE (slots 2,3)**

### Intrinsics

<code>Rd=sfmin(Rs,Rt)</code>	<code>Word32 Q6_R_sfmin_RR(Word32 Rs, Word32 Rt)</code>
<code>Rdd=dfmin(Rss,Rtt)</code>	<code>Word64 Q6_P_dfmin_PP(Word64 Rss, Word64 Rtt)</code>

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			d5							
1	1	1	0	1	0	0	0	1	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	1	d	d	d	d	d	Rdd=dfmin(Rss,Rtt)
1	1	1	0	1	0	1	1	1	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	d	d	d	d	d	Rd=sfmin(Rs,Rt)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t



## Floating point multiply

Multiply two floating point values.

Syntax	Behavior
<code>Rd=sfmpy(Rs,Rt)</code>	<code>Rd=Rs*Rt;</code>
<code>Rdd=dfmpyfix(Rss,Rtt)</code>	<pre>if (is_denormal(Rss) &amp;&amp; (df_exponent(Rtt) &gt;= 512) &amp;&amp; is_normal(Rtt)) Rdd = Rss * 0x1.0p52; else if (is_denormal(Rtt) &amp;&amp; (df_exponent(Rss) &gt;= 512) &amp;&amp; is_normal(Rss)) Rdd = Rss * 0x1.0p-52; else Rdd = Rss;</pre>
<code>Rdd=dfmpyll(Rss,Rtt)</code>	<pre>prod = (Rss.uw[0] * Rtt.uw[0]); Rdd = (prod &gt;&gt; 32) &lt;&lt; 1; if (prod.uw[0] != 0) Rdd.0 = 1;</pre>

### Class: XTYPE (slots 2,3)

#### Intrinsics

<code>Rd=sfmpy(Rs,Rt)</code>	Word32 Q6_R_sfmpy_RR(Word32 Rs, Word32 Rt)
<code>Rdd=dfmpyfix(Rss,Rtt)</code>	Word64 Q6_P_dfmpyfix_PP(Word64 Rss, Word64 Rtt)
<code>Rdd=dfmpyll(Rss,Rtt)</code>	Word64 Q6_P_dfmpyll_PP(Word64 Rss, Word64 Rtt)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp				s5					Parse		t5					MinOp			d5					
1	1	1	0	1	0	0	0	0	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	1	d	d	d	d	d	Rdd=dfmpyfix(Rss,Rtt)
1	1	1	0	1	0	0	0	1	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	1	d	d	d	d	d	Rdd=dfmpyll(Rss,Rtt)
1	1	1	0	1	0	1	1	0	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	d	d	d	d	d	Rd=sfmpy(Rs,Rt)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Floating point reciprocal approximation

Provides an approximation of the reciprocal of the denominator (Rt), if combined with the appropriate fixup instructions. Certain values (such as infinities or zeros) in the numerator or denominator may yield values that are not reciprocal approximations, but yield the correct answer when combined with fixup instructions and the appropriate routines.

For compatibility, exact results of these instructions cannot be relied on. The precision of the approximation for this architecture and later is at least 6.6 bits.

Syntax	Behavior
<code>Rd, Pe=sfrecipa (Rs, Rt)</code>	<pre> if ((Rs, Rt, Rd, adjust)=recip_common(Rs, Rt)) {     Pe = adjust;     idx = (Rt &gt;&gt; 16) &amp; 0x7f;     mant = (recip_lut[idx] &lt;&lt; 15)   1;     exp = 127 - (exponent(Rt) - 127) - 1;     Rd = -1**Rt.31 * 1.MANT * 2**(exp- BIAS); } </pre>

### Class: XTYPE (slots 2,3)

#### Notes

- This instruction provides a certain amount of accuracy. In future versions the accuracy may increase. For future compatibility, avoid dependence on exact values.
- The predicate generated by this instruction cannot be used as a .new predicate, nor can it be automatically ANDed with another predicate.

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp					s5					Parse		t5					e2		d5						
1	1	1	0	1	0	1	1	1	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	e	e	d	d	d	d	d	Rd, Pe=sfrecipa(Rs, Rt)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
e2	Field to encode register e
s5	Field to encode register s
t5	Field to encode register t

## Floating point subtraction

Subtract two floating point values.

Syntax	Behavior
$Rd = sfsub(Rs, Rt)$	$Rd = Rs - Rt;$
$Rdd = dfsb(Rss, Rtt)$	$Rdd = Rss - Rtt;$

**Class: XTYPE (slots 2,3)**

### Intrinsics

$Rd = sfsub(Rs, Rt)$	Word32 Q6_R_sfsub_RR(Word32 Rs, Word32 Rt)
$Rdd = dfsb(Rss, Rtt)$	Word64 Q6_P_dfsb_PP(Word64 Rss, Word64 Rtt)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			d5							
1	1	1	0	1	0	0	0	1	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	1	d	d	d	d	d	Rdd=dfsbsub(Rss,Rtt)
1	1	1	0	1	0	1	1	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	d	d	d	d	d	Rd=sfsub(Rs,Rt)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

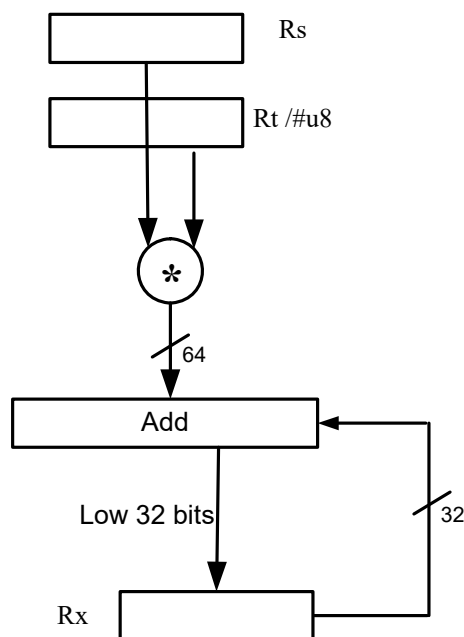
## 11.10.5 XTYPE MPY

The XTYPE MPY instruction subclass includes instructions which perform multiplication.

### Multiply and use lower result

Multiply the signed 32-bit integer in Rs by either the signed 32-bit integer in Rt or an unsigned immediate value. The 64-bit result optionally accumulates with the 32-bit destination, or is added to an immediate. The least-significant 32-bits of the result are written to the single destination register.

This multiply produces the correct results for the ANSI C multiplication of two signed or unsigned integers with an integer result.



#### Syntax

#### Behavior

<code>Rd=+mpyi (Rs, #u8)</code>	<code>apply_extension(#u);</code> <code>Rd=Rs*#u;</code>
<code>Rd=-mpyi (Rs, #u8)</code>	<code>Rd=Rs*-#u;</code>
<code>Rd=add(#u6,mpyi (Rs, #U6))</code>	<code>apply_extension(#u);</code> <code>Rd = #u + Rs*#U;</code>
<code>Rd=add(#u6,mpyi (Rs, Rt))</code>	<code>apply_extension(#u);</code> <code>Rd = #u + Rs*Rt;</code>
<code>Rd=add(Ru,mpyi (#u6:2, Rs))</code>	<code>Rd = Ru + Rs*#u;</code>
<code>Rd=add(Ru,mpyi (Rs, #u6))</code>	<code>apply_extension(#u);</code> <code>Rd = Ru + Rs*#u;</code>

Syntax	Behavior
<code>Rd=mpyi (Rs, #m9)</code>	<pre>if ("((#m9&lt;0) &amp;&amp; (#m9&gt;-256))") {     Assembler mapped to: "Rd=-mpyi (Rs, #m9*(-1))"; } else {     Assembler mapped to: "Rd+=mpyi (Rs, #m9)"; }</pre>
<code>Rd=mpyi (Rs, Rt)</code>	<code>Rd=Rs*Rt;</code>
<code>Rd=mpyui (Rs, Rt)</code>	Assembler mapped to: <code>"Rd=mpyi (Rs, Rt) "</code>
<code>Rx+=mpyi (Rs, #u8)</code>	<pre>apply_extension(#u); Rx=Rx + (Rs*#u);</pre>
<code>Rx+=mpyi (Rs, Rt)</code>	<code>Rx=Rx + Rs*Rt;</code>
<code>Rx-=mpyi (Rs, #u8)</code>	<pre>apply_extension(#u); Rx=Rx - (Rs*#u);</pre>
<code>Rx-=mpyi (Rs, Rt)</code>	<code>Rx=Rx - Rs*Rt;</code>
<code>Ry=add (Ru, mpyi (Ry, Rs))</code>	<code>Ry = Ru + Rs*Ry;</code>

**Class: XTYPE (slots 2,3)****Intrinsics**

Rd=add(#u6,mpyi(Rs,#U6))	Word32 Q6_R_add_mpyi_IRI(Word32 Iu6, Word32 Rs, Word32 IU6)
Rd=add(#u6,mpyi(Rs,Rt))	Word32 Q6_R_add_mpyi_IRR(Word32 Iu6, Word32 Rs, Word32 Rt)
Rd=add(Ru,mpyi(#u6:2,Rs))	Word32 Q6_R_add_mpyi_RIR(Word32 Ru, Word32 Iu6_2, Word32 Rs)
Rd=add(Ru,mpyi(Rs,#u6))	Word32 Q6_R_add_mpyi_RRI(Word32 Ru, Word32 Rs, Word32 Iu6)
Rd=mpyi(Rs,#m9)	Word32 Q6_R_mpyi_RI(Word32 Rs, Word32 Im9)
Rd=mpyi(Rs,Rt)	Word32 Q6_R_mpyi_RR(Word32 Rs, Word32 Rt)
Rd=mpyui(Rs,Rt)	Word32 Q6_R_mpyui_RR(Word32 Rs, Word32 Rt)
Rx+=mpyi(Rs,#u8)	Word32 Q6_R_mpyiacc_RI(Word32 Rx, Word32 Rs, Word32 Iu8)
Rx+=mpyi(Rs,Rt)	Word32 Q6_R_mpyiacc_RR(Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpyi(Rs,#u8)	Word32 Q6_R_mpyinac_RI(Word32 Rx, Word32 Rs, Word32 Iu8)
Rx-=mpyi(Rs,Rt)	Word32 Q6_R_mpyinac_RR(Word32 Rx, Word32 Rs, Word32 Rt)
Ry=add(Ru,mpyi(Ry,Rs))	Word32 Q6_R_add_mpyi_RRR(Word32 Ru, Word32 Ry, Word32 Rs)

**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType			s5					Parse		t5					MinOp		d5													
1	1	0	1	0	1	1	1	0	i	i	s	s	s	s	s	P	P	i	t	t	t	t	t	i	i	i	d	d	d	d	d	Rd=add(#u6,mpyi(Rs,Rt))
ICLASS		RegType			s5					Parse		d5																				
1	1	0	1	1	0	0	0	1	i	i	s	s	s	s	s	P	P	i	d	d	d	d	d	i	i	i	l	l	l	l	l	Rd=add(#u6,mpyi(Rs,#U6))
ICLASS		RegType			s5					Parse		d5					u5															
1	1	0	1	1	1	1	1	0	i	i	s	s	s	s	s	P	P	i	d	d	d	d	d	i	i	i	u	u	u	u	u	Rd=add(Ru,mpyi(#u6:2,Rs))
1	1	0	1	1	1	1	1	1	i	i	s	s	s	s	s	P	P	i	d	d	d	d	d	i	i	i	u	u	u	u	u	Rd=add(Ru,mpyi(Rs,#u6))
ICLASS		RegType			MajOp		s5					Parse		y5					u5													
1	1	1	0	0	0	1	1	0	0	0	s	s	s	s	s	P	P	-	y	y	y	y	y	-	-	-	u	u	u	u	u	Ry=add(Ru,mpyi(Ry,Rs))
ICLASS		RegType			MajOp		s5					Parse		MinOp					d5													
1	1	1	0	0	0	0	0	0	-	-	s	s	s	s	s	P	P	0	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd+=mpyi(Rs,#u8)
1	1	1	0	0	0	0	0	1	-	-	s	s	s	s	s	P	P	0	i	i	i	i	i	i	i	i	d	d	d	d	d	Rd-=mpyi(Rs,#u8)
ICLASS		RegType			MajOp		s5					Parse		MinOp					x5													
1	1	1	0	0	0	0	1	0	-	-	s	s	s	s	s	P	P	0	i	i	i	i	i	i	i	i	x	x	x	x	x	Rx+=mpyi(Rs,#u8)
1	1	1	0	0	0	0	1	1	-	-	s	s	s	s	s	P	P	0	i	i	i	i	i	i	i	i	x	x	x	x	x	Rx-=mpyi(Rs,#u8)

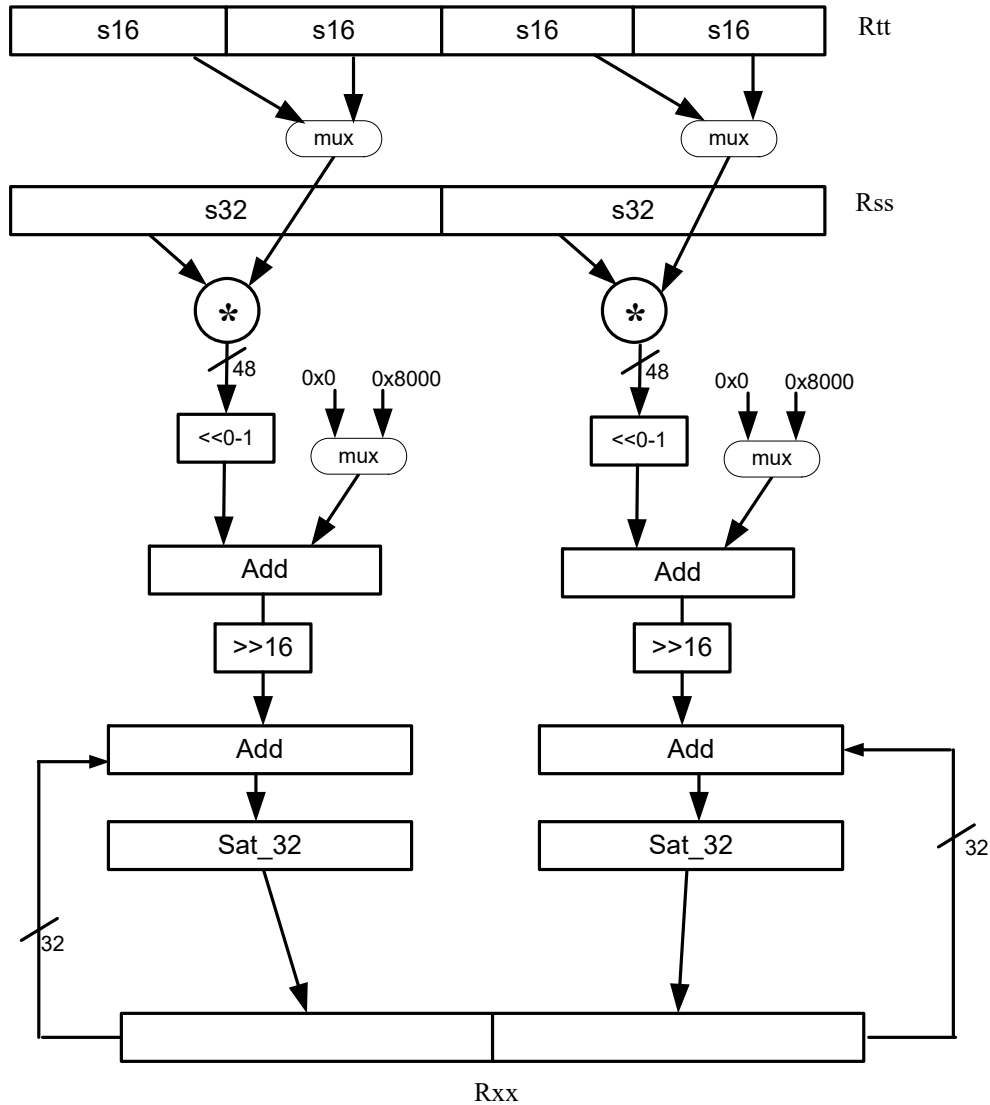
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			d5							
1	1	1	0	1	1	0	1	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	d	d	d	d	d	Rd=mpyi(Rs,Rt)
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			x5							
1	1	1	0	1	1	1	1	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	x	x	x	x	x	Rx+=mpyi(Rs,Rt)
1	1	1	0	1	1	1	1	1	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	x	x	x	x	x	Rx-=mpyi(Rs,Rt)

Field name	Description
RegType	Register type
MajOp	Major opcode
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
u5	Field to encode register u
x5	Field to encode register x
y5	Field to encode register y

## Vector multiply word by signed half (32 × 16)

Perform mixed precision vector multiply operations. A 32-bit word from vector Rss is multiplied by a 16-bit halfword (either even or odd) from vector Rtt. The multiplication is performed as a signed 32 × 16, which produces a 48-bit result. This result is optionally scaled left by one bit. This result is then shifted right by 16 bits, optionally accumulated and then saturated to 32 bits.

This operation is available in vector form (vmpyweh/vmpywoh) and non-vector form (multiply and use upper result).



### Syntax

```
Rdd=vmpyweh (Rss, Rtt) [:<<1] :rnd:sat
```

### Behavior

```
Rdd.w[1]=sat32((Rss.w[1] *  
Rtt.h[2]) [<<1]+0x8000)>>16);  
Rdd.w[0]=sat32((Rss.w[0] *  
Rtt.h[0]) [<<1]+0x8000)>>16);
```



Syntax	Behavior
<code>Rdd=vmpyweh (Rss, Rtt) [:&lt;&lt;1]:sat</code>	<code>Rdd.w[1]=sat<sub>32</sub>((Rss.w[1] * Rtt.h[2]) [&lt;&lt;1]&gt;&gt;16); Rdd.w[0]=sat<sub>32</sub>((Rss.w[0] * Rtt.h[0]) [&lt;&lt;1]&gt;&gt;16);</code>
<code>Rdd=vmpywoh (Rss, Rtt) [:&lt;&lt;1]:rnd:sat</code>	<code>Rdd.w[1]=sat<sub>32</sub>((Rss.w[1] * Rtt.h[3]) [&lt;&lt;1]+0x8000)&gt;&gt;16); Rdd.w[0]=sat<sub>32</sub>((Rss.w[0] * Rtt.h[1]) [&lt;&lt;1]+0x8000)&gt;&gt;16);</code>
<code>Rdd=vmpywoh (Rss, Rtt) [:&lt;&lt;1]:sat</code>	<code>Rdd.w[1]=sat<sub>32</sub>((Rss.w[1] * Rtt.h[3]) [&lt;&lt;1]&gt;&gt;16); Rdd.w[0]=sat<sub>32</sub>((Rss.w[0] * Rtt.h[1]) [&lt;&lt;1]&gt;&gt;16);</code>
<code>Rxx+=vmpyweh (Rss, Rtt) [:&lt;&lt;1]:rnd:sat</code>	<code>Rxx.w[1]=sat<sub>32</sub>(Rxx.w[1] + ((Rss.w[1] * Rtt.h[2]) [&lt;&lt;1]+0x8000)&gt;&gt;16)); Rxx.w[0]=sat<sub>32</sub>(Rxx.w[0] + ((Rss.w[0] * Rtt.h[0]) [&lt;&lt;1]+0x8000)&gt;&gt;16));</code>
<code>Rxx+=vmpyweh (Rss, Rtt) [:&lt;&lt;1]:sat</code>	<code>Rxx.w[1]=sat<sub>32</sub>(Rxx.w[1] + ((Rss.w[1] * Rtt.h[2]) [&lt;&lt;1]&gt;&gt;16)); Rxx.w[0]=sat<sub>32</sub>(Rxx.w[0] + ((Rss.w[0] * Rtt.h[0]) [&lt;&lt;1]&gt;&gt;16));</code>
<code>Rxx+=vmpywoh (Rss, Rtt) [:&lt;&lt;1]:rnd:sat</code>	<code>Rxx.w[1]=sat<sub>32</sub>(Rxx.w[1] + ((Rss.w[1] * Rtt.h[3]) [&lt;&lt;1]+0x8000)&gt;&gt;16)); Rxx.w[0]=sat<sub>32</sub>(Rxx.w[0] + ((Rss.w[0] * Rtt.h[1]) [&lt;&lt;1]+0x8000)&gt;&gt;16));</code>
<code>Rxx+=vmpywoh (Rss, Rtt) [:&lt;&lt;1]:sat</code>	<code>Rxx.w[1]=sat<sub>32</sub>(Rxx.w[1] + ((Rss.w[1] * Rtt.h[3]) [&lt;&lt;1]&gt;&gt;16)); Rxx.w[0]=sat<sub>32</sub>(Rxx.w[0] + ((Rss.w[0] * Rtt.h[1]) [&lt;&lt;1]&gt;&gt;16));</code>

**Class: XTYPE (slots 2,3)****Notes**

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

**Intrinsics**

<code>Rdd=vmpyweh (Rss, Rtt) :&lt;&lt;1:rnd:sat</code>	<code>Word64 Q6_P_vmpyweh_PP_s1_rnd_sat (Word64 Rss, Word64 Rtt)</code>
<code>Rdd=vmpyweh (Rss, Rtt) :&lt;&lt;1:sat</code>	<code>Word64 Q6_P_vmpyweh_PP_s1_sat (Word64 Rss, Word64 Rtt)</code>
<code>Rdd=vmpyweh (Rss, Rtt) :rnd:sat</code>	<code>Word64 Q6_P_vmpyweh_PP_rnd_sat (Word64 Rss, Word64 Rtt)</code>
<code>Rdd=vmpyweh (Rss, Rtt) :sat</code>	<code>Word64 Q6_P_vmpyweh_PP_sat (Word64 Rss, Word64 Rtt)</code>

Rdd=vmpywoh(Rss,Rtt):<<1:rnd:sat	Word64 Q6_P_vmpywoh_PP_s1_rnd_sat(Word64 Rss, Word64 Rtt)
Rdd=vmpywoh(Rss,Rtt):<<1:sat	Word64 Q6_P_vmpywoh_PP_s1_sat(Word64 Rss, Word64 Rtt)
Rdd=vmpywoh(Rss,Rtt):rnd:sat	Word64 Q6_P_vmpywoh_PP_rnd_sat(Word64 Rss, Word64 Rtt)
Rdd=vmpywoh(Rss,Rtt):sat	Word64 Q6_P_vmpywoh_PP_sat(Word64 Rss, Word64 Rtt)
Rxx+=vmpyweh(Rss,Rtt):<<1:rnd:sat	Word64 Q6_P_vmpywehacc_PP_s1_rnd_sat(Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=vmpyweh(Rss,Rtt):<<1:sat	Word64 Q6_P_vmpywehacc_PP_s1_sat(Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=vmpyweh(Rss,Rtt):rnd:sat	Word64 Q6_P_vmpywehacc_PP_rnd_sat(Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=vmpyweh(Rss,Rtt):sat	Word64 Q6_P_vmpywehacc_PP_sat(Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=vmpywoh(Rss,Rtt):<<1:rnd:sat	Word64 Q6_P_vmpywohacc_PP_s1_rnd_sat(Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=vmpywoh(Rss,Rtt):<<1:sat	Word64 Q6_P_vmpywohacc_PP_s1_sat(Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=vmpywoh(Rss,Rtt):rnd:sat	Word64 Q6_P_vmpywohacc_PP_rnd_sat(Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=vmpywoh(Rss,Rtt):sat	Word64 Q6_P_vmpywohacc_PP_sat(Word64 Rxx, Word64 Rss, Word64 Rtt)

### Encoding

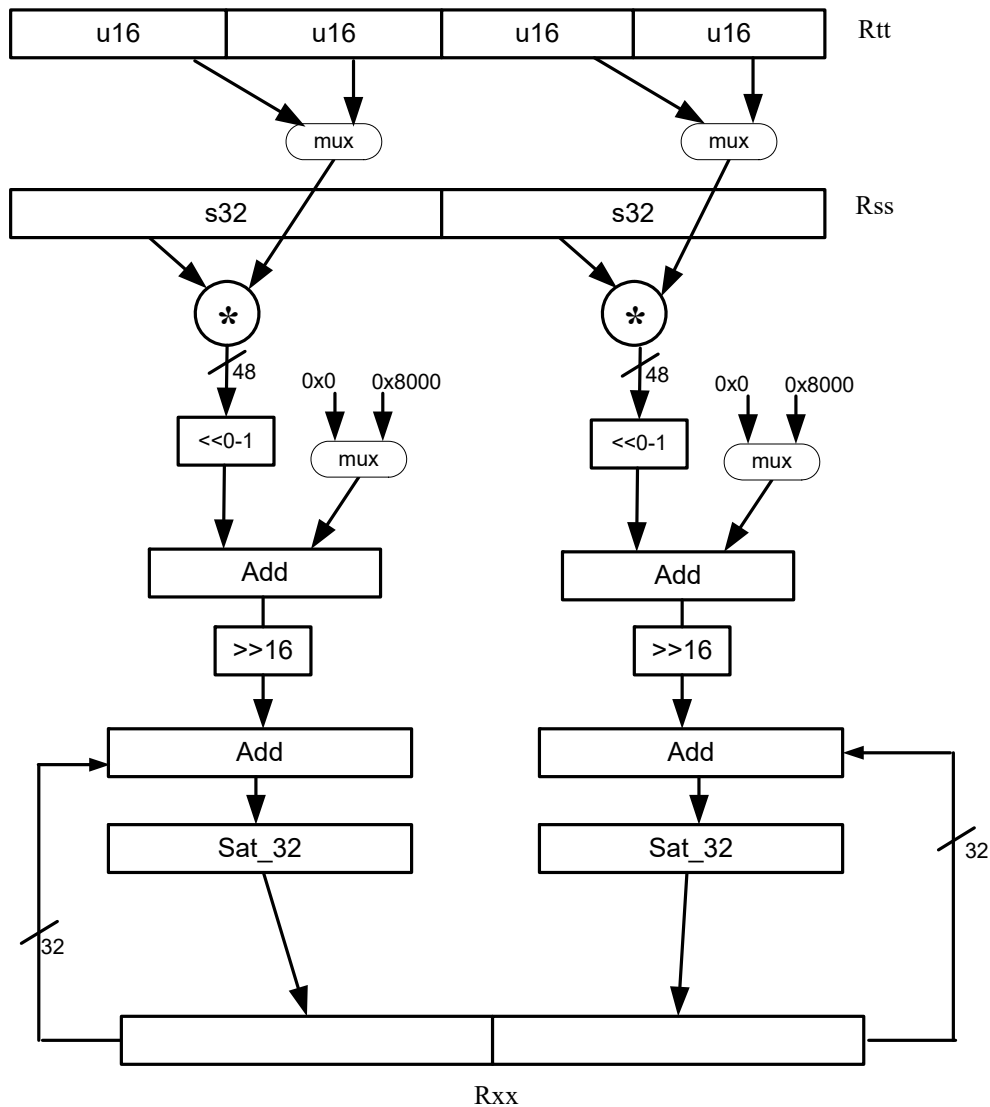
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType			MajOp		s5					Parse		t5					MinOp		d5											
1	1	1	0	1	0	0	0	N	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	1	d	d	d	d	d	Rdd=vmpyweh(Rss,Rtt):<<N]:sat
1	1	1	0	1	0	0	0	N	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	1	d	d	d	d	d	Rdd=vmpywoh(Rss,Rtt):<<N]:sat
1	1	1	0	1	0	0	0	N	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	1	d	d	d	d	d	Rdd=vmpyweh(Rss,Rtt):<<N]:rnd:sat
1	1	1	0	1	0	0	0	N	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	1	d	d	d	d	d	Rdd=vmpywoh(Rss,Rtt):<<N]:rnd:sat
ICLASS		RegType			MajOp		s5					Parse		t5					MinOp		x5											
1	1	1	0	1	0	1	0	N	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	1	x	x	x	x	x	Rxx+=vmpyweh(Rss,Rtt):<<N]:sat
1	1	1	0	1	0	1	0	N	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	1	x	x	x	x	x	Rxx+=vmpywoh(Rss,Rtt):<<N]:sat
1	1	1	0	1	0	1	0	N	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	1	x	x	x	x	x	Rxx+=vmpyweh(Rss,Rtt):<<N]:rnd:sat
1	1	1	0	1	0	1	0	N	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	1	x	x	x	x	x	Rxx+=vmpywoh(Rss,Rtt):<<N]:rnd:sat

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode

<b>Field name</b>	<b>Description</b>
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Vector multiply word by unsigned half ( $32 \times 16$ )

Perform mixed precision vector multiply operations. A 32-bit signed word from vector  $R_{ss}$  is multiplied by a 16-bit unsigned halfword (either odd or even) from vector  $R_{tt}$ . This multiplication produces a 48-bit result. This result is optionally scaled left by one bit, and then a rounding constant is optionally added to the lower 16 bits. This result is then shifted right by 16 bits, optionally accumulated and then saturated to 32 bits. This is a dual vector operation and is performed for both high and low word of  $R_{ss}$ .



### Syntax

```
Rdd=vmpyweuh (Rss, Rtt) [ :<<1] :rnd:s  
at
```

### Behavior

```
Rdd.w[1]=sat32((Rss.w[1] *  
Rtt.uh[2]) [<<1]+0x8000)>>16);  
Rdd.w[0]=sat32((Rss.w[0] *  
Rtt.uh[0]) [<<1]+0x8000)>>16);
```

Syntax	Behavior
<code>Rdd=vmpyweuh (Rss,Rtt) [ :&lt;&lt;1 ] :sat</code>	<code>Rdd.w[1]=sat<sub>32</sub>((Rss.w[1] * Rtt.uh[2]) [&lt;&lt;1]&gt;&gt;16); Rdd.w[0]=sat<sub>32</sub>((Rss.w[0] * Rtt.uh[0]) [&lt;&lt;1]&gt;&gt;16);</code>
<code>Rdd=vmpywouh (Rss,Rtt) [ :&lt;&lt;1 ] :rnd:sa t</code>	<code>Rdd.w[1]=sat<sub>32</sub>((Rss.w[1] * Rtt.uh[3]) [&lt;&lt;1]+0x8000)&gt;&gt;16); Rdd.w[0]=sat<sub>32</sub>((Rss.w[0] * Rtt.uh[1]) [&lt;&lt;1]+0x8000)&gt;&gt;16);</code>
<code>Rdd=vmpywouh (Rss,Rtt) [ :&lt;&lt;1 ] :sat</code>	<code>Rdd.w[1]=sat<sub>32</sub>((Rss.w[1] * Rtt.uh[3]) [&lt;&lt;1]&gt;&gt;16); Rdd.w[0]=sat<sub>32</sub>((Rss.w[0] * Rtt.uh[1]) [&lt;&lt;1]&gt;&gt;16);</code>
<code>Rxx+=vmpyweuh (Rss,Rtt) [ :&lt;&lt;1 ] :rnd:sa t</code>	<code>Rxx.w[1]=sat<sub>32</sub>(Rxx.w[1] + ((Rss.w[1] * Rtt.uh[2]) [&lt;&lt;1]+0x8000)&gt;&gt;16)); Rxx.w[0]=sat<sub>32</sub>(Rxx.w[0] + ((Rss.w[0] * Rtt.uh[0]) [&lt;&lt;1]+0x8000)&gt;&gt;16));</code>
<code>Rxx+=vmpyweuh (Rss,Rtt) [ :&lt;&lt;1 ] :sat</code>	<code>Rxx.w[1]=sat<sub>32</sub>(Rxx.w[1] + ((Rss.w[1] * Rtt.uh[2]) [&lt;&lt;1]&gt;&gt;16)); Rxx.w[0]=sat<sub>32</sub>(Rxx.w[0] + ((Rss.w[0] * Rtt.uh[0]) [&lt;&lt;1]&gt;&gt;16));</code>
<code>Rxx+=vmpywouh (Rss,Rtt) [ :&lt;&lt;1 ] :rnd:sa t</code>	<code>Rxx.w[1]=sat<sub>32</sub>(Rxx.w[1] + ((Rss.w[1] * Rtt.uh[3]) [&lt;&lt;1]+0x8000)&gt;&gt;16)); Rxx.w[0]=sat<sub>32</sub>(Rxx.w[0] + ((Rss.w[0] * Rtt.uh[1]) [&lt;&lt;1]+0x8000)&gt;&gt;16));</code>
<code>Rxx+=vmpywouh (Rss,Rtt) [ :&lt;&lt;1 ] :sat</code>	<code>Rxx.w[1]=sat<sub>32</sub>(Rxx.w[1] + ((Rss.w[1] * Rtt.uh[3]) [&lt;&lt;1]&gt;&gt;16)); Rxx.w[0]=sat<sub>32</sub>(Rxx.w[0] + ((Rss.w[0] * Rtt.uh[1]) [&lt;&lt;1]&gt;&gt;16));</code>

### Class: XTYPE (slots 2,3)

#### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

#### Intrinsics

<code>Rdd=vmpyweuh (Rss,Rtt) :&lt;&lt;1:rnd:sa t</code>	Word64 Q6_P_vmpyweuh_PP_sl_rnd_sat (Word64 Rss, Word64 Rtt)
<code>Rdd=vmpyweuh (Rss,Rtt) :&lt;&lt;1:sat</code>	Word64 Q6_P_vmpyweuh_PP_sl_sat (Word64 Rss, Word64 Rtt)
<code>Rdd=vmpyweuh (Rss,Rtt) :rnd:sat</code>	Word64 Q6_P_vmpyweuh_PP_rnd_sat (Word64 Rss, Word64 Rtt)

Rdd=vmpyweuh (Rss,Rtt) :sat	Word64 Q6_P_vmpyweuh_PP_sat (Word64 Rss, Word64 Rtt)
Rdd=vmpywouh (Rss,Rtt) :<<1:rnd:sat	Word64 Q6_P_vmpywouh_PP_sl_rnd_sat (Word64 Rss, Word64 Rtt)
Rdd=vmpywouh (Rss,Rtt) :<<1:sat	Word64 Q6_P_vmpywouh_PP_sl_sat (Word64 Rss, Word64 Rtt)
Rdd=vmpywouh (Rss,Rtt) :rnd:sat	Word64 Q6_P_vmpywouh_PP_rnd_sat (Word64 Rss, Word64 Rtt)
Rdd=vmpywouh (Rss,Rtt) :sat	Word64 Q6_P_vmpywouh_PP_sat (Word64 Rss, Word64 Rtt)
Rxx+=vmpyweuh (Rss,Rtt) :<<1:rnd:sat	Word64 Q6_P_vmpyweuhacc_PP_sl_rnd_sat (Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=vmpyweuh (Rss,Rtt) :<<1:sat	Word64 Q6_P_vmpyweuhacc_PP_sl_sat (Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=vmpyweuh (Rss,Rtt) :rnd:sat	Word64 Q6_P_vmpyweuhacc_PP_rnd_sat (Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=vmpyweuh (Rss,Rtt) :sat	Word64 Q6_P_vmpyweuhacc_PP_sat (Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=vmpywouh (Rss,Rtt) :<<1:rnd:sat	Word64 Q6_P_vmpywouhacc_PP_sl_rnd_sat (Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=vmpywouh (Rss,Rtt) :<<1:sat	Word64 Q6_P_vmpywouhacc_PP_sl_sat (Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=vmpywouh (Rss,Rtt) :rnd:sat	Word64 Q6_P_vmpywouhacc_PP_rnd_sat (Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=vmpywouh (Rss,Rtt) :sat	Word64 Q6_P_vmpywouhacc_PP_sat (Word64 Rxx, Word64 Rss, Word64 Rtt)

## Encoding

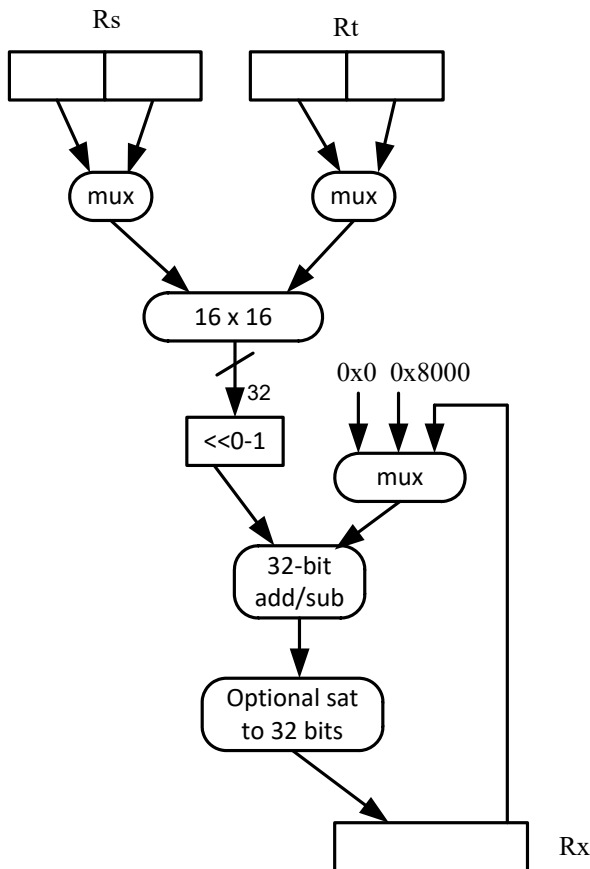
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse		t5					MinOp		d5									
1	1	1	0	1	0	0	0	N	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	1	d	d	d	d	d	Rdd=vmpyweuh(Rss,Rtt):<<N]:sat
1	1	1	0	1	0	0	0	N	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	1	d	d	d	d	d	Rdd=vmpywouh(Rss,Rtt):<<N]:sat
1	1	1	0	1	0	0	0	N	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	1	d	d	d	d	d	Rdd=vmpyweuh(Rss,Rtt):<<N]:rnd:sat
1	1	1	0	1	0	0	0	N	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	1	d	d	d	d	d	Rdd=vmpywouh(Rss,Rtt):<<N]:rnd:sat
ICLASS			RegType				MajOp		s5					Parse		t5					MinOp		x5									
1	1	1	0	1	0	1	0	N	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	1	x	x	x	x	x	Rxx+=vmpyweuh(Rss,Rtt):<<N]:sat
1	1	1	0	1	0	1	0	N	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	1	x	x	x	x	x	Rxx+=vmpywouh(Rss,Rtt):<<N]:sat
1	1	1	0	1	0	1	0	N	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	1	x	x	x	x	x	Rxx+=vmpyweuh(Rss,Rtt):<<N]:rnd:sat
1	1	1	0	1	0	1	0	N	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	1	x	x	x	x	x	Rxx+=vmpywouh(Rss,Rtt):<<N]:rnd:sat

<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

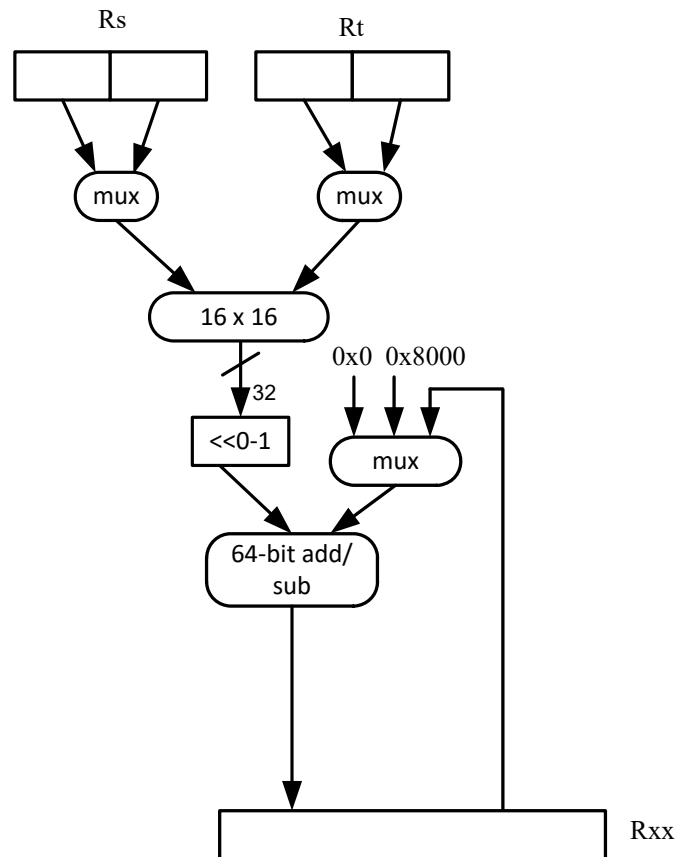
## Multiply signed halfwords

Multiply two signed halfwords. Optionally shift the multiplier result by 1 bit. This result can be accumulated or rounded. The destination/accumulator is either 32 or 64 bits. For 32-bit results, saturation is optional.

$Rx += mpy(Rs.[HL], Rt.[HL])[<<1][:sat]$   
 $Rd = mpy(Rs.[HL], Rt.[HL])[<<1][:rnd][:sat]$



$Rxx += mpy(Rs.[HL], Rt.[HL])[<<1]$   
 $Rdd = mpy(Rs.[HL], Rt.[HL])[<<1][:rnd]$



### Syntax

$Rd = mpy(Rs.[HL], Rt.[HL])[<<1][:rnd][:sat]$
$Rdd = mpy(Rs.[HL], Rt.[HL])[<<1][:rnd]$
$Rx += mpy(Rs.[HL], Rt.[HL])[<<1][:sat]$
$Rx -= mpy(Rs.[HL], Rt.[HL])[<<1][:sat]$
$Rxx += mpy(Rs.[HL], Rt.[HL])[<<1]$
$Rxx -= mpy(Rs.[HL], Rt.[HL])[<<1]$

### Behavior

$Rd = [sat_{32}] ([round] ((Rs.h[01] * Rt.h[01]) [<<1]));$
$Rdd = [round] ((Rs.h[01] * Rt.h[01]) [<<1]);$
$Rx = [sat_{32}] (Rx + (Rs.h[01] * Rt.h[01]) [<<1]);$
$Rx = [sat_{32}] (Rx - (Rs.h[01] * Rt.h[01]) [<<1]);$
$Rxx = Rxx + (Rs.h[01] * Rt.h[01]) [<<1];$
$Rxx = Rxx - (Rs.h[01] * Rt.h[01]) [<<1];$



**Class: XTYPE (slots 2,3)****Notes**

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

**Intrinsics**

Rd=mpy(Rs.H,Rt.H)	Word32 Q6_R_mpy_RhRh(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.H,Rt.H) <<1	Word32 Q6_R_mpy_RhRh_s1(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.H,Rt.H) <<1:rnd	Word32 Q6_R_mpy_RhRh_s1_rnd(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.H,Rt.H) <<1:rnd:sat	Word32 Q6_R_mpy_RhRh_s1_rnd_sat(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.H,Rt.H) <<1:sat	Word32 Q6_R_mpy_RhRh_s1_sat(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.H,Rt.H) :rnd	Word32 Q6_R_mpy_RhRh_rnd(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.H,Rt.H) :rnd:sat	Word32 Q6_R_mpy_RhRh_rnd_sat(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.H,Rt.H) :sat	Word32 Q6_R_mpy_RhRh_sat(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.H,Rt.L)	Word32 Q6_R_mpy_RhRl(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.H,Rt.L) <<1	Word32 Q6_R_mpy_RhRl_s1(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.H,Rt.L) <<1:rnd	Word32 Q6_R_mpy_RhRl_s1_rnd(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.H,Rt.L) <<1:rnd:sat	Word32 Q6_R_mpy_RhRl_s1_rnd_sat(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.H,Rt.L) <<1:sat	Word32 Q6_R_mpy_RhRl_s1_sat(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.H,Rt.L) :rnd	Word32 Q6_R_mpy_RhRl_rnd(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.H,Rt.L) :rnd:sat	Word32 Q6_R_mpy_RhRl_rnd_sat(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.H,Rt.L) :sat	Word32 Q6_R_mpy_RhRl_sat(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.L,Rt.H)	Word32 Q6_R_mpy_RlRh(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.L,Rt.H) <<1	Word32 Q6_R_mpy_RlRh_s1(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.L,Rt.H) <<1:rnd	Word32 Q6_R_mpy_RlRh_s1_rnd(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.L,Rt.H) <<1:rnd:sat	Word32 Q6_R_mpy_RlRh_s1_rnd_sat(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.L,Rt.H) <<1:sat	Word32 Q6_R_mpy_RlRh_s1_sat(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.L,Rt.H) :rnd	Word32 Q6_R_mpy_RlRh_rnd(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.L,Rt.H) :rnd:sat	Word32 Q6_R_mpy_RlRh_rnd_sat(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.L,Rt.H) :sat	Word32 Q6_R_mpy_RlRh_sat(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.L,Rt.L)	Word32 Q6_R_mpy_RlRl(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.L,Rt.L) <<1	Word32 Q6_R_mpy_RlRl_s1(Word32 Rs, Word32 Rt)
Rd=mpy(Rs.L,Rt.L) <<1:rnd	Word32 Q6_R_mpy_RlRl_s1_rnd(Word32 Rs, Word32 Rt)

Rd=mpy (Rs.L,Rt.L) :<<1:rnd:sat	Word32 Q6_R_mpy_RlRl_s1_rnd_sat (Word32 Rs, Word32 Rt)
Rd=mpy (Rs.L,Rt.L) :<<1:sat	Word32 Q6_R_mpy_RlRl_s1_sat (Word32 Rs, Word32 Rt)
Rd=mpy (Rs.L,Rt.L) :rnd	Word32 Q6_R_mpy_RlRl_rnd (Word32 Rs, Word32 Rt)
Rd=mpy (Rs.L,Rt.L) :rnd:sat	Word32 Q6_R_mpy_RlRl_rnd_sat (Word32 Rs, Word32 Rt)
Rd=mpy (Rs.L,Rt.L) :sat	Word32 Q6_R_mpy_RlRl_sat (Word32 Rs, Word32 Rt)
Rdd=mpy (Rs.H,Rt.H)	Word64 Q6_P_mpy_RhRh (Word32 Rs, Word32 Rt)
Rdd=mpy (Rs.H,Rt.H) :<<1	Word64 Q6_P_mpy_RhRh_s1 (Word32 Rs, Word32 Rt)
Rdd=mpy (Rs.H,Rt.H) :<<1:rnd	Word64 Q6_P_mpy_RhRh_s1_rnd (Word32 Rs, Word32 Rt)
Rdd=mpy (Rs.H,Rt.H) :rnd	Word64 Q6_P_mpy_RhRh_rnd (Word32 Rs, Word32 Rt)
Rdd=mpy (Rs.H,Rt.L)	Word64 Q6_P_mpy_RhRl (Word32 Rs, Word32 Rt)
Rdd=mpy (Rs.H,Rt.L) :<<1	Word64 Q6_P_mpy_RhRl_s1 (Word32 Rs, Word32 Rt)
Rdd=mpy (Rs.H,Rt.L) :<<1:rnd	Word64 Q6_P_mpy_RhRl_s1_rnd (Word32 Rs, Word32 Rt)
Rdd=mpy (Rs.H,Rt.L) :rnd	Word64 Q6_P_mpy_RhRl_rnd (Word32 Rs, Word32 Rt)
Rdd=mpy (Rs.L,Rt.H)	Word64 Q6_P_mpy_RlRh (Word32 Rs, Word32 Rt)
Rdd=mpy (Rs.L,Rt.H) :<<1	Word64 Q6_P_mpy_RlRh_s1 (Word32 Rs, Word32 Rt)
Rdd=mpy (Rs.L,Rt.H) :<<1:rnd	Word64 Q6_P_mpy_RlRh_s1_rnd (Word32 Rs, Word32 Rt)
Rdd=mpy (Rs.L,Rt.H) :rnd	Word64 Q6_P_mpy_RlRh_rnd (Word32 Rs, Word32 Rt)
Rdd=mpy (Rs.L,Rt.L)	Word64 Q6_P_mpy_RlRl (Word32 Rs, Word32 Rt)
Rdd=mpy (Rs.L,Rt.L) :<<1	Word64 Q6_P_mpy_RlRl_s1 (Word32 Rs, Word32 Rt)
Rdd=mpy (Rs.L,Rt.L) :<<1:rnd	Word64 Q6_P_mpy_RlRl_s1_rnd (Word32 Rs, Word32 Rt)
Rdd=mpy (Rs.L,Rt.L) :rnd	Word64 Q6_P_mpy_RlRl_rnd (Word32 Rs, Word32 Rt)
Rx+=mpy (Rs.H,Rt.H)	Word32 Q6_R_mpyacc_RhRh (Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpy (Rs.H,Rt.H) :<<1	Word32 Q6_R_mpyacc_RhRh_s1 (Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpy (Rs.H,Rt.H) :<<1:sat	Word32 Q6_R_mpyacc_RhRh_s1_sat (Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpy (Rs.H,Rt.H) :sat	Word32 Q6_R_mpyacc_RhRh_sat (Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpy (Rs.H,Rt.L)	Word32 Q6_R_mpyacc_RhRl (Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpy (Rs.H,Rt.L) :<<1	Word32 Q6_R_mpyacc_RhRl_s1 (Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpy (Rs.H,Rt.L) :<<1:sat	Word32 Q6_R_mpyacc_RhRl_s1_sat (Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpy (Rs.H,Rt.L) :sat	Word32 Q6_R_mpyacc_RhRl_sat (Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpy (Rs.L,Rt.H)	Word32 Q6_R_mpyacc_RlRh (Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpy (Rs.L,Rt.H) :<<1	Word32 Q6_R_mpyacc_RlRh_s1 (Word32 Rx, Word32 Rs, Word32 Rt)

Rx+=mpy (Rs.L,Rt.H) :<<1:sat	Word32 Q6_R_mpyacc_RlRh_s1_sat (Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpy (Rs.L,Rt.H) :sat	Word32 Q6_R_mpyacc_RlRh_sat (Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpy (Rs.L,Rt.L)	Word32 Q6_R_mpyacc_RlRl (Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpy (Rs.L,Rt.L) :<<1	Word32 Q6_R_mpyacc_RlRl_s1 (Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpy (Rs.L,Rt.L) :<<1:sat	Word32 Q6_R_mpyacc_RlRl_s1_sat (Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpy (Rs.L,Rt.L) :sat	Word32 Q6_R_mpyacc_RlRl_sat (Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpy (Rs.H,Rt.H)	Word32 Q6_R_mpynac_RhRh (Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpy (Rs.H,Rt.H) :<<1	Word32 Q6_R_mpynac_RhRh_s1 (Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpy (Rs.H,Rt.H) :<<1:sat	Word32 Q6_R_mpynac_RhRh_s1_sat (Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpy (Rs.H,Rt.H) :sat	Word32 Q6_R_mpynac_RhRh_sat (Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpy (Rs.H,Rt.L)	Word32 Q6_R_mpynac_RhRl (Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpy (Rs.H,Rt.L) :<<1	Word32 Q6_R_mpynac_RhRl_s1 (Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpy (Rs.H,Rt.L) :<<1:sat	Word32 Q6_R_mpynac_RhRl_s1_sat (Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpy (Rs.H,Rt.L) :sat	Word32 Q6_R_mpynac_RhRl_sat (Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpy (Rs.L,Rt.H)	Word32 Q6_R_mpynac_RlRh (Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpy (Rs.L,Rt.H) :<<1	Word32 Q6_R_mpynac_RlRh_s1 (Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpy (Rs.L,Rt.H) :<<1:sat	Word32 Q6_R_mpynac_RlRh_s1_sat (Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpy (Rs.L,Rt.H) :sat	Word32 Q6_R_mpynac_RlRh_sat (Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpy (Rs.L,Rt.L)	Word32 Q6_R_mpynac_RlRl (Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpy (Rs.L,Rt.L) :<<1	Word32 Q6_R_mpynac_RlRl_s1 (Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpy (Rs.L,Rt.L) :<<1:sat	Word32 Q6_R_mpynac_RlRl_s1_sat (Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpy (Rs.L,Rt.L) :sat	Word32 Q6_R_mpynac_RlRl_sat (Word32 Rx, Word32 Rs, Word32 Rt)
Rxx+=mpy (Rs.H,Rt.H)	Word64 Q6_P_mpyacc_RhRh (Word64 Rxx, Word32 Rs, Word32 Rt)

Rxx+=mpy (Rs.H,Rt.H) :<<1	Word64 Q6_P_mpyacc_RhRh_s1 (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx+=mpy (Rs.H,Rt.L)	Word64 Q6_P_mpyacc_RhRl (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx+=mpy (Rs.H,Rt.L) :<<1	Word64 Q6_P_mpyacc_RhRl_s1 (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx+=mpy (Rs.L,Rt.H)	Word64 Q6_P_mpyacc_RlRh (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx+=mpy (Rs.L,Rt.H) :<<1	Word64 Q6_P_mpyacc_RlRh_s1 (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx+=mpy (Rs.L,Rt.L)	Word64 Q6_P_mpyacc_RlRl (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx+=mpy (Rs.L,Rt.L) :<<1	Word64 Q6_P_mpyacc_RlRl_s1 (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx-=mpy (Rs.H,Rt.H)	Word64 Q6_P_mpynac_RhRh (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx-=mpy (Rs.H,Rt.H) :<<1	Word64 Q6_P_mpynac_RhRh_s1 (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx-=mpy (Rs.H,Rt.L)	Word64 Q6_P_mpynac_RhRl (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx-=mpy (Rs.H,Rt.L) :<<1	Word64 Q6_P_mpynac_RhRl_s1 (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx-=mpy (Rs.L,Rt.H)	Word64 Q6_P_mpynac_RlRh (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx-=mpy (Rs.L,Rt.H) :<<1	Word64 Q6_P_mpynac_RlRh_s1 (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx-=mpy (Rs.L,Rt.L)	Word64 Q6_P_mpynac_RlRl (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx-=mpy (Rs.L,Rt.L) :<<1	Word64 Q6_P_mpynac_RlRl_s1 (Word64 Rxx, Word32 Rs, Word32 Rt)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp				s5				Parse		t5				MinOp		d5								
1	1	1	0	0	1	0	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	0	0	d	d	d	d	d	Rdd=mpy(Rs.L,Rt.L)[:<<N]
1	1	1	0	0	1	0	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	0	1	d	d	d	d	d	Rdd=mpy(Rs.L,Rt.H)[:<<N]
1	1	1	0	0	1	0	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	1	0	d	d	d	d	d	Rdd=mpy(Rs.H,Rt.L)[:<<N]
1	1	1	0	0	1	0	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	1	1	d	d	d	d	d	Rdd=mpy(Rs.H,Rt.H)[:<<N]
1	1	1	0	0	1	0	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	0	0	d	d	d	d	d	Rdd=mpy(Rs.L,Rt.L)[:<<N]:rnd
1	1	1	0	0	1	0	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	0	1	d	d	d	d	d	Rdd=mpy(Rs.L,Rt.H)[:<<N]:rnd
1	1	1	0	0	1	0	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	1	0	d	d	d	d	d	Rdd=mpy(Rs.H,Rt.L)[:<<N]:rnd
1	1	1	0	0	1	0	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	1	1	d	d	d	d	d	Rdd=mpy(Rs.H,Rt.H)[:<<N]:rnd
ICLASS				RegType				MajOp				s5				Parse		t5				MinOp		x5								

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
1	1	1	0	0	1	1	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	x	x	x	x	x	Rxx+=mpy(Rs.L,Rt.L):<<N]
1	1	1	0	0	1	1	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	x	x	x	x	x	Rxx+=mpy(Rs.L,Rt.H):<<N]
1	1	1	0	0	1	1	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	x	x	x	x	x	Rxx+=mpy(Rs.H,Rt.L):<<N]
1	1	1	0	0	1	1	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	x	x	x	x	x	Rxx+=mpy(Rs.H,Rt.H):<<N]
1	1	1	0	0	1	1	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	x	x	x	x	x	Rxx-=mpy(Rs.L,Rt.L):<<N]
1	1	1	0	0	1	1	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	x	x	x	x	x	Rxx-=mpy(Rs.L,Rt.H):<<N]
1	1	1	0	0	1	1	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	x	x	x	x	x	Rxx-=mpy(Rs.H,Rt.L):<<N]
1	1	1	0	0	1	1	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	x	x	x	x	x	Rxx-=mpy(Rs.H,Rt.H):<<N]
IClass		RegType		MajOp		s5					Parse		t5			MinOp		d5														
1	1	1	0	1	1	0	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	d	d	d	d	d	Rd=mpy(Rs.L,Rt.L):<<N]
1	1	1	0	1	1	0	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	d	d	d	d	d	Rd=mpy(Rs.L,Rt.H):<<N]
1	1	1	0	1	1	0	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	d	d	d	d	d	Rd=mpy(Rs.H,Rt.L):<<N]
1	1	1	0	1	1	0	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	d	d	d	d	d	Rd=mpy(Rs.H,Rt.H):<<N]
1	1	1	0	1	1	0	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	0	d	d	d	d	d	Rd=mpy(Rs.L,Rt.L):<<N]:sat
1	1	1	0	1	1	0	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	1	d	d	d	d	d	Rd=mpy(Rs.L,Rt.H):<<N]:sat
1	1	1	0	1	1	0	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	0	d	d	d	d	d	Rd=mpy(Rs.H,Rt.L):<<N]:sat
1	1	1	0	1	1	0	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	1	d	d	d	d	d	Rd=mpy(Rs.H,Rt.H):<<N]:sat
1	1	1	0	1	1	0	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	d	d	d	d	d	Rd=mpy(Rs.L,Rt.L):<<N]:rnd
1	1	1	0	1	1	0	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	d	d	d	d	d	Rd=mpy(Rs.L,Rt.H):<<N]:rnd
1	1	1	0	1	1	0	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	d	d	d	d	d	Rd=mpy(Rs.H,Rt.L):<<N]:rnd
1	1	1	0	1	1	0	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	d	d	d	d	d	Rd=mpy(Rs.H,Rt.H):<<N]:rnd
1	1	1	0	1	1	0	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	0	d	d	d	d	d	Rd=mpy(Rs.L,Rt.L):<<N]:rnd:sat
1	1	1	0	1	1	0	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	1	d	d	d	d	d	Rd=mpy(Rs.L,Rt.H):<<N]:rnd:sat
1	1	1	0	1	1	0	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	0	d	d	d	d	d	Rd=mpy(Rs.H,Rt.L):<<N]:rnd:sat
1	1	1	0	1	1	0	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	1	d	d	d	d	d	Rd=mpy(Rs.H,Rt.H):<<N]:rnd:sat
IClass		RegType		MajOp		s5					Parse		t5			MinOp		x5														
1	1	1	0	1	1	1	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	x	x	x	x	x	Rx+=mpy(Rs.L,Rt.L):<<N]
1	1	1	0	1	1	1	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	x	x	x	x	x	Rx+=mpy(Rs.L,Rt.H):<<N]
1	1	1	0	1	1	1	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	x	x	x	x	x	Rx+=mpy(Rs.H,Rt.L):<<N]
1	1	1	0	1	1	1	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	x	x	x	x	x	Rx+=mpy(Rs.H,Rt.H):<<N]
1	1	1	0	1	1	1	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	0	x	x	x	x	x	Rx+=mpy(Rs.L,Rt.L):<<N]:sat
1	1	1	0	1	1	1	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	1	x	x	x	x	x	Rx+=mpy(Rs.L,Rt.H):<<N]:sat
1	1	1	0	1	1	1	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	0	x	x	x	x	x	Rx+=mpy(Rs.H,Rt.L):<<N]:sat
1	1	1	0	1	1	1	0	N	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	1	x	x	x	x	x	Rx+=mpy(Rs.H,Rt.H):<<N]:sat
1	1	1	0	1	1	1	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	x	x	x	x	x	Rx-=mpy(Rs.L,Rt.L):<<N]
1	1	1	0	1	1	1	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	x	x	x	x	x	Rx-=mpy(Rs.L,Rt.H):<<N]
1	1	1	0	1	1	1	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	x	x	x	x	x	Rx-=mpy(Rs.H,Rt.L):<<N]

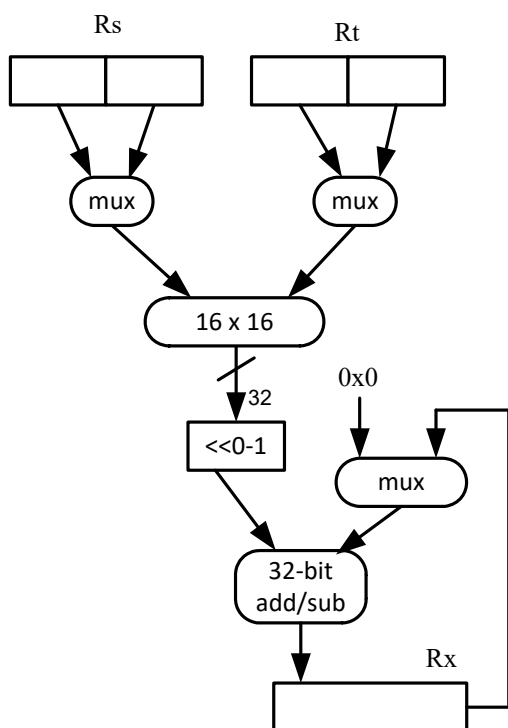
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
1	1	1	0	1	1	1	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	x	x	x	x	x	Rx= =mpy(Rs.H,Rt.H)[:<<N]
1	1	1	0	1	1	1	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	0	x	x	x	x	x	Rx= =mpy(Rs.L,Rt.L)[:<<N]:sat
1	1	1	0	1	1	1	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	1	x	x	x	x	x	Rx= =mpy(Rs.L,Rt.H)[:<<N]:sat
1	1	1	0	1	1	1	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	0	x	x	x	x	x	Rx= =mpy(Rs.H,Rt.L)[:<<N]:sat
1	1	1	0	1	1	1	0	N	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	1	x	x	x	x	x	Rx= =mpy(Rs.H,Rt.H)[:<<N]:sat

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
sH	Rs is high
tH	Rt is high
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

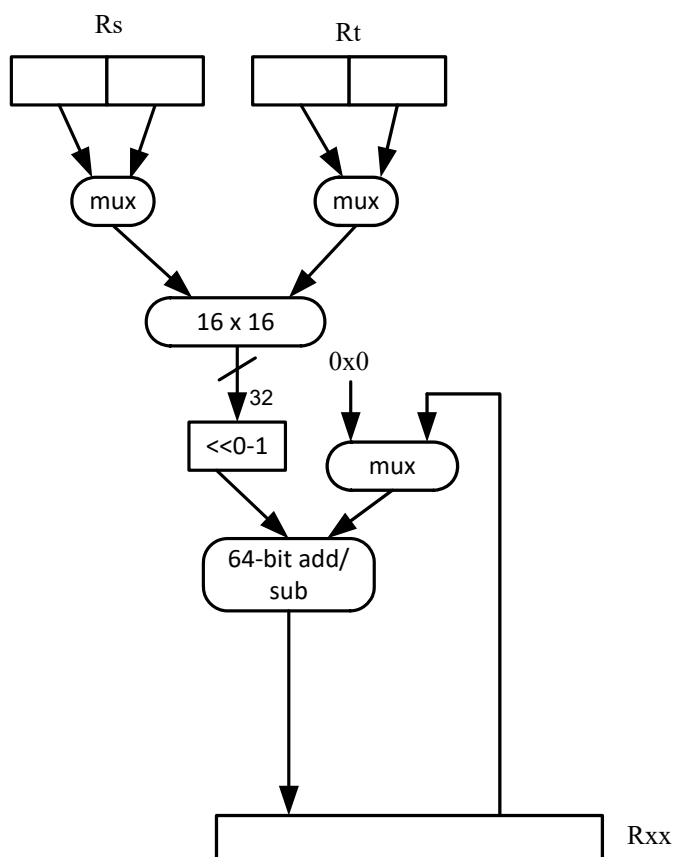
## Multiply unsigned halfwords

Multiply two unsigned halfwords. Scale the result by 0-3 bits. Optionally, add or subtract the result from the accumulator.

$Rx += mpyu(Rs.[HL], Rt.[HL])[:<<1]$   
 $Rd = mpyu(Rs.[HL], Rt.[HL])[:<<1]$



$Rxx += mpyu(Rs.[HL], Rt.[HL])[:<<1]$   
 $Rdd = mpyu(Rs.[HL], Rt.[HL])[:<<1]$



### Syntax

$Rd = mpyu(Rs.[HL], Rt.[HL])[:<<1]$

$Rdd = mpyu(Rs.[HL], Rt.[HL])[:<<1]$

$Rx += mpyu(Rs.[HL], Rt.[HL])[:<<1]$

$Rx -$   
 $= mpyu(Rs.[HL], Rt.[HL])[:<<1]$

$Rxx += mpyu(Rs.[HL], Rt.[HL])[:<<1]$

$Rxx -$   
 $= mpyu(Rs.[HL], Rt.[HL])[:<<1]$

### Behavior

$Rd = (Rs.uh[01] * Rt.uh[01])[:<<1];$

$Rdd = (Rs.uh[01] * Rt.uh[01])[:<<1];$

$Rx = Rx + (Rs.uh[01] * Rt.uh[01])[:<<1];$

$Rx = Rx - (Rs.uh[01] * Rt.uh[01])[:<<1];$

$Rxx = Rxx + (Rs.uh[01] * Rt.uh[01])[:<<1];$

$Rxx = Rxx - (Rs.uh[01] * Rt.uh[01])[:<<1];$

**Class: XTYPE (slots 2,3)****Intrinsics**

Rd=mpyu (Rs.H,Rt.H)	UWord32 Q6_R_mpyu_RhRh(Word32 Rs, Word32 Rt)
Rd=mpyu (Rs.H,Rt.H) :<<1	UWord32 Q6_R_mpyu_RhRh_s1(Word32 Rs, Word32 Rt)
Rd=mpyu (Rs.H,Rt.L)	UWord32 Q6_R_mpyu_RhRl(Word32 Rs, Word32 Rt)
Rd=mpyu (Rs.H,Rt.L) :<<1	UWord32 Q6_R_mpyu_RhRl_s1(Word32 Rs, Word32 Rt)
Rd=mpyu (Rs.L,Rt.H)	UWord32 Q6_R_mpyu_RlRh(Word32 Rs, Word32 Rt)
Rd=mpyu (Rs.L,Rt.H) :<<1	UWord32 Q6_R_mpyu_RlRh_s1(Word32 Rs, Word32 Rt)
Rd=mpyu (Rs.L,Rt.L)	UWord32 Q6_R_mpyu_RlRl(Word32 Rs, Word32 Rt)
Rd=mpyu (Rs.L,Rt.L) :<<1	UWord32 Q6_R_mpyu_RlRl_s1(Word32 Rs, Word32 Rt)
Rdd=mpyu (Rs.H,Rt.H)	UWord64 Q6_P_mpyu_RhRh(Word32 Rs, Word32 Rt)
Rdd=mpyu (Rs.H,Rt.H) :<<1	UWord64 Q6_P_mpyu_RhRh_s1(Word32 Rs, Word32 Rt)
Rdd=mpyu (Rs.H,Rt.L)	UWord64 Q6_P_mpyu_RhRl(Word32 Rs, Word32 Rt)
Rdd=mpyu (Rs.H,Rt.L) :<<1	UWord64 Q6_P_mpyu_RhRl_s1(Word32 Rs, Word32 Rt)
Rdd=mpyu (Rs.L,Rt.H)	UWord64 Q6_P_mpyu_RlRh(Word32 Rs, Word32 Rt)
Rdd=mpyu (Rs.L,Rt.H) :<<1	UWord64 Q6_P_mpyu_RlRh_s1(Word32 Rs, Word32 Rt)
Rdd=mpyu (Rs.L,Rt.L)	UWord64 Q6_P_mpyu_RlRl(Word32 Rs, Word32 Rt)
Rdd=mpyu (Rs.L,Rt.L) :<<1	UWord64 Q6_P_mpyu_RlRl_s1(Word32 Rs, Word32 Rt)
Rx+=mpyu (Rs.H,Rt.H)	Word32 Q6_R_mpyuacc_RhRh(Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpyu (Rs.H,Rt.H) :<<1	Word32 Q6_R_mpyuacc_RhRh_s1(Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpyu (Rs.H,Rt.L)	Word32 Q6_R_mpyuacc_RhRl(Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpyu (Rs.H,Rt.L) :<<1	Word32 Q6_R_mpyuacc_RhRl_s1(Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpyu (Rs.L,Rt.H)	Word32 Q6_R_mpyuacc_RlRh(Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpyu (Rs.L,Rt.H) :<<1	Word32 Q6_R_mpyuacc_RlRh_s1(Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpyu (Rs.L,Rt.L)	Word32 Q6_R_mpyuacc_RlRl(Word32 Rx, Word32 Rs, Word32 Rt)
Rx+=mpyu (Rs.L,Rt.L) :<<1	Word32 Q6_R_mpyuacc_RlRl_s1(Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpyu (Rs.H,Rt.H)	Word32 Q6_R_mpyunac_RhRh(Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpyu (Rs.H,Rt.H) :<<1	Word32 Q6_R_mpyunac_RhRh_s1(Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpyu (Rs.H,Rt.L)	Word32 Q6_R_mpyunac_RhRl(Word32 Rx, Word32 Rs, Word32 Rt)



Rx==mpyu (Rs.H,Rt.L) :<<1	Word32 Q6_R_mpyunac_RhRl_s1 (Word32 Rx, Word32 Rs, Word32 Rt)
Rx==mpyu (Rs.L,Rt.H)	Word32 Q6_R_mpyunac_RlRh (Word32 Rx, Word32 Rs, Word32 Rt)
Rx==mpyu (Rs.L,Rt.H) :<<1	Word32 Q6_R_mpyunac_RlRh_s1 (Word32 Rx, Word32 Rs, Word32 Rt)
Rx==mpyu (Rs.L,Rt.L)	Word32 Q6_R_mpyunac_RlRl (Word32 Rx, Word32 Rs, Word32 Rt)
Rx==mpyu (Rs.L,Rt.L) :<<1	Word32 Q6_R_mpyunac_RlRl_s1 (Word32 Rx, Word32 Rs, Word32 Rt)
Rxx+=mpyu (Rs.H,Rt.H)	Word64 Q6_P_mpyuacc_RhRh (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx+=mpyu (Rs.H,Rt.H) :<<1	Word64 Q6_P_mpyuacc_RhRh_s1 (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx+=mpyu (Rs.H,Rt.L)	Word64 Q6_P_mpyuacc_RhRl (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx+=mpyu (Rs.H,Rt.L) :<<1	Word64 Q6_P_mpyuacc_RhRl_s1 (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx+=mpyu (Rs.L,Rt.H)	Word64 Q6_P_mpyuacc_RlRh (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx+=mpyu (Rs.L,Rt.H) :<<1	Word64 Q6_P_mpyuacc_RlRh_s1 (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx+=mpyu (Rs.L,Rt.L)	Word64 Q6_P_mpyuacc_RlRl (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx+=mpyu (Rs.L,Rt.L) :<<1	Word64 Q6_P_mpyuacc_RlRl_s1 (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx-=mpyu (Rs.H,Rt.H)	Word64 Q6_P_mpyunac_RhRh (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx-=mpyu (Rs.H,Rt.H) :<<1	Word64 Q6_P_mpyunac_RhRh_s1 (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx-=mpyu (Rs.H,Rt.L)	Word64 Q6_P_mpyunac_RhRl (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx-=mpyu (Rs.H,Rt.L) :<<1	Word64 Q6_P_mpyunac_RhRl_s1 (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx-=mpyu (Rs.L,Rt.H)	Word64 Q6_P_mpyunac_RlRh (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx-=mpyu (Rs.L,Rt.H) :<<1	Word64 Q6_P_mpyunac_RlRh_s1 (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx-=mpyu (Rs.L,Rt.L)	Word64 Q6_P_mpyunac_RlRl (Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx-=mpyu (Rs.L,Rt.L) :<<1	Word64 Q6_P_mpyunac_RlRl_s1 (Word64 Rxx, Word32 Rs, Word32 Rt)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp			s5					Parse		t5					MinOp		d5							
1	1	1	0	0	1	0	0	N	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	0	0	d	d	d	d	d	Rdd=mpyu(Rs.L,Rt.L)[:<<N]
1	1	1	0	0	1	0	0	N	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	0	1	d	d	d	d	d	Rdd=mpyu(Rs.L,Rt.H)[:<<N]
1	1	1	0	0	1	0	0	N	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	1	0	d	d	d	d	d	Rdd=mpyu(Rs.H,Rt.L)[:<<N]
1	1	1	0	0	1	0	0	N	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	1	1	d	d	d	d	d	Rdd=mpyu(Rs.H,Rt.H)[:<<N]
ICLASS				RegType				MajOp			s5					Parse		t5					MinOp		x5							
1	1	1	0	0	1	1	0	N	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	x	x	x	x	x	Rxx+=mpyu(Rs.L,Rt.L)[:<<N]
1	1	1	0	0	1	1	0	N	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	x	x	x	x	x	Rxx+=mpyu(Rs.L,Rt.H)[:<<N]
1	1	1	0	0	1	1	0	N	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	x	x	x	x	x	Rxx+=mpyu(Rs.H,Rt.L)[:<<N]
1	1	1	0	0	1	1	0	N	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	x	x	x	x	x	Rxx+=mpyu(Rs.H,Rt.H)[:<<N]
1	1	1	0	0	1	1	0	N	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	x	x	x	x	x	Rxx-=mpyu(Rs.L,Rt.L)[:<<N]
1	1	1	0	0	1	1	0	N	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	x	x	x	x	x	Rxx-=mpyu(Rs.L,Rt.H)[:<<N]
1	1	1	0	0	1	1	0	N	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	x	x	x	x	x	Rxx-=mpyu(Rs.H,Rt.L)[:<<N]
1	1	1	0	0	1	1	0	N	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	x	x	x	x	x	Rxx-=mpyu(Rs.H,Rt.H)[:<<N]
ICLASS				RegType				MajOp			s5					Parse		t5					MinOp		d5							
1	1	1	0	1	1	0	0	N	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	d	d	d	d	d	Rd=mpyu(Rs.L,Rt.L)[:<<N]
1	1	1	0	1	1	0	0	N	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	d	d	d	d	d	Rd=mpyu(Rs.L,Rt.H)[:<<N]
1	1	1	0	1	1	0	0	N	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	d	d	d	d	d	Rd=mpyu(Rs.H,Rt.L)[:<<N]
1	1	1	0	1	1	0	0	N	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	d	d	d	d	d	Rd=mpyu(Rs.H,Rt.H)[:<<N]
ICLASS				RegType				MajOp			s5					Parse		t5					MinOp		x5							
1	1	1	0	1	1	1	0	N	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	x	x	x	x	x	Rx+=mpyu(Rs.L,Rt.L)[:<<N]
1	1	1	0	1	1	1	0	N	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	x	x	x	x	x	Rx+=mpyu(Rs.L,Rt.H)[:<<N]
1	1	1	0	1	1	1	0	N	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	x	x	x	x	x	Rx+=mpyu(Rs.H,Rt.L)[:<<N]
1	1	1	0	1	1	1	0	N	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	x	x	x	x	x	Rx+=mpyu(Rs.H,Rt.H)[:<<N]
1	1	1	0	1	1	1	0	N	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	x	x	x	x	x	Rx-=mpyu(Rs.L,Rt.L)[:<<N]
1	1	1	0	1	1	1	0	N	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	1	x	x	x	x	x	Rx-=mpyu(Rs.L,Rt.H)[:<<N]
1	1	1	0	1	1	1	0	N	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	x	x	x	x	x	Rx-=mpyu(Rs.H,Rt.L)[:<<N]
1	1	1	0	1	1	1	0	N	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	x	x	x	x	x	Rx-=mpyu(Rs.H,Rt.H)[:<<N]

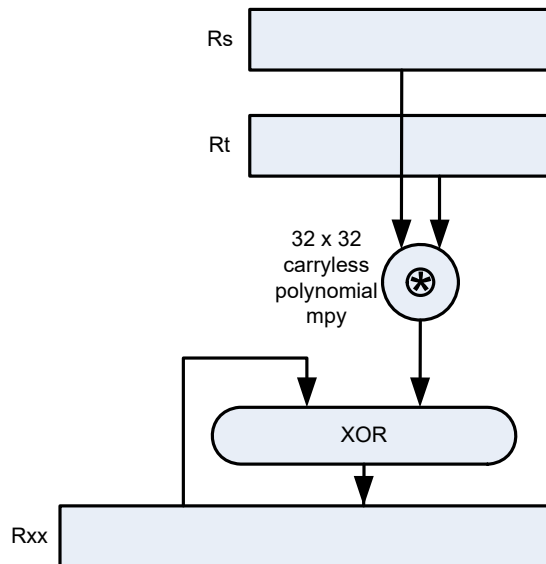
<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
sH	Rs is high
tH	Rt is high

<b>Field name</b>	<b>Description</b>
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Polynomial multiply words

Perform a  $32 \times 32$  carryless polynomial multiply using 32-bit source registers Rs and Rt. The 64-bit result is optionally accumulated (XORed) with the destination register. Finite field multiply instructions are useful for many algorithms including scramble code generation, cryptographic algorithms, convolutional, and Reed Solomon codes.

`Rxx += pmpyw(Rs,Rt)`



### Syntax

`Rdd=pmpyw (Rs, Rt)`

```
x = Rs.uw[0];
y = Rt.uw[0];
prod = 0;
for(i=0; i < 32; i++) {
    if((y >> i) & 1) prod ^= (x << i);
}
Rdd = prod;
```

### Behavior

`Rxx^=pmpyw (Rs, Rt)`

```
x = Rs.uw[0];
y = Rt.uw[0];
prod = 0;
for(i=0; i < 32; i++) {
    if((y >> i) & 1) prod ^= (x << i);
}
Rxx ^= prod;
```

**Class: XTYPE (slots 2,3)****Intrinsics**

Rdd=pmpyw (Rs, Rt)

Word64 Q6\_P\_pmpyw\_RR (Word32 Rs, Word32 Rt)

Rxx^=pmpyw (Rs, Rt)

Word64 Q6\_P\_pmpywxacc\_RR (Word64 Rxx, Word32 Rs, Word32 Rt)

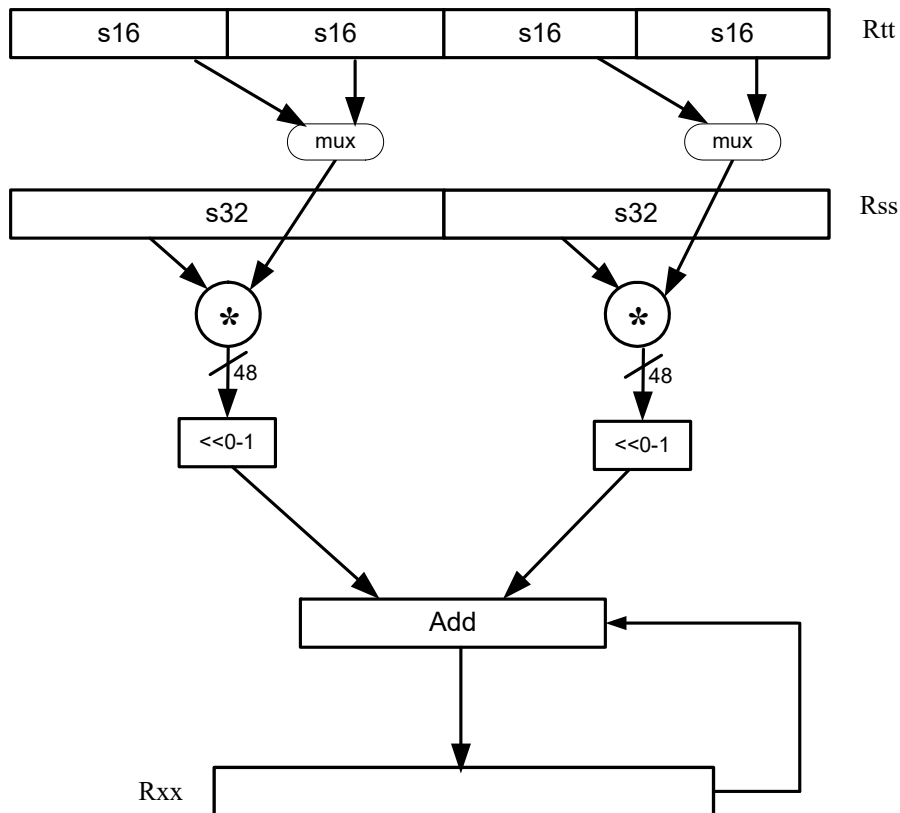
**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			d5							
1	1	1	0	0	1	0	1	0	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	1	d	d	d	d	d	Rdd=pmpyw(Rs,Rt)
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			x5							
1	1	1	0	0	1	1	1	0	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	1	x	x	x	x	x	Rxx^=pmpyw(Rs,Rt)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Vector reduce multiply word by signed half (32 × 16)

Perform mixed precision vector multiply operations and accumulate the results. A 32-bit word from vector *Rss* is multiplied by a 16-bit halfword (either even or odd) from vector *Rtt*. The multiplication is performed as a signed 32 × 16, which produces a 48-bit result. This result is optionally scaled left by one bit. A similar operation is performed for both words in *Rss*, and accumulates the two results. The final result optionally accumulates with *Rxx*.



Syntax	Behavior
<code>Rdd=vrmpyweh (Rss,Rtt) [:&lt;&lt;1]</code>	$Rdd = (Rss.w[1] * Rtt.h[2]) [ <<1 ] + (Rss.w[0] * Rtt.h[0]) [ <<1 ] ;$
<code>Rdd=vrmpywoh (Rss,Rtt) [:&lt;&lt;1]</code>	$Rdd = (Rss.w[1] * Rtt.h[3]) [ <<1 ] + (Rss.w[0] * Rtt.h[1]) [ <<1 ] ;$
<code>Rxx+=vrmpyweh (Rss,Rtt) [:&lt;&lt;1]</code>	$Rxx += (Rss.w[1] * Rtt.h[2]) [ <<1 ] + (Rss.w[0] * Rtt.h[0]) [ <<1 ] ;$
<code>Rxx+=vrmpywoh (Rss,Rtt) [:&lt;&lt;1]</code>	$Rxx += (Rss.w[1] * Rtt.h[3]) [ <<1 ] + (Rss.w[0] * Rtt.h[1]) [ <<1 ] ;$

**Class: XTYPE (slots 2,3)****Intrinsics**

Rdd=vrmpyweh (Rss, Rtt)	Word64 Q6_P_vrmpyweh_PP (Word64 Rss, Word64 Rtt)
Rdd=vrmpyweh (Rss, Rtt) :<< 1	Word64 Q6_P_vrmpyweh_PP_s1 (Word64 Rss, Word64 Rtt)
Rdd=vrmpywoh (Rss, Rtt)	Word64 Q6_P_vrmpywoh_PP (Word64 Rss, Word64 Rtt)
Rdd=vrmpywoh (Rss, Rtt) :<< 1	Word64 Q6_P_vrmpywoh_PP_s1 (Word64 Rss, Word64 Rtt)
Rxx+=vrmpyweh (Rss, Rtt)	Word64 Q6_P_vrmpywehacc_PP (Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=vrmpyweh (Rss, Rtt) :< <1	Word64 Q6_P_vrmpywehacc_PP_s1 (Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=vrmpywoh (Rss, Rtt)	Word64 Q6_P_vrmpywohacc_PP (Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=vrmpywoh (Rss, Rtt) :< <1	Word64 Q6_P_vrmpywohacc_PP_s1 (Word64 Rxx, Word64 Rss, Word64 Rtt)

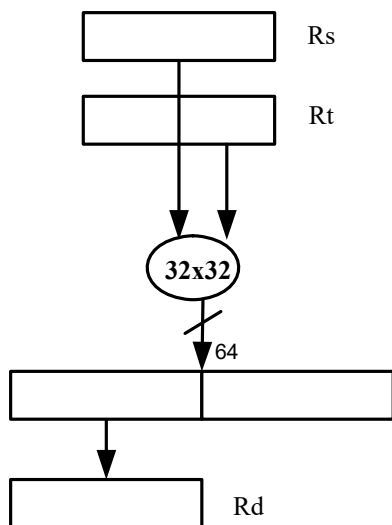
**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse		t5					MinOp		d5									
1	1	1	0	1	0	0	0	N	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	0	d	d	d	d	d	Rdd=vrmpywoh(Rss,Rtt)[:<N]
1	1	1	0	1	0	0	0	N	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	0	d	d	d	d	d	Rdd=vrmpyweh(Rss,Rtt)[:<N]
ICLASS			RegType				MajOp		s5					Parse		t5					MinOp		x5									
1	1	1	0	1	0	1	0	N	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	0	x	x	x	x	x	Rxx+=vrmpyweh(Rss,Rtt)[:<N]
1	1	1	0	1	0	1	0	N	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	0	x	x	x	x	x	Rxx+=vrmpywoh(Rss,Rtt)[:<N]

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Multiply and use upper result

Multiply two signed or unsigned 32-bit words. Store the upper 32-bits of this result to a single destination register. Optional rounding is available.



Syntax	Behavior
$Rd = \text{mpy}(Rs, Rt.H) : \ll 1 : \text{rnd} : \text{sat}$	$Rd = \text{sat}_{32}(((Rs * Rt.h[1]) \ll 1 + 0x8000) \gg 16);$
$Rd = \text{mpy}(Rs, Rt.H) : \ll 1 : \text{sat}$	$Rd = \text{sat}_{32}(((Rs * Rt.h[1]) \ll 1) \gg 16);$
$Rd = \text{mpy}(Rs, Rt.L) : \ll 1 : \text{rnd} : \text{sat}$	$Rd = \text{sat}_{32}(((Rs * Rt.h[0]) \ll 1 + 0x8000) \gg 16);$
$Rd = \text{mpy}(Rs, Rt.L) : \ll 1 : \text{sat}$	$Rd = \text{sat}_{32}(((Rs * Rt.h[0]) \ll 1) \gg 16);$
$Rd = \text{mpy}(Rs, Rt)$	$Rd = (Rs * Rt) \gg 32;$
$Rd = \text{mpy}(Rs, Rt) : \ll 1$	$Rd = (Rs * Rt) \gg 31;$
$Rd = \text{mpy}(Rs, Rt) : \ll 1 : \text{sat}$	$Rd = \text{sat}_{32}((Rs * Rt) \gg 31);$
$Rd = \text{mpy}(Rs, Rt) : \text{rnd}$	$Rd = ((Rs * Rt) + 0x80000000) \gg 32;$
$Rd = \text{mpysu}(Rs, Rt)$	$Rd = (Rs * Rt.uw[0]) \gg 32;$
$Rd = \text{mpyu}(Rs, Rt)$	$Rd = (Rs.uw[0] * Rt.uw[0]) \gg 32;$
$Rx += \text{mpy}(Rs, Rt) : \ll 1 : \text{sat}$	$Rx = \text{sat}_{32}((Rx) + ((Rs * Rt) \gg 31));$
$Rx -= \text{mpy}(Rs, Rt) : \ll 1 : \text{sat}$	$Rx = \text{sat}_{32}((Rx) - ((Rs * Rt) \gg 31));$

### Class: XTYPE (slots 2,3)

#### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.



## Intrinsics

Rd=mpy(Rs,Rt.H):<<1:rnd:sat	Word32 Q6_R_mpy_RRh_s1_rnd_sat(Word32 Rs, Word32 Rt)
Rd=mpy(Rs,Rt.H):<<1:sat	Word32 Q6_R_mpy_RRh_s1_sat(Word32 Rs, Word32 Rt)
Rd=mpy(Rs,Rt.L):<<1:rnd:sat	Word32 Q6_R_mpy_RRl_s1_rnd_sat(Word32 Rs, Word32 Rt)
Rd=mpy(Rs,Rt.L):<<1:sat	Word32 Q6_R_mpy_RRl_s1_sat(Word32 Rs, Word32 Rt)
Rd=mpy(Rs,Rt)	Word32 Q6_R_mpy_RR(Word32 Rs, Word32 Rt)
Rd=mpy(Rs,Rt):<<1	Word32 Q6_R_mpy_RR_s1(Word32 Rs, Word32 Rt)
Rd=mpy(Rs,Rt):<<1:sat	Word32 Q6_R_mpy_RR_s1_sat(Word32 Rs, Word32 Rt)
Rd=mpy(Rs,Rt):rnd	Word32 Q6_R_mpy_RR_rnd(Word32 Rs, Word32 Rt)
Rd=mpysu(Rs,Rt)	Word32 Q6_R_mpysu_RR(Word32 Rs, Word32 Rt)
Rd=mpyu(Rs,Rt)	UWord32 Q6_R_mpyu_RR(Word32 Rs, Word32 Rt)
Rx+=mpy(Rs,Rt):<<1:sat	Word32 Q6_R_mpyacc_RR_s1_sat(Word32 Rx, Word32 Rs, Word32 Rt)
Rx-=mpy(Rs,Rt):<<1:sat	Word32 Q6_R_mpynac_RR_s1_sat(Word32 Rx, Word32 Rs, Word32 Rt)

## Encoding

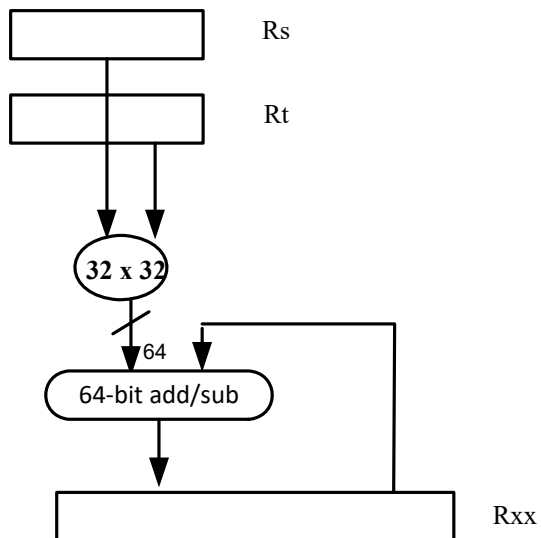
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse		t5					MinOp			d5								
1	1	1	0	1	1	0	1	0	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	d	d	d	d	d	Rd=mpy(Rs,Rt):rnd
1	1	1	0	1	1	0	1	0	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	d	d	d	d	d	Rd=mpyu(Rs,Rt)
1	1	1	0	1	1	0	1	0	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	d	d	d	d	d	Rd=mpysu(Rs,Rt)
1	1	1	0	1	1	0	1	1	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	d	d	d	d	d	Rd=mpy(Rs,Rt.H):<<1:sat
1	1	1	0	1	1	0	1	1	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	d	d	d	d	d	Rd=mpy(Rs,Rt.L):<<1:sat
1	1	1	0	1	1	0	1	1	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	0	d	d	d	d	d	Rd=mpy(Rs,Rt.H):<<1:rnd:sat
1	1	1	0	1	1	0	1	1	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	d	d	d	d	d	Rd=mpy(Rs,Rt):<<1:sat
1	1	1	0	1	1	0	1	1	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	0	d	d	d	d	d	Rd=mpy(Rs,Rt.L):<<1:rnd:sat
1	1	1	0	1	1	0	1	N	0	N	s	s	s	s	s	P	P	0	t	t	t	t	t	0	N	N	d	d	d	d	d	Rd=mpy(Rs,Rt)[:<<N]
ICLASS			RegType				MajOp		s5					Parse		t5					MinOp			x5								
1	1	1	0	1	1	1	1	0	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	x	x	x	x	x	Rx+=mpy(Rs,Rt):<<1:sat
1	1	1	0	1	1	1	1	0	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	x	x	x	x	x	Rx-=mpy(Rs,Rt):<<1:sat

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d

<b>Field name</b>	<b>Description</b>
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Multiply and use full result

Multiply two signed or unsigned 32-bit words. Optionally, add or subtract this value from the 64-bit accumulator. The result is a full-precision 64-bit value.



Syntax	Behavior
<code>Rdd=mpy (Rs, Rt)</code>	<code>Rdd = (Rs * Rt);</code>
<code>Rdd=mpyu (Rs, Rt)</code>	<code>Rdd = (Rs.uw[0] * Rt.uw[0]);</code>
<code>Rxx[+-]=mpy (Rs, Rt)</code>	<code>Rxx = Rxx [+-] (Rs * Rt);</code>
<code>Rxx[+-]=mpyu (Rs, Rt)</code>	<code>Rxx = Rxx [+-] (Rs.uw[0] * Rt.uw[0]);</code>

### Class: XTYPE (slots 2,3)

#### Intrinsics

<code>Rdd=mpy (Rs, Rt)</code>	<code>Word64 Q6_P_mpy_RR (Word32 Rs, Word32 Rt)</code>
<code>Rdd=mpyu (Rs, Rt)</code>	<code>UWord64 Q6_P_mpyu_RR (Word32 Rs, Word32 Rt)</code>
<code>Rxx+=mpy (Rs, Rt)</code>	<code>Word64 Q6_P_mpyacc_RR (Word64 Rxx, Word32 Rs, Word32 Rt)</code>
<code>Rxx+=mpyu (Rs, Rt)</code>	<code>Word64 Q6_P_mpyuacc_RR (Word64 Rxx, Word32 Rs, Word32 Rt)</code>
<code>Rxx-=mpy (Rs, Rt)</code>	<code>Word64 Q6_P_mpynac_RR (Word64 Rxx, Word32 Rs, Word32 Rt)</code>
<code>Rxx-=mpyu (Rs, Rt)</code>	<code>Word64 Q6_P_mpyunac_RR (Word64 Rxx, Word32 Rs, Word32 Rt)</code>

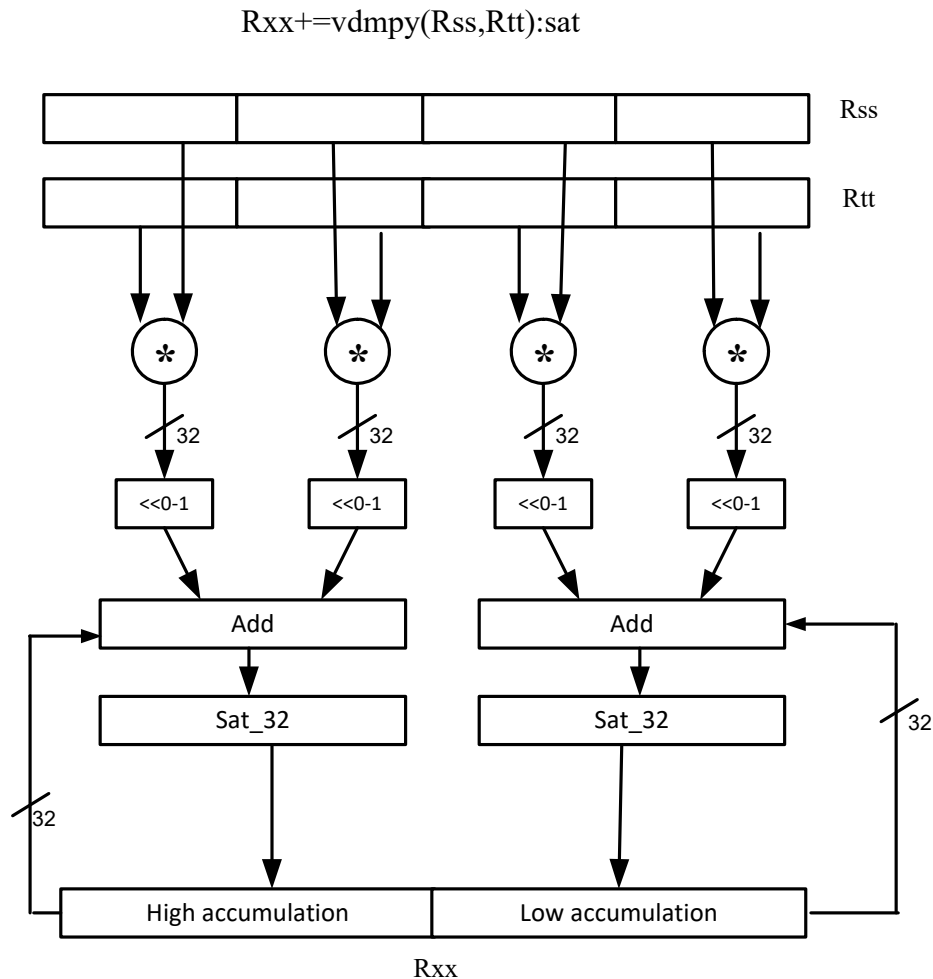
## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			d5							
1	1	1	0	0	1	0	1	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	d	d	d	d	d	Rdd=mpy(Rs,Rt)
1	1	1	0	0	1	0	1	0	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	d	d	d	d	d	Rdd=mpyu(Rs,Rt)
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			x5							
1	1	1	0	0	1	1	1	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	x	x	x	x	x	Rxx+=mpy(Rs,Rt)
1	1	1	0	0	1	1	1	0	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	x	x	x	x	x	Rxx+=mpy(Rs,Rt)
1	1	1	0	0	1	1	1	0	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	x	x	x	x	x	Rxx+=mpyu(Rs,Rt)
1	1	1	0	0	1	1	1	0	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	x	x	x	x	x	Rxx-=mpyu(Rs,Rt)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Vector dual multiply

Multiply four 16-bit halfwords in *Rss* by the corresponding 16-bit halfwords in *Rtt*. Add and scale the two lower results. Optionally add the lower word of the accumulator. Saturate this result to 32-bits and store in the lower word of the accumulator. Perform the same operation on the upper two products using the upper word of the accumulator.



### Syntax

```
Rdd=vdmpy(Rss,Rtt):<<1:sat
```

```
Rdd=vdmpy(Rss,Rtt):sat
```

### Behavior

```
Rdd.w[0]=sat32((Rss.h[0] * Rtt.h[0])<<1 + (Rss.h[1] * Rtt.h[1])<<1);
Rdd.w[1]=sat32((Rss.h[2] * Rtt.h[2])<<1 + (Rss.h[3] * Rtt.h[3])<<1);
```

```
Rdd.w[0]=sat32((Rss.h[0] * Rtt.h[0])<<0 + (Rss.h[1] * Rtt.h[1])<<0);
Rdd.w[1]=sat32((Rss.h[2] * Rtt.h[2])<<0 + (Rss.h[3] * Rtt.h[3])<<0);
```

Syntax	Behavior
$Rxx += \text{vdmpy}(Rss, Rtt) : \ll 1 : \text{sat}$	$Rxx.w[0] = \text{sat}_{32}(Rxx.w[0] + (Rss.h[0] * Rtt.h[0]) \ll 1 + (Rss.h[1] * Rtt.h[1]) \ll 1);$ $Rxx.w[1] = \text{sat}_{32}(Rxx.w[1] + (Rss.h[2] * Rtt.h[2]) \ll 1 + (Rss.h[3] * Rtt.h[3]) \ll 1);$
$Rxx += \text{vdmpy}(Rss, Rtt) : \text{sat}$	$Rxx.w[0] = \text{sat}_{32}(Rxx.w[0] + (Rss.h[0] * Rtt.h[0]) \ll 0 + (Rss.h[1] * Rtt.h[1]) \ll 0);$ $Rxx.w[1] = \text{sat}_{32}(Rxx.w[1] + (Rss.h[2] * Rtt.h[2]) \ll 0 + (Rss.h[3] * Rtt.h[3]) \ll 0);$

### Class: XTYPE (slots 2,3)

#### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

#### Intrinsics

$Rdd = \text{vdmpy}(Rss, Rtt) : \ll 1 : \text{sat}$	Word64 Q6_P_vdmpy_PP_s1_sat(Word64 Rss, Word64 Rtt)
$Rdd = \text{vdmpy}(Rss, Rtt) : \text{sat}$	Word64 Q6_P_vdmpy_PP_sat(Word64 Rss, Word64 Rtt)
$Rxx += \text{vdmpy}(Rss, Rtt) : \ll 1 : \text{sat}$	Word64 Q6_P_vdmpyacc_PP_s1_sat(Word64 Rxx, Word64 Rss, Word64 Rtt)
$Rxx += \text{vdmpy}(Rss, Rtt) : \text{sat}$	Word64 Q6_P_vdmpyacc_PP_sat(Word64 Rxx, Word64 Rss, Word64 Rtt)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			d5							
1	1	1	0	1	0	0	0	N	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	0	d	d	d	d	d	Rdd=vdmpy(Rss,Rtt):<<N]:sat
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			x5							
1	1	1	0	1	0	1	0	N	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	0	x	x	x	x	x	Rxx+=vdmpy(Rss,Rtt):<<N]:sat

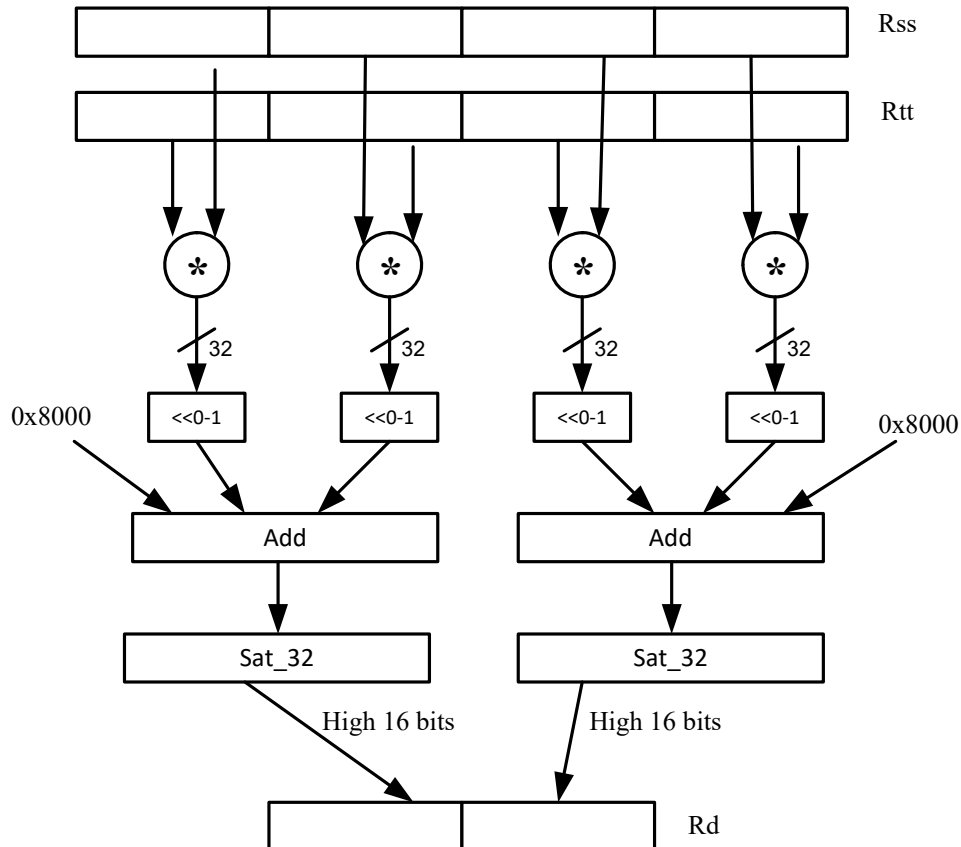
Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d

<b>Field name</b>	<b>Description</b>
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Vector dual multiply with round and pack

Multiply four 16-bit halfwords in *Rss* by the corresponding 16-bit halfwords in *Rtt*. Scale and add together the two lower results with a rounding constant. Saturate this result to 32-bits, and store the upper 16-bits of this result in the lower 16-bits of the destination register. The same operation is performed on the upper two products and the result is stored in the upper 16-bit halfword of the destination.

$Rd = \text{vdmpy}(Rss, Rtt) : \text{rnd} : \text{sat}$



### Syntax

```
Rd=vdmpy(Rss,Rtt)[:<<1]:rnd:
sat
```

### Behavior

```
Rd.h[0]=(sat32((Rss.h[0] * Rtt.h[0]) [<<1] +
(Rss.h[1] * Rtt.h[1]) [<<1] + 0x8000)).h[1];
Rd.h[1]=(sat32((Rss.h[2] * Rtt.h[2]) [<<1] +
(Rss.h[3] * Rtt.h[3]) [<<1] + 0x8000)).h[1];
```

**Class:** XTYPE (slots 2,3)

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.



## Intrinsics

```
Rd=vdmpy(Rss,Rtt):<<1:rnd:sa Word32 Q6_R_vdmpy_PP_s1_rnd_sat(Word64 Rss, Word64
t Rtt)
Rd=vdmpy(Rss,Rtt):rnd:sat Word32 Q6_R_vdmpy_PP_rnd_sat(Word64 Rss, Word64
Rtt)
```

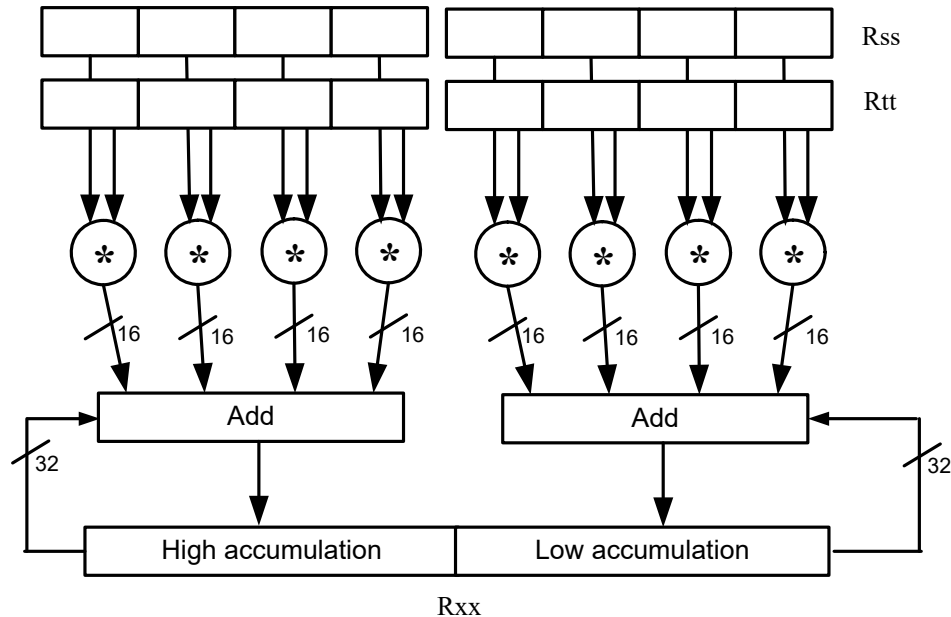
## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse		t5				MinOp			d5								
1	1	1	0	1	0	0	1	N	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	d	d	d	d	d	Rd=vdmpy(Rss,Rtt):<<N]:r nd:sat

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector reduce multiply bytes

Multiply eight 8-bit bytes in Rss by the corresponding 8-bit bytes in Rtt. Accumulate the four lower results. Optionally add the lower word of the accumulator. Store this result in the lower 32-bits of the accumulator. The same operation is performed on the upper four products using the upper word of the accumulator. The eight bytes of Rss are treated as either signed or unsigned.



Syntax	Behavior
<code>Rdd=vrmpybsu (Rss, Rtt)</code>	$\begin{aligned} Rdd.w[0] &= (Rss.b[0] * Rtt.ub[0]) + (Rss.b[1] * Rtt.ub[1]) \\ &+ (Rss.b[2] * Rtt.ub[2]) + (Rss.b[3] * Rtt.ub[3]); \\ Rdd.w[1] &= (Rss.b[4] * Rtt.ub[4]) + (Rss.b[5] * Rtt.ub[5]) \\ &+ (Rss.b[6] * Rtt.ub[6]) + (Rss.b[7] * Rtt.ub[7]); \end{aligned}$
<code>Rdd=vrmpybu (Rss, Rtt)</code>	$\begin{aligned} Rdd.w[0] &= (Rss.ub[0] * Rtt.ub[0]) + (Rss.ub[1] * \\ &Rtt.ub[1]) + (Rss.ub[2] * Rtt.ub[2]) + (Rss.ub[3] * \\ &Rtt.ub[3]); \\ Rdd.w[1] &= (Rss.ub[4] * Rtt.ub[4]) + (Rss.ub[5] * \\ &Rtt.ub[5]) + (Rss.ub[6] * Rtt.ub[6]) + (Rss.ub[7] * \\ &Rtt.ub[7]); \end{aligned}$
<code>Rxx+=vrmpybsu (Rss, Rtt)</code>	$\begin{aligned} Rxx.w[0] &= (Rxx.w[0] + (Rss.b[0] * Rtt.ub[0]) + (Rss.b[1] * \\ &Rtt.ub[1]) + (Rss.b[2] * Rtt.ub[2]) + (Rss.b[3] * \\ &Rtt.ub[3])); \\ Rxx.w[1] &= (Rxx.w[1] + (Rss.b[4] * Rtt.ub[4]) + (Rss.b[5] * \\ &Rtt.ub[5]) + (Rss.b[6] * Rtt.ub[6]) + (Rss.b[7] * \\ &Rtt.ub[7])); \end{aligned}$
<code>Rxx+=vrmpybu (Rss, Rtt)</code>	$\begin{aligned} Rxx.w[0] &= (Rxx.w[0] + (Rss.ub[0] * Rtt.ub[0]) + (Rss.ub[1] * \\ &Rtt.ub[1]) + (Rss.ub[2] * Rtt.ub[2]) + (Rss.ub[3] * \\ &Rtt.ub[3])); \\ Rxx.w[1] &= (Rxx.w[1] + (Rss.ub[4] * Rtt.ub[4]) + (Rss.ub[5] * \\ &Rtt.ub[5]) + (Rss.ub[6] * Rtt.ub[6]) + (Rss.ub[7] * \\ &Rtt.ub[7])); \end{aligned}$

**Class: XTYPE (slots 2,3)****Intrinsics**

Rdd=vrmpybsu (Rss, Rtt)	Word64 Q6_P_vrmpybsu_PP(Word64 Rss, Word64 Rtt)
Rdd=vrmpybu (Rss, Rtt)	Word64 Q6_P_vrmpybu_PP(Word64 Rss, Word64 Rtt)
Rxx+=vrmpybsu (Rss, Rtt)	Word64 Q6_P_vrmpybsuacc_PP(Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=vrmpybu (Rss, Rtt)	Word64 Q6_P_vrmpybuacc_PP(Word64 Rxx, Word64 Rss, Word64 Rtt)

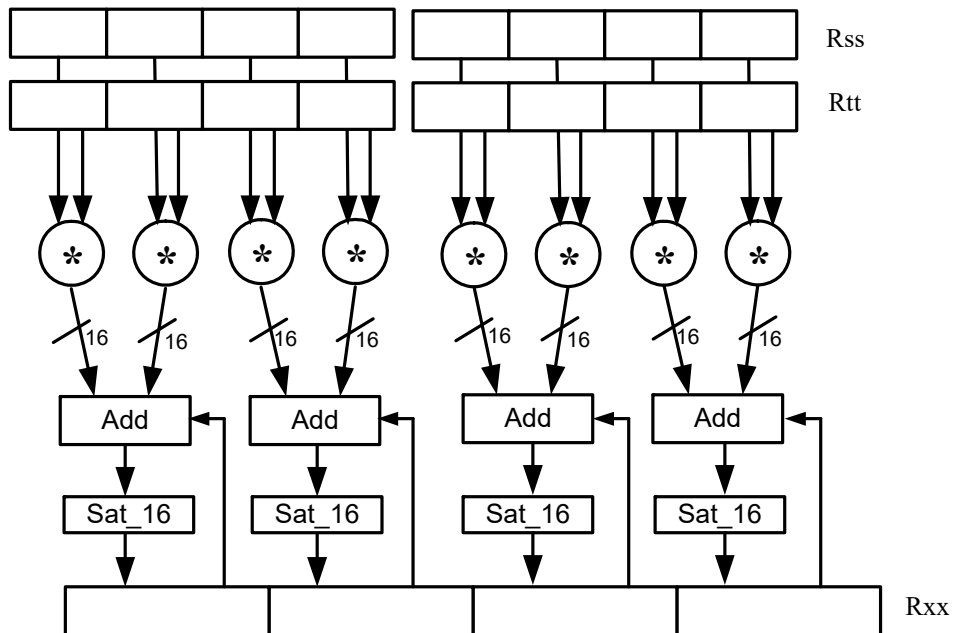
**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			d5							
1	1	1	0	1	0	0	0	1	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	d	d	d	d	d	Rdd=vrmpybu(Rss,Rtt)
1	1	1	0	1	0	0	0	1	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	d	d	d	d	d	Rdd=vrmpybsu(Rss,Rtt)
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			x5							
1	1	1	0	1	0	1	0	1	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	x	x	x	x	x	Rxx+=vrmpybu(Rss,Rtt)
1	1	1	0	1	0	1	0	1	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	x	x	x	x	x	Rxx+=vrmpybsu(Rss,Rtt)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Vector dual multiply signed by unsigned bytes

Multiply eight 8-bit signed bytes in *Rss* by the corresponding 8-bit unsigned bytes in *Rtt*. Add the results in pairs, and optionally add the accumulator. Saturate the results to signed 16 bits and store in the four halfwords of the destination register.



### Syntax

```
Rdd=vdmpybsu(Rss,Rtt):sat
```

```
Rxx+=vdmpybsu(Rss,Rtt):sat
```

### Behavior

```
Rdd.h[0]=sat16((Rss.b[0] * Rtt.ub[0]) + (Rss.b[1] * Rtt.ub[1]));
Rdd.h[1]=sat16((Rss.b[2] * Rtt.ub[2]) + (Rss.b[3] * Rtt.ub[3]));
Rdd.h[2]=sat16((Rss.b[4] * Rtt.ub[4]) + (Rss.b[5] * Rtt.ub[5]));
Rdd.h[3]=sat16((Rss.b[6] * Rtt.ub[6]) + (Rss.b[7] * Rtt.ub[7]));
```

```
Rxx.h[0]=sat16((Rxx.h[0] + (Rss.b[0] * Rtt.ub[0]) + (Rss.b[1] * Rtt.ub[1]));
Rxx.h[1]=sat16((Rxx.h[1] + (Rss.b[2] * Rtt.ub[2]) + (Rss.b[3] * Rtt.ub[3]));
Rxx.h[2]=sat16((Rxx.h[2] + (Rss.b[4] * Rtt.ub[4]) + (Rss.b[5] * Rtt.ub[5]));
Rxx.h[3]=sat16((Rxx.h[3] + (Rss.b[6] * Rtt.ub[6]) + (Rss.b[7] * Rtt.ub[7]));
```

**Class: XTYPE (slots 2,3)****Notes**

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

**Intrinsics**

```
Rdd=vdmpybsu(Rss,Rtt):sat Word64 Q6_P_vdmpybsu_PP_sat(Word64 Rss,
Word64 Rtt)
```

```
Rxx+=vdmpybsu(Rss,Rtt):sat Word64 Q6_P_vdmpybsuacc_PP_sat(Word64 Rxx,
Word64 Rss, Word64 Rtt)
```

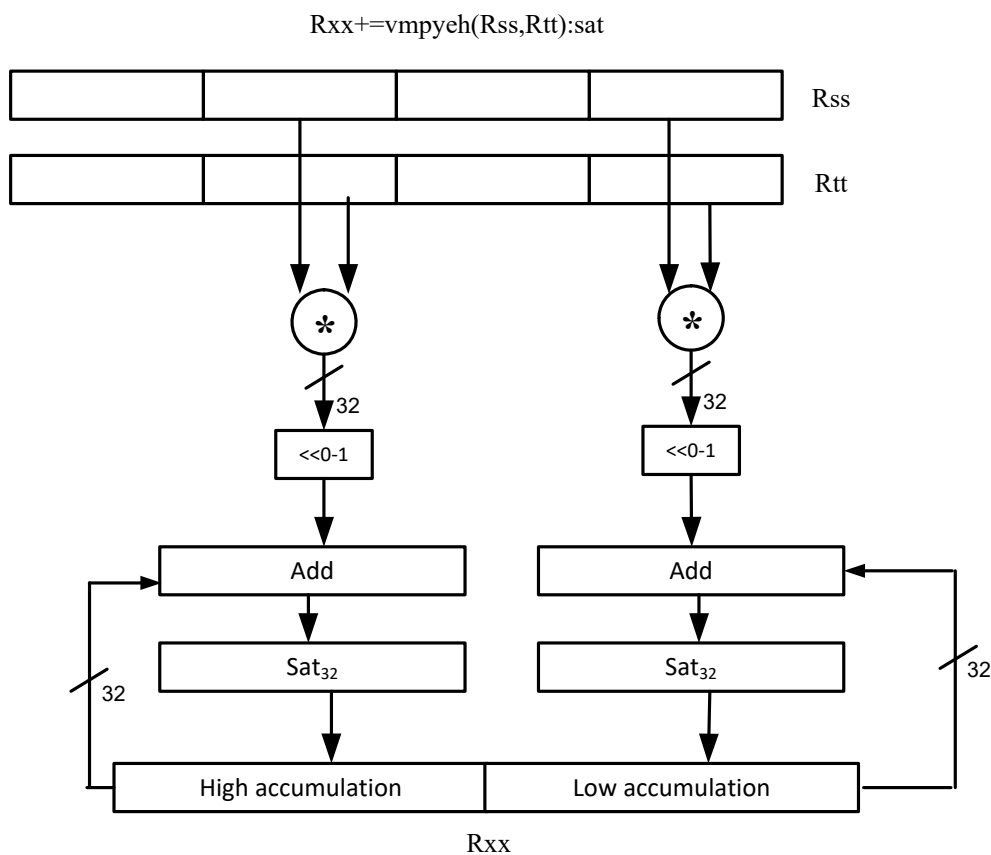
**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			d5							
1	1	1	0	1	0	0	0	1	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	d	d	d	d	d	Rdd=vdmpybsu(Rss,Rtt):sat
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			x5							
1	1	1	0	1	0	1	0	0	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	x	x	x	x	x	Rxx+=vdmpybsu(Rss,Rtt):sat

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Vector multiply even halfwords

Multiply the even 16-bit halfwords from Rss and Rtt separately. Optionally accumulate with the low and high words of the destination register pair and optionally saturate.



Syntax	Behavior
$Rdd = \text{vmpyeh}(Rss, Rtt) : \ll 1 : \text{sat}$	$Rdd.w[0] = \text{sat}_{32}((Rss.h[0] * Rtt.h[0]) \ll 1);$ $Rdd.w[1] = \text{sat}_{32}((Rss.h[2] * Rtt.h[2]) \ll 1);$
$Rdd = \text{vmpyeh}(Rss, Rtt) : \text{sat}$	$Rdd.w[0] = \text{sat}_{32}((Rss.h[0] * Rtt.h[0]) \ll 0);$ $Rdd.w[1] = \text{sat}_{32}((Rss.h[2] * Rtt.h[2]) \ll 0);$
$Rxx += \text{vmpyeh}(Rss, Rtt)$	$Rxx.w[0] = Rxx.w[0] + (Rss.h[0] * Rtt.h[0]);$ $Rxx.w[1] = Rxx.w[1] + (Rss.h[2] * Rtt.h[2]);$
$Rxx += \text{vmpyeh}(Rss, Rtt) : \ll 1 : \text{sat}$	$Rxx.w[0] = \text{sat}_{32}(Rxx.w[0] + (Rss.h[0] * Rtt.h[0]) \ll 1);$ $Rxx.w[1] = \text{sat}_{32}(Rxx.w[1] + (Rss.h[2] * Rtt.h[2]) \ll 1);$
$Rxx += \text{vmpyeh}(Rss, Rtt) : \text{sat}$	$Rxx.w[0] = \text{sat}_{32}(Rxx.w[0] + (Rss.h[0] * Rtt.h[0]) \ll 0);$ $Rxx.w[1] = \text{sat}_{32}(Rxx.w[1] + (Rss.h[2] * Rtt.h[2]) \ll 0);$

**Class: XTYPE (slots 2,3)****Notes**

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

**Intrinsics**

Rdd=vmpyeh(Rss,Rtt):<<1:sat	Word64 Q6_P_vmpyeh_PP_s1_sat(Word64 Rss, Word64 Rtt)
Rdd=vmpyeh(Rss,Rtt):sat	Word64 Q6_P_vmpyeh_PP_sat(Word64 Rss, Word64 Rtt)
Rxx+=vmpyeh(Rss,Rtt)	Word64 Q6_P_vmpyehacc_PP(Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=vmpyeh(Rss,Rtt):<<1:sat	Word64 Q6_P_vmpyehacc_PP_s1_sat(Word64 Rxx, Word64 Rss, Word64 Rtt)
Rxx+=vmpyeh(Rss,Rtt):sat	Word64 Q6_P_vmpyehacc_PP_sat(Word64 Rxx, Word64 Rss, Word64 Rtt)

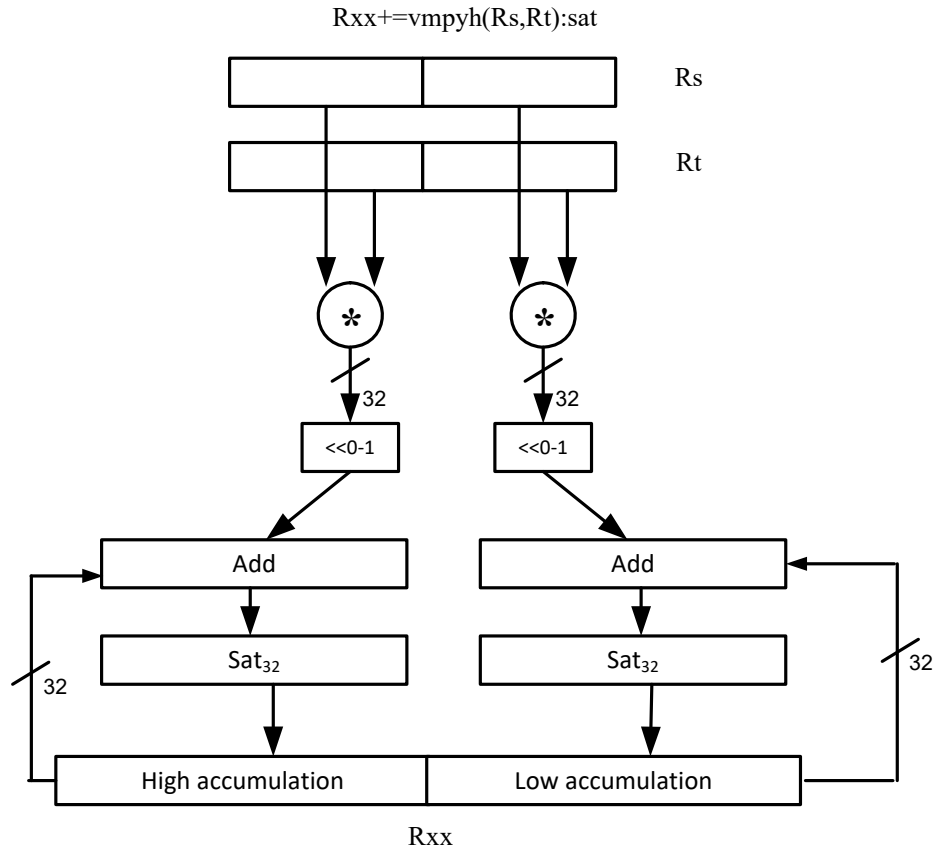
**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			d5							
1	1	1	0	1	0	0	0	N	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	0	d	d	d	d	d	Rdd=vmpyeh(Rss,Rtt):<<N]:sat
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			x5							
1	1	1	0	1	0	1	0	0	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	0	x	x	x	x	x	Rxx+=vmpyeh(Rss,Rtt)
1	1	1	0	1	0	1	0	N	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	0	x	x	x	x	x	Rxx+=vmpyeh(Rss,Rtt):<<N]:sat

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Vector multiply halfwords

Multiply two 16-bit halfwords separately, and optionally accumulate with the low and high words of the destination. Optionally saturate, and store the results back to the destination register pair.



Syntax	Behavior
$R_{dd} = \text{vmpyh}(R_s, R_t) [ : \ll 1 ] : \text{sat}$	$R_{dd}.w[0] = \text{sat}_{32}((R_s.h[0] * R_t.h[0]) [\ll 1]);$ $R_{dd}.w[1] = \text{sat}_{32}((R_s.h[1] * R_t.h[1]) [\ll 1]);$
$R_{xx} += \text{vmpyh}(R_s, R_t)$	$R_{xx}.w[0] = R_{xx}.w[0] + (R_s.h[0] * R_t.h[0]);$ $R_{xx}.w[1] = R_{xx}.w[1] + (R_s.h[1] * R_t.h[1]);$
$R_{xx} += \text{vmpyh}(R_s, R_t) [ : \ll 1 ] : \text{sat}$	$R_{xx}.w[0] = \text{sat}_{32}(R_{xx}.w[0] + (R_s.h[0] * R_t.h[0]) [\ll 1]);$ $R_{xx}.w[1] = \text{sat}_{32}(R_{xx}.w[1] + (R_s.h[1] * R_t.h[1]) [\ll 1]);$

**Class: XTYPE (slots 2,3)**

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.



## Intrinsics

Rdd=vmpyh(Rs,Rt):<<1:sat	Word64 Q6_P_vmpyh_RR_s1_sat(Word32 Rs, Word32 Rt)
Rdd=vmpyh(Rs,Rt):sat	Word64 Q6_P_vmpyh_RR_sat(Word32 Rs, Word32 Rt)
Rxx+=vmpyh(Rs,Rt)	Word64 Q6_P_vmpyhacc_RR(Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx+=vmpyh(Rs,Rt):<<1:sat	Word64 Q6_P_vmpyhacc_RR_s1_sat(Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx+=vmpyh(Rs,Rt):sat	Word64 Q6_P_vmpyhacc_RR_sat(Word64 Rxx, Word32 Rs, Word32 Rt)

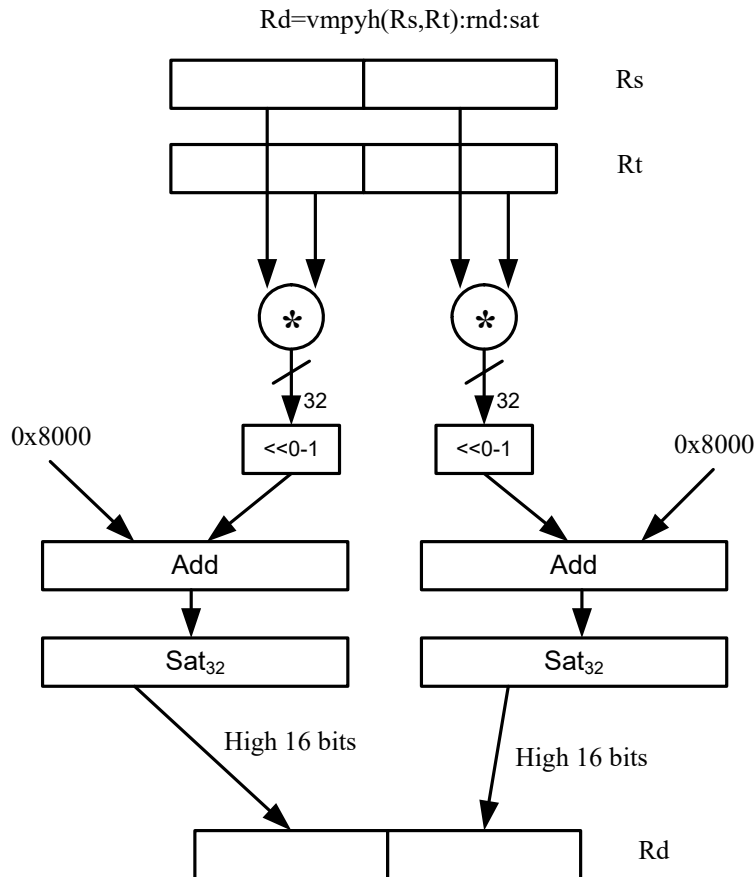
## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			d5							
1	1	1	0	0	1	0	1	N	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	1	d	d	d	d	d	Rdd=vmpyh(Rs,Rt):<<N]:sat
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			x5							
1	1	1	0	0	1	1	1	0	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	x	x	x	x	x	Rxx+=vmpyh(Rs,Rt)
1	1	1	0	0	1	1	1	N	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	1	x	x	x	x	x	Rxx+=vmpyh(Rs,Rt):<<N]:sat

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Vector multiply halfwords with round and pack

Multiply two 16-bit halfwords separately. Round the results, and store the high halfwords packed in a single register destination.



### Syntax

```
Rd=vmpyh(Rs,Rt)[:<<1]:rnd:sat
```

### Behavior

```
Rd.h[1]=(sat32((Rs.h[1] * Rt.h[1])<<1) + 0x8000).h[1];
Rd.h[0]=(sat32((Rs.h[0] * Rt.h[0])<<1) + 0x8000).h[1];
```

**Class:** XTYPE (slots 2,3)

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

## Intrinsics

$Rd = \text{vmpyh}(Rs, Rt) : \ll 1 : \text{rnd} : \text{sat}$	Word32 Q6_R_vmpyh_RR_s1_rnd_sat(Word32 Rs, Word32 Rt)
$Rd = \text{vmpyh}(Rs, Rt) : \text{rnd} : \text{sat}$	Word32 Q6_R_vmpyh_RR_rnd_sat(Word32 Rs, Word32 Rt)

## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse		t5					MinOp			d5								
1	1	1	0	1	1	0	1	N	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	1	d	d	d	d	d	Rd=vmpyh(Rs,Rt):<<N>>:rnd:sat

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector multiply halfwords signed by unsigned

Multiply two 16-bit halfwords. Rs is considered signed, Ru unsigned.

Syntax	Behavior
$Rdd = \text{vmpyhsu}(Rs, Rt) [ : \ll 1 ] : \text{sat}$	$Rdd.w[0] = \text{sat}_{32} (Rs.h[0] * Rt.uh[0]) [ \ll 1 ] ;$ $Rdd.w[1] = \text{sat}_{32} (Rs.h[1] * Rt.uh[1]) [ \ll 1 ] ;$
$Rxx += \text{vmpyhsu}(Rs, Rt) [ : \ll 1 ] : \text{sat}$	$Rxx.w[0] = \text{sat}_{32} (Rxx.w[0] + (Rs.h[0] * Rt.uh[0]) [ \ll 1 ] ) ;$ $Rxx.w[1] = \text{sat}_{32} (Rxx.w[1] + (Rs.h[1] * Rt.uh[1]) [ \ll 1 ] ) ;$

### Class: XTYPE (slots 2,3)

#### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

#### Intrinsics

$Rdd = \text{vmpyhsu}(Rs, Rt) [ : \ll 1 ] : \text{sat}$	Word64 Q6_P_vmpyhsu_RR_s1_sat (Word32 Rs, Word32 Rt)
$Rdd = \text{vmpyhsu}(Rs, Rt) : \text{sat}$	Word64 Q6_P_vmpyhsu_RR_sat (Word32 Rs, Word32 Rt)
$Rxx += \text{vmpyhsu}(Rs, Rt) [ : \ll 1 ] : \text{sat}$	Word64 Q6_P_vmpyhsuacc_RR_s1_sat (Word64 Rxx, Word32 Rs, Word32 Rt)
$Rxx += \text{vmpyhsu}(Rs, Rt) : \text{sat}$	Word64 Q6_P_vmpyhsuacc_RR_sat (Word64 Rxx, Word32 Rs, Word32 Rt)

#### Encoding

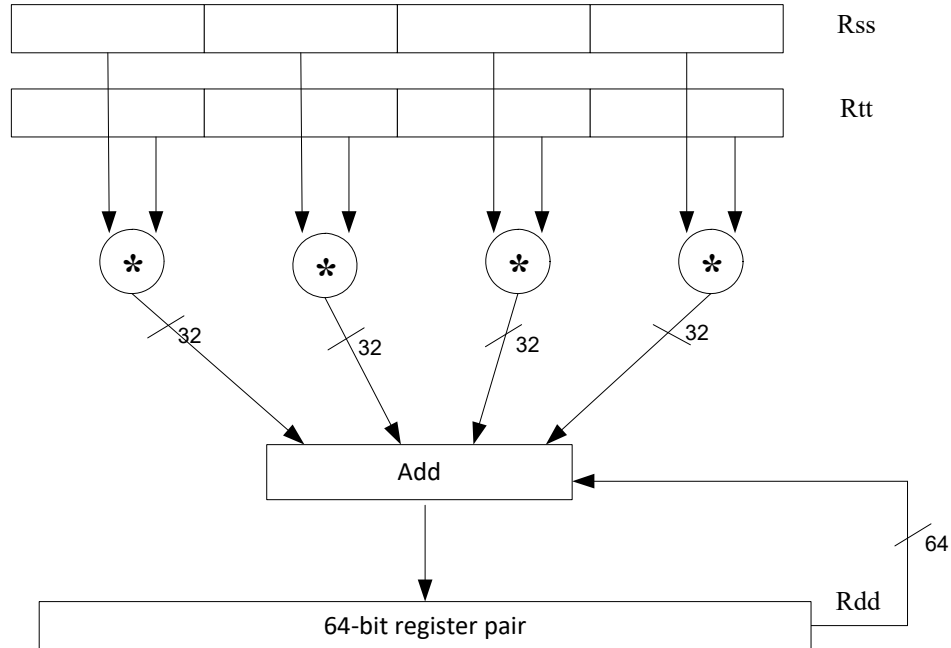
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				MajOp			s5					Parse		t5					MinOp			d5								
1	1	1	0	0	1	0	1	N	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	1	d	d	d	d	d	$Rdd = \text{vmpyhsu}(Rs, Rt) [ : \ll N ] : \text{sat}$
ICLASS		RegType				MajOp			s5					Parse		t5					MinOp			x5								
1	1	1	0	0	1	1	1	N	1	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	1	x	x	x	x	x	$Rxx += \text{vmpyhsu}(Rs, Rt) [ : \ll N ] : \text{sat}$

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d

<b>Field name</b>	<b>Description</b>
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Vector reduce multiply halfwords

Multiply each halfword of *Rss* by the corresponding halfword in *Rtt*. Add the intermediate products together and then optionally add the accumulator. Store the full 64-bit result in the destination register pair.



### Syntax

```
Rdd=vrmpyh (Rss, Rtt)
```

```
Rxx+=vrmpyh (Rss, Rtt)
```

### Behavior

```
Rdd = (Rss.h[0] * Rtt.h[0]) + (Rss.h[1] * Rtt.h[1]) +  
(Rss.h[2] * Rtt.h[2]) + (Rss.h[3] * Rtt.h[3]);
```

```
Rxx = Rxx + (Rss.h[0] * Rtt.h[0]) + (Rss.h[1] *  
Rtt.h[1]) + (Rss.h[2] * Rtt.h[2]) + (Rss.h[3] *  
Rtt.h[3]);
```

**Class: XTYPE (slots 2,3)**

### Intrinsics

```
Rdd=vrmpyh (Rss, Rtt)
```

```
Word64 Q6_P_vrmpyh_PP (Word64 Rss, Word64 Rtt)
```

```
Rxx+=vrmpyh (Rss, Rtt)
```

```
Word64 Q6_P_vrmpyhacc_PP (Word64 Rxx, Word64 Rss,  
Word64 Rtt)
```

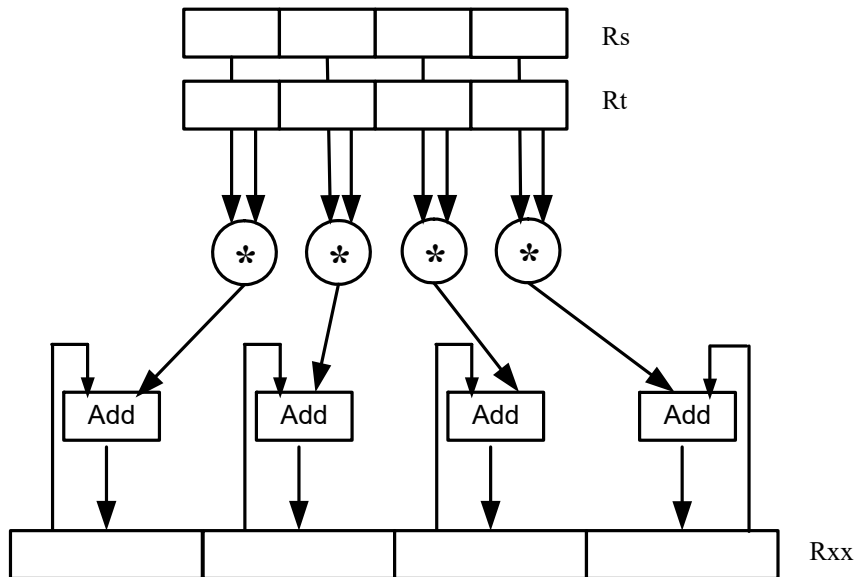
## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp				s5					Parse		t5					MinOp			d5					
1	1	1	0	1	0	0	0	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	0	d	d	d	d	d	Rdd=vrmphyh(Rss,Rtt)
ICLASS				RegType				MajOp				s5					Parse		t5					MinOp			x5					
1	1	1	0	1	0	1	0	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	0	x	x	x	x	x	Rxx+=vrmphyh(Rss,Rtt)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Vector multiply bytes

Multiply four 8-bit bytes from register *Rs* by four 8-bit bytes from *Rt*. Optionally accumulate the product with the 16-bit value from the destination register. Pack the 16-bit results in the destination register pair. The bytes of *Rs* are treated as either signed or unsigned.



### Syntax

### Behavior

<code>Rdd=vmpybsu (Rs, Rt)</code>	<pre>Rdd.h[0] = (Rs.b[0] * Rt.ub[0]); Rdd.h[1] = (Rs.b[1] * Rt.ub[1]); Rdd.h[2] = (Rs.b[2] * Rt.ub[2]); Rdd.h[3] = (Rs.b[3] * Rt.ub[3]);</pre>
<code>Rdd=vmpybu (Rs, Rt)</code>	<pre>Rdd.h[0] = (Rs.ub[0] * Rt.ub[0]); Rdd.h[1] = (Rs.ub[1] * Rt.ub[1]); Rdd.h[2] = (Rs.ub[2] * Rt.ub[2]); Rdd.h[3] = (Rs.ub[3] * Rt.ub[3]);</pre>
<code>Rxx+=vmpybsu (Rs, Rt)</code>	<pre>Rxx.h[0] = (Rxx.h[0] + (Rs.b[0] * Rt.ub[0])); Rxx.h[1] = (Rxx.h[1] + (Rs.b[1] * Rt.ub[1])); Rxx.h[2] = (Rxx.h[2] + (Rs.b[2] * Rt.ub[2])); Rxx.h[3] = (Rxx.h[3] + (Rs.b[3] * Rt.ub[3]));</pre>
<code>Rxx+=vmpybu (Rs, Rt)</code>	<pre>Rxx.h[0] = (Rxx.h[0] + (Rs.ub[0] * Rt.ub[0])); Rxx.h[1] = (Rxx.h[1] + (Rs.ub[1] * Rt.ub[1])); Rxx.h[2] = (Rxx.h[2] + (Rs.ub[2] * Rt.ub[2])); Rxx.h[3] = (Rxx.h[3] + (Rs.ub[3] * Rt.ub[3]));</pre>



**Class: XTYPE (slots 2,3)****Intrinsics**

Rdd=vmpybsu (Rs, Rt)	Word64 Q6_P_vmpybsu_RR(Word32 Rs, Word32 Rt)
Rdd=vmpybu (Rs, Rt)	Word64 Q6_P_vmpybu_RR(Word32 Rs, Word32 Rt)
Rxx+=vmpybsu (Rs, Rt)	Word64 Q6_P_vmpybsuacc_RR(Word64 Rxx, Word32 Rs, Word32 Rt)
Rxx+=vmpybu (Rs, Rt)	Word64 Q6_P_vmpybuacc_RR(Word64 Rxx, Word32 Rs, Word32 Rt)

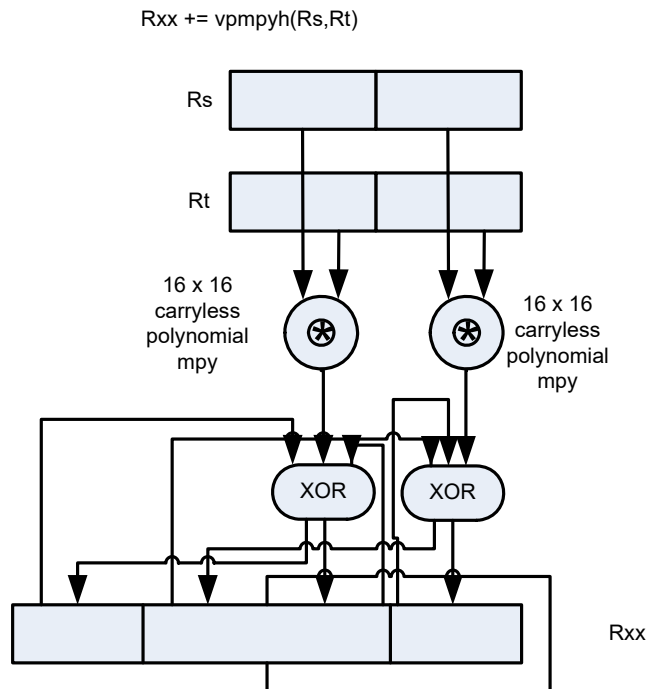
**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			d5							
1	1	1	0	0	1	0	1	0	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	d	d	d	d	d	Rdd=vmpybsu(Rs,Rt)
1	1	1	0	0	1	0	1	1	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	d	d	d	d	d	Rdd=vmpybu(Rs,Rt)
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			x5							
1	1	1	0	0	1	1	1	1	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	x	x	x	x	x	Rxx+=vmpybu(Rs,Rt)
1	1	1	0	0	1	1	1	1	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	x	x	x	x	x	Rxx+=vmpybsu(Rs,Rt)

Field name	Description
ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## Vector polynomial multiply halfwords

Perform a vector  $16 \times 16$  carryless polynomial multiply using 32-bit source registers Rs and Rt. Store the 64-bit result in packed H,H,L,L format in the destination register. The destination register can also be optionally accumulated (XORed). Finite field multiply instructions are useful for many algorithms including scramble code generation, cryptographic algorithms, convolutional, and Reed Solomon codes.



### Syntax

```
Rdd=vpmpyh (Rs, Rt)
```

### Behavior

```
x0 = Rs.uh[0];
x1 = Rs.uh[1];
y0 = Rt.uh[0];
y1 = Rt.uh[1];
prod0 = prod1 = 0;
for(i=0; i < 16; i++) {
    if((y0 >> i) & 1) prod0 ^= (x0 << i);
    if((y1 >> i) & 1) prod1 ^= (x1 << i);
}
Rdd.h[0]=prod0.uh[0];
Rdd.h[1]=prod1.uh[0];
Rdd.h[2]=prod0.uh[1];
Rdd.h[3]=prod1.uh[1];
```

**Syntax**

```
Rxx^=vpmpyh (Rs, Rt)
```

**Behavior**

```
x0 = Rs.uh[0];
x1 = Rs.uh[1];
y0 = Rt.uh[0];
y1 = Rt.uh[1];
prod0 = prod1 = 0;
for(i=0; i < 16; i++) {
    if((y0 >> i) & 1) prod0 ^= (x0 << i);
    if((y1 >> i) & 1) prod1 ^= (x1 << i);
}
Rxx.h[0]=Rxx.uh[0] ^ prod0.uh[0];
Rxx.h[1]=Rxx.uh[1] ^ prod1.uh[0];
Rxx.h[2]=Rxx.uh[2] ^ prod0.uh[1];
Rxx.h[3]=Rxx.uh[3] ^ prod1.uh[1];
```

**Class: XTYPE (slots 2,3)****Intrinsics**

```
Rdd=vpmpyh (Rs, Rt)
```

```
Word64 Q6_P_vpmpyh_RR(Word32 Rs, Word32 Rt)
```

```
Rxx^=vpmpyh (Rs, Rt)
```

```
Word64 Q6_P_vpmpyh_xacc_RR(Word64 Rxx, Word32 Rs,
Word32 Rt)
```

**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			d5							
1	1	1	0	0	1	0	1	1	1	0	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	1	d	d	d	d	d	Rdd=vpmpyh(Rs,Rt)
ICLASS			RegType				MajOp			s5					Parse		t5					MinOp			x5							
1	1	1	0	0	1	1	1	1	0	1	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	1	x	x	x	x	x	Rxx^=vpmpyh(Rs,Rt)

**Field name****Description**

ICLASS	Instruction class
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x

## 11.10.6 XTYPE PERM

The XTYPE PERM instruction subclass includes instructions that perform permutations.

### CABAC decode bin

This is a special-purpose instruction to support H.264 Context Adaptive Binary Arithmetic Coding (CABAC).

Syntax	Behavior
Rdd=decbin(Rss,Rtt)	<pre> state = Rtt.w[1][5:0]; valMPS = Rtt.w[1][8:8]; bitpos = Rtt.w[0][4:0]; range = Rss.w[0]; offset = Rss.w[1]; range &lt;&lt;= bitpos; offset &lt;&lt;= bitpos; rLPS = rLPS_table_64x4[state][ (range &gt;&gt;29)&amp;3]; rLPS = rLPS &lt;&lt; 23; rMPS= (range&amp;0xff800000) - rLPS; if (offset &lt; rMPS) {     Rdd = AC_next_state_MPS_64[state];     Rdd[8:8]=valMPS;     Rdd[31:23]=(rMPS&gt;&gt;23);     Rdd.w[1]=offset;     P0=valMPS; } else {     Rdd = AC_next_state_LPS_64[state];     Rdd[8:8]=(!state)?(1- valMPS):(valMPS));     Rdd[31:23]=(rLPS&gt;&gt;23);     Rdd.w[1]=(offset-rMPS);     P0=(valMPS^1); } </pre>

**Class: XTYPE (slots 2,3)**

### Notes

- The predicate generated by this instruction cannot be used as a .new predicate, nor can it be automatically ANDed with another predicate.

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				Maj		s5					Parse		t5				Min		d5										
1	1	0	0	0	0	0	1	1	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	-	d	d	d	d	d	Rdd=decbin(Rss,Rtt)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d

<b>Field name</b>	<b>Description</b>
s5	Field to encode register s
t5	Field to encode register t
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Saturate

Saturate a single scalar value.

The `sath` instruction saturates a signed 32-bit number to a signed 16-bit number, which is sign-extended back to 32 bits and placed in the destination register. The minimum negative value of the result is `0xffff8000` and the maximum positive value is `0x00007fff`.

The `satuh` instruction saturates a signed 32-bit number to an unsigned 16-bit number, which is zero-extended back to 32 bits and placed in the destination register. The minimum value of the result is 0 and the maximum value is `0x0000ffff`.

The `satb` instruction saturates a signed 32-bit number to a signed 8-bit number, which is sign-extended back to 32 bits and placed in the destination register. The minimum value of the result is `0xfffff80` and the maximum value is `0x0000007f`.

The `satub` instruction saturates a signed 32-bit number to an unsigned 8-bit number, which is zero-extended back to 32 bits and placed in the destination register. The minimum value of the result is 0 and the maximum value is `0x000000ff`.

Syntax	Behavior
<code>Rd=sat(Rss)</code>	<code>Rd = sat<sub>32</sub>(Rss);</code>
<code>Rd=satb(Rs)</code>	<code>Rd = sat<sub>8</sub>(Rs);</code>
<code>Rd=sath(Rs)</code>	<code>Rd = sat<sub>16</sub>(Rs);</code>
<code>Rd=satub(Rs)</code>	<code>Rd = usat<sub>8</sub>(Rs);</code>
<code>Rd=satuh(Rs)</code>	<code>Rd = usat<sub>16</sub>(Rs);</code>

### Class: XTYPE (slots 2,3)

#### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

#### Intrinsics

<code>Rd=sat(Rss)</code>	<code>Word32 Q6_R_sat_P(Word64 Rss)</code>
<code>Rd=satb(Rs)</code>	<code>Word32 Q6_R_satb_R(Word32 Rs)</code>
<code>Rd=sath(Rs)</code>	<code>Word32 Q6_R_sath_R(Word32 Rs)</code>
<code>Rd=satub(Rs)</code>	<code>Word32 Q6_R_satub_R(Word32 Rs)</code>
<code>Rd=satuh(Rs)</code>	<code>Word32 Q6_R_satuh_R(Word32 Rs)</code>

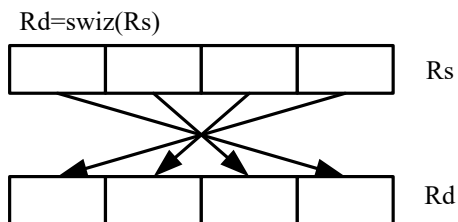
## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp		s5					Parse				MinOp			d5										
1	0	0	0	1	0	0	0	1	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	0	d	d	d	d	d	Rd=sat(Rss)
1	0	0	0	1	1	0	0	1	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	0	d	d	d	d	d	Rd=sath(Rs)
1	0	0	0	1	1	0	0	1	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	1	d	d	d	d	d	Rd=satuh(Rs)
1	0	0	0	1	1	0	0	1	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	1	0	d	d	d	d	d	Rd=satub(Rs)
1	0	0	0	1	1	0	0	1	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	1	1	d	d	d	d	d	Rd=satb(Rs)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Swizzle bytes

Swizzle the bytes of a word. This instruction is useful in converting between little and big endian formats.



### Syntax

```
Rd=swiz (Rs)
```

### Behavior

```
Rd.b[0]=Rs.b[3];
Rd.b[1]=Rs.b[2];
Rd.b[2]=Rs.b[1];
Rd.b[3]=Rs.b[0];
```

**Class: XTYPE (slots 2,3)**

### Intrinsics

```
Rd=swiz (Rs)
```

```
Word32 Q6_R_swiz_R(Word32 Rs)
```

### Encoding

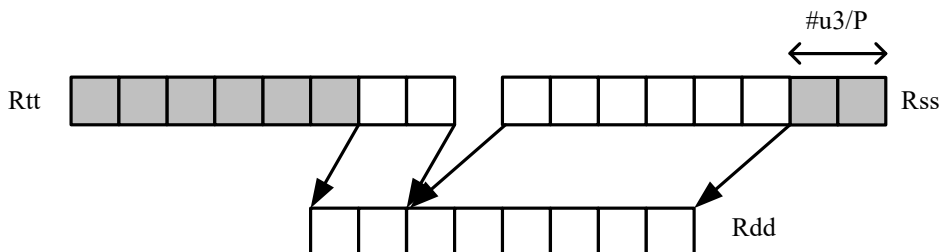
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp				s5					Parse		MinOp			d5										
1	0	0	0	1	1	0	0	1	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	1	1	d	d	d	d	d	Rd=swiz(Rs)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type



## Vector align

Align a vector. Use the immediate amount, or the least significant three bits of a predicate register as the number of bytes to align. Shift the Rss register pair right by this number of bytes. Fill the vacated positions with the least significant elements from Rtt.



### Syntax

`Rdd=valignb(Rtt,Rss,#u3)`

`Rdd=valignb(Rtt,Rss,Pu)`

### Behavior

`Rdd = (Rss >>> #u*8) | (Rtt << ((8-#u)*8));`

`PREDUSE_TIMING;`  
`Rdd = Rss >>> (Pu&0x7)*8 | (Rtt << (8-(Pu&0x7))*8);`

## Class: XTYPE (slots 2,3)

### Intrinsics

`Rdd=valignb(Rtt,Rss,#u3)`

Word64 Q6\_P\_valignb\_PPI(Word64 Rtt, Word64 Rss, Word32 Iu3)

`Rdd=valignb(Rtt,Rss,Pu)`

Word64 Q6\_P\_valignb\_PPp(Word64 Rtt, Word64 Rss, Byte Pu)

### Encoding

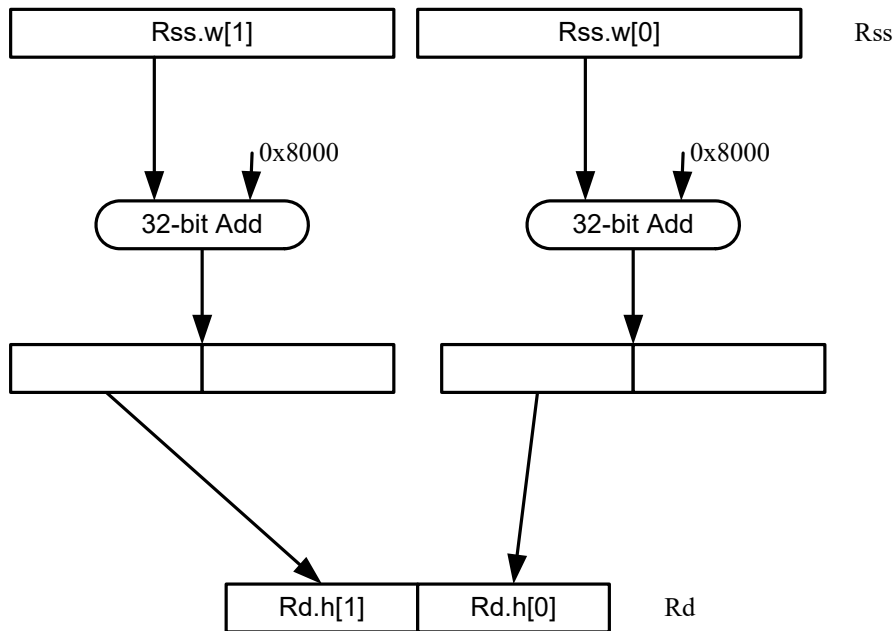
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				Maj		s5					Parse		t5					Min		d5										
1	1	0	0	0	0	0	0	0	-	-	s	s	s	s	s	P	P	-	t	t	t	t	t	i	i	i	d	d	d	d	d	Rdd=valignb(Rtt,Rss,#u3)
ICLASS		RegType				Maj		s5					Parse		t5					u2		d5										
1	1	0	0	0	0	1	0	0	-	-	s	s	s	s	s	P	P	-	t	t	t	t	t	-	u	u	d	d	d	d	d	Rdd=valignb(Rtt,Rss,Pu)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
u2	Field to encode register u

<b>Field name</b>	<b>Description</b>
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Vector round and pack

Add the constant 0x00008000 to each word in the 64-bit source vector Rss. Optionally saturate this addition to 32 bits. Pack the high halfwords of the result into the corresponding halfword of the 32-bit destination register.



Syntax	Behavior
<code>Rd=vrndwh(Rss)</code>	<pre>for (i=0;i&lt;2;i++) {   Rd.h[i]=(Rss.w[i]+0x08000).h[1]; }</pre>
<code>Rd=vrndwh(Rss):sat</code>	<pre>for (i=0;i&lt;2;i++) {   Rd.h[i]=sat<sub>32</sub>(Rss.w[i]+0x08000).h[1]; }</pre>

**Class: XTYPE (slots 2,3)**

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

### Intrinsics

<code>Rd=vrndwh(Rss)</code>	<code>Word32 Q6_R_vrndwh_P(Word64 Rss)</code>
<code>Rd=vrndwh(Rss):sat</code>	<code>Word32 Q6_R_vrndwh_P_sat(Word64 Rss)</code>

## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp			s5					Parse				MinOp			d5									
1	0	0	0	1	0	0	0	1	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	0	d	d	d	d	d	Rd=vrndwh(Rss)
1	0	0	0	1	0	0	0	1	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	1	0	d	d	d	d	d	Rd=vrndwh(Rss):sat

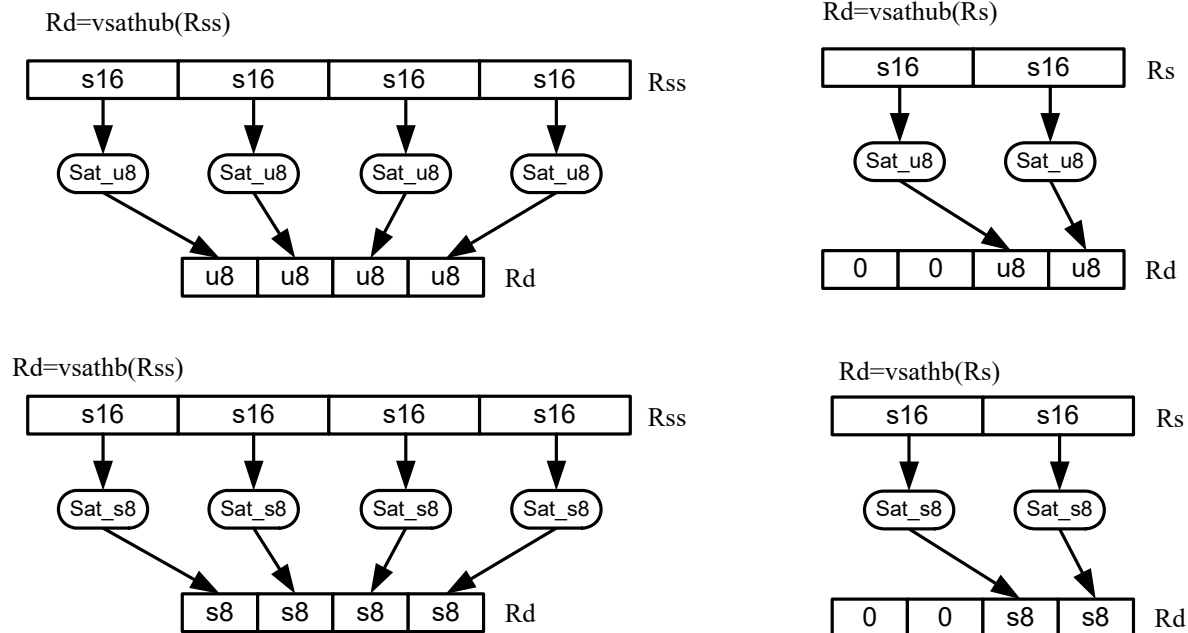
Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Vector saturate and pack

For each element in the vector, saturate the value to the next smaller size.

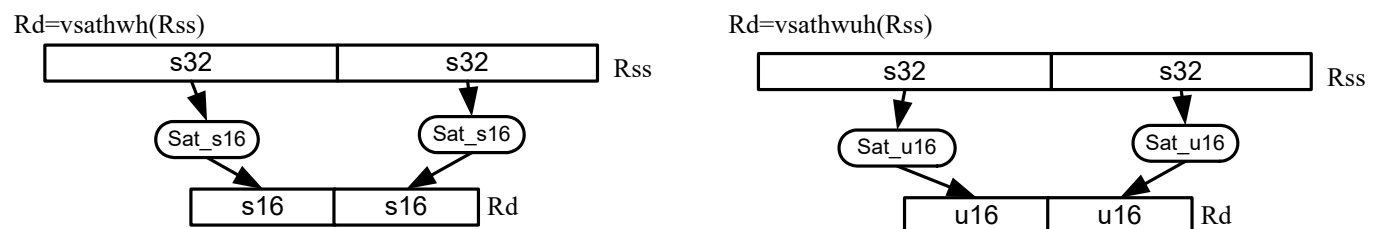
The vsathub instruction saturates signed halfwords to unsigned bytes,

The vsathb instruction saturates signed halfwords to signed bytes.



The vsatwh instruction saturates signed words to signed halfwords.

The vsathwuh instruction saturates signed words to unsigned halfwords. The resulting values are packed together into the destination register Rd.



Syntax	Behavior
Rd=vsathb (Rs)	<pre>Rd.b[0]=sat<sub>8</sub>(Rs.h[0]); Rd.b[1]=sat<sub>8</sub>(Rs.h[1]); Rd.b[2]=0; Rd.b[3]=0;</pre>
Rd=vsathb (Rss)	<pre>for (i=0;i&lt;4;i++) {     Rd.b[i]=sat<sub>8</sub>(Rss.h[i]); }</pre>
Rd=vsathub (Rs)	<pre>Rd.b[0]=usat<sub>8</sub>(Rs.h[0]); Rd.b[1]=usat<sub>8</sub>(Rs.h[1]); Rd.b[2]=0; Rd.b[3]=0;</pre>
Rd=vsathub (Rss)	<pre>for (i=0;i&lt;4;i++) {     Rd.b[i]=usat<sub>8</sub>(Rss.h[i]); }</pre>
Rd=vsatwh (Rss)	<pre>for (i=0;i&lt;2;i++) {     Rd.h[i]=sat<sub>16</sub>(Rss.w[i]); }</pre>
Rd=vsatwuh (Rss)	<pre>for (i=0;i&lt;2;i++) {     Rd.h[i]=usat<sub>16</sub>(Rss.w[i]); }</pre>

**Class: XTYPE (slots 2,3)****Notes**

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

**Intrinsics**

Rd=vsathb (Rs)	Word32 Q6_R_vsathb_R(Word32 Rs)
Rd=vsathb (Rss)	Word32 Q6_R_vsathb_P(Word64 Rss)
Rd=vsathub (Rs)	Word32 Q6_R_vsathub_R(Word32 Rs)
Rd=vsathub (Rss)	Word32 Q6_R_vsathub_P(Word64 Rss)
Rd=vsatwh (Rss)	Word32 Q6_R_vsatwh_P(Word64 Rss)
Rd=vsatwuh (Rss)	Word32 Q6_R_vsatwuh_P(Word64 Rss)

## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp				s5					Parse				MinOp				d5							
1	0	0	0	1	0	0	0	0	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	0	d	d	d	d	d	Rd=vsathub(Rss)
1	0	0	0	1	0	0	0	0	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	0	1	0	d	d	d	d	d	Rd=vsatwh(Rss)
1	0	0	0	1	0	0	0	0	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	0	d	d	d	d	d	Rd=vsatwuh(Rss)
1	0	0	0	1	0	0	0	0	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	1	0	d	d	d	d	d	Rd=vsathb(Rss)
1	0	0	0	1	1	0	0	1	0	-	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	-	d	d	d	d	d	Rd=vsathb(Rs)
1	0	0	0	1	1	0	0	1	0	-	s	s	s	s	s	P	P	-	-	-	-	-	-	0	1	-	d	d	d	d	d	Rd=vsathub(Rs)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Vector saturate without pack

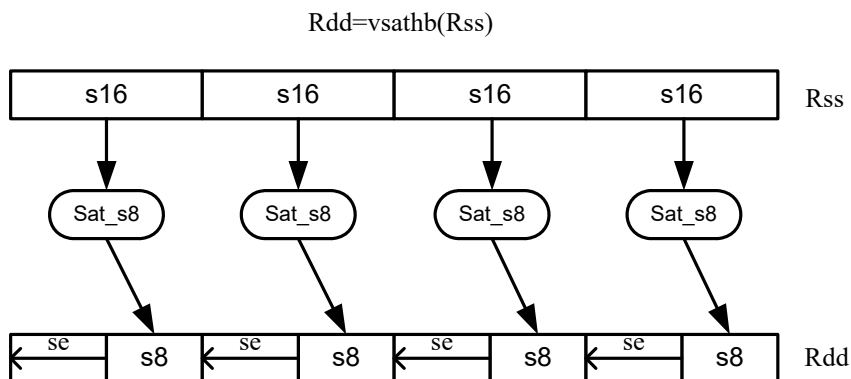
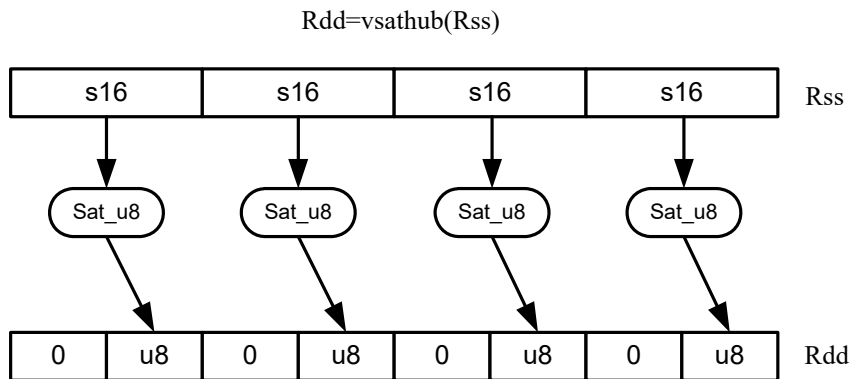
Saturate each element of source vector Rss to the next smaller size.

The vsathub instruction saturates signed halfwords to unsigned bytes.

The vsatwh instruction saturates signed words to signed halfwords.

The vsatwuh instruction saturates signed words to unsigned halfwords.

The resulting values are placed in destination register Rdd in unpacked form.





Syntax	Behavior
Rdd=vsathb(Rss)	for (i=0;i<4;i++) { Rdd.h[i]=sat <sub>8</sub> (Rss.h[i]); }
Rdd=vsathub(Rss)	for (i=0;i<4;i++) { Rdd.h[i]=usat <sub>8</sub> (Rss.h[i]); }
Rdd=vsatwh(Rss)	for (i=0;i<2;i++) { Rdd.w[i]=sat <sub>16</sub> (Rss.w[i]); }
Rdd=vsatwuh(Rss)	for (i=0;i<2;i++) { Rdd.w[i]=usat <sub>16</sub> (Rss.w[i]); }

### Class: XTYPE (slots 2,3)

#### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

#### Intrinsics

Rdd=vsathb(Rss)	Word64 Q6_P_vsathb_P(Word64 Rss)
Rdd=vsathub(Rss)	Word64 Q6_P_vsathub_P(Word64 Rss)
Rdd=vsatwh(Rss)	Word64 Q6_P_vsatwh_P(Word64 Rss)
Rdd=vsatwuh(Rss)	Word64 Q6_P_vsatwuh_P(Word64 Rss)

#### Encoding

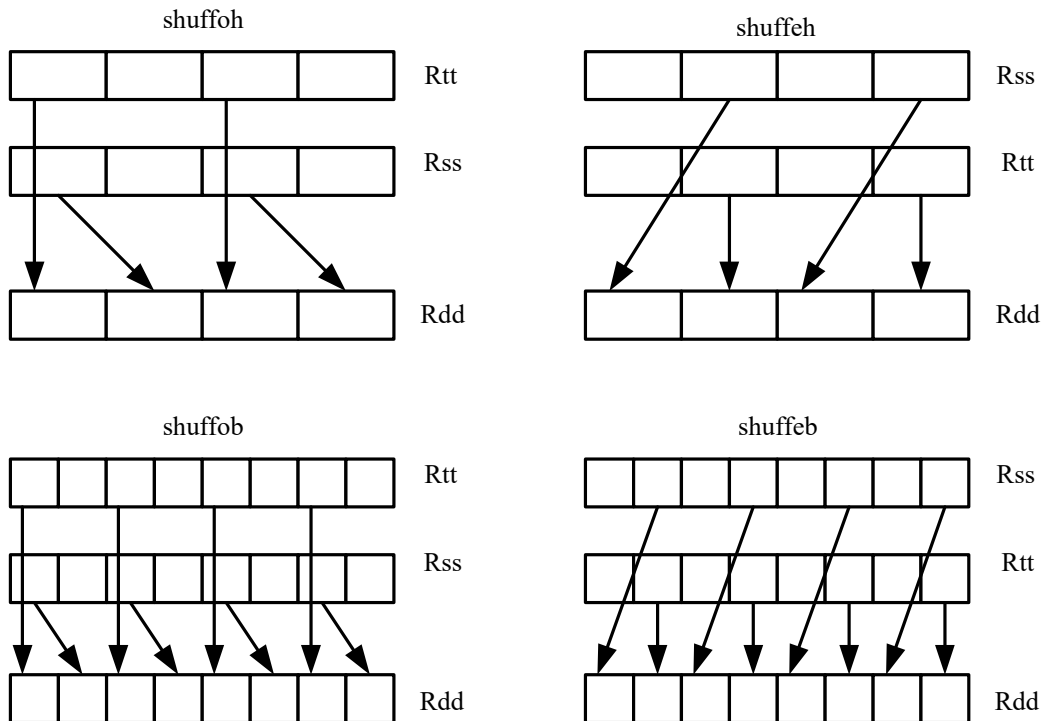
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				MajOp				s5					Parse		MinOp			d5												
1	0	0	0	0	0	0	0	0	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	0	d	d	d	d	d	Rdd=vsathub(Rss)
1	0	0	0	0	0	0	0	0	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	1	d	d	d	d	d	Rdd=vsatwuh(Rss)
1	0	0	0	0	0	0	0	0	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	1	0	d	d	d	d	d	Rdd=vsatwh(Rss)
1	0	0	0	0	0	0	0	0	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	1	1	d	d	d	d	d	Rdd=vsathb(Rss)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s

<b>Field name</b>	<b>Description</b>
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Vector shuffle

Shuffle odd halfwords (shuffoh) takes the odd halfwords from Rtt and the odd halfwords from Rss and merges them together into vector Rdd. Shuffle even halfwords (shuffeh) performs the same operation on every even halfword in Rss and Rtt. The same operation is available for odd and even bytes.



### Syntax

Rdd=shuffeb(Rss,Rtt)

Rdd=shuffeh(Rss,Rtt)

Rdd=shuffob(Rtt,Rss)

Rdd=shuffoh(Rtt,Rss)

### Behavior

```
for (i=0;i<4;i++) {
  Rdd.b[i*2]=Rtt.b[i*2];
  Rdd.b[i*2+1]=Rss.b[i*2];
}
```

```
for (i=0;i<2;i++) {
  Rdd.h[i*2]=Rtt.h[i*2];
  Rdd.h[i*2+1]=Rss.h[i*2];
}
```

```
for (i=0;i<4;i++) {
  Rdd.b[i*2]=Rss.b[i*2+1];
  Rdd.b[i*2+1]=Rtt.b[i*2+1];
}
```

```
for (i=0;i<2;i++) {
  Rdd.h[i*2]=Rss.h[i*2+1];
  Rdd.h[i*2+1]=Rtt.h[i*2+1];
}
```

**Class: XTYPE (slots 2,3)****Intrinsics**

Rdd=shuffeb(Rss,Rtt)	Word64 Q6_P_shuffeb_PP(Word64 Rss, Word64 Rtt)
Rdd=shuffeh(Rss,Rtt)	Word64 Q6_P_shuffeh_PP(Word64 Rss, Word64 Rtt)
Rdd=shuffob(Rtt,Rss)	Word64 Q6_P_shuffob_PP(Word64 Rtt, Word64 Rss)
Rdd=shuffoh(Rtt,Rss)	Word64 Q6_P_shuffoh_PP(Word64 Rtt, Word64 Rss)

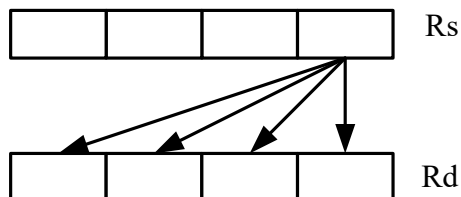
**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				Maj		s5					Parse		t5					Min		d5									
1	1	0	0	0	0	0	1	0	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	-	d	d	d	d	d	Rdd=shuffeb(Rss,Rtt)
1	1	0	0	0	0	0	1	0	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	-	d	d	d	d	d	Rdd=shuffob(Rtt,Rss)
1	1	0	0	0	0	0	1	0	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	-	d	d	d	d	d	Rdd=shuffeh(Rss,Rtt)
1	1	0	0	0	0	0	1	1	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	d	d	d	d	d	Rdd=shuffoh(Rtt,Rss)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Vector splat bytes

Replicate the low 8-bits from register Rs into each of the four bytes of destination register Rd.

$$Rd = vsplatb(Rs)$$


Syntax	Behavior
$Rd = vsplatb(Rs)$	<pre>for (i=0;i&lt;4;i++) {   Rd.b[i]=Rs.b[0]; }</pre>
$Rdd = vsplatb(Rs)$	<pre>for (i=0;i&lt;8;i++) {   Rdd.b[i]=Rs.b[0]; }</pre>

**Class: XTYPE (slots 2,3)**

### Intrinsics

$Rd = vsplatb(Rs)$       Word32 Q6\_R\_vsplatb\_R(Word32 Rs)

$Rdd = vsplatb(Rs)$       Word64 Q6\_P\_vsplatb\_R(Word32 Rs)

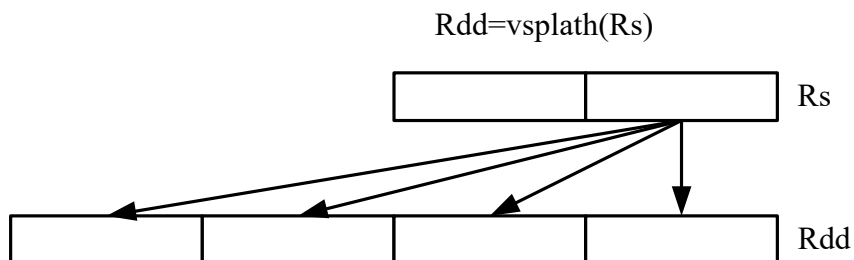
### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp				s5					Parse		MinOp			d5										
1	0	0	0	0	1	0	0	0	1	-	s	s	s	s	s	P	P	-	-	-	-	-	-	1	0	-	d	d	d	d	d	Rdd=vsplatb(Rs)
1	0	0	0	1	1	0	0	0	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	1	1	1	d	d	d	d	d	Rd=vsplatb(Rs)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Vector splat halfwords

Replicate the low 16-bits from register Rs into each of the four halfwords of destination Rdd.



### Syntax

```
Rdd=vsplath(Rs)
```

### Behavior

```
for (i=0;i<4;i++) {
    Rdd.h[i]=Rs.h[0];
}
```

**Class: XTYPE (slots 2,3)**

### Intrinsics

```
Rdd=vsplath(Rs)
```

```
Word64 Q6_P_vsplath_R(Word32 Rs)
```

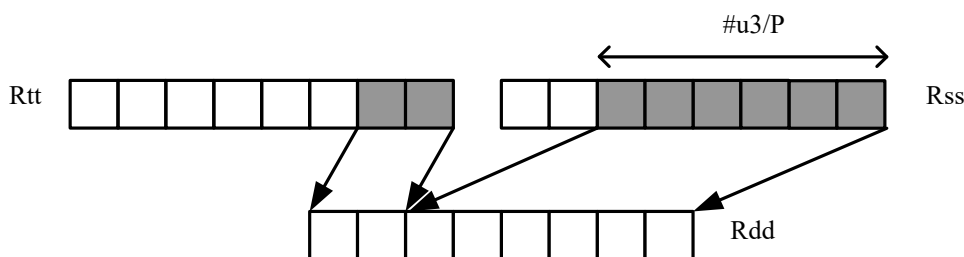
### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp				s5					Parse		MinOp				d5									
1	0	0	0	0	1	0	0	0	1	-	s	s	s	s	s	P	P	-	-	-	-	-	-	0	1	-	d	d	d	d	d	Rdd=vsplath(Rs)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Vector splice

Concatenate the low (8-N) bytes of vector Rtt with the low N bytes of vector Rss. This instruction is helpful to vectorize unaligned stores.



### Syntax

```
Rdd=vspliceb(Rss,Rtt,#u3)
```

```
Rdd=vspliceb(Rss,Rtt,Pu)
```

### Behavior

```
Rdd = Rtt << #u*8 | zxt_{#u*8->64}(Rss);
```

```
PREDUSE_TIMING;
```

```
Rdd = Rtt << (Pu&7)*8 | zxt_{(Pu&7)*8->64}(Rss);
```

## Class: XTYPE (slots 2,3)

### Intrinsics

```
Rdd=vspliceb(Rss,Rtt,#u3)
```

```
Word64 Q6_P_vspliceb_PPI(Word64 Rss, Word64  
Rtt, Word32 Iu3)
```

```
Rdd=vspliceb(Rss,Rtt,Pu)
```

```
Word64 Q6_P_vspliceb_PPp(Word64 Rss, Word64  
Rtt, Byte Pu)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				Maj		s5					Parse		t5					Min		d5										
1	1	0	0	0	0	0	0	1	-	-	s	s	s	s	s	P	P	-	t	t	t	t	t	i	i	i	d	d	d	d	d	Rdd=vspliceb(Rss,Rtt,#u3)
ICLASS		RegType				Maj		s5					Parse		t5					u2		d5										
1	1	0	0	0	0	1	0	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	u	u	d	d	d	d	d	Rdd=vspliceb(Rss,Rtt,Pu)

#### Field name

#### Description

ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
u2	Field to encode register u

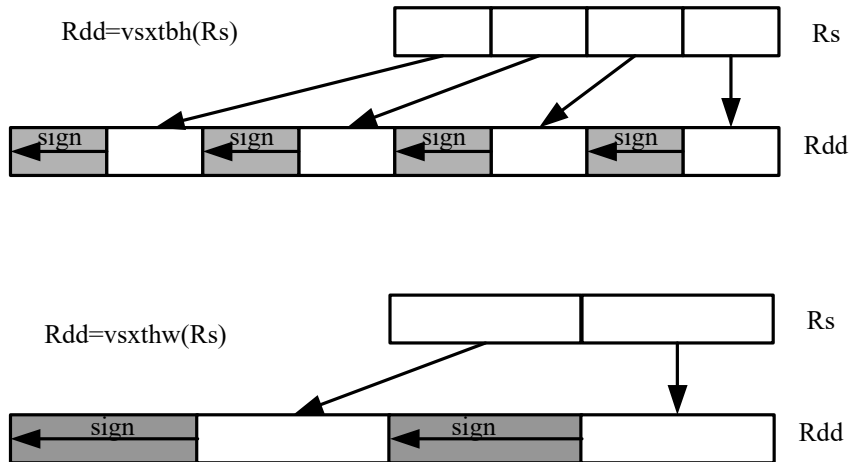
<b>Field name</b>	<b>Description</b>
Maj	Major opcode
Min	Minor opcode
RegType	Register type



## Vector sign extend

The `vsxtbh` instruction sign-extends each byte of a single register source to halfwords, and places the result in the destination register pair.

The `vsxthw` instruction sign-extends each halfword of a single register source to words, and places the result in the destination register pair.



### Syntax

`Rdd=vsxtbh (Rs)`

`Rdd=vsxthw (Rs)`

### Behavior

```
for (i=0;i<4;i++) {
    Rdd.h[i]=Rs.b[i];
}
```

```
for (i=0;i<2;i++) {
    Rdd.w[i]=Rs.h[i];
}
```

## Class: XTYPE (slots 2,3)

### Intrinsics

`Rdd=vsxtbh (Rs)`

`Word64 Q6_P_vsxtbh_R(Word32 Rs)`

`Rdd=vsxthw (Rs)`

`Word64 Q6_P_vsxthw_R(Word32 Rs)`

## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp				s5					Parse				MinOp			d5								
1	0	0	0	0	1	0	0	0	0	-	s	s	s	s	s	P	P	-	-	-	-	-	0	0	-	d	d	d	d	d	Rdd=vsxtbh(Rs)	
1	0	0	0	0	1	0	0	0	0	-	s	s	s	s	s	P	P	-	-	-	-	-	1	0	-	d	d	d	d	d	Rdd=vsxthw(Rs)	

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d

<b>Field name</b>	<b>Description</b>
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

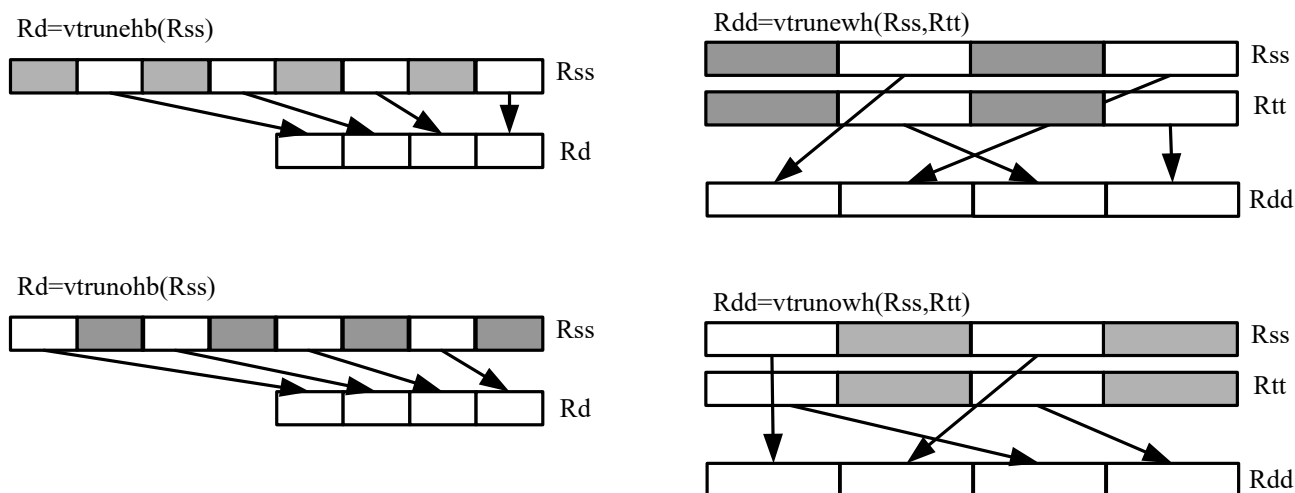
## Vector truncate

In the `vtrunehb` instruction, for each halfword in a vector, take the even (lower) byte and ignore the other byte. Pack the resulting values into destination register `Rd`.

The `vtrunohb` instruction takes each odd byte of the source vector.

The `vtrunewh` instruction uses two source register pairs, `Rss` and `Rtt`. Pack the even (lower) halfwords of `Rss` in the upper word of `Rdd`, and pack the lower halfwords of `Rtt` in the lower word of `Rdd`.

The `vtrunowh` instruction performs the same operation as the `vtrunewh` instruction, but uses the odd (upper) halfwords of the source vectors instead.



Syntax	Behavior
<code>Rd=vtrunehb(Rss)</code>	<pre>for (i=0;i&lt;4;i++) {   Rd.b[i]=Rss.b[i*2]; }</pre>
<code>Rd=vtrunohb(Rss)</code>	<pre>for (i=0;i&lt;4;i++) {   Rd.b[i]=Rss.b[i*2+1]; }</pre>
<code>Rdd=vtrunehb(Rss,Rtt)</code>	<pre>for (i=0;i&lt;4;i++) {   Rdd.b[i]=Rtt.b[i*2];   Rdd.b[i+4]=Rss.b[i*2]; }</pre>
<code>Rdd=vtrunewh(Rss,Rtt)</code>	<pre>Rdd.h[0]=Rtt.h[0]; Rdd.h[1]=Rtt.h[2]; Rdd.h[2]=Rss.h[0]; Rdd.h[3]=Rss.h[2];</pre>
<code>Rdd=vtrunowh(Rss,Rtt)</code>	<pre>for (i=0;i&lt;4;i++) {   Rdd.b[i]=Rtt.b[i*2+1];   Rdd.b[i+4]=Rss.b[i*2+1]; }</pre>

**Syntax**

```
Rdd=vtrunowh(Rss,Rtt)
```

**Behavior**

```
Rdd.h[0]=Rtt.h[1];
Rdd.h[1]=Rtt.h[3];
Rdd.h[2]=Rss.h[1];
Rdd.h[3]=Rss.h[3];
```

**Class: XTYPE (slots 2,3)****Intrinsics**

Rd=vtrunehb(Rss)	Word32 Q6_R_vtrunehb_P(Word64 Rss)
Rd=vtrunohb(Rss)	Word32 Q6_R_vtrunohb_P(Word64 Rss)
Rdd=vtrunehb(Rss,Rtt)	Word64 Q6_P_vtrunehb_PP(Word64 Rss, Word64 Rtt)
Rdd=vtrunewh(Rss,Rtt)	Word64 Q6_P_vtrunewh_PP(Word64 Rss, Word64 Rtt)
Rdd=vtrunohb(Rss,Rtt)	Word64 Q6_P_vtrunohb_PP(Word64 Rss, Word64 Rtt)
Rdd=vtrunowh(Rss,Rtt)	Word64 Q6_P_vtrunowh_PP(Word64 Rss, Word64 Rtt)

**Encoding**

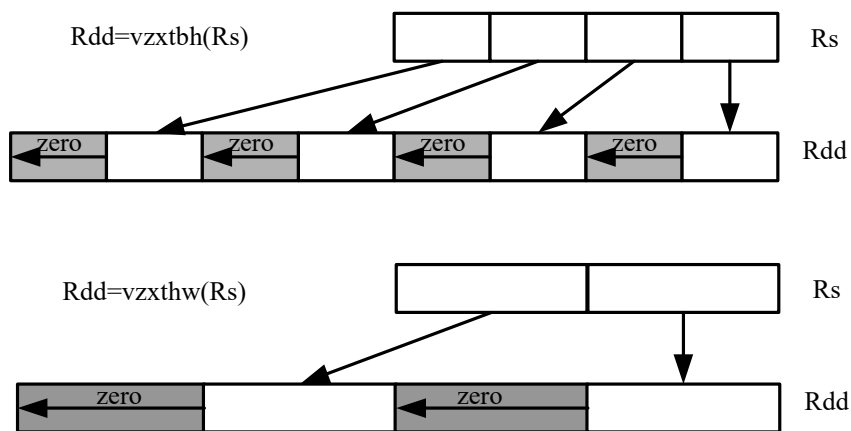
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp			s5					Parse			MinOp			d5											
1	0	0	0	1	0	0	0	1	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	0	0	0	d	d	d	d	d	Rd=vtrunohb(Rss)
1	0	0	0	1	0	0	0	1	0	0	s	s	s	s	s	P	P	-	-	-	-	-	-	0	1	0	d	d	d	d	d	Rd=vtrunehb(Rss)
ICLASS			RegType				Maj			s5					Parse			t5			Min			d5								
1	1	0	0	0	0	0	1	1	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	d	d	d	d	d	Rdd=vtrunewh(Rss,Rtt)
1	1	0	0	0	0	0	1	1	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	d	d	d	d	d	Rdd=vtrunehb(Rss,Rtt)
1	1	0	0	0	0	0	1	1	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	0	d	d	d	d	d	Rdd=vtrunowh(Rss,Rtt)
1	1	0	0	0	0	0	1	1	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	1	d	d	d	d	d	Rdd=vtrunohb(Rss,Rtt)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
MajOp	Major opcode
MinOp	Minor opcode
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Vector zero extend

The `vzxtbh` instruction zero-extends each byte of a single register source to halfwords, and places the result in the destination register pair.

The `vzxtbw` instruction zero-extends each halfword of a single register source to words, and places the result in the destination register pair.



Syntax	Behavior
<code>Rdd=vzxtbh (Rs)</code>	<pre>for (i=0;i&lt;4;i++) {   Rdd.h[i]=Rs.ub[i]; }</pre>
<code>Rdd=vzxtbw (Rs)</code>	<pre>for (i=0;i&lt;2;i++) {   Rdd.w[i]=Rs.uh[i]; }</pre>

### Class: XTYPE (slots 2,3)

#### Intrinsics

`Rdd=vzxtbh (Rs)`      `Word64 Q6_P_vzxtbh_R(Word32 Rs)`

`Rdd=vzxtbw (Rs)`      `Word64 Q6_P_vzxtbw_R(Word32 Rs)`

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				MajOp				s5					Parse				MinOp			d5										
1	0	0	0	0	1	0	0	0	0	-	s	s	s	s	s	P	P	-	-	-	-	-	0	1	-	d	d	d	d	d	Rdd=vzxtbh(Rs)	
1	0	0	0	0	1	0	0	0	0	-	s	s	s	s	s	P	P	-	-	-	-	-	1	1	-	d	d	d	d	d	Rdd=vzxtbw(Rs)	

<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## 11.10.7 XTYPE PRED

The XTYPE PRED instruction subclass includes instructions that perform miscellaneous operations on predicates, including mask generation, predicate transfers, and the Viterbi pack operation.

### Bounds check

Determine if Rs falls in the range defined by Rtt.

The user sets Rtt.w0 to the lower bound, and Rtt.w1 to the upper bound.

All bits of the destination predicate are set if the value falls within the range, or all cleared otherwise.

Syntax	Behavior
Pd=boundscheck(Rs,Rtt)	<pre>if ("Rs &amp; 1") {     Assembler mapped to:     "Pd=boundscheck(Rss,Rtt):raw:hi"; } else {     Assembler mapped to:     "Pd=boundscheck(Rss,Rtt):raw:lo"; }</pre>
Pd=boundscheck(Rss,Rtt):raw:hi	<pre>src = Rss.uw[1]; Pd = (src.uw[0] &gt;= Rtt.uw[0]) &amp;&amp; (src.uw[0] &lt; Rtt.uw[1]) ? 0xff : 0x00;</pre>
Pd=boundscheck(Rss,Rtt):raw:lo	<pre>src = Rss.uw[0]; Pd = (src.uw[0] &gt;= Rtt.uw[0]) &amp;&amp; (src.uw[0] &lt; Rtt.uw[1]) ? 0xff : 0x00;</pre>

### Class: XTYPE (slots 2,3)

#### Intrinsics

Pd=boundscheck(Rs,Rtt) Byte Q6\_p\_boundscheck\_RP(Word32 Rs, Word64 Rtt)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5					MinOp			d2										
1	1	0	1	0	0	1	0	0	-	-	s	s	s	s	s	P	P	1	t	t	t	t	t	1	0	0	-	-	-	d	d	Pd=boundscheck(Rss,Rtt):raw:lo
1	1	0	1	0	0	1	0	0	-	-	s	s	s	s	s	P	P	1	t	t	t	t	t	1	0	1	-	-	-	d	d	Pd=boundscheck(Rss,Rtt):raw:hi

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits

<b>Field name</b>	<b>Description</b>
d2	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t



## Compare byte

These instructions sign- or zero-extend the low 8 bits of the source registers and perform 32-bit comparisons on the result. When there is an extended 32-bit immediate operand, the full 32 immediate bits are used for the comparison.

Syntax	Behavior
Pd=cmpb.eq(Rs,#u8)	Pd=Rs.ub[0] == #u ? 0xff : 0x00;
Pd=cmpb.eq(Rs,Rt)	Pd=Rs.b[0] == Rt.b[0] ? 0xff : 0x00;
Pd=cmpb.gt(Rs,#s8)	Pd=Rs.b[0] > #s ? 0xff : 0x00;
Pd=cmpb.gt(Rs,Rt)	Pd=Rs.b[0] > Rt.b[0] ? 0xff : 0x00;
Pd=cmpb.gtu(Rs,#u7)	apply_extension(#u); Pd=Rs.ub[0] > #u.uw[0] ? 0xff : 0x00;
Pd=cmpb.gtu(Rs,Rt)	Pd=Rs.ub[0] > Rt.ub[0] ? 0xff : 0x00;

### Class: XTYPE (slots 2,3)

#### Intrinsics

Pd=cmpb.eq(Rs,#u8)	Byte Q6_p_cmpb_eq_RI(Word32 Rs, Word32 Iu8)
Pd=cmpb.eq(Rs,Rt)	Byte Q6_p_cmpb_eq_RR(Word32 Rs, Word32 Rt)
Pd=cmpb.gt(Rs,#s8)	Byte Q6_p_cmpb_gt_RI(Word32 Rs, Word32 Is8)
Pd=cmpb.gt(Rs,Rt)	Byte Q6_p_cmpb_gt_RR(Word32 Rs, Word32 Rt)
Pd=cmpb.gtu(Rs,#u7)	Byte Q6_p_cmpb_gtu_RI(Word32 Rs, Word32 Iu7)
Pd=cmpb.gtu(Rs,Rt)	Byte Q6_p_cmpb_gtu_RR(Word32 Rs, Word32 Rt)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				Maj		s5					Parse		t5					Min		d2									
1	1	0	0	0	1	1	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	-	-	-	d	d	Pd=cmpb.gt(Rs,Rt)
1	1	0	0	0	1	1	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	0	-	-	-	d	d	Pd=cmpb.eq(Rs,Rt)
1	1	0	0	0	1	1	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	1	-	-	-	d	d	Pd=cmpb.gtu(Rs,Rt)
ICLASS			RegType				s5					Parse															d2					
1	1	0	1	1	1	0	1	-	0	0	s	s	s	s	s	P	P	-	i	i	i	i	i	i	i	i	0	0	-	d	d	Pd=cmpb.eq(Rs,#u8)
1	1	0	1	1	1	0	1	-	0	1	s	s	s	s	s	P	P	-	i	i	i	i	i	i	i	i	0	0	-	d	d	Pd=cmpb.gt(Rs,#s8)
1	1	0	1	1	1	0	1	-	1	0	s	s	s	s	s	P	P	-	0	i	i	i	i	i	i	i	0	0	-	d	d	Pd=cmpb.gtu(Rs,#u7)

Field name	Description
RegType	Register type
MajOp	Major opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d2	Field to encode register d
s5	Field to encode register s

<b>Field name</b>	<b>Description</b>
t5	Field to encode register t
Maj	Major opcode
Min	Minor opcode

## Compare half

These instructions sign- or zero-extend the low 16 bits of the source registers and perform 32-bit comparisons on the result. When there is an extended 32-bit immediate operand, the full 32 immediate bits are used for the comparison.

Syntax	Behavior
Pd=cmph.eq(Rs,#s8)	apply_extension(#s); Pd=Rs.h[0] == #s ? 0xff : 0x00;
Pd=cmph.eq(Rs,Rt)	Pd=Rs.h[0] == Rt.h[0] ? 0xff : 0x00;
Pd=cmph.gt(Rs,#s8)	apply_extension(#s); Pd=Rs.h[0] > #s ? 0xff : 0x00;
Pd=cmph.gt(Rs,Rt)	Pd=Rs.h[0] > Rt.h[0] ? 0xff : 0x00;
Pd=cmph.gtu(Rs,#u7)	apply_extension(#u); Pd=Rs.uh[0] > #u.uw[0] ? 0xff : 0x00;
Pd=cmph.gtu(Rs,Rt)	Pd=Rs.uh[0] > Rt.uh[0] ? 0xff : 0x00;

### Class: XTYPE (slots 2,3)

#### Intrinsics

Pd=cmph.eq(Rs,#s8)	Byte Q6_p_cmph_eq_RI(Word32 Rs, Word32 Is8)
Pd=cmph.eq(Rs,Rt)	Byte Q6_p_cmph_eq_RR(Word32 Rs, Word32 Rt)
Pd=cmph.gt(Rs,#s8)	Byte Q6_p_cmph_gt_RI(Word32 Rs, Word32 Is8)
Pd=cmph.gt(Rs,Rt)	Byte Q6_p_cmph_gt_RR(Word32 Rs, Word32 Rt)
Pd=cmph.gtu(Rs,#u7)	Byte Q6_p_cmph_gtu_RI(Word32 Rs, Word32 Iu7)
Pd=cmph.gtu(Rs,Rt)	Byte Q6_p_cmph_gtu_RR(Word32 Rs, Word32 Rt)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS		RegType			Maj		s5					Parse		t5				Min			d2												
1	1	0	0	0	1	1	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	1	-	-	-	d	d	Pd=cmph.eq(Rs,Rt)	
1	1	0	0	0	1	1	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	0	-	-	-	d	d	Pd=cmph.gt(Rs,Rt)	
1	1	0	0	0	1	1	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	1	-	-	-	d	d	Pd=cmph.gtu(Rs,Rt)	
ICLASS		RegType			s5					Parse																d2							
1	1	0	1	1	1	0	1	-	0	0	s	s	s	s	s	P	P	-	i	i	i	i	i	i	i	i	i	0	1	-	d	d	Pd=cmph.eq(Rs,#s8)
1	1	0	1	1	1	0	1	-	0	1	s	s	s	s	s	P	P	-	i	i	i	i	i	i	i	i	i	0	1	-	d	d	Pd=cmph.gt(Rs,#s8)
1	1	0	1	1	1	0	1	-	1	0	s	s	s	s	s	P	P	-	0	i	i	i	i	i	i	i	0	1	-	d	d	Pd=cmph.gtu(Rs,#u7)	

Field name	Description
RegType	Register type
MajOp	Major opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d2	Field to encode register d

<b>Field name</b>	<b>Description</b>
s5	Field to encode register s
t5	Field to encode register t
Maj	Major opcode
Min	Minor opcode

## Compare doublewords

Compare two 64-bit register pairs for unsigned greater than, greater than, or equal. The 8-bit predicate register Pd is set to all 1s or all 0s, depending on the result.

Syntax	Behavior
<code>Pd=cmp.eq(Rss,Rtt)</code>	<code>Pd=Rss==Rtt ? 0xff : 0x00;</code>
<code>Pd=cmp.gt(Rss,Rtt)</code>	<code>Pd=Rss&gt;Rtt ? 0xff : 0x00;</code>
<code>Pd=cmp.gtu(Rss,Rtt)</code>	<code>Pd=Rss.u64&gt;Rtt.u64 ? 0xff : 0x00;</code>

### Class: XTYPE (slots 2,3)

#### Intrinsics

<code>Pd=cmp.eq(Rss,Rtt)</code>	Byte <code>Q6_p_cmp_eq_PP(Word64 Rss, Word64 Rtt)</code>
<code>Pd=cmp.gt(Rss,Rtt)</code>	Byte <code>Q6_p_cmp_gt_PP(Word64 Rss, Word64 Rtt)</code>
<code>Pd=cmp.gtu(Rss,Rtt)</code>	Byte <code>Q6_p_cmp_gtu_PP(Word64 Rss, Word64 Rtt)</code>

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5					MinOp			d2										
1	1	0	1	0	0	1	0	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	0	-	-	-	d	d	Pd=cmp.eq(Rss,Rtt)
1	1	0	1	0	0	1	0	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	-	-	-	d	d	Pd=cmp.gt(Rss,Rtt)
1	1	0	1	0	0	1	0	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	0	-	-	-	d	d	Pd=cmp.gtu(Rss,Rtt)

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d2	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Compare bit mask

If all the bits in the mask in Rt or a short immediate are set (bitsset) or clear (bitsclear) in Rs, set the Pd to true. Otherwise, set the bits in Pd to false.

Syntax	Behavior
Pd=[!]bitsclr(Rs, #u6)	Pd=(Rs&#u) [!]=0 ? 0xff : 0x00;
Pd=[!]bitsclr(Rs, Rt)	Pd=(Rs&Rt) [!]=0 ? 0xff : 0x00;
Pd=[!]bitsset(Rs, Rt)	Pd=(Rs&Rt) [!]=Rt ? 0xff : 0x00;

### Class: XTYPE (slots 2,3)

#### Intrinsics

Pd=!bitsclr(Rs, #u6)	Byte Q6_p_not_bitsclr_RI(Word32 Rs, Word32 Iu6)
Pd=!bitsclr(Rs, Rt)	Byte Q6_p_not_bitsclr_RR(Word32 Rs, Word32 Rt)
Pd=!bitsset(Rs, Rt)	Byte Q6_p_not_bitsset_RR(Word32 Rs, Word32 Rt)
Pd=bitsclr(Rs, #u6)	Byte Q6_p_bitsclr_RI(Word32 Rs, Word32 Iu6)
Pd=bitsclr(Rs, Rt)	Byte Q6_p_bitsclr_RR(Word32 Rs, Word32 Rt)
Pd=bitsset(Rs, Rt)	Byte Q6_p_bitsset_RR(Word32 Rs, Word32 Rt)

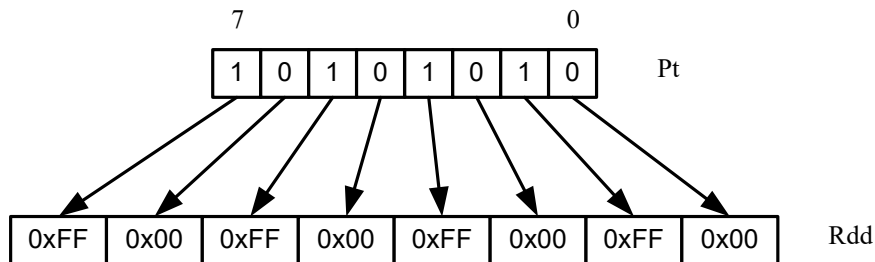
#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse										d2								
1	0	0	0	0	1	0	1	1	0	0	s	s	s	s	s	P	P	i	i	i	i	i	i	-	-	-	-	-	-	d	d	Pd=bitsclr(Rs,#u6)
1	0	0	0	0	1	0	1	1	0	1	s	s	s	s	s	P	P	i	i	i	i	i	i	-	-	-	-	-	-	d	d	Pd=!bitsclr(Rs,#u6)
ICLASS			RegType				Maj		s5					Parse					t5					d2								
1	1	0	0	0	1	1	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	-	-	-	d	d	Pd=bitsset(Rs,Rt)
1	1	0	0	0	1	1	1	0	1	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	-	-	-	d	d	Pd=!bitsset(Rs,Rt)
1	1	0	0	0	1	1	1	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	-	-	-	d	d	Pd=bitsclr(Rs,Rt)
1	1	0	0	0	1	1	1	1	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	-	-	-	d	d	Pd=!bitsclr(Rs,Rt)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d2	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
MajOp	Major opcode
Maj	Major opcode
RegType	Register type

## Mask generate from predicate

For each of the low 8 bits in predicate register Pt, when the bit is set, set the corresponding byte in 64-bit register pair Rdd to 0xff, otherwise, set the corresponding byte to 0x00.



### Syntax

```
Rdd=mask (Pt)
```

### Behavior

```
PREDUSE_TIMING;
for (i = 0; i < 8; i++) {
    Rdd.b[i]=(Pt.i?(0xff):(0x00));
}
```

## Class: XTYPE (slots 2,3)

### Intrinsics

```
Rdd=mask (Pt)
```

```
Word64 Q6_P_mask_p(Byte Pt)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				Parse								t2		d5														
1	0	0	0	0	1	1	0	-	-	-	-	-	-	-	-	P	P	-	-	-	-	t	t	-	-	-	d	d	d	d	d	Rdd=mask(Pt)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
t2	Field to encode register t
RegType	Register type

## Check for TLB match

Determine if the TLB entry in Rss matches the ASID:PPN in Rt.

### Syntax

```
Pd=tlbmatch(Rss,Rt)
```

### Behavior

```
MASK = 0x07ffffff;
TLBLO = Rss.uw[0];
TLBHI = Rss.uw[1];
SIZE = min(6, count_leading_ones(~reverse_bits(TLBLO)));
MASK &= (0xffffffff << 2*SIZE);
Pd = TLBHI.31 && ((TLBHI & MASK) == (Rt & MASK)) ? 0xff
: 0x00;
```

### Class: XTYPE (slots 2,3)

### Notes

- The predicate generated by this instruction cannot be used as a .new predicate, nor can it be automatically ANDed with another predicate.

### Intrinsics

```
Pd=tlbmatch(Rss,Rt)
```

```
Byte Q6_p_tlbmatch_PR(Word64 Rss, Word32
Rt)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5			MinOp			d2												
1	1	0	1	0	0	1	0	0	-	-	s	s	s	s	s	P	P	1	t	t	t	t	t	0	1	1	-	-	-	d	d	Pd=tlbmatch(Rss,Rt)

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d2	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t



## Predicate transfer

Pd=Rs transfers a predicate to the eight least-significant bits of a general register and zeros the other bits.

Rd=Ps transfers the eight least-significant bits of a general register to a predicate.

Syntax	Behavior
Pd=Rs	<code>Pd = Rs.ub[0];</code>
Rd=Ps	<code>PREDUSE_TIMING;</code> <code>Rd = zxt<sub>8-&gt;32</sub>(Ps);</code>

### Class: XTYPE (slots 2,3)

#### Intrinsics

Pd=Rs `Byte Q6_p_equals_R(Word32 Rs)`

Rd=Ps `Word32 Q6_R_equals_p(Byte Ps)`

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp		s5					Parse					d2												
1	0	0	0	0	1	0	1	0	1	0	s	s	s	s	s	P	P	-	-	-	-	-	-	-	-	-	-	-	-	d	d	Pd=Rs
ICLASS				RegType				MajOp		s2					Parse					d5												
1	0	0	0	1	0	0	1	-	1	-	-	-	-	s	s	P	P	-	-	-	-	-	-	-	-	-	d	d	d	d	Rd=Ps	

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d2	Field to encode register d
d5	Field to encode register d
s2	Field to encode register s
s5	Field to encode register s
MajOp	Major opcode
RegType	Register type

## Test bit

Extract a bit from a register. If the bit is true (1), set all the bits of the predicate register destination to 1. If the bit is false (0), set all the bits of the predicate register destination to 0. The bit to test can be indicated using an immediate or register value.

If a register is used to indicate the bit to test, and the value specified is out of range, the predicate result is zero.

Syntax	Behavior
<code>Pd=[!]tstbit(Rs,#u5)</code>	$Pd = (Rs \& (1 \ll \#u)) == 0 ? 0xff : 0x00;$
<code>Pd=[!]tstbit(Rs,Rt)</code>	$Pd = (zxt_{32 \rightarrow 64}(Rs) \& (sxt_{7 \rightarrow 32}(Rt) > 0) ? (zxt_{32 \rightarrow 64}(1) \ll sxt_{7 \rightarrow 32}(Rt)) : (zxt_{32 \rightarrow 64}(1) \gg sxt_{7 \rightarrow 32}(Rt))) == 0 ? 0xff : 0x00;$

### Class: XTYPE (slots 2,3)

#### Intrinsics

<code>Pd=!tstbit(Rs,#u5)</code>	Byte Q6_p_not_tstbit_RI(Word32 Rs, Word32 Iu5)
<code>Pd=!tstbit(Rs,Rt)</code>	Byte Q6_p_not_tstbit_RR(Word32 Rs, Word32 Rt)
<code>Pd=tstbit(Rs,#u5)</code>	Byte Q6_p_tstbit_RI(Word32 Rs, Word32 Iu5)
<code>Pd=tstbit(Rs,Rt)</code>	Byte Q6_p_tstbit_RR(Word32 Rs, Word32 Rt)

#### Encoding

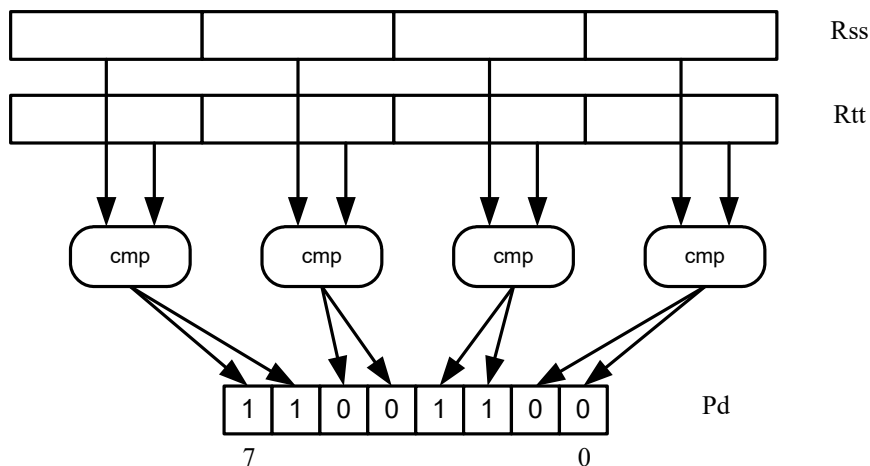
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse					d2													
1	0	0	0	0	1	0	1	0	0	0	s	s	s	s	s	P	P	0	i	i	i	i	i	-	-	-	-	-	-	d	d	Pd=tstbit(Rs,#u5)
1	0	0	0	0	1	0	1	0	0	1	s	s	s	s	s	P	P	0	i	i	i	i	i	-	-	-	-	-	-	d	d	Pd=!tstbit(Rs,#u5)
ICLASS			RegType				Maj		s5					Parse					t5					d2								
1	1	0	0	0	1	1	1	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	-	-	-	d	d	Pd=tstbit(Rs,Rt)
1	1	0	0	0	1	1	1	0	0	1	s	s	s	s	s	P	P	-	t	t	t	t	t	-	-	-	-	-	-	d	d	Pd=!tstbit(Rs,Rt)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d2	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
MajOp	Major opcode
Maj	Major opcode
RegType	Register type

## Vector compare halfwords

Compare each of four 16-bit halfwords in two 64-bit vectors and set the corresponding bits in a predicate destination to '11' if true, '00' if false.

Halfword comparisons are for equal, signed greater than, or unsigned greater than.



### Syntax

### Behavior

<code>Pd=vcmph.eq(Rss,#s8)</code>	<pre>for (i = 0; i &lt; 4; i++) {   Pd.i*2 = (Rss.h[i] == #s);   Pd.i*2+1 = (Rss.h[i] == #s); }</pre>
<code>Pd=vcmph.eq(Rss,Rtt)</code>	<pre>for (i = 0; i &lt; 4; i++) {   Pd.i*2 = (Rss.h[i] == Rtt.h[i]);   Pd.i*2+1 = (Rss.h[i] == Rtt.h[i]); }</pre>
<code>Pd=vcmph.gt(Rss,#s8)</code>	<pre>for (i = 0; i &lt; 4; i++) {   Pd.i*2 = (Rss.h[i] &gt; #s);   Pd.i*2+1 = (Rss.h[i] &gt; #s); }</pre>
<code>Pd=vcmph.gt(Rss,Rtt)</code>	<pre>for (i = 0; i &lt; 4; i++) {   Pd.i*2 = (Rss.h[i] &gt; Rtt.h[i]);   Pd.i*2+1 = (Rss.h[i] &gt; Rtt.h[i]); }</pre>
<code>Pd=vcmph.gtu(Rss,#u7)</code>	<pre>for (i = 0; i &lt; 4; i++) {   Pd.i*2 = (Rss.uh[i] &gt; #u);   Pd.i*2+1 = (Rss.uh[i] &gt; #u); }</pre>
<code>Pd=vcmph.gtu(Rss,Rtt)</code>	<pre>for (i = 0; i &lt; 4; i++) {   Pd.i*2 = (Rss.uh[i] &gt; Rtt.uh[i]);   Pd.i*2+1 = (Rss.uh[i] &gt; Rtt.uh[i]); }</pre>

**Class: XTYPE (slots 2,3)****Intrinsics**

Pd=vcmph.eq(Rss, #s8)	Byte Q6_p_vcmph_eq_PI(Word64 Rss, Word32 Is8)
Pd=vcmph.eq(Rss, Rtt)	Byte Q6_p_vcmph_eq_PP(Word64 Rss, Word64 Rtt)
Pd=vcmph.gt(Rss, #s8)	Byte Q6_p_vcmph_gt_PI(Word64 Rss, Word32 Is8)
Pd=vcmph.gt(Rss, Rtt)	Byte Q6_p_vcmph_gt_PP(Word64 Rss, Word64 Rtt)
Pd=vcmph.gtu(Rss, #u7)	Byte Q6_p_vcmph_gtu_PI(Word64 Rss, Word32 Iu7)
Pd=vcmph.gtu(Rss, Rtt)	Byte Q6_p_vcmph_gtu_PP(Word64 Rss, Word64 Rtt)

**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				s5					Parse		t5					MinOp			d2											
1	1	0	1	0	0	1	0	0	-	-	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	1	-	-	-	d	d	Pd=vcmph.eq(Rss,Rtt)
1	1	0	1	0	0	1	0	0	-	-	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	0	-	-	-	d	d	Pd=vcmph.gt(Rss,Rtt)
1	1	0	1	0	0	1	0	0	-	-	s	s	s	s	s	P	P	0	t	t	t	t	t	1	0	1	-	-	-	d	d	Pd=vcmph.gtu(Rss,Rtt)
ICLASS		RegType				s5					Parse													d2								
1	1	0	1	1	1	0	0	0	0	0	s	s	s	s	s	P	P	-	i	i	i	i	i	i	i	i	0	1	-	d	d	Pd=vcmph.eq(Rss,#s8)
1	1	0	1	1	1	0	0	0	0	1	s	s	s	s	s	P	P	-	i	i	i	i	i	i	i	i	0	1	-	d	d	Pd=vcmph.gt(Rss,#s8)
1	1	0	1	1	1	0	0	0	1	0	s	s	s	s	s	P	P	-	0	i	i	i	i	i	i	i	0	1	-	d	d	Pd=vcmph.gtu(Rss,#u7)

Field name	Description
RegType	Register type
MajOp	Major opcode
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d2	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector compare bytes for any match

Compare each byte in two 64-bit source vectors and set a predicate if any of the eight bytes are equal.

This instruction can quickly find the null terminator in a string.

Syntax	Behavior
<code>Pd=!any8(vcmpb.eq(Rss,Rtt))</code>	<pre>Pd = 0; for (i = 0; i &lt; 8; i++) {     if (Rss.b[i] == Rtt.b[i]) Pd = 0xff; } Pd = ~Pd;</pre>
<code>Pd=any8(vcmpb.eq(Rss,Rtt))</code>	<pre>Pd = 0; for (i = 0; i &lt; 8; i++) {     if (Rss.b[i] == Rtt.b[i]) Pd = 0xff; }</pre>

### Class: XTYPE (slots 2,3)

#### Intrinsics

`Pd=!any8(vcmpb.eq(Rss,Rtt))` Byte Q6\_p\_not\_any8\_vcmpb\_eq\_PP(Word64 Rss, Word64 Rtt)

`Pd=any8(vcmpb.eq(Rss,Rtt))` Byte Q6\_p\_any8\_vcmpb\_eq\_PP(Word64 Rss, Word64 Rtt)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				s5					Parse		t5					MinOp			d2										
1	1	0	1	0	0	1	0	0	-	-	s	s	s	s	s	P	P	1	t	t	t	t	t	0	0	0	-	-	-	d	d	Pd=any8(vcmpb.eq(Rss,Rtt))
1	1	0	1	0	0	1	0	0	-	-	s	s	s	s	s	P	P	1	t	t	t	t	t	0	0	1	-	-	-	d	d	Pd=!any8(vcmpb.eq(Rss,Rtt))

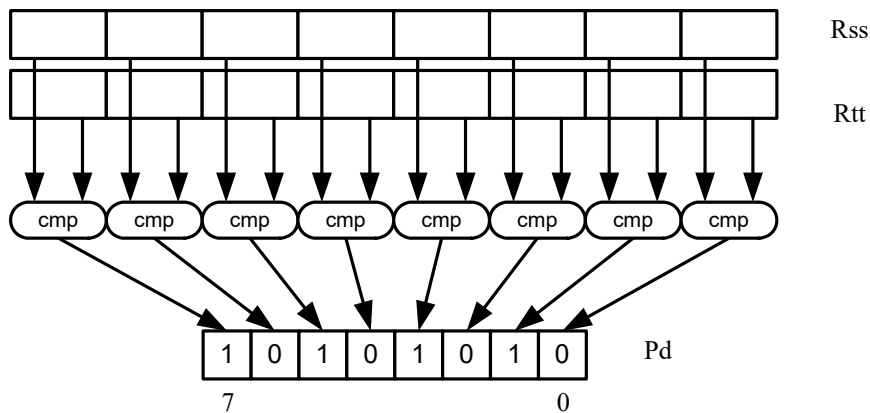
Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d2	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector compare bytes

Compare each of eight bytes in two 64-bit vectors and set the corresponding bit in a predicate destination to 1 if true, 0 if false.

Byte comparisons are for equal or for unsigned greater than.

In the following example, every other comparison is true.



Syntax	Behavior
<code>Pd=vcmpb.eq(Rss, #u8)</code>	<pre>for (i = 0; i &lt; 8; i++) {     Pd.i = (Rss.ub[i] == #u); }</pre>
<code>Pd=vcmpb.eq(Rss, Rtt)</code>	<pre>for (i = 0; i &lt; 8; i++) {     Pd.i = (Rss.b[i] == Rtt.b[i]); }</pre>
<code>Pd=vcmpb.gt(Rss, #s8)</code>	<pre>for (i = 0; i &lt; 8; i++) {     Pd.i = (Rss.b[i] &gt; #s); }</pre>
<code>Pd=vcmpb.gt(Rss, Rtt)</code>	<pre>for (i = 0; i &lt; 8; i++) {     Pd.i = (Rss.b[i] &gt; Rtt.b[i]); }</pre>
<code>Pd=vcmpb.gtu(Rss, #u7)</code>	<pre>for (i = 0; i &lt; 8; i++) {     Pd.i = (Rss.ub[i] &gt; #u); }</pre>
<code>Pd=vcmpb.gtu(Rss, Rtt)</code>	<pre>for (i = 0; i &lt; 8; i++) {     Pd.i = (Rss.ub[i] &gt; Rtt.ub[i]); }</pre>

**Class: XTYPE (slots 2,3)****Intrinsics**

Pd=vcmpb.eq(Rss, #u8)	Byte Q6_p_vcmpb_eq_PI(Word64 Rss, Word32 Iu8)
Pd=vcmpb.eq(Rss, Rtt)	Byte Q6_p_vcmpb_eq_PP(Word64 Rss, Word64 Rtt)
Pd=vcmpb.gt(Rss, #s8)	Byte Q6_p_vcmpb_gt_PI(Word64 Rss, Word32 Is8)
Pd=vcmpb.gt(Rss, Rtt)	Byte Q6_p_vcmpb_gt_PP(Word64 Rss, Word64 Rtt)
Pd=vcmpb.gtu(Rss, #u7)	Byte Q6_p_vcmpb_gtu_PI(Word64 Rss, Word32 Iu7)
Pd=vcmpb.gtu(Rss, Rtt)	Byte Q6_p_vcmpb_gtu_PP(Word64 Rss, Word64 Rtt)

**Encoding**

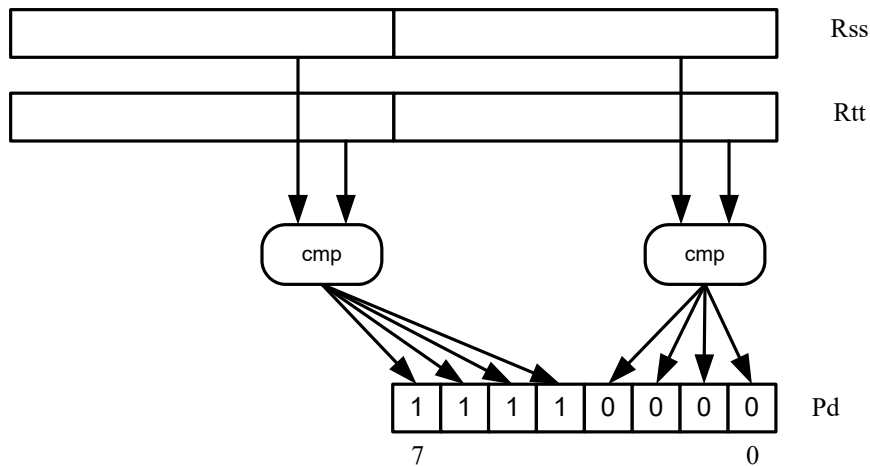
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
ICLASS		RegType				s5					Parse		t5					MinOp			d2												
1	1	0	1	0	0	1	0	0	-	-	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	0	-	-	-	d	d	Pd=vcmpb.eq(Rss,Rtt)	
1	1	0	1	0	0	1	0	0	-	-	s	s	s	s	s	P	P	0	t	t	t	t	t	1	1	1	-	-	-	d	d	Pd=vcmpb.gtu(Rss,Rtt)	
1	1	0	1	0	0	1	0	0	-	-	s	s	s	s	s	P	P	1	t	t	t	t	t	0	1	0	-	-	-	d	d	Pd=vcmpb.gt(Rss,Rtt)	
ICLASS		RegType				s5					Parse																	d2					
1	1	0	1	1	1	0	0	0	0	0	s	s	s	s	s	P	P	-	i	i	i	i	i	i	i	i	i	0	0	-	d	d	Pd=vcmpb.eq(Rss,#u8)
1	1	0	1	1	1	0	0	0	0	1	s	s	s	s	s	P	P	-	i	i	i	i	i	i	i	i	0	0	-	d	d	Pd=vcmpb.gt(Rss,#s8)	
1	1	0	1	1	1	0	0	0	1	0	s	s	s	s	s	P	P	-	0	i	i	i	i	i	i	i	0	0	-	d	d	Pd=vcmpb.gtu(Rss,#u7)	

Field name	Description
RegType	Register type
MajOp	Major opcode
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d2	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Vector compare words

Compare each of two 32-bit words in two 64-bit vectors and set the corresponding bits in a predicate destination to '1111' if true, '0000' if false.

Word comparisons are for equal, signed greater than, or unsigned greater than.



Syntax	Behavior
<code>Pd=vcmpw.eq(Rss, #s8)</code>	<code>Pd[3:0] = (Rss.w[0]==#s);</code> <code>Pd[7:4] = (Rss.w[1]==#s);</code>
<code>Pd=vcmpw.eq(Rss, Rtt)</code>	<code>Pd[3:0] = (Rss.w[0]==Rtt.w[0]);</code> <code>Pd[7:4] = (Rss.w[1]==Rtt.w[1]);</code>
<code>Pd=vcmpw.gt(Rss, #s8)</code>	<code>Pd[3:0] = (Rss.w[0]&gt;#s);</code> <code>Pd[7:4] = (Rss.w[1]&gt;#s);</code>
<code>Pd=vcmpw.gt(Rss, Rtt)</code>	<code>Pd[3:0] = (Rss.w[0]&gt;Rtt.w[0]);</code> <code>Pd[7:4] = (Rss.w[1]&gt;Rtt.w[1]);</code>
<code>Pd=vcmpw.gtu(Rss, #u7)</code>	<code>Pd[3:0] = (Rss.uw[0]&gt;#u.uw[0]);</code> <code>Pd[7:4] = (Rss.uw[1]&gt;#u.uw[0]);</code>
<code>Pd=vcmpw.gtu(Rss, Rtt)</code>	<code>Pd[3:0] = (Rss.uw[0]&gt;Rtt.uw[0]);</code> <code>Pd[7:4] = (Rss.uw[1]&gt;Rtt.uw[1]);</code>



**Class: XTYPE (slots 2,3)****Intrinsics**

Pd=vcmpw.eq(Rss, #s8)	Byte Q6_p_vcmpw_eq_PI(Word64 Rss, Word32 Is8)
Pd=vcmpw.eq(Rss, Rtt)	Byte Q6_p_vcmpw_eq_PP(Word64 Rss, Word64 Rtt)
Pd=vcmpw.gt(Rss, #s8)	Byte Q6_p_vcmpw_gt_PI(Word64 Rss, Word32 Is8)
Pd=vcmpw.gt(Rss, Rtt)	Byte Q6_p_vcmpw_gt_PP(Word64 Rss, Word64 Rtt)
Pd=vcmpw.gtu(Rss, #u7)	Byte Q6_p_vcmpw_gtu_PI(Word64 Rss, Word32 Iu7)
Pd=vcmpw.gtu(Rss, Rtt)	Byte Q6_p_vcmpw_gtu_PP(Word64 Rss, Word64 Rtt)

**Encoding**

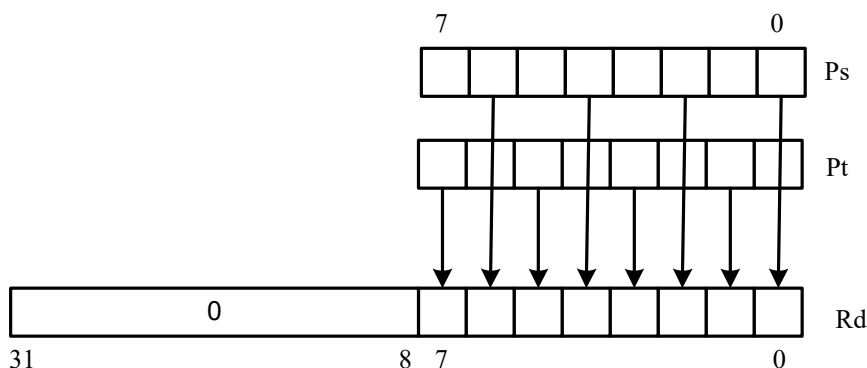
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				s5					Parse		t5					MinOp			d2											
1	1	0	1	0	0	1	0	0	-	-	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	0	-	-	-	d	d	Pd=vcmpw.eq(Rss,Rtt)
1	1	0	1	0	0	1	0	0	-	-	s	s	s	s	s	P	P	0	t	t	t	t	t	0	0	1	-	-	-	d	d	Pd=vcmpw.gt(Rss,Rtt)
1	1	0	1	0	0	1	0	0	-	-	s	s	s	s	s	P	P	0	t	t	t	t	t	0	1	0	-	-	-	d	d	Pd=vcmpw.gtu(Rss,Rtt)
ICLASS		RegType				s5					Parse												d2									
1	1	0	1	1	1	0	0	0	0	0	s	s	s	s	s	P	P	-	i	i	i	i	i	i	i	i	1	0	-	d	d	Pd=vcmpw.eq(Rss,#s8)
1	1	0	1	1	1	0	0	0	0	1	s	s	s	s	s	P	P	-	i	i	i	i	i	i	i	i	1	0	-	d	d	Pd=vcmpw.gt(Rss,#s8)
1	1	0	1	1	1	0	0	0	1	0	s	s	s	s	s	P	P	-	0	i	i	i	i	i	i	i	1	0	-	d	d	Pd=vcmpw.gtu(Rss,#u7)

Field name	Description
RegType	Register type
MajOp	Major opcode
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d2	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t

## Viterbi pack even and odd predicate bits

Pack the even and odd bits of two predicate registers into a single destination register. A variant of this instruction is R3:2 |= vitpack(P1,P0), which places the packed predicate bits into the lower eight bits of the register pair, which is preshifted by eight bits.

This instruction is useful in Viterbi decoding. Repeated use of the push version enables a history storage for traceback, purposes.



### Syntax

Rd=vitpack(Ps,Pt)

### Behavior

PREDUSE\_TIMING;  
Rd = (Ps&0x55) | (Pt&0xAA);

## Class: XTYPE (slots 2,3)

### Intrinsics

Rd=vitpack(Ps,Pt)

Word32 Q6\_R\_vitpack\_pp(Byte Ps, Byte Pt)

### Encoding

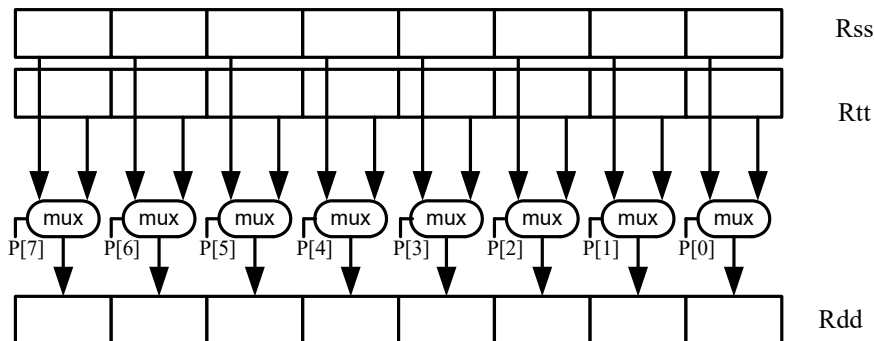
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp				s2		Parse				t2				d5										
1	0	0	0	1	0	0	1	-	0	0	-	-	-	s	s	P	P	-	-	-	-	t	t	-	-	-	d	d	d	d	d	Rd=vitpack(Ps,Pt)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s2	Field to encode register s
t2	Field to encode register t
MajOp	Major opcode
RegType	Register type

## Vector mux

Perform an element-wise byte selection between two vectors.

For each of the low eight bits of predicate register Pu, if the bit is set, the corresponding byte in Rdd is set to the corresponding byte from Rss. Otherwise, set the byte in Rdd to the byte from Rtt.



### Syntax

Rdd=vmux (Pu, Rss, Rtt)

### Behavior

```

PREDUSE_TIMING;
for (i = 0; i < 8; i++) {
    Rdd.b[i] = (Pu.i ? (Rss.b[i]) : (Rtt.b[i]));
}
    
```

## Class: XTYPE (slots 2,3)

### Intrinsics

Rdd=vmux (Pu, Rss, Rtt)

Word64 Q6\_P\_vmux\_ppp(Byte Pu, Word64 Rss, Word64 Rtt)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				s5					Parse		t5				u2		d5													
1	1	0	1	0	0	0	1	-	-	-	s	s	s	s	s	P	P	-	t	t	t	t	t	-	u	u	d	d	d	d	d	Rdd=vmux(Pu,Rss,Rtt)

Field name	Description
RegType	Register type
MinOp	Minor opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
u2	Field to encode register u

## 11.10.8 XTYPE SHIFT

The XTYPE SHIFT instruction subclass includes instructions that perform shifts.

### Mask generate from immediate

Generate a mask from two immediate values.

#### Syntax

```
Rd=mask(#u5, #U5)
```

#### Behavior

```
Rd = ((1<<#u)-1) << #U;
```

### Class: XTYPE (slots 2,3)

#### Intrinsics

```
Rd=mask(#u5, #U5)
```

```
Word32 Q6_R_mask_II(Word32 Iu5, Word32 IU5)
```

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType			MajOp			Parse				MinOp			d5															
1	0	0	0	1	1	0	1	0	I	I	-	-	-	-	-	P	P	1	i	i	i	i	i	I	I	I	d	d	d	d	d	Rd=mask(#u5,#U5)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

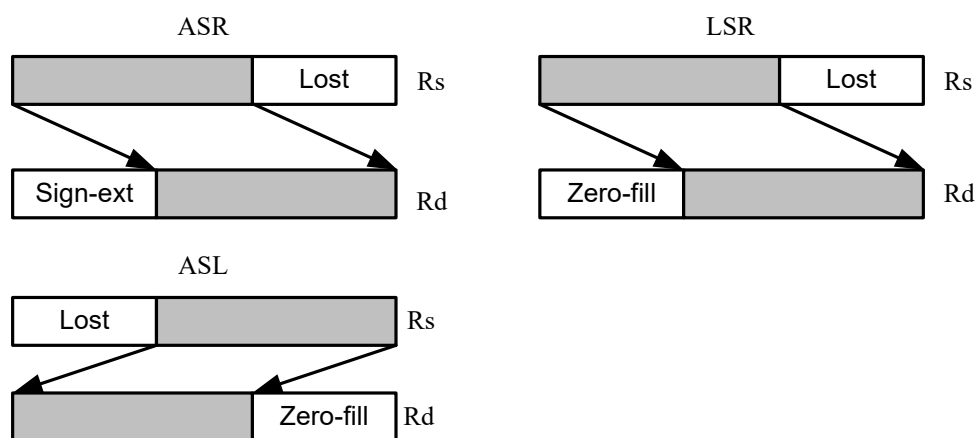
## Shift by immediate

Shift the source register value right or left based on the type of instruction. In these instructions, the shift amount is contained in an unsigned immediate (five bits for 32-bit shifts, six bits for 64-bit shifts) and the shift instruction gives the shift direction.

Arithmetic right shifts place the sign bit of the source value in the vacated positions.

Logical right shifts place zeros in the vacated positions.

Left shifts always zero-fill the vacated bits.



Syntax	Behavior
$Rd=asl(Rs, \#u5)$	$Rd = Rs \ll \#u;$
$Rd=asr(Rs, \#u5)$	$Rd = Rs \gg \#u;$
$Rd=lsr(Rs, \#u5)$	$Rd = Rs \ggg \#u;$
$Rd=rol(Rs, \#u5)$	$Rd = Rs \ll_R \#u;$
$Rdd=asl(Rss, \#u6)$	$Rdd = Rss \ll \#u;$
$Rdd=asr(Rss, \#u6)$	$Rdd = Rss \gg \#u;$
$Rdd=lsr(Rss, \#u6)$	$Rdd = Rss \ggg \#u;$
$Rdd=rol(Rss, \#u6)$	$Rdd = Rss \ll_R \#u;$

### Class: XTYPE (slots 2,3)

#### Intrinsics

$Rd=asl(Rs, \#u5)$	<code>Word32 Q6_R_asl_RI(Word32 Rs, Word32 Iu5)</code>
$Rd=asr(Rs, \#u5)$	<code>Word32 Q6_R_asr_RI(Word32 Rs, Word32 Iu5)</code>
$Rd=lsr(Rs, \#u5)$	<code>Word32 Q6_R_lsr_RI(Word32 Rs, Word32 Iu5)</code>
$Rd=rol(Rs, \#u5)$	<code>Word32 Q6_R_rol_RI(Word32 Rs, Word32 Iu5)</code>

Rdd=asl(Rss,#u6)	Word64 Q6_P_asl_PI(Word64 Rss, Word32 Iu6)
Rdd=asr(Rss,#u6)	Word64 Q6_P_asr_PI(Word64 Rss, Word32 Iu6)
Rdd=lsr(Rss,#u6)	Word64 Q6_P_lsr_PI(Word64 Rss, Word32 Iu6)
Rdd=rol(Rss,#u6)	Word64 Q6_P_rol_PI(Word64 Rss, Word32 Iu6)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp				s5					Parse			MinOp			d5										
1	0	0	0	0	0	0	0	0	0	0	s	s	s	s	s	P	P	i	i	i	i	i	i	0	0	0	d	d	d	d	d	Rdd=asr(Rss,#u6)
1	0	0	0	0	0	0	0	0	0	0	s	s	s	s	s	P	P	i	i	i	i	i	i	0	0	1	d	d	d	d	d	Rdd=lsr(Rss,#u6)
1	0	0	0	0	0	0	0	0	0	0	s	s	s	s	s	P	P	i	i	i	i	i	i	0	1	0	d	d	d	d	d	Rdd=asl(Rss,#u6)
1	0	0	0	0	0	0	0	0	0	0	s	s	s	s	s	P	P	i	i	i	i	i	i	0	1	1	d	d	d	d	d	Rdd=rol(Rss,#u6)
1	0	0	0	1	1	0	0	0	0	0	s	s	s	s	s	P	P	0	i	i	i	i	i	0	0	0	d	d	d	d	d	Rd=asr(Rs,#u5)
1	0	0	0	1	1	0	0	0	0	0	s	s	s	s	s	P	P	0	i	i	i	i	i	0	0	1	d	d	d	d	d	Rd=lsr(Rs,#u5)
1	0	0	0	1	1	0	0	0	0	0	s	s	s	s	s	P	P	0	i	i	i	i	i	0	1	0	d	d	d	d	d	Rd=asl(Rs,#u5)
1	0	0	0	1	1	0	0	0	0	0	s	s	s	s	s	P	P	0	i	i	i	i	i	0	1	1	d	d	d	d	d	Rd=rol(Rs,#u5)

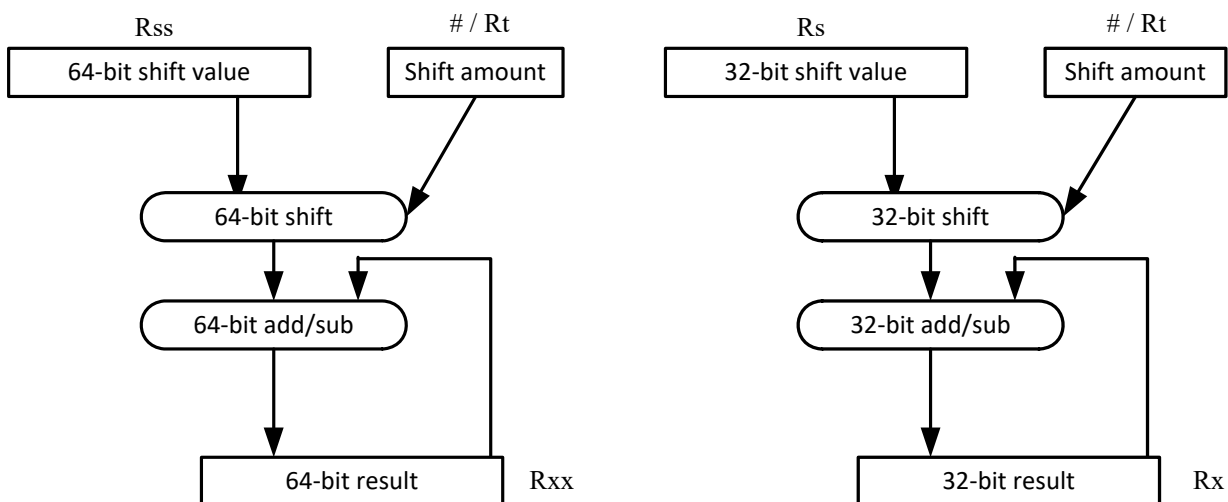
Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Shift by immediate and accumulate

Shift the source register value right or left, based on the type of instruction. In these instructions, an unsigned immediate (5 bits for 32-bit shifts, 6 bits for 64-bit shifts) contains the shift amount, and the shift instruction gives the shift direction.

Arithmetic right shifts place the sign bit of the source value in the vacated positions. Logical right shifts place zeros in the vacated positions. Left shifts always zero-fill the vacated bits.

After shifting, add or subtract the shifted value from the destination register or register pair.



### Syntax

```
Rx=add(#u8,as1(Rx,#U5))
```

```
Rx=add(#u8,lsr(Rx,#U5))
```

```
Rx=sub(#u8,as1(Rx,#U5))
```

```
Rx=sub(#u8,lsr(Rx,#U5))
```

```
Rx[+-]=as1(Rs,#u5)
```

```
Rx[+-]=asr(Rs,#u5)
```

```
Rx[+-]=lsr(Rs,#u5)
```

```
Rx[+-]=rol(Rs,#u5)
```

```
Rxx[+-]=as1(Rss,#u6)
```

```
Rxx[+-]=asr(Rss,#u6)
```

```
Rxx[+-]=lsr(Rss,#u6)
```

```
Rxx[+-]=rol(Rss,#u6)
```

### Behavior

```
Rx=apply_extension(#u)+(Rx<<#U);
```

```
Rx=apply_extension(#u)+(((unsigned int)Rx)>>#U);
```

```
Rx=apply_extension(#u)-(Rx<<#U);
```

```
Rx=apply_extension(#u)-(((unsigned int)Rx)>>#U);
```

```
Rx = Rx [+-] Rs << #u;
```

```
Rx = Rx [+-] Rs >> #u;
```

```
Rx = Rx [+-] Rs >>> #u;
```

```
Rx = Rx [+-] Rs <<R #u;
```

```
Rxx = Rxx [+-] Rss << #u;
```

```
Rxx = Rxx [+-] Rss >> #u;
```

```
Rxx = Rxx [+-] Rss >>> #u;
```

```
Rxx = Rxx [+-] Rss <<R #u;
```

**Class: XTYPE (slots 2,3)****Intrinsics**

Rx+=asl(Rs, #u5)	Word32 Q6_R_aslacc_RI(Word32 Rx, Word32 Rs, Word32 Iu5)
Rx+=asr(Rs, #u5)	Word32 Q6_R_asracc_RI(Word32 Rx, Word32 Rs, Word32 Iu5)
Rx+=lsr(Rs, #u5)	Word32 Q6_R_lsracc_RI(Word32 Rx, Word32 Rs, Word32 Iu5)
Rx+=rol(Rs, #u5)	Word32 Q6_R_rolacc_RI(Word32 Rx, Word32 Rs, Word32 Iu5)
Rx-=asl(Rs, #u5)	Word32 Q6_R_aslnac_RI(Word32 Rx, Word32 Rs, Word32 Iu5)
Rx-=asr(Rs, #u5)	Word32 Q6_R_asrnac_RI(Word32 Rx, Word32 Rs, Word32 Iu5)
Rx-=lsr(Rs, #u5)	Word32 Q6_R_lsrnac_RI(Word32 Rx, Word32 Rs, Word32 Iu5)
Rx-=rol(Rs, #u5)	Word32 Q6_R_rolnac_RI(Word32 Rx, Word32 Rs, Word32 Iu5)
Rx=add(#u8, asl(Rx, #U5))	Word32 Q6_R_add_asl_IRI(Word32 Iu8, Word32 Rx, Word32 IU5)
Rx=add(#u8, lsr(Rx, #U5))	Word32 Q6_R_add_lsr_IRI(Word32 Iu8, Word32 Rx, Word32 IU5)
Rx=sub(#u8, asl(Rx, #U5))	Word32 Q6_R_sub_asl_IRI(Word32 Iu8, Word32 Rx, Word32 IU5)
Rx=sub(#u8, lsr(Rx, #U5))	Word32 Q6_R_sub_lsr_IRI(Word32 Iu8, Word32 Rx, Word32 IU5)
Rxx+=asl(Rss, #u6)	Word64 Q6_P_aslacc_PI(Word64 Rxx, Word64 Rss, Word32 Iu6)
Rxx+=asr(Rss, #u6)	Word64 Q6_P_asracc_PI(Word64 Rxx, Word64 Rss, Word32 Iu6)
Rxx+=lsr(Rss, #u6)	Word64 Q6_P_lsracc_PI(Word64 Rxx, Word64 Rss, Word32 Iu6)
Rxx+=rol(Rss, #u6)	Word64 Q6_P_rolacc_PI(Word64 Rxx, Word64 Rss, Word32 Iu6)
Rxx-=asl(Rss, #u6)	Word64 Q6_P_aslnac_PI(Word64 Rxx, Word64 Rss, Word32 Iu6)
Rxx-=asr(Rss, #u6)	Word64 Q6_P_asrnac_PI(Word64 Rxx, Word64 Rss, Word32 Iu6)
Rxx-=lsr(Rss, #u6)	Word64 Q6_P_lsrnac_PI(Word64 Rxx, Word64 Rss, Word32 Iu6)
Rxx-=rol(Rss, #u6)	Word64 Q6_P_rolnac_PI(Word64 Rxx, Word64 Rss, Word32 Iu6)



### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse				MinOp			x5											
1	0	0	0	0	0	1	0	0	0	-	s	s	s	s	s	P	P	i	i	i	i	i	i	0	0	0	x	x	x	x	x	Rxx-=asr(Rss,#u6)
1	0	0	0	0	0	1	0	0	0	-	s	s	s	s	s	P	P	i	i	i	i	i	i	0	0	1	x	x	x	x	x	Rxx-=lsl(Rss,#u6)
1	0	0	0	0	0	1	0	0	0	-	s	s	s	s	s	P	P	i	i	i	i	i	i	0	1	0	x	x	x	x	x	Rxx-=asl(Rss,#u6)
1	0	0	0	0	0	1	0	0	0	-	s	s	s	s	s	P	P	i	i	i	i	i	i	0	1	1	x	x	x	x	x	Rxx-=rol(Rss,#u6)
1	0	0	0	0	0	1	0	0	0	-	s	s	s	s	s	P	P	i	i	i	i	i	i	1	0	0	x	x	x	x	x	Rxx+=asr(Rss,#u6)
1	0	0	0	0	0	1	0	0	0	-	s	s	s	s	s	P	P	i	i	i	i	i	i	1	0	1	x	x	x	x	x	Rxx+=lsl(Rss,#u6)
1	0	0	0	0	0	1	0	0	0	-	s	s	s	s	s	P	P	i	i	i	i	i	i	1	1	1	x	x	x	x	x	Rxx+=asl(Rss,#u6)
1	0	0	0	1	1	1	0	0	0	-	s	s	s	s	s	P	P	0	i	i	i	i	i	0	0	0	x	x	x	x	x	Rx-=asr(Rs,#u5)
1	0	0	0	1	1	1	0	0	0	-	s	s	s	s	s	P	P	0	i	i	i	i	i	0	0	1	x	x	x	x	x	Rx-=lsl(Rs,#u5)
1	0	0	0	1	1	1	0	0	0	-	s	s	s	s	s	P	P	0	i	i	i	i	i	0	1	0	x	x	x	x	x	Rx-=asl(Rs,#u5)
1	0	0	0	1	1	1	0	0	0	-	s	s	s	s	s	P	P	0	i	i	i	i	i	0	1	1	x	x	x	x	x	Rx-=rol(Rs,#u5)
1	0	0	0	1	1	1	0	0	0	-	s	s	s	s	s	P	P	0	i	i	i	i	i	1	0	0	x	x	x	x	x	Rx+=asr(Rs,#u5)
1	0	0	0	1	1	1	0	0	0	-	s	s	s	s	s	P	P	0	i	i	i	i	i	1	0	1	x	x	x	x	x	Rx+=lsl(Rs,#u5)
1	0	0	0	1	1	1	0	0	0	-	s	s	s	s	s	P	P	0	i	i	i	i	i	1	1	0	x	x	x	x	x	Rx+=asl(Rs,#u5)
1	0	0	0	1	1	1	0	0	0	-	s	s	s	s	s	P	P	0	i	i	i	i	i	1	1	1	x	x	x	x	x	Rx+=rol(Rs,#u5)
ICLASS			RegType				x5					Parse				MajOp																
1	1	0	1	1	1	1	0	i	i	i	x	x	x	x	x	P	P	i	l	l	l	l	l	i	i	i	0	i	1	0	-	Rx=add(#u8,asl(Rx,#U5))
1	1	0	1	1	1	1	0	i	i	i	x	x	x	x	x	P	P	i	l	l	l	l	l	i	i	i	0	i	1	1	-	Rx=sub(#u8,asl(Rx,#U5))
1	1	0	1	1	1	1	0	i	i	i	x	x	x	x	x	P	P	i	l	l	l	l	l	i	i	i	1	i	1	0	-	Rx=add(#u8,lsl(Rx,#U5))
1	1	0	1	1	1	1	0	i	i	i	x	x	x	x	x	P	P	i	l	l	l	l	l	i	i	i	1	i	1	1	-	Rx=sub(#u8,lsl(Rx,#U5))

<b>Field name</b>	<b>Description</b>
RegType	Register type
MajOp	Major opcode
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s
x5	Field to encode register x
MinOp	Minor opcode

## Shift by immediate and add

Shift Rs left by 0-7 bits, add to Rt, and place the result in Rd.

This instruction is useful for calculating array pointers, where destruction of the base pointer is undesirable.

### Syntax

```
Rd=addasl(Rt,Rs,#u3)
```

### Behavior

```
Rd = Rt + Rs << #u;
```

### Class: XTYPE (slots 2,3)

### Intrinsics

```
Rd=addasl(Rt,Rs,#u3)
```

```
Word32 Q6_R_addasl_RRI(Word32 Rt, Word32
Rs, Word32 Iu3)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				Maj		s5					Parse		t5				Min		d5										
1	1	0	0	0	1	0	0	0	0	0	s	s	s	s	s	P	P	0	t	t	t	t	t	i	i	i	d	d	d	d	d	Rd=addasl(Rt,Rs,#u3)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
Maj	Major opcode
Min	Minor opcode
RegType	Register type

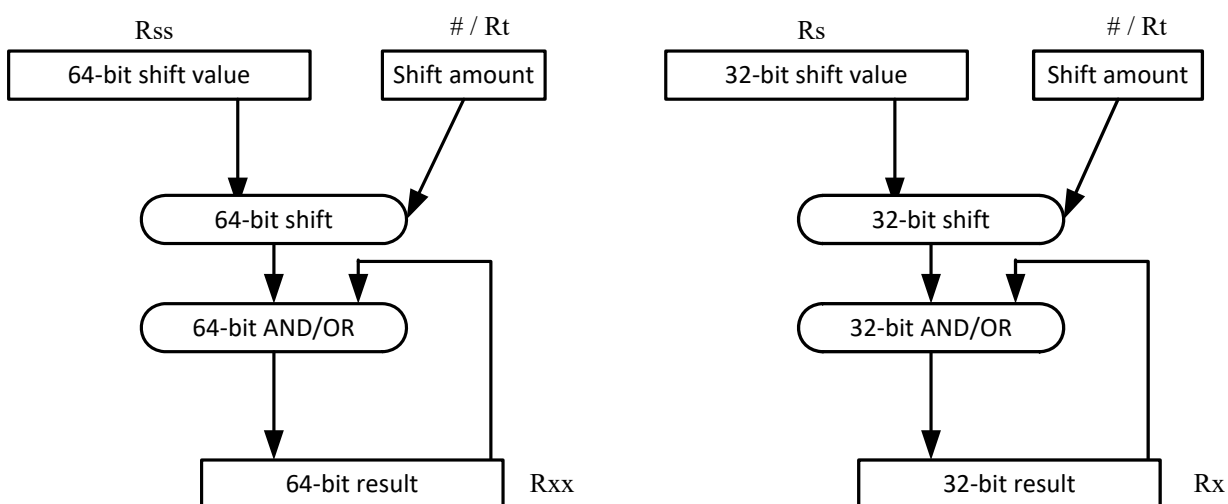
## Shift by immediate and logical

Shift the source register value right or left based on the type of instruction. In these instructions, an unsigned immediate (five bits for 32-bit shifts, six bits for 64-bit shifts) contains the shift amount and the shift instruction gives the shift direction.

Arithmetic right shifts place the sign bit of the source value in the vacated positions. Logical right shifts place zeros in the vacated positions. Left shifts always zero-fill the vacated bits.

After shifting, take the logical AND, OR, or XOR of the shifted amount and the destination register or register pair, and place the result back in the destination register or register pair.

Saturation is not available for these instructions.



Syntax	Behavior
<code>Rx=and(#u8,asl(Rx,#U5))</code>	<code>Rx=apply_extension(#u) &amp; (Rx&lt;&lt;#U);</code>
<code>Rx=and(#u8,lsr(Rx,#U5))</code>	<code>Rx=apply_extension(#u) &amp; (((unsigned int)Rx)&gt;&gt;#U);</code>
<code>Rx=or(#u8,asl(Rx,#U5))</code>	<code>Rx=apply_extension(#u)   (Rx&lt;&lt;#U);</code>
<code>Rx=or(#u8,lsr(Rx,#U5))</code>	<code>Rx=apply_extension(#u)   (((unsigned int)Rx)&gt;&gt;#U);</code>
<code>Rx[&amp; =asl(Rs,#u5)</code>	<code>Rx = Rx [ &amp;] Rs &lt;&lt; #u;</code>
<code>Rx[&amp; =asr(Rs,#u5)</code>	<code>Rx = Rx [ &amp;] Rs &gt;&gt; #u;</code>
<code>Rx[&amp; =lsr(Rs,#u5)</code>	<code>Rx = Rx [ &amp;] Rs &gt;&gt;&gt; #u;</code>
<code>Rx[&amp; =rol(Rs,#u5)</code>	<code>Rx = Rx [ &amp;] Rs &lt;&lt;<sub>R</sub> #u;</code>
<code>Rx^=asl(Rs,#u5)</code>	<code>Rx = Rx ^ Rs &lt;&lt; #u;</code>
<code>Rx^=lsr(Rs,#u5)</code>	<code>Rx = Rx ^ Rs &gt;&gt;&gt; #u;</code>
<code>Rx^=rol(Rs,#u5)</code>	<code>Rx = Rx ^ Rs &lt;&lt;<sub>R</sub> #u;</code>
<code>Rxx[&amp; =asl(Rss,#u6)</code>	<code>Rxx = Rxx [ &amp;] Rss &lt;&lt; #u;</code>

Syntax	Behavior
$Rxx[ \&   ] = asr(Rss, \#u6)$	$Rxx = Rxx [ \&   ] Rss \gg \#u;$
$Rxx[ \&   ] = lsr(Rss, \#u6)$	$Rxx = Rxx [ \&   ] Rss \ggg \#u;$
$Rxx[ \&   ] = rol(Rss, \#u6)$	$Rxx = Rxx [ \&   ] Rss \ll_R \#u;$
$Rxx^{\wedge} = asl(Rss, \#u6)$	$Rxx = Rxx \wedge Rss \ll \#u;$
$Rxx^{\wedge} = lsr(Rss, \#u6)$	$Rxx = Rxx \wedge Rss \ggg \#u;$
$Rxx^{\wedge} = rol(Rss, \#u6)$	$Rxx = Rxx \wedge Rss \ll_R \#u;$

### Class: XTYPE (slots 2,3)

#### Intrinsics

$Rx\&=asl(Rs, \#u5)$	Word32 Q6_R_asland_RI (Word32 Rx, Word32 Rs, Word32 Iu5)
$Rx\&=asr(Rs, \#u5)$	Word32 Q6_R_asrand_RI (Word32 Rx, Word32 Rs, Word32 Iu5)
$Rx\&=lsr(Rs, \#u5)$	Word32 Q6_R_lsrland_RI (Word32 Rx, Word32 Rs, Word32 Iu5)
$Rx\&=rol(Rs, \#u5)$	Word32 Q6_R_roland_RI (Word32 Rx, Word32 Rs, Word32 Iu5)
$Rx=and(\#u8, asl(Rx, \#U5))$	Word32 Q6_R_and_asl_IRI (Word32 Iu8, Word32 Rx, Word32 IU5)
$Rx=and(\#u8, lsr(Rx, \#U5))$	Word32 Q6_R_and_lsr_IRI (Word32 Iu8, Word32 Rx, Word32 IU5)
$Rx=or(\#u8, asl(Rx, \#U5))$	Word32 Q6_R_or_asl_IRI (Word32 Iu8, Word32 Rx, Word32 IU5)
$Rx=or(\#u8, lsr(Rx, \#U5))$	Word32 Q6_R_or_lsr_IRI (Word32 Iu8, Word32 Rx, Word32 IU5)
$Rx^{\wedge}=asl(Rs, \#u5)$	Word32 Q6_R_aslxacc_RI (Word32 Rx, Word32 Rs, Word32 Iu5)
$Rx^{\wedge}=lsr(Rs, \#u5)$	Word32 Q6_R_lsrxacc_RI (Word32 Rx, Word32 Rs, Word32 Iu5)
$Rx^{\wedge}=rol(Rs, \#u5)$	Word32 Q6_R_rolxacc_RI (Word32 Rx, Word32 Rs, Word32 Iu5)
$Rx =asl(Rs, \#u5)$	Word32 Q6_R_aslor_RI (Word32 Rx, Word32 Rs, Word32 Iu5)
$Rx =asr(Rs, \#u5)$	Word32 Q6_R_asror_RI (Word32 Rx, Word32 Rs, Word32 Iu5)
$Rx =lsr(Rs, \#u5)$	Word32 Q6_R_lsr_ror_RI (Word32 Rx, Word32 Rs, Word32 Iu5)
$Rx =rol(Rs, \#u5)$	Word32 Q6_R_rol_ror_RI (Word32 Rx, Word32 Rs, Word32 Iu5)
$Rxx\&=asl(Rss, \#u6)$	Word64 Q6_P_asland_PI (Word64 Rxx, Word64 Rss, Word32 Iu6)
$Rxx\&=asr(Rss, \#u6)$	Word64 Q6_P_asrand_PI (Word64 Rxx, Word64 Rss, Word32 Iu6)
$Rxx\&=lsr(Rss, \#u6)$	Word64 Q6_P_lsrland_PI (Word64 Rxx, Word64 Rss, Word32 Iu6)
$Rxx\&=rol(Rss, \#u6)$	Word64 Q6_P_roland_PI (Word64 Rxx, Word64 Rss, Word32 Iu6)
$Rxx^{\wedge}=asl(Rss, \#u6)$	Word64 Q6_P_aslxacc_PI (Word64 Rxx, Word64 Rss, Word32 Iu6)
$Rxx^{\wedge}=lsr(Rss, \#u6)$	Word64 Q6_P_lsrxacc_PI (Word64 Rxx, Word64 Rss, Word32 Iu6)

Rxx^=rol(Rss,#u6)	Word64 Q6_P_rolxacc_PI(Word64 Rxx, Word64 Rss, Word32 Iu6)
Rxx =asl(Rss,#u6)	Word64 Q6_P_aslor_PI(Word64 Rxx, Word64 Rss, Word32 Iu6)
Rxx =asr(Rss,#u6)	Word64 Q6_P_asror_PI(Word64 Rxx, Word64 Rss, Word32 Iu6)
Rxx =lsr(Rss,#u6)	Word64 Q6_P_lsror_PI(Word64 Rxx, Word64 Rss, Word32 Iu6)
Rxx =rol(Rss,#u6)	Word64 Q6_P_rolor_PI(Word64 Rxx, Word64 Rss, Word32 Iu6)

### Encoding

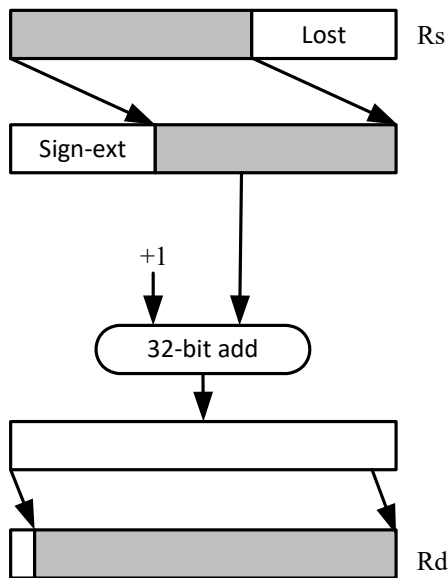
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse					MinOp			x5										
1	0	0	0	0	0	1	0	0	1	-	s	s	s	s	s	P	P	i	i	i	i	i	i	0	0	0	x	x	x	x	x	Rxx&=asr(Rss,#u6)
1	0	0	0	0	0	1	0	0	1	-	s	s	s	s	s	P	P	i	i	i	i	i	i	0	0	1	x	x	x	x	x	Rxx&=lsr(Rss,#u6)
1	0	0	0	0	0	1	0	0	1	-	s	s	s	s	s	P	P	i	i	i	i	i	i	0	1	0	x	x	x	x	x	Rxx&=asl(Rss,#u6)
1	0	0	0	0	0	1	0	0	1	-	s	s	s	s	s	P	P	i	i	i	i	i	i	0	1	1	x	x	x	x	x	Rxx&=rol(Rss,#u6)
1	0	0	0	0	0	1	0	0	1	-	s	s	s	s	s	P	P	i	i	i	i	i	i	1	0	0	x	x	x	x	x	Rxx =asr(Rss,#u6)
1	0	0	0	0	0	1	0	0	1	-	s	s	s	s	s	P	P	i	i	i	i	i	i	1	0	1	x	x	x	x	x	Rxx =lsr(Rss,#u6)
1	0	0	0	0	0	1	0	0	1	-	s	s	s	s	s	P	P	i	i	i	i	i	i	1	1	0	x	x	x	x	x	Rxx =asl(Rss,#u6)
1	0	0	0	0	0	1	0	0	1	-	s	s	s	s	s	P	P	i	i	i	i	i	i	1	1	1	x	x	x	x	x	Rxx =rol(Rss,#u6)
1	0	0	0	0	0	1	0	1	0	-	s	s	s	s	s	P	P	i	i	i	i	i	i	0	0	1	x	x	x	x	x	Rxx^=lsr(Rss,#u6)
1	0	0	0	0	0	1	0	1	0	-	s	s	s	s	s	P	P	i	i	i	i	i	i	0	1	0	x	x	x	x	x	Rxx^=asl(Rss,#u6)
1	0	0	0	0	0	1	0	1	0	-	s	s	s	s	s	P	P	i	i	i	i	i	i	0	1	1	x	x	x	x	x	Rxx^=rol(Rss,#u6)
1	0	0	0	1	1	1	0	0	1	-	s	s	s	s	s	P	P	0	i	i	i	i	i	0	0	0	x	x	x	x	x	Rx&=asr(Rs,#u5)
1	0	0	0	1	1	1	0	0	1	-	s	s	s	s	s	P	P	0	i	i	i	i	i	0	0	1	x	x	x	x	x	Rx&=lsr(Rs,#u5)
1	0	0	0	1	1	1	0	0	1	-	s	s	s	s	s	P	P	0	i	i	i	i	i	0	1	0	x	x	x	x	x	Rx&=asl(Rs,#u5)
1	0	0	0	1	1	1	0	0	1	-	s	s	s	s	s	P	P	0	i	i	i	i	i	0	1	1	x	x	x	x	x	Rx&=rol(Rs,#u5)
1	0	0	0	1	1	1	0	0	1	-	s	s	s	s	s	P	P	0	i	i	i	i	i	1	0	1	x	x	x	x	x	Rx =lsr(Rs,#u5)
1	0	0	0	1	1	1	0	0	1	-	s	s	s	s	s	P	P	0	i	i	i	i	i	1	1	0	x	x	x	x	x	Rx =asl(Rs,#u5)
1	0	0	0	1	1	1	0	0	1	-	s	s	s	s	s	P	P	0	i	i	i	i	i	1	1	1	x	x	x	x	x	Rx =rol(Rs,#u5)
1	0	0	0	1	1	1	0	1	0	-	s	s	s	s	s	P	P	0	i	i	i	i	i	0	0	1	x	x	x	x	x	Rx^=lsr(Rs,#u5)
1	0	0	0	1	1	1	0	1	0	-	s	s	s	s	s	P	P	0	i	i	i	i	i	0	1	0	x	x	x	x	x	Rx^=asl(Rs,#u5)
1	0	0	0	1	1	1	0	1	0	-	s	s	s	s	s	P	P	0	i	i	i	i	i	0	1	1	x	x	x	x	x	Rx^=rol(Rs,#u5)
ICLASS	RegType			x5					Parse					MajOp																		
1	1	0	1	1	1	1	0	i	i	i	x	x	x	x	x	P	P	i	l	l	l	l	l	i	i	i	0	i	0	0	-	Rx=and(#u8,asl(Rx,#U5))
1	1	0	1	1	1	1	0	i	i	i	x	x	x	x	x	P	P	i	l	l	l	l	l	i	i	i	0	i	0	1	-	Rx=or(#u8,asl(Rx,#U5))
1	1	0	1	1	1	1	0	i	i	i	x	x	x	x	x	P	P	i	l	l	l	l	l	i	i	i	1	i	0	0	-	Rx=and(#u8,lsr(Rx,#U5))
1	1	0	1	1	1	1	0	i	i	i	x	x	x	x	x	P	P	i	l	l	l	l	l	i	i	i	1	i	0	1	-	Rx=or(#u8,lsr(Rx,#U5))

<b>Field name</b>	<b>Description</b>
RegType	Register type
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s
x5	Field to encode register x
MajOp	Major opcode
MinOp	Minor opcode

## Shift right by immediate with rounding

Perform an arithmetic right shift by an immediate amount, and then round the result. This instruction works by first shifting right, then adding the value +1 to the result, and finally shifting right again by one bit. The right shifts always inserts the sign-bit in the vacated position.

When using the `asrrnd` instruction, the assembler adjusts the immediate appropriately.



Syntax	Behavior
<code>Rd=asr(Rs, #u5) :rnd</code>	<code>Rd = ((Rs &gt;&gt; #u)+1) &gt;&gt; 1;</code>
<code>Rd=asrrnd(Rs, #u5)</code>	<pre>if ("#u5==0") {     Assembler mapped to: "Rd=Rs"; } else {     Assembler mapped to: "Rd=asr(Rs, #u5-1) :rnd"; }</pre>
<code>Rdd=asr(Rss, #u6) :rnd</code>	<pre>tmp = Rss &gt;&gt; #u; rnd = tmp &amp; 1; Rdd = tmp &gt;&gt; 1 + rnd;</pre>
<code>Rdd=asrrnd(Rss, #u6)</code>	<pre>if ("#u6==0") {     Assembler mapped to: "Rdd=Rss"; } else {     Assembler mapped to: "Rdd=asr(Rss, #u6-1) :rnd"; }</pre>

**Class: XTYPE (slots 2,3)****Intrinsics**

Rd=asr(Rs,#u5):rnd	Word32 Q6_R_asr_RI_rnd(Word32 Rs, Word32 Iu5)
Rd=asrrnd(Rs,#u5)	Word32 Q6_R_asrrnd_RI(Word32 Rs, Word32 Iu5)
Rdd=asr(Rss,#u6):rnd	Word64 Q6_P_asr_PI_rnd(Word64 Rss, Word32 Iu6)
Rdd=asrrnd(Rss,#u6)	Word64 Q6_P_asrrnd_PI(Word64 Rss, Word32 Iu6)

**Encoding**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse		MinOp			d5													
1	0	0	0	0	0	0	0	1	1	0	s	s	s	s	s	P	P	i	i	i	i	i	i	1	1	1	d	d	d	d	d	Rdd=asr(Rss,#u6):rnd
1	0	0	0	1	1	0	0	0	1	0	s	s	s	s	s	P	P	0	i	i	i	i	i	0	0	0	d	d	d	d	d	Rd=asr(Rs,#u5):rnd

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Shift left by immediate with saturation

Perform a left shift of the 32-bit source register value by an immediate amount and saturate.

Saturation works by first sign-extending the 32-bit Rs register to 64 bits. It is then left-shifted by the immediate amount. If this 64-bit value cannot fit in a signed 32-bit number (the upper word is not the sign-extension of bit 31), saturation is performed based on the sign of the original value. Saturation clamps the 32-bit result to the range 0x8000\_0000 to 0x7fff\_ffff.

### Syntax

```
Rd=asl(Rs, #u5):sat
```

### Behavior

```
Rd = sat32(sxt32->64(Rs) << #u);
```

### Class: XTYPE (slots 2,3)

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

### Intrinsics

```
Rd=asl(Rs, #u5):sat
```

```
Word32 Q6_R_asl_RI_sat(Word32 Rs, Word32  
Iu5)
```

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp		s5					Parse		MinOp			d5												
1	0	0	0	1	1	0	0	0	1	0	s	s	s	s	s	P	P	0	i	i	i	i	i	0	1	0	d	d	d	d	d	Rd=asl(Rs,#u5):sat

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

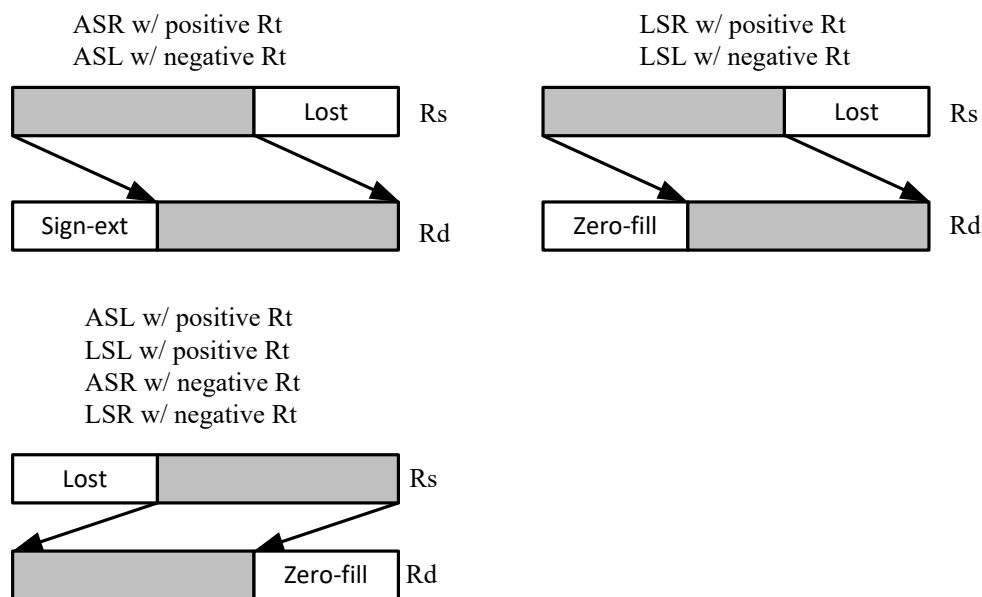


## Shift by register

The shift amount is the least significant seven bits of  $R_t$ , treated as a two's complement value. If the shift amount is negative (bit 6 of  $R_t$  is set), reverse the direction of the shift indicated in the opcode (see Figure).

The source data to shift is always performed as a 64-bit shift. When the  $R_s$  source register is a 32-bit register, this register is first sign or zero-extended to 64-bits. Arithmetic shifts sign-extend the 32-bit source to 64-bits, whereas logical shifts zero extend.

The 64-bit source value is then right or left shifted based on the shift amount and the type of instruction. Arithmetic right shifts place the sign bit of the source value in the vacated positions. Logical right shifts place zeros in the vacated positions.



Syntax	Behavior
$Rd=asl(Rs, Rt)$	$shamt=sxt_{7 \rightarrow 32}(Rt);$ $Rd = (shamt > 0) ? (sxt_{32 \rightarrow 64}(Rs) \ll shamt) : (sxt_{32 \rightarrow 64}(Rs) \gg shamt);$
$Rd=asr(Rs, Rt)$	$shamt=sxt_{7 \rightarrow 32}(Rt);$ $Rd = (shamt > 0) ? (sxt_{32 \rightarrow 64}(Rs) \gg shamt) : (sxt_{32 \rightarrow 64}(Rs) \ll shamt);$
$Rd=lsl(\#s6, Rt)$	$shamt = sxt_{7 \rightarrow 32}(Rt);$ $Rd = (shamt > 0) ? (zxt_{32 \rightarrow 64}(\#s) \ll shamt) : (zxt_{32 \rightarrow 64}(\#s) \gg shamt);$
$Rd=lsr(Rs, Rt)$	$shamt=sxt_{7 \rightarrow 32}(Rt);$ $Rd = (shamt > 0) ? (zxt_{32 \rightarrow 64}(Rs) \ll shamt) : (zxt_{32 \rightarrow 64}(Rs) \gg shamt);$

Syntax	Behavior
Rd=lsr(Rs,Rt)	shamt=sxt <sub>7-&gt;32</sub> (Rt); Rd = (shamt>0)?(zxt <sub>32-&gt;64</sub> (Rs)>>>shamt):(zxt <sub>32-&gt;64</sub> (Rs)<<<shamt);
Rdd=asl(Rss,Rt)	shamt=sxt <sub>7-&gt;32</sub> (Rt); Rdd = (shamt>0)?(Rss<<<shamt):(Rss>>>shamt);
Rdd=asr(Rss,Rt)	shamt=sxt <sub>7-&gt;32</sub> (Rt); Rdd = (shamt>0)?(Rss>>>shamt):(Rss<<<shamt);
Rdd=lsl(Rss,Rt)	shamt=sxt <sub>7-&gt;32</sub> (Rt); Rdd = (shamt>0)?(Rss<<<shamt):(Rss>>>shamt);
Rdd=lsr(Rss,Rt)	shamt=sxt <sub>7-&gt;32</sub> (Rt); Rdd = (shamt>0)?(Rss>>>shamt):(Rss<<<shamt);

### Class: XTYPE (slots 2,3)

#### Intrinsics

Rd=asl(Rs,Rt)	Word32 Q6_R_asl_RR(Word32 Rs, Word32 Rt)
Rd=asr(Rs,Rt)	Word32 Q6_R_asr_RR(Word32 Rs, Word32 Rt)
Rd=lsl(#s6,Rt)	Word32 Q6_R_lsl_IR(Word32 Is6, Word32 Rt)
Rd=lsl(Rs,Rt)	Word32 Q6_R_lsl_RR(Word32 Rs, Word32 Rt)
Rd=lsr(Rs,Rt)	Word32 Q6_R_lsr_RR(Word32 Rs, Word32 Rt)
Rdd=asl(Rss,Rt)	Word64 Q6_P_asl_PR(Word64 Rss, Word32 Rt)
Rdd=asr(Rss,Rt)	Word64 Q6_P_asr_PR(Word64 Rss, Word32 Rt)
Rdd=lsl(Rss,Rt)	Word64 Q6_P_lsl_PR(Word64 Rss, Word32 Rt)
Rdd=lsr(Rss,Rt)	Word64 Q6_P_lsr_PR(Word64 Rss, Word32 Rt)

#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				Maj		s5					Parse		t5				Min		d5											
1	1	0	0	0	0	1	1	1	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	d	d	d	d	d	Rdd=asr(Rss,Rt)
1	1	0	0	0	0	1	1	1	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	-	d	d	d	d	d	Rdd=lsr(Rss,Rt)
1	1	0	0	0	0	1	1	1	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	-	d	d	d	d	d	Rdd=asl(Rss,Rt)
1	1	0	0	0	0	1	1	1	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	-	d	d	d	d	d	Rdd=lsl(Rss,Rt)
1	1	0	0	0	1	1	0	0	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	d	d	d	d	d	Rd=asr(Rs,Rt)
1	1	0	0	0	1	1	0	0	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	-	d	d	d	d	d	Rd=lsr(Rs,Rt)
1	1	0	0	0	1	1	0	0	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	-	d	d	d	d	d	Rd=asl(Rs,Rt)
1	1	0	0	0	1	1	0	0	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	-	d	d	d	d	d	Rd=lsl(Rs,Rt)
ICLASS		RegType				Maj							Parse		t5				Min		d5											
1	1	0	0	0	1	1	0	1	0	-	i	i	i	i	i	P	P	-	t	t	t	t	t	1	1	i	d	d	d	d	d	Rd=lsl(#s6,Rt)

<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
Maj	Major opcode
Min	Minor opcode
RegType	Register type

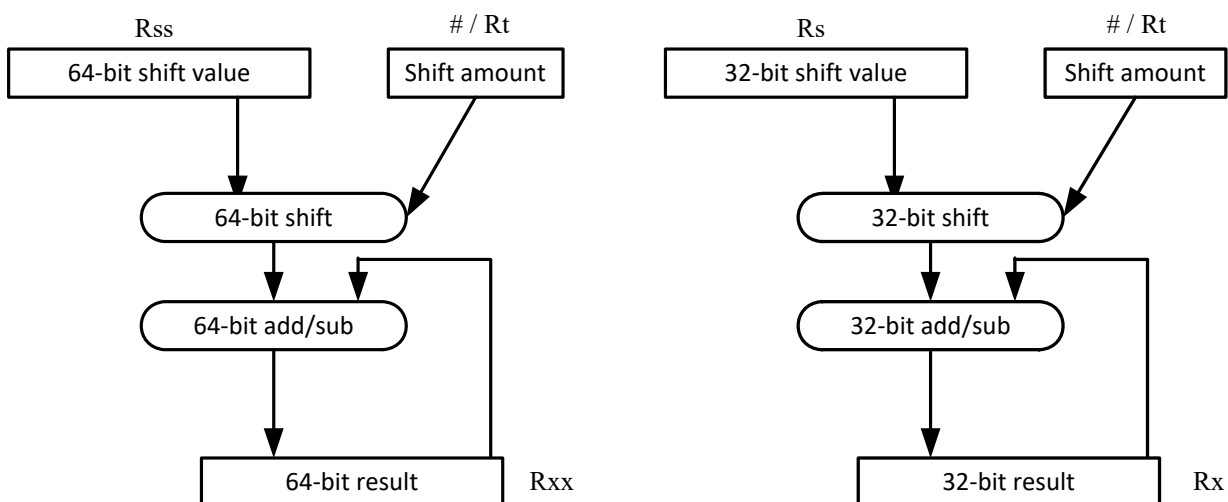
## Shift by register and accumulate

The shift amount is the least significant seven bits of  $R_t$ , treated as a two's complement value. When the shift amount is negative (bit 6 of  $R_t$  is set), reverse the direction of the shift indicated in the opcode.

Shift the source register value right or left based on the shift amount and the type of instruction. Arithmetic right shifts place the sign bit of the source value in the vacated positions. Logical right shifts place zeros in the vacated positions.

The shift operation is always performed as a 64-bit shift. When  $R_s$  is a 32-bit register, this register is first sign- or zero-extended to 64-bits. Arithmetic shifts sign-extend the 32-bit source to 64-bits, whereas logical shifts zero extend.

After shifting, add or subtract the 64-bit shifted amount from the destination register or register pair.



### Syntax

Syntax	Behavior
$R_x[+-]=asl(R_s, R_t)$	$shamt = sxt_{7 \rightarrow 32}(R_t);$ $R_x = R_x[+-](shamt > 0) ? (sxt_{32 \rightarrow 64}(R_s) \ll shamt) : (sxt_{32 \rightarrow 64}(R_s) \gg shamt);$
$R_x[+-]=asr(R_s, R_t)$	$shamt = sxt_{7 \rightarrow 32}(R_t);$ $R_x = R_x[+-](shamt > 0) ? (sxt_{32 \rightarrow 64}(R_s) \gg shamt) : (sxt_{32 \rightarrow 64}(R_s) \ll shamt);$
$R_x[+-]=lsl(R_s, R_t)$	$shamt = sxt_{7 \rightarrow 32}(R_t);$ $R_x = R_x[+-](shamt > 0) ? (zxt_{32 \rightarrow 64}(R_s) \ll shamt) : (zxt_{32 \rightarrow 64}(R_s) \gg shamt);$
$R_x[+-]=lsr(R_s, R_t)$	$shamt = sxt_{7 \rightarrow 32}(R_t);$ $R_x = R_x[+-](shamt > 0) ? (zxt_{32 \rightarrow 64}(R_s) \gg shamt) : (zxt_{32 \rightarrow 64}(R_s) \ll shamt);$
$R_{xx}[+-]=asl(R_{ss}, R_t)$	$shamt = sxt_{7 \rightarrow 32}(R_t);$ $R_{xx} = R_{xx}[+-](shamt > 0) ? (R_{ss} \ll shamt) : (R_{ss} \gg shamt);$

Syntax	Behavior
<code>Rxx[+-] ]=asr(Rss,Rt)</code>	<code>shamt=sxt<sub>7-&gt;32</sub>(Rt); Rxx = Rxx [+-] (shamt&gt;0)?(Rss&gt;&gt;shamt):(Rss&lt;&lt;shamt);</code>
<code>Rxx[+-] ]=lsl(Rss,Rt)</code>	<code>shamt=sxt<sub>7-&gt;32</sub>(Rt); Rxx = Rxx [+-] (shamt&gt;0)?(Rss&lt;&lt;shamt):(Rss&gt;&gt;&gt;shamt);</code>
<code>Rxx[+-] ]=lsr(Rss,Rt)</code>	<code>shamt=sxt<sub>7-&gt;32</sub>(Rt); Rxx = Rxx [+-] (shamt&gt;0)?(Rss&gt;&gt;&gt;shamt):(Rss&lt;&lt;shamt);</code>

## Class: XTYPE (slots 2,3)

### Intrinsics

<code>Rx+=asl(Rs,Rt)</code>	<code>Word32 Q6_R_aslacc_RR(Word32 Rx, Word32 Rs, Word32 Rt)</code>
<code>Rx+=asr(Rs,Rt)</code>	<code>Word32 Q6_R_asracc_RR(Word32 Rx, Word32 Rs, Word32 Rt)</code>
<code>Rx+=lsl(Rs,Rt)</code>	<code>Word32 Q6_R_lslacc_RR(Word32 Rx, Word32 Rs, Word32 Rt)</code>
<code>Rx+=lsr(Rs,Rt)</code>	<code>Word32 Q6_R_lsracc_RR(Word32 Rx, Word32 Rs, Word32 Rt)</code>
<code>Rx-=asl(Rs,Rt)</code>	<code>Word32 Q6_R_aslnac_RR(Word32 Rx, Word32 Rs, Word32 Rt)</code>
<code>Rx-=asr(Rs,Rt)</code>	<code>Word32 Q6_R_asrnac_RR(Word32 Rx, Word32 Rs, Word32 Rt)</code>
<code>Rx-=lsl(Rs,Rt)</code>	<code>Word32 Q6_R_lslnac_RR(Word32 Rx, Word32 Rs, Word32 Rt)</code>
<code>Rx-=lsr(Rs,Rt)</code>	<code>Word32 Q6_R_lsrnac_RR(Word32 Rx, Word32 Rs, Word32 Rt)</code>
<code>Rxx+=asl(Rss,Rt)</code>	<code>Word64 Q6_P_aslacc_PR(Word64 Rxx, Word64 Rss, Word32 Rt)</code>
<code>Rxx+=asr(Rss,Rt)</code>	<code>Word64 Q6_P_asracc_PR(Word64 Rxx, Word64 Rss, Word32 Rt)</code>
<code>Rxx+=lsl(Rss,Rt)</code>	<code>Word64 Q6_P_lslacc_PR(Word64 Rxx, Word64 Rss, Word32 Rt)</code>
<code>Rxx+=lsr(Rss,Rt)</code>	<code>Word64 Q6_P_lsracc_PR(Word64 Rxx, Word64 Rss, Word32 Rt)</code>
<code>Rxx-=asl(Rss,Rt)</code>	<code>Word64 Q6_P_aslnac_PR(Word64 Rxx, Word64 Rss, Word32 Rt)</code>
<code>Rxx-=asr(Rss,Rt)</code>	<code>Word64 Q6_P_asrnac_PR(Word64 Rxx, Word64 Rss, Word32 Rt)</code>
<code>Rxx-=lsl(Rss,Rt)</code>	<code>Word64 Q6_P_lslnac_PR(Word64 Rxx, Word64 Rss, Word32 Rt)</code>
<code>Rxx-=lsr(Rss,Rt)</code>	<code>Word64 Q6_P_lsrnac_PR(Word64 Rxx, Word64 Rss, Word32 Rt)</code>

## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				Maj		s5					Parse		t5					Min		x5									
1	1	0	0	1	0	1	1	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	x	x	x	x	x	Rxx-=asr(Rss,Rt)
1	1	0	0	1	0	1	1	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	-	x	x	x	x	x	Rxx-=lsr(Rss,Rt)
1	1	0	0	1	0	1	1	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	-	x	x	x	x	x	Rxx-=asl(Rss,Rt)
1	1	0	0	1	0	1	1	1	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	-	x	x	x	x	x	Rxx-=lsl(Rss,Rt)
1	1	0	0	1	0	1	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	x	x	x	x	x	Rxx+=asr(Rss,Rt)
1	1	0	0	1	0	1	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	-	x	x	x	x	x	Rxx+=lsr(Rss,Rt)
1	1	0	0	1	0	1	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	-	x	x	x	x	x	Rxx+=asl(Rss,Rt)
1	1	0	0	1	0	1	1	1	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	-	x	x	x	x	x	Rxx+=lsl(Rss,Rt)
1	1	0	0	1	1	0	0	1	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	x	x	x	x	x	Rx-=asr(Rs,Rt)
1	1	0	0	1	1	0	0	1	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	-	x	x	x	x	x	Rx-=lsr(Rs,Rt)
1	1	0	0	1	1	0	0	1	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	-	x	x	x	x	x	Rx-=asl(Rs,Rt)
1	1	0	0	1	1	0	0	1	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	-	x	x	x	x	x	Rx-=lsl(Rs,Rt)
1	1	0	0	1	1	0	0	1	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	x	x	x	x	x	Rx+=asr(Rs,Rt)
1	1	0	0	1	1	0	0	1	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	-	x	x	x	x	x	Rx+=lsr(Rs,Rt)
1	1	0	0	1	1	0	0	1	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	-	x	x	x	x	x	Rx+=asl(Rs,Rt)
1	1	0	0	1	1	0	0	1	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	-	x	x	x	x	x	Rx+=lsl(Rs,Rt)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Shift by register and logical

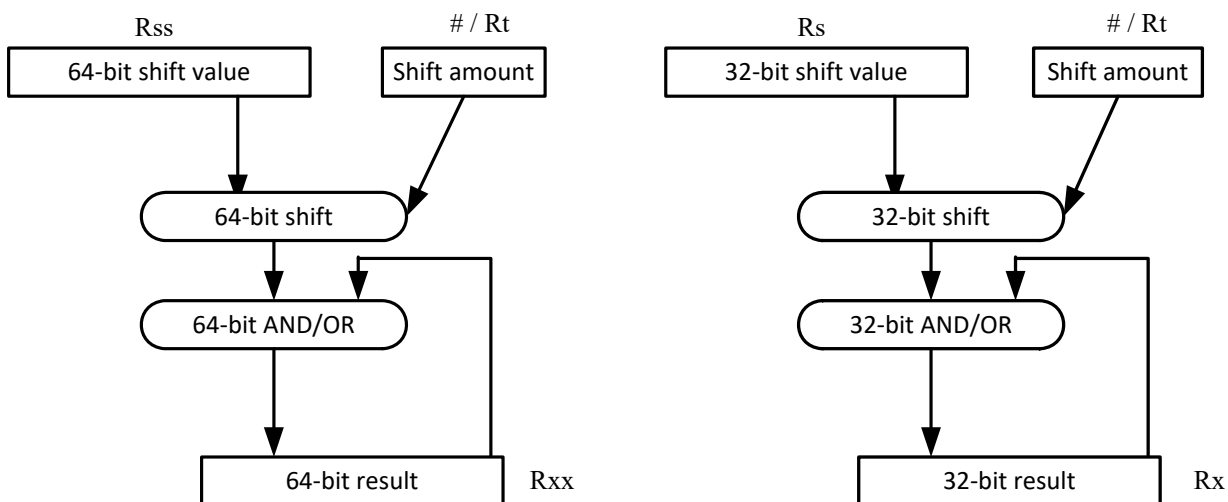
The shift amount is the least significant seven bits of  $R_t$ , treated as a two's complement value. When the shift amount is negative (bit 6 of  $R_t$  is set), reverse the direction of the shift indicated in the opcode.

Shift the source register value right or left based on the shift amount and the type of instruction. Arithmetic right shifts place the sign bit of the source value in the vacated positions. Logical right shifts place zeros in the vacated positions.

The shift operation is always performed as a 64-bit shift. When the  $R_s$  source register is a 32-bit register, this register is first sign or zero-extended to 64-bits. Arithmetic shifts sign-extend the 32-bit source to 64-bits, whereas logical shifts zero extend.

After shifting, take the logical AND or OR of the shifted amount and the destination register or register pair, and place the result back in the destination register or register pair.

Saturation is not available for these instructions.



### Syntax

$Rx[ \& | ] = asl (Rs, Rt)$

```
shamt=sxt7->32 (Rt);
Rx = Rx [ | & ] (shamt>0) ? (sxt32->64 (Rs) <<shamt) : (sxt32->64 (Rs) >>shamt);
```

$Rx[ \& | ] = asr (Rs, Rt)$

```
shamt=sxt7->32 (Rt);
Rx = Rx [ | & ] (shamt>0) ? (sxt32->64 (Rs) >>shamt) : (sxt32->64 (Rs) <<shamt);
```

$Rx[ \& | ] = lsl (Rs, Rt)$

```
shamt=sxt7->32 (Rt);
Rx = Rx [ | & ] (shamt>0) ? (zxt32->64 (Rs) <<shamt) : (zxt32->64 (Rs) >>>shamt);
```

$Rx[ \& | ] = lsr (Rs, Rt)$

```
shamt=sxt7->32 (Rt);
Rx = Rx [ | & ] (shamt>0) ? (zxt32->64 (Rs) >>>shamt) : (zxt32->64 (Rs) <<shamt);
```

### Behavior

Syntax	Behavior
<code>Rxx[&amp; ]=asl(Rss,Rt)</code>	<code>shamt=sxt<sub>7-&gt;32</sub>(Rt);</code> <code>Rxx = Rxx [ &amp;] (shamt&gt;0)?(Rss&lt;&lt;shamt):(Rss&gt;&gt;shamt);</code>
<code>Rxx[&amp; ]=asr(Rss,Rt)</code>	<code>shamt=sxt<sub>7-&gt;32</sub>(Rt);</code> <code>Rxx = Rxx [ &amp;] (shamt&gt;0)?(Rss&gt;&gt;shamt):(Rss&lt;&lt;shamt);</code>
<code>Rxx[&amp; ]=lsl(Rss,Rt)</code>	<code>shamt=sxt<sub>7-&gt;32</sub>(Rt);</code> <code>Rxx = Rxx [ &amp;] (shamt&gt;0)?(Rss&lt;&lt;shamt):(Rss&gt;&gt;&gt;shamt);</code>
<code>Rxx[&amp; ]=lsr(Rss,Rt)</code>	<code>shamt=sxt<sub>7-&gt;32</sub>(Rt);</code> <code>Rxx = Rxx [ &amp;] (shamt&gt;0)?(Rss&gt;&gt;&gt;shamt):(Rss&lt;&lt;shamt);</code>
<code>Rxx^=asl(Rss,Rt)</code>	<code>shamt=sxt<sub>7-&gt;32</sub>(Rt);</code> <code>Rxx = Rxx ^ (shamt&gt;0)?(Rss&lt;&lt;shamt):(Rss&gt;&gt;shamt);</code>
<code>Rxx^=asr(Rss,Rt)</code>	<code>shamt=sxt<sub>7-&gt;32</sub>(Rt);</code> <code>Rxx = Rxx ^ (shamt&gt;0)?(Rss&gt;&gt;shamt):(Rss&lt;&lt;shamt);</code>
<code>Rxx^=lsl(Rss,Rt)</code>	<code>shamt=sxt<sub>7-&gt;32</sub>(Rt);</code> <code>Rxx = Rxx ^ (shamt&gt;0)?(Rss&lt;&lt;shamt):(Rss&gt;&gt;&gt;shamt);</code>
<code>Rxx^=lsr(Rss,Rt)</code>	<code>shamt=sxt<sub>7-&gt;32</sub>(Rt);</code> <code>Rxx = Rxx ^ (shamt&gt;0)?(Rss&gt;&gt;&gt;shamt):(Rss&lt;&lt;shamt);</code>

**Class: XTYPE (slots 2,3)****Intrinsics**

<code>Rx&amp;=asl(Rs,Rt)</code>	<code>Word32 Q6_R_asland_RR(Word32 Rx, Word32 Rs, Word32 Rt)</code>
<code>Rx&amp;=asr(Rs,Rt)</code>	<code>Word32 Q6_R_asrand_RR(Word32 Rx, Word32 Rs, Word32 Rt)</code>
<code>Rx&amp;=lsl(Rs,Rt)</code>	<code>Word32 Q6_R_lsland_RR(Word32 Rx, Word32 Rs, Word32 Rt)</code>
<code>Rx&amp;=lsr(Rs,Rt)</code>	<code>Word32 Q6_R_lsrnd_RR(Word32 Rx, Word32 Rs, Word32 Rt)</code>
<code>Rx =asl(Rs,Rt)</code>	<code>Word32 Q6_R_aslor_RR(Word32 Rx, Word32 Rs, Word32 Rt)</code>
<code>Rx =asr(Rs,Rt)</code>	<code>Word32 Q6_R_asror_RR(Word32 Rx, Word32 Rs, Word32 Rt)</code>
<code>Rx =lsl(Rs,Rt)</code>	<code>Word32 Q6_R_lslor_RR(Word32 Rx, Word32 Rs, Word32 Rt)</code>
<code>Rx =lsr(Rs,Rt)</code>	<code>Word32 Q6_R_lsrer_RR(Word32 Rx, Word32 Rs, Word32 Rt)</code>
<code>Rxx&amp;=asl(Rss,Rt)</code>	<code>Word64 Q6_P_asland_PR(Word64 Rxx, Word64 Rss, Word32 Rt)</code>
<code>Rxx&amp;=asr(Rss,Rt)</code>	<code>Word64 Q6_P_asrand_PR(Word64 Rxx, Word64 Rss, Word32 Rt)</code>
<code>Rxx&amp;=lsl(Rss,Rt)</code>	<code>Word64 Q6_P_lsland_PR(Word64 Rxx, Word64 Rss, Word32 Rt)</code>
<code>Rxx&amp;=lsr(Rss,Rt)</code>	<code>Word64 Q6_P_lsrnd_PR(Word64 Rxx, Word64 Rss, Word32 Rt)</code>
<code>Rxx^=asl(Rss,Rt)</code>	<code>Word64 Q6_P_aslxacc_PR(Word64 Rxx, Word64 Rss, Word32 Rt)</code>
<code>Rxx^=asr(Rss,Rt)</code>	<code>Word64 Q6_P_asrxacc_PR(Word64 Rxx, Word64 Rss, Word32 Rt)</code>
<code>Rxx^=lsl(Rss,Rt)</code>	<code>Word64 Q6_P_lslxacc_PR(Word64 Rxx, Word64 Rss, Word32 Rt)</code>



Rxx <sup>^</sup> =lsr(Rss,Rt)	Word64 Q6_P_lsr <sub>xacc</sub> _PR(Word64 Rxx, Word64 Rss, Word32 Rt)
Rxx =asl(Rss,Rt)	Word64 Q6_P_asl <sub>or</sub> _PR(Word64 Rxx, Word64 Rss, Word32 Rt)
Rxx =asr(Rss,Rt)	Word64 Q6_P_asr <sub>or</sub> _PR(Word64 Rxx, Word64 Rss, Word32 Rt)
Rxx =lsl(Rss,Rt)	Word64 Q6_P_lsl <sub>or</sub> _PR(Word64 Rxx, Word64 Rss, Word32 Rt)
Rxx =lsr(Rss,Rt)	Word64 Q6_P_lsr <sub>or</sub> _PR(Word64 Rxx, Word64 Rss, Word32 Rt)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				Maj		s5					Parse		t5					Min		x5									
1	1	0	0	1	0	1	1	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	x	x	x	x	x	Rxx =asr(Rss,Rt)
1	1	0	0	1	0	1	1	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	-	x	x	x	x	x	Rxx =lsr(Rss,Rt)
1	1	0	0	1	0	1	1	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	-	x	x	x	x	x	Rxx =asl(Rss,Rt)
1	1	0	0	1	0	1	1	0	0	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	-	x	x	x	x	x	Rxx =lsl(Rss,Rt)
1	1	0	0	1	0	1	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	x	x	x	x	x	Rxx&=asr(Rss,Rt)
1	1	0	0	1	0	1	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	-	x	x	x	x	x	Rxx&=lsr(Rss,Rt)
1	1	0	0	1	0	1	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	-	x	x	x	x	x	Rxx&=asl(Rss,Rt)
1	1	0	0	1	0	1	1	0	1	0	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	-	x	x	x	x	x	Rxx&=lsl(Rss,Rt)
1	1	0	0	1	1	0	0	0	0	0	-	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	x	x	x	x	x	Rx =asr(Rs,Rt)
1	1	0	0	1	1	0	0	0	0	0	-	s	s	s	s	P	P	-	t	t	t	t	t	0	1	-	x	x	x	x	x	Rx =lsr(Rs,Rt)
1	1	0	0	1	1	0	0	0	0	0	-	s	s	s	s	P	P	-	t	t	t	t	t	1	0	-	x	x	x	x	x	Rx =asl(Rs,Rt)
1	1	0	0	1	1	0	0	0	0	0	-	s	s	s	s	P	P	-	t	t	t	t	t	1	1	-	x	x	x	x	x	Rx =lsl(Rs,Rt)
1	1	0	0	1	1	0	0	0	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	x	x	x	x	x	Rx&=asr(Rs,Rt)
1	1	0	0	1	1	0	0	0	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	-	x	x	x	x	x	Rx&=lsr(Rs,Rt)
1	1	0	0	1	1	0	0	0	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	-	x	x	x	x	x	Rx&=asl(Rs,Rt)
1	1	0	0	1	1	0	0	0	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	-	x	x	x	x	x	Rx&=lsl(Rs,Rt)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
s5	Field to encode register s
t5	Field to encode register t
x5	Field to encode register x
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Shift by register with saturation

The shift amount is the least significant seven bits of *Rt*, treated as a two's complement value. When the shift amount is negative (bit 6 of *Rt* is set), reverse the direction of the shift indicated in the opcode.

Saturation is available for 32-bit arithmetic left shifts: either an ASL instruction with positive *Rt*, or an ASR instruction with negative *Rt*. Saturation works by first sign-extending the 32-bit *Rs* register to 64 bits. It is then shifted by the shift amount. If this 64-bit value cannot fit in a signed 32-bit number (the upper word is not the sign-extension of bit 31), saturation is performed based on the sign of the original value. Saturation clamps the 32-bit result to the range 0x80000000 to 0x7fffffff.

Syntax	Behavior
<code>Rd=asl (Rs,Rt) :sat</code>	<code>shamt=sxt7-&gt;32 (Rt) ;</code> <code>Rd = bidir_shifl (Rs,shamt) ;</code>
<code>Rd=asr (Rs,Rt) :sat</code>	<code>shamt=sxt7-&gt;32 (Rt) ;</code> <code>Rd = bidir_shiftr (Rs,shamt) ;</code>

### Class: XTYPE (slots 2,3)

#### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

#### Intrinsics

`Rd=asl (Rs,Rt) :sat`      `Word32 Q6_R_asl_RR_sat (Word32 Rs, Word32 Rt)`

`Rd=asr (Rs,Rt) :sat`      `Word32 Q6_R_asr_RR_sat (Word32 Rs, Word32 Rt)`

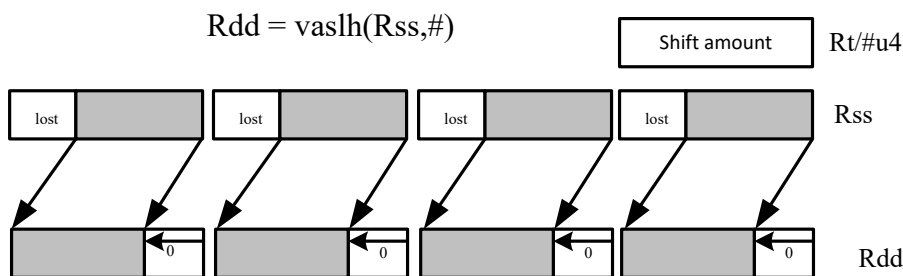
#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				Maj		s5					Parse		t5					Min		d5										
1	1	0	0	0	1	1	0	0	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	d	d	d	d	d	Rd=asr(Rs,Rt):sat
1	1	0	0	0	1	1	0	0	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	-	d	d	d	d	d	Rd=asl(Rs,Rt):sat

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Vector shift halfwords by immediate

Shift individual halfwords of the source vector. Arithmetic right shifts place the sign bit of the source values in the vacated positions. Logical right shifts place zeros in the vacated positions.



Syntax	Behavior
$Rdd = \text{vaslh}(Rss, \#u4)$	<pre>for (i=0;i&lt;4;i++) {     Rdd.h[i] = (Rss.h[i] &lt;&lt; #u); }</pre>
$Rdd = \text{vasrh}(Rss, \#u4)$	<pre>for (i=0;i&lt;4;i++) {     Rdd.h[i] = (Rss.h[i] &gt;&gt; #u); }</pre>
$Rdd = \text{vlsrh}(Rss, \#u4)$	<pre>for (i=0;i&lt;4;i++) {     Rdd.h[i] = (Rss.uh[i] &gt;&gt; #u); }</pre>

### Class: XTYPE (slots 2,3)

#### Intrinsics

$Rdd = \text{vaslh}(Rss, \#u4)$	Word64 Q6_P_vaslh_PI (Word64 Rss, Word32 Iu4)
$Rdd = \text{vasrh}(Rss, \#u4)$	Word64 Q6_P_vasrh_PI (Word64 Rss, Word32 Iu4)
$Rdd = \text{vlsrh}(Rss, \#u4)$	Word64 Q6_P_vlsrh_PI (Word64 Rss, Word32 Iu4)

#### Encoding

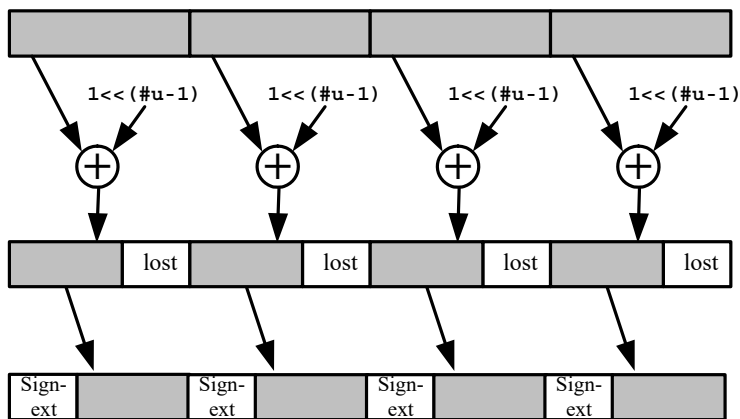
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp				s5					Parse		MinOp			d5										
1	0	0	0	0	0	0	0	1	0	0	s	s	s	s	s	P	P	0	0	i	i	i	i	0	0	0	d	d	d	d	d	Rdd=vasrh(Rss,#u4)
1	0	0	0	0	0	0	0	1	0	0	s	s	s	s	s	P	P	0	0	i	i	i	i	0	0	1	d	d	d	d	d	Rdd=vlsrh(Rss,#u4)
1	0	0	0	0	0	0	0	1	0	0	s	s	s	s	s	P	P	0	0	i	i	i	i	0	1	0	d	d	d	d	d	Rdd=vaslh(Rss,#u4)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Vector arithmetic shift halfwords with round

For each halfword in the vector, round then arithmetic shift right by an immediate amount. The results are stored in the destination register.

$$Rdd = \text{vasrh}(Rss, \#u):rnd$$



### Syntax

`Rdd=vasrh(Rss,#u4):raw`

`Rdd=vasrh(Rss,#u4):rnd`

### Behavior

```
for (i=0;i<4;i++) {
    Rdd.h[i]=((Rss.h[i] >> #u)+1)>>1;
}
```

```
if ("#u4==0") {
    Assembler mapped to: "Rdd=Rss";
} else {
    Assembler mapped to: "Rdd=vasrh(Rss,#u4-1):raw";
}
```

## Class: XTYPE (slots 2,3)

### Intrinsics

`Rdd=vasrh(Rss,#u4):rnd` `Word64 Q6_P_vasrh_PI_rnd(Word64 Rss, Word32 Iu4)`

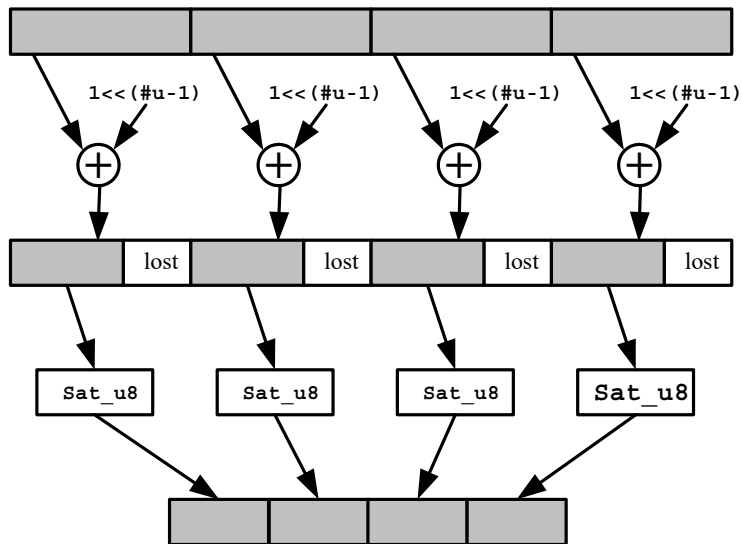
### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				MajOp				s5					Parse		MinOp			d5												
1	0	0	0	0	0	0	0	0	0	1	s	s	s	s	s	P	P	0	0	i	i	i	i	0	0	0	d	d	d	d	d	Rdd=vasrh(Rss,#u4):raw

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Vector arithmetic shift halfwords with saturate and pack

For each halfword in the vector, optionally round, then arithmetic shift right by an immediate amount. Saturate the results to unsigned [0-255] and then pack in the destination register.

$$Rd = \text{vasrhub}(Rss, \#u) : \text{rnd} : \text{sat}$$


### Syntax

$$Rd = \text{vasrhub}(Rss, \#u4) : \text{raw}$$

$$Rd = \text{vasrhub}(Rss, \#u4) : \text{rnd} : \text{sat}$$

$$Rd = \text{vasrhub}(Rss, \#u4) : \text{sat}$$

### Behavior

```
for (i=0; i<4; i++) {
    Rd.b[i] = usat8(((Rss.h[i] >> #u) + 1) >> 1);
}
```

```
if ("#u4==0") {
    Assembler mapped to: "Rd=vsathub(Rss)";
} else {
    Assembler mapped to: "Rd=vasrhub(Rss, #u4-1):raw";
}
```

```
for (i=0; i<4; i++) {
    Rd.b[i] = usat8(Rss.h[i] >> #u);
}
```

**Class: XTYPE (slots 2,3)**

### Notes

- If saturation occurs during execution of this instruction (a result is clamped to either maximum or minimum values), the OVF bit in the status register is set. OVF remains set until explicitly cleared by a transfer to the status register.

## Intrinsics

Rd=vasrhub(Rss,#u4):rnd:sat	Word32 Q6_R_vasrhub_PI_rnd_sat(Word64 Rss, Word32 Iu4)
Rd=vasrhub(Rss,#u4):sat	Word32 Q6_R_vasrhub_PI_sat(Word64 Rss, Word32 Iu4)

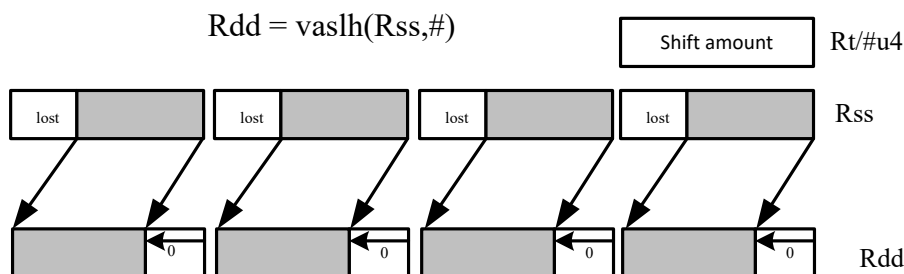
## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				MajOp		s5					Parse				MinOp				d5									
1	0	0	0	1	0	0	0	0	1	1	s	s	s	s	s	P	P	0	0	i	i	i	i	1	0	0	d	d	d	d	d	Rd=vasrhub(Rss,#u4):raw
1	0	0	0	1	0	0	0	0	1	1	s	s	s	s	s	P	P	0	0	i	i	i	i	1	0	1	d	d	d	d	d	Rd=vasrhub(Rss,#u4):sat

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Vector shift halfwords by register

The shift amount is the least significant seven bits of  $Rt$ , treated as a two's complement value. If the shift amount is negative, reverse the direction of the shift. Shift the source values right or left based on the shift amount and the type of instruction. Arithmetic right shifts place the sign bit of the source value in the vacated positions. Logical right shifts place zeros in the vacated positions.



### Syntax

### Behavior

$Rdd = \text{vaslh}(Rss, Rt)$	<pre>for (i=0; i&lt;4; i++) {     Rdd.h[i] = (sxt7-&gt;32(Rt) &gt; 0) ? (sxt16-&gt;64(Rss.h[i]) &lt;&lt; sxt7-&gt;32(Rt)) : (sxt16-&gt;64(Rss.h[i]) &gt;&gt; sxt7-&gt;32(Rt)); }</pre>
$Rdd = \text{vasrh}(Rss, Rt)$	<pre>for (i=0; i&lt;4; i++) {     Rdd.h[i] = (sxt7-&gt;32(Rt) &gt; 0) ? (sxt16-&gt;64(Rss.h[i]) &gt;&gt; sxt7-&gt;32(Rt)) : (sxt16-&gt;64(Rss.h[i]) &lt;&lt; sxt7-&gt;32(Rt)); }</pre>
$Rdd = \text{vslh}(Rss, Rt)$	<pre>for (i=0; i&lt;4; i++) {     Rdd.h[i] = (sxt7-&gt;32(Rt) &gt; 0) ? (zxt16-&gt;64(Rss.uh[i]) &lt;&lt; sxt7-&gt;32(Rt)) : (zxt16-&gt;64(Rss.uh[i]) &gt;&gt;&gt; sxt7-&gt;32(Rt)); }</pre>
$Rdd = \text{vlsrh}(Rss, Rt)$	<pre>for (i=0; i&lt;4; i++) {     Rdd.h[i] = (sxt7-&gt;32(Rt) &gt; 0) ? (zxt16-&gt;64(Rss.uh[i]) &gt;&gt;&gt; sxt7-&gt;32(Rt)) : (zxt16-&gt;64(Rss.uh[i]) &lt;&lt; sxt7-&gt;32(Rt)); }</pre>

### Class: XTYPE (slots 2,3)

### Notes

- If the number of bits to shift is greater than the width of the vector element, the result is either all sign-bits (for arithmetic right shifts) or all zeros for logical and left shifts.

## Intrinsics

Rdd=vaslh(Rss,Rt)	Word64 Q6_P_vaslh_PR(Word64 Rss, Word32 Rt)
Rdd=vasrh(Rss,Rt)	Word64 Q6_P_vasrh_PR(Word64 Rss, Word32 Rt)
Rdd=vlslh(Rss,Rt)	Word64 Q6_P_vlslh_PR(Word64 Rss, Word32 Rt)
Rdd=vlsrh(Rss,Rt)	Word64 Q6_P_vlsrh_PR(Word64 Rss, Word32 Rt)

## Encoding

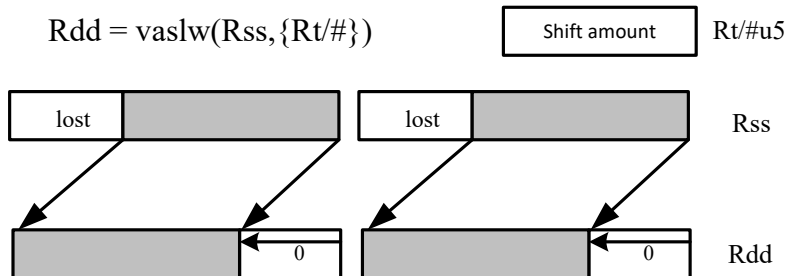
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				Maj		s5					Parse		t5					Min		d5									
1	1	0	0	0	0	1	1	0	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	d	d	d	d	d	Rdd=vasrh(Rss,Rt)
1	1	0	0	0	0	1	1	0	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	-	d	d	d	d	d	Rdd=vlsrh(Rss,Rt)
1	1	0	0	0	0	1	1	0	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	-	d	d	d	d	d	Rdd=vaslh(Rss,Rt)
1	1	0	0	0	0	1	1	0	1	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	-	d	d	d	d	d	Rdd=vlslh(Rss,Rt)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
Maj	Major opcode
Min	Minor opcode
RegType	Register type



## Vector shift words by immediate

Shift individual words of the source vector. Arithmetic right shifts place the sign bit of the source values in the vacated positions. Logical right shifts place zeros in the vacated positions.



Syntax	Behavior
$Rdd = \text{vaslw}(Rss, \#u5)$	<pre>for (i=0;i&lt;2;i++) {   Rdd.w[i] = (Rss.w[i] &lt;&lt; #u); }</pre>
$Rdd = \text{vasrw}(Rss, \#u5)$	<pre>for (i=0;i&lt;2;i++) {   Rdd.w[i] = (Rss.w[i] &gt;&gt; #u); }</pre>
$Rdd = \text{vlsrw}(Rss, \#u5)$	<pre>for (i=0;i&lt;2;i++) {   Rdd.w[i] = (Rss.uw[i] &gt;&gt; #u); }</pre>

### Class: XTYPE (slots 2,3)

#### Intrinsics

$Rdd = \text{vaslw}(Rss, \#u5)$	Word64 Q6_P_vaslw_PI(Word64 Rss, Word32 Iu5)
$Rdd = \text{vasrw}(Rss, \#u5)$	Word64 Q6_P_vasrw_PI(Word64 Rss, Word32 Iu5)
$Rdd = \text{vlsrw}(Rss, \#u5)$	Word64 Q6_P_vlsrw_PI(Word64 Rss, Word32 Iu5)

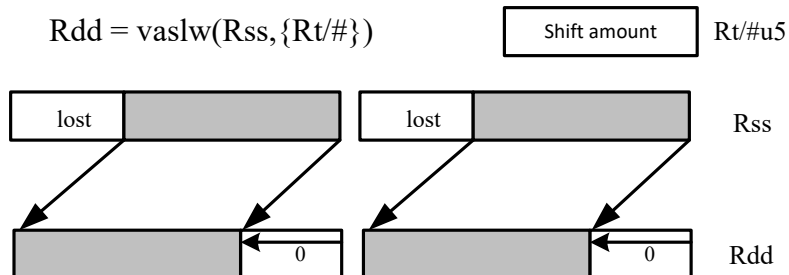
#### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS		RegType				MajOp				s5				Parse				MinOp				d5										
1	0	0	0	0	0	0	0	0	1	0	s	s	s	s	s	P	P	0	i	i	i	i	i	0	0	0	d	d	d	d	d	$Rdd = \text{vasrw}(Rss, \#u5)$
1	0	0	0	0	0	0	0	0	1	0	s	s	s	s	s	P	P	0	i	i	i	i	i	0	0	1	d	d	d	d	d	$Rdd = \text{vlsrw}(Rss, \#u5)$
1	0	0	0	0	0	0	0	0	1	0	s	s	s	s	s	P	P	0	i	i	i	i	i	0	1	0	d	d	d	d	d	$Rdd = \text{vaslw}(Rss, \#u5)$

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
MajOp	Major opcode
MinOp	Minor opcode
RegType	Register type

## Vector shift words by register

The shift amount is the least significant seven bits of  $Rt$ , treated as a two's complement value. If the shift amount is negative, reverse the direction of the shift. Shift the source values right or left based on the shift amount and the type of instruction. Arithmetic right shifts place the sign bit of the source value in the vacated positions. Logical right shifts place zeros in the vacated positions.



Syntax	Behavior
$Rdd = \text{vaslw}(Rss, Rt)$	<pre>for (i=0; i&lt;2; i++) {     Rdd.w[i] = (sxt7-&gt;32(Rt) &gt; 0) ? (sxt32-&gt;64(Rss.w[i]) &lt;&lt; sxt7-&gt;32(Rt)) : (sxt32-&gt;64(Rss.w[i]) &gt;&gt; sxt7-&gt;32(Rt)); }</pre>
$Rdd = \text{vasrw}(Rss, Rt)$	<pre>for (i=0; i&lt;2; i++) {     Rdd.w[i] = (sxt7-&gt;32(Rt) &gt; 0) ? (sxt32-&gt;64(Rss.w[i]) &gt;&gt; sxt7-&gt;32(Rt)) : (sxt32-&gt;64(Rss.w[i]) &lt;&lt; sxt7-&gt;32(Rt)); }</pre>
$Rdd = \text{vlslw}(Rss, Rt)$	<pre>for (i=0; i&lt;2; i++) {     Rdd.w[i] = (sxt7-&gt;32(Rt) &gt; 0) ? (zxt32-&gt;64(Rss.uw[i]) &lt;&lt; sxt7-&gt;32(Rt)) : (zxt32-&gt;64(Rss.uw[i]) &gt;&gt;&gt; sxt7-&gt;32(Rt)); }</pre>
$Rdd = \text{vlsrw}(Rss, Rt)$	<pre>for (i=0; i&lt;2; i++) {     Rdd.w[i] = (sxt7-&gt;32(Rt) &gt; 0) ? (zxt32-&gt;64(Rss.uw[i]) &gt;&gt;&gt; sxt7-&gt;32(Rt)) : (zxt32-&gt;64(Rss.uw[i]) &lt;&lt; sxt7-&gt;32(Rt)); }</pre>

### Class: XTYPE (slots 2,3)

#### Notes

- If the number of bits to shift is greater than the width of the vector element, the result is either all sign-bits (for arithmetic right shifts) or all zeros for logical and left shifts.

#### Intrinsics

$Rdd = \text{vaslw}(Rss, Rt)$	Word64 Q6_P_vaslw_PR(Word64 Rss, Word32 Rt)
$Rdd = \text{vasrw}(Rss, Rt)$	Word64 Q6_P_vasrw_PR(Word64 Rss, Word32 Rt)
$Rdd = \text{vlslw}(Rss, Rt)$	Word64 Q6_P_vlslw_PR(Word64 Rss, Word32 Rt)
$Rdd = \text{vlsrw}(Rss, Rt)$	Word64 Q6_P_vlsrw_PR(Word64 Rss, Word32 Rt)

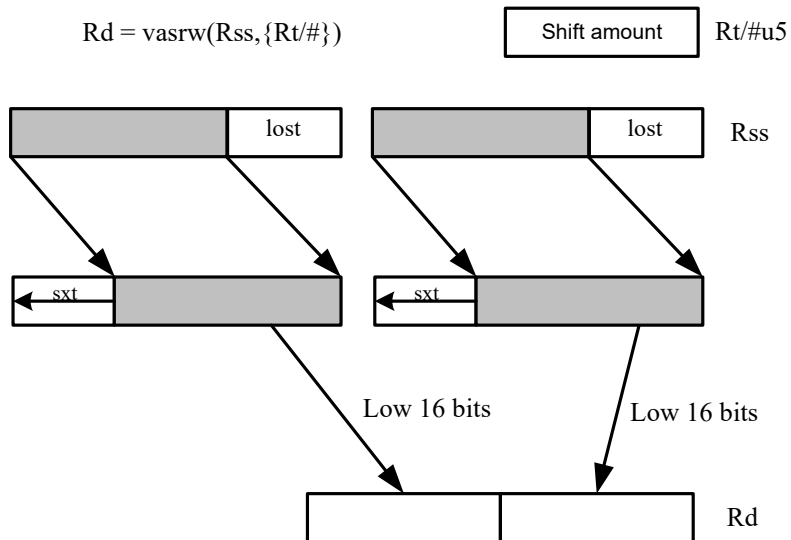
## Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS				RegType				Maj		s5					Parse		t5					Min		d5								
1	1	0	0	0	0	1	1	0	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	0	-	d	d	d	d	d	Rdd=vasrw(Rss,Rt)
1	1	0	0	0	0	1	1	0	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	-	d	d	d	d	d	Rdd=visrw(Rss,Rt)
1	1	0	0	0	0	1	1	0	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	0	-	d	d	d	d	d	Rdd=vaslw(Rss,Rt)
1	1	0	0	0	0	1	1	0	0	-	s	s	s	s	s	P	P	-	t	t	t	t	t	1	1	-	d	d	d	d	d	Rdd=vislw(Rss,Rt)

Field name	Description
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
Maj	Major opcode
Min	Minor opcode
RegType	Register type

## Vector shift words with truncate and pack

Shift individual words of the source vector Rss right by a register or immediate amount. The low 16-bits of each word are packed into destination register Rd.



### Syntax

`Rd=vasrw(Rss, #u5)`

```
for (i=0; i<2; i++) {
    Rd.h[i] = (Rss.w[i] >> #u) .h[0];
}
```

`Rd=vasrw(Rss, Rt)`

```
for (i=0; i<2; i++) {
    Rd.h[i] = (sxt7->32(Rt) > 0) ? (sxt32->64(Rss.w[i]) >> sxt7->
    >32(Rt)) : (sxt32->64(Rss.w[i]) << sxt7->32(Rt)) .h[0];
}
```

### Behavior

## Class: XTYPE (slots 2,3)

### Intrinsics

`Rd=vasrw(Rss, #u5)`

Word32 Q6\_R\_vasrw\_PI(Word64 Rss, Word32 Iu5)

`Rd=vasrw(Rss, Rt)`

Word32 Q6\_R\_vasrw\_PR(Word64 Rss, Word32 Rt)

### Encoding

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
ICLASS			RegType				MajOp		s5					Parse		MinOp					d5											
1	0	0	0	1	0	0	0	1	1	0	s	s	s	s	s	P	P	0	i	i	i	i	i	0	1	0	d	d	d	d	d	Rd=vasrw(Rss,#u5)
ICLASS			RegType				s5					Parse		t5					Min		d5											
1	1	0	0	0	1	0	1	-	-	-	s	s	s	s	s	P	P	-	t	t	t	t	t	0	1	0	d	d	d	d	d	Rd=vasrw(Rss,Rt)

<b>Field name</b>	<b>Description</b>
ICLASS	Instruction class
Parse	Packet/loop parse bits
d5	Field to encode register d
s5	Field to encode register s
t5	Field to encode register t
MajOp	Major opcode
MinOp	Minor opcode
Min	Minor opcode
RegType	Register type

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    Rdd=cmpy (Rs, Rt*) [ :<<1 ] :sat 433
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    Rxx+=cmpy (Rs, Rt*) [ :<<1 ] :sat 434
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    Rd=cmpyiw (Rss, Rtt) :<<1:sat 443
    Rd=cmpyiw (Rss, Rtt*) :<<1:rnd:sat 443
    Rd=cmpyiw (Rss, Rtt*) :<<1:sat 443
    Rdd=cmpyiw (Rss, Rtt) 444
    Rdd=cmpyiw (Rss, Rtt*) 444
    Rxx+=cmpyiw (Rss, Rtt) 444
    Rxx+=cmpyiw (Rss, Rtt*) 444

cmpyiw
    Rd=cmpyiw (Rss, Rtt) :<<1:rnd:sat 441
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    Rdd=cmpyrw (Rss, Rtt*) 444
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    Rxx+=cmpyrw (Rss, Rtt*) 444

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    Rd=cmpyrwh (Rss, Rt) :<<1:rnd:sat 441
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Rd=convert\_sf2w(Rs) 469  
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```

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## memb

if ([!]Pt [.new]) Rd=memb (#u6) 229  
 if ([!]Pt [.new]) Rd=memb (Rs+#u6:0) 229  
 if ([!]Pt [.new]) Rd=memb (Rx++#s4:0) 229  
 if ([!]Pv [.new]) memb (#u6)=Nt.new 278  
 if ([!]Pv [.new]) memb (#u6)=Rt 295  
 if ([!]Pv [.new]) memb (Rs+#u6:0)=#S6 295  
 if ([!]Pv [.new]) memb (Rs+#u6:0)=Nt.new 278  
 if ([!]Pv [.new]) memb (Rs+#u6:0)=Rt 295  
 if ([!]Pv [.new]) memb (Rs+Ru<<#u2)=Nt.new 278  
 if ([!]Pv [.new]) memb (Rs+Ru<<#u2)=Rt 295  
 if ([!]Pv [.new]) memb (Rx++#s4:0)=Nt.new 278  
 if ([!]Pv [.new]) memb (Rx++#s4:0)=Rt 295  
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 memb (gp+#u16:0)=Rt 293  
 memb (Re=#U6)=Nt.new 276  
 memb (Re=#U6)=Rt 293  
 memb (Rs+#s11:0)=Nt.new 276  
 memb (Rs+#s11:0)=Rt 293  
 memb (Rs+#u6:0) [+]=#U5 268  
 memb (Rs+#u6:0) [+|&]=Rt 268  
 memb (Rs+#u6:0)=#S8 293  
 memb (Rs+Ru<<#u2)=Nt.new 276  
 memb (Rs+Ru<<#u2)=Rt 293  
 memb (Ru<<#u2+#U6)=Nt.new 276  
 memb (Ru<<#u2+#U6)=Rt 293  
 memb (Rx++#s4:0:circ (Mu))=Nt.new 276  
 memb (Rx++#s4:0:circ (Mu))=Rt 293  
 memb (Rx++#s4:0)=Nt.new 276  
 memb (Rx++#s4:0)=Rt 293  
 memb (Rx++I:circ (Mu))=Nt.new 276  
 memb (Rx++I:circ (Mu))=Rt 293  
 memb (Rx++Mu:brev)=Nt.new 276  
 memb (Rx++Mu:brev)=Rt 293  
 memb (Rx++Mu)=Nt.new 276  
 memb (Rx++Mu)=Rt 293  
 Rd=memb (gp+#u16:0) 227  
 Rd=memb (Re=#U6) 227  
 Rd=memb (Rs+#s11:0) 227  
 Rd=memb (Rs+Rt<<#u2) 227  
 Rd=memb (Rt<<#u2+#U6) 227  
 Rd=memb (Rx++#s4:0:circ (Mu)) 227  
 Rd=memb (Rx++#s4:0) 227  
 Rd=memb (Rx++I:circ (Mu)) 227  
 Rd=memb (Rx++Mu:brev) 227  
 Rd=memb (Rx++Mu) 227

## memb\_fifo

Ryy=memb\_fifo (Re=#U6) 231  
 Ryy=memb\_fifo (Rs) 231  
 Ryy=memb\_fifo (Rs+#s11:0) 231  
 Ryy=memb\_fifo (Rt<<#u2+#U6) 231  
 Ryy=memb\_fifo (Rx++#s4:0:circ (Mu)) 231  
 Ryy=memb\_fifo (Rx++#s4:0) 231  
 Ryy=memb\_fifo (Rx++I:circ (Mu)) 232  
 Ryy=memb\_fifo (Rx++Mu:brev) 232  
 Ryy=memb\_fifo (Rx++Mu) 232



## membh

Rd=membh (Re=#U6) 260  
 Rd=membh (Rs) 260  
 Rd=membh (Rs+#s11:1) 260  
 Rd=membh (Rt<<#u2+#U6) 260  
 Rd=membh (Rx++#s4:1:circ (Mu)) 261  
 Rd=membh (Rx++#s4:1) 261  
 Rd=membh (Rx++I:circ (Mu)) 261  
 Rd=membh (Rx++Mu:brev) 261  
 Rd=membh (Rx++Mu) 261  
 Rdd=membh (Re=#U6) 263  
 Rdd=membh (Rs) 263  
 Rdd=membh (Rs+#s11:2) 263  
 Rdd=membh (Rt<<#u2+#U6) 263  
 Rdd=membh (Rx++#s4:2:circ (Mu)) 263  
 Rdd=membh (Rx++#s4:2) 263  
 Rdd=membh (Rx++I:circ (Mu)) 264  
 Rdd=membh (Rx++Mu:brev) 264  
 Rdd=membh (Rx++Mu) 264

## memd

if ([!]Pt [.new]) Rdd=memd (#u6) 225  
 if ([!]Pt [.new]) Rdd=memd (Rs+#u6:3) 225  
 if ([!]Pt [.new]) Rdd=memd (Rx++#s4:3) 225  
 if ([!]Pv [.new]) memd (#u6)=Rtt 291  
 if ([!]Pv [.new]) memd (Rs+#u6:3)=Rtt 291  
 if ([!]Pv [.new]) memd (Rs+Ru<<#u2)=Rtt 291  
 if ([!]Pv [.new]) memd (Rx++#s4:3)=Rtt 291  
 if ([!]Pv [.new]) Rdd=memd (Rs+Rt<<#u2) 225  
 memd (gp+#u16:3)=Rtt 288  
 memd (Re=#U6)=Rtt 288  
 memd (Rs+#s11:3)=Rtt 288  
 memd (Rs+Ru<<#u2)=Rtt 288  
 memd (Ru<<#u2+#U6)=Rtt 288  
 memd (Rx++#s4:3:circ (Mu))=Rtt 288  
 memd (Rx++#s4:3)=Rtt 288  
 memd (Rx++I:circ (Mu))=Rtt 288  
 memd (Rx++Mu:brev)=Rtt 288  
 memd (Rx++Mu)=Rtt 288  
 Rdd=memd (gp+#u16:3) 222  
 Rdd=memd (Re=#U6) 222  
 Rdd=memd (Rs+#s11:3) 222  
 Rdd=memd (Rs+Rt<<#u2) 222  
 Rdd=memd (Rt<<#u2+#U6) 222  
 Rdd=memd (Rx++#s4:3:circ (Mu)) 222  
 Rdd=memd (Rx++#s4:3) 222  
 Rdd=memd (Rx++I:circ (Mu)) 222  
 Rdd=memd (Rx++Mu:brev) 222  
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if ([!]Pt[.new]) Rd=memh(Rs+#u6:1) 239
if ([!]Pt[.new]) Rd=memh(Rx++#s4:1) 239
if ([!]Pv[.new]) memh(#u6)=Nt.new 282
if ([!]Pv[.new]) memh(#u6)=Rt 301
if ([!]Pv[.new]) memh(#u6)=Rt.H 301
if ([!]Pv[.new]) memh(Rs+#u6:1)=#S6 301
if ([!]Pv[.new]) memh(Rs+#u6:1)=Nt.new 282
if ([!]Pv[.new]) memh(Rs+#u6:1)=Rt 301
if ([!]Pv[.new]) memh(Rs+#u6:1)=Rt.H 301
if ([!]Pv[.new]) memh(Rs+Ru<<#u2)=Nt.new 282
if ([!]Pv[.new]) memh(Rs+Ru<<#u2)=Rt 302
if ([!]Pv[.new]) memh(Rs+Ru<<#u2)=Rt.H 301
if ([!]Pv[.new]) memh(Rx++#s4:1)=Nt.new 282
if ([!]Pv[.new]) memh(Rx++#s4:1)=Rt 302
if ([!]Pv[.new]) memh(Rx++#s4:1)=Rt.H 302
if ([!]Pv[.new]) Rd=memh(Rs+Rt<<#u2) 239
memh(gp+#u16:1)=Nt.new 280
memh(gp+#u16:1)=Rt 299
memh(gp+#u16:1)=Rt.H 299
memh(Re=#U6)=Nt.new 280
memh(Re=#U6)=Rt 298
memh(Re=#U6)=Rt.H 298
memh(Rs+#s11:1)=Nt.new 280
memh(Rs+#s11:1)=Rt 298
memh(Rs+#s11:1)=Rt.H 298
memh(Rs+#u6:1)[+-]=#U5 270
memh(Rs+#u6:1)[+ - | &]=Rt 270
memh(Rs+#u6:1)=#S8 298
memh(Rs+Ru<<#u2)=Nt.new 280
memh(Rs+Ru<<#u2)=Rt 298
memh(Rs+Ru<<#u2)=Rt.H 298
memh(Ru<<#u2+#U6)=Nt.new 280
memh(Ru<<#u2+#U6)=Rt 298
memh(Ru<<#u2+#U6)=Rt.H 298
memh(Rx++#s4:1:circ(Mu))=Nt.new 280
memh(Rx++#s4:1:circ(Mu))=Rt 298
memh(Rx++#s4:1:circ(Mu))=Rt.H 298
memh(Rx++#s4:1)=Nt.new 280
memh(Rx++#s4:1)=Rt 298
memh(Rx++#s4:1)=Rt.H 298
memh(Rx++I:circ(Mu))=Nt.new 280
memh(Rx++I:circ(Mu))=Rt 298
memh(Rx++I:circ(Mu))=Rt.H 298
memh(Rx++Mu:brev)=Nt.new 280
memh(Rx++Mu:brev)=Rt 299
memh(Rx++Mu:brev)=Rt.H 299
memh(Rx++Mu)=Nt.new 280
memh(Rx++Mu)=Rt 299
memh(Rx++Mu)=Rt.H 299
Rd=memh(gp+#u16:1) 237
Rd=memh(Re=#U6) 237
Rd=memh(Rs+#s11:1) 237
Rd=memh(Rs+Rt<<#u2) 237
Rd=memh(Rt<<#u2+#U6) 237
Rd=memh(Rx++#s4:1:circ(Mu)) 237
Rd=memh(Rx++#s4:1) 237
Rd=memh(Rx++I:circ(Mu)) 237
Rd=memh(Rx++Mu:brev) 237
Rd=memh(Rx++Mu) 237

```

## memh\_fifo

Ryy=memh\_fifo (Re=#U6) 234  
 Ryy=memh\_fifo (Rs) 234  
 Ryy=memh\_fifo (Rs+#s11:1) 234  
 Ryy=memh\_fifo (Rt<<#u2+#U6) 234  
 Ryy=memh\_fifo (Rx++#s4:1:circ (Mu)) 234  
 Ryy=memh\_fifo (Rx++#s4:1) 234  
 Ryy=memh\_fifo (Rx++I:circ (Mu)) 235  
 Ryy=memh\_fifo (Rx++Mu:brev) 235  
 Ryy=memh\_fifo (Rx++Mu) 235

## memub

if ([!]Pt [.new]) Rd=memub (#u6) 245  
 if ([!]Pt [.new]) Rd=memub (Rs+#u6:0) 245  
 if ([!]Pt [.new]) Rd=memub (Rx++#s4:0) 245  
 if ([!]Pv [.new]) Rd=memub (Rs+Rt<<#u2) 245  
 Rd=memub (gp+#u16:0) 243  
 Rd=memub (Re=#U6) 243  
 Rd=memub (Rs+#s11:0) 243  
 Rd=memub (Rs+Rt<<#u2) 243  
 Rd=memub (Rt<<#u2+#U6) 243  
 Rd=memub (Rx++#s4:0:circ (Mu)) 243  
 Rd=memub (Rx++#s4:0) 243  
 Rd=memub (Rx++I:circ (Mu)) 243  
 Rd=memub (Rx++Mu:brev) 243  
 Rd=memub (Rx++Mu) 243

## memubh

Rd=memubh (Re=#U6) 261  
 Rd=memubh (Rs+#s11:1) 262  
 Rd=memubh (Rt<<#u2+#U6) 262  
 Rd=memubh (Rx++#s4:1:circ (Mu)) 262  
 Rd=memubh (Rx++#s4:1) 262  
 Rd=memubh (Rx++I:circ (Mu)) 262  
 Rd=memubh (Rx++Mu:brev) 263  
 Rd=memubh (Rx++Mu) 262  
 Rdd=memubh (Re=#U6) 264  
 Rdd=memubh (Rs+#s11:2) 264  
 Rdd=memubh (Rt<<#u2+#U6) 264  
 Rdd=memubh (Rx++#s4:2:circ (Mu)) 265  
 Rdd=memubh (Rx++#s4:2) 265  
 Rdd=memubh (Rx++I:circ (Mu)) 265  
 Rdd=memubh (Rx++Mu:brev) 265  
 Rdd=memubh (Rx++Mu) 265

## memuh

if ([!]Pt [.new]) Rd=memuh (#u6) 249  
 if ([!]Pt [.new]) Rd=memuh (Rs+#u6:1) 249  
 if ([!]Pt [.new]) Rd=memuh (Rx++#s4:1) 249  
 if ([!]Pv [.new]) Rd=memuh (Rs+Rt<<#u2) 249  
 Rd=memuh (gp+#u16:1) 247  
 Rd=memuh (Re=#U6) 247  
 Rd=memuh (Rs+#s11:1) 247  
 Rd=memuh (Rs+Rt<<#u2) 247  
 Rd=memuh (Rt<<#u2+#U6) 247  
 Rd=memuh (Rx++#s4:1:circ (Mu)) 247  
 Rd=memuh (Rx++#s4:1) 247  
 Rd=memuh (Rx++I:circ (Mu)) 247  
 Rd=memuh (Rx++Mu:brev) 247  
 Rd=memuh (Rx++Mu) 247

```

memw
  if ([!]Pt[.new]) Rd=memw(#u6) 254
  if ([!]Pt[.new]) Rd=memw(Rs+#u6:2) 254
  if ([!]Pt[.new]) Rd=memw(Rx++#s4:2) 254
  if ([!]Pv[.new]) memw(#u6)=Nt.new 286
  if ([!]Pv[.new]) memw(#u6)=Rt 309
  if ([!]Pv[.new]) memw(Rs+#u6:2)=#S6 309
  if ([!]Pv[.new]) memw(Rs+#u6:2)=Nt.new 286
  if ([!]Pv[.new]) memw(Rs+#u6:2)=Rt 309
  if ([!]Pv[.new]) memw(Rs+Ru<<#u2)=Nt.new 286
  if ([!]Pv[.new]) memw(Rs+Ru<<#u2)=Rt 309
  if ([!]Pv[.new]) memw(Rx++#s4:2)=Nt.new 286
  if ([!]Pv[.new]) memw(Rx++#s4:2)=Rt 309
  if ([!]Pv[.new]) Rd=memw(Rs+Rt<<#u2) 254
  memw(gp+#u16:2)=Nt.new 284
  memw(gp+#u16:2)=Rt 306
  memw(Re=#U6)=Nt.new 284
  memw(Re=#U6)=Rt 306
  memw(Rs+#s11:2)=Nt.new 284
  memw(Rs+#s11:2)=Rt 306
  memw(Rs+#u6:2)[+-]=#U5 271
  memw(Rs+#u6:2)[+&]=Rt 271
  memw(Rs+#u6:2)=#S8 306
  memw(Rs+Ru<<#u2)=Nt.new 284
  memw(Rs+Ru<<#u2)=Rt 306
  memw(Ru<<#u2+#U6)=Nt.new 284
  memw(Ru<<#u2+#U6)=Rt 306
  memw(Rx++#s4:2:circ(Mu))=Nt.new 284
  memw(Rx++#s4:2:circ(Mu))=Rt 306
  memw(Rx++#s4:2)=Nt.new 284
  memw(Rx++#s4:2)=Rt 306
  memw(Rx++I:circ(Mu))=Nt.new 284
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  Pd=cmpb.gtu(Rs,Rt)
    Byte Q6_p_cmpb_gtu_RR(Word32 Rs, Word32 Rt) 569

cmph.eq
  Pd=cmph.eq(Rs,#s8)
    Byte Q6_p_cmph_eq_RI(Word32 Rs, Word32 Is8) 571
  Pd=cmph.eq(Rs,Rt)
    Byte Q6_p_cmph_eq_RR(Word32 Rs, Word32 Rt) 571

cmph.gt
  Pd=cmph.gt(Rs,#s8)
    Byte Q6_p_cmph_gt_RI(Word32 Rs, Word32 Is8) 571
  Pd=cmph.gt(Rs,Rt)
    Byte Q6_p_cmph_gt_RR(Word32 Rs, Word32 Rt) 571

cmph.gtu
  Pd=cmph.gtu(Rs,#u7)
    Byte Q6_p_cmph_gtu_RI(Word32 Rs, Word32 Iu7) 571
  Pd=cmph.gtu(Rs,Rt)
    Byte Q6_p_cmph_gtu_RR(Word32 Rs, Word32 Rt) 571

cmpy
  Rd=cmpy(Rs,Rt):<<1:rnd:sat

```



```

    Word32 Q6_R_cmpy_RR_sl_rnd_sat(Word32 Rs, Word32 Rt) 440
Rd=cmpy(Rs,Rt):rnd:sat
    Word32 Q6_R_cmpy_RR_rnd_sat(Word32 Rs, Word32 Rt) 440
Rd=cmpy(Rs,Rt*):<<1:rnd:sat
    Word32 Q6_R_cmpy_RR_conj_sl_rnd_sat(Word32 Rs, Word32 Rt) 440
Rd=cmpy(Rs,Rt*):rnd:sat
    Word32 Q6_R_cmpy_RR_conj_rnd_sat(Word32 Rs, Word32 Rt) 440
Rdd=cmpy(Rs,Rt):<<1:sat
    Word64 Q6_P_cmpy_RR_sl_sat(Word32 Rs, Word32 Rt) 435
Rdd=cmpy(Rs,Rt):sat
    Word64 Q6_P_cmpy_RR_sat(Word32 Rs, Word32 Rt) 435
Rdd=cmpy(Rs,Rt*):<<1:sat
    Word64 Q6_P_cmpy_RR_conj_sl_sat(Word32 Rs, Word32 Rt) 435
Rdd=cmpy(Rs,Rt*):sat
    Word64 Q6_P_cmpy_RR_conj_sat(Word32 Rs, Word32 Rt) 435
Rxx+=cmpy(Rs,Rt):<<1:sat
    Word64 Q6_P_cmpyacc_RR_sl_sat(Word64 Rxx, Word32 Rs, Word32 Rt) 435
Rxx+=cmpy(Rs,Rt):sat
    Word64 Q6_P_cmpyacc_RR_sat(Word64 Rxx, Word32 Rs, Word32 Rt) 435
Rxx+=cmpy(Rs,Rt*):<<1:sat
    Word64 Q6_P_cmpyacc_RR_conj_sl_sat(Word64 Rxx, Word32 Rs, Word32 Rt) 435
Rxx+=cmpy(Rs,Rt*):sat
    Word64 Q6_P_cmpyacc_RR_conj_sat(Word64 Rxx, Word32 Rs, Word32 Rt) 435
Rxx-=cmpy(Rs,Rt):<<1:sat
    Word64 Q6_P_cmpynac_RR_sl_sat(Word64 Rxx, Word32 Rs, Word32 Rt) 435
Rxx-=cmpy(Rs,Rt):sat
    Word64 Q6_P_cmpynac_RR_sat(Word64 Rxx, Word32 Rs, Word32 Rt) 435
Rxx-=cmpy(Rs,Rt*):<<1:sat
    Word64 Q6_P_cmpynac_RR_conj_sl_sat(Word64 Rxx, Word32 Rs, Word32 Rt) 435
Rxx-=cmpy(Rs,Rt*):sat
    Word64 Q6_P_cmpynac_RR_conj_sat(Word64 Rxx, Word32 Rs, Word32 Rt) 435

cmpyi
    Rdd=cmpyi(Rs,Rt)
    Word64 Q6_P_cmpyi_RR(Word32 Rs, Word32 Rt) 438
    Rxx+=cmpyi(Rs,Rt)
    Word64 Q6_P_cmpyiacc_RR(Word64 Rxx, Word32 Rs, Word32 Rt) 438

cmpyiw
    Rd=cmpyiw(Rss,Rtt):<<1:rnd:sat
    Word32 Q6_R_cmpyiw_PP_sl_rnd_sat(Word64 Rss, Word64 Rtt) 445
    Rd=cmpyiw(Rss,Rtt):<<1:sat
    Word32 Q6_R_cmpyiw_PP_sl_sat(Word64 Rss, Word64 Rtt) 445
    Rd=cmpyiw(Rss,Rtt*):<<1:rnd:sat
    Word32 Q6_R_cmpyiw_PP_conj_sl_rnd_sat(Word64 Rss, Word64 Rtt) 445
    Rd=cmpyiw(Rss,Rtt*):<<1:sat
    Word32 Q6_R_cmpyiw_PP_conj_sl_sat(Word64 Rss, Word64 Rtt) 445
    Rdd=cmpyiw(Rss,Rtt)
    Word64 Q6_P_cmpyiw_PP(Word64 Rss, Word64 Rtt) 445
    Rdd=cmpyiw(Rss,Rtt*)
    Word64 Q6_P_cmpyiw_PP_conj(Word64 Rss, Word64 Rtt) 445
    Rxx+=cmpyiw(Rss,Rtt)
    Word64 Q6_P_cmpyiwacc_PP(Word64 Rxx, Word64 Rss, Word64 Rtt) 445
    Rxx+=cmpyiw(Rss,Rtt*)
    Word64 Q6_P_cmpyiwacc_PP_conj(Word64 Rxx, Word64 Rss, Word64 Rtt) 445

cmpyihw
    Rd=cmpyihw(Rss,Rt):<<1:rnd:sat

```

```

    Word32 Q6_R_cmpyiwH_PR_sl_rnd_sat(Word64 Rss, Word32 Rt) 442
    Rd=cmpyiwH(Rss,Rt*):<<1:rnd:sat
    Word32 Q6_R_cmpyiwH_PR_conj_sl_rnd_sat(Word64 Rss, Word32 Rt) 442

cmpyr
    Rdd=cmpyr(Rs,Rt)
    Word64 Q6_P_cmpyr_RR(Word32 Rs, Word32 Rt) 438
    Rxx+=cmpyr(Rs,Rt)
    Word64 Q6_P_cmpyracc_RR(Word64 Rxx, Word32 Rs, Word32 Rt) 438

cmpyrw
    Rd=cmpyrw(Rss,Rtt):<<1:rnd:sat
    Word32 Q6_R_cmpyrw_PP_sl_rnd_sat(Word64 Rss, Word64 Rtt) 445
    Rd=cmpyrw(Rss,Rtt):<<1:sat
    Word32 Q6_R_cmpyrw_PP_sl_sat(Word64 Rss, Word64 Rtt) 445
    Rd=cmpyrw(Rss,Rtt*):<<1:rnd:sat
    Word32 Q6_R_cmpyrw_PP_conj_sl_rnd_sat(Word64 Rss, Word64 Rtt) 445
    Rd=cmpyrw(Rss,Rtt*):<<1:sat
    Word32 Q6_R_cmpyrw_PP_conj_sl_sat(Word64 Rss, Word64 Rtt) 445
    Rdd=cmpyrw(Rss,Rtt)
    Word64 Q6_P_cmpyrw_PP(Word64 Rss, Word64 Rtt) 445
    Rdd=cmpyrw(Rss,Rtt*)
    Word64 Q6_P_cmpyrw_PP_conj(Word64 Rss, Word64 Rtt) 445
    Rxx+=cmpyrw(Rss,Rtt)
    Word64 Q6_P_cmpyrwacc_PP(Word64 Rxx, Word64 Rss, Word64 Rtt) 445
    Rxx+=cmpyrw(Rss,Rtt*)
    Word64 Q6_P_cmpyrwacc_PP_conj(Word64 Rxx, Word64 Rss, Word64 Rtt) 445

cmpyrwh
    Rd=cmpyrwh(Rss,Rt):<<1:rnd:sat
    Word32 Q6_R_cmpyrwh_PR_sl_rnd_sat(Word64 Rss, Word32 Rt) 442
    Rd=cmpyrwh(Rss,Rt*):<<1:rnd:sat
    Word32 Q6_R_cmpyrwh_PR_conj_sl_rnd_sat(Word64 Rss, Word32 Rt) 442

combine
    Rd=combine(Rt.H,Rs.H)
    Word32 Q6_R_combine_RhRh(Word32 Rt, Word32 Rs) 173
    Rd=combine(Rt.H,Rs.L)
    Word32 Q6_R_combine_RhRl(Word32 Rt, Word32 Rs) 173
    Rd=combine(Rt.L,Rs.H)
    Word32 Q6_R_combine_RlRh(Word32 Rt, Word32 Rs) 173
    Rd=combine(Rt.L,Rs.L)
    Word32 Q6_R_combine_RlRl(Word32 Rt, Word32 Rs) 173
    Rdd=combine(#s8,#s8)
    Word64 Q6_P_combine_II(Word32 Is8, Word32 IS8) 173
    Rdd=combine(#s8,Rs)
    Word64 Q6_P_combine_IR(Word32 Is8, Word32 Rs) 173
    Rdd=combine(Rs,#s8)
    Word64 Q6_P_combine_RI(Word32 Rs, Word32 Is8) 173
    Rdd=combine(Rs,Rt)
    Word64 Q6_P_combine_RR(Word32 Rs, Word32 Rt) 173

convert_d2df
    Rdd=convert_d2df(Rss)
    Word64 Q6_P_convert_d2df_P(Word64 Rss) 467

convert_d2sf
    Rd=convert_d2sf(Rss)
    Word32 Q6_R_convert_d2sf_P(Word64 Rss) 467

```

```
convert_df2d
  Rdd=convert_df2d(Rss)
    Word64 Q6_P_convert_df2d_P(Word64 Rss) 470
  Rdd=convert_df2d(Rss):chop
    Word64 Q6_P_convert_df2d_P_chop(Word64 Rss) 470

convert_df2sf
  Rd=convert_df2sf(Rss)
    Word32 Q6_R_convert_df2sf_P(Word64 Rss) 466

convert_df2ud
  Rdd=convert_df2ud(Rss)
    Word64 Q6_P_convert_df2ud_P(Word64 Rss) 470
  Rdd=convert_df2ud(Rss):chop
    Word64 Q6_P_convert_df2ud_P_chop(Word64 Rss) 470

convert_df2uw
  Rd=convert_df2uw(Rss)
    Word32 Q6_R_convert_df2uw_P(Word64 Rss) 470
  Rd=convert_df2uw(Rss):chop
    Word32 Q6_R_convert_df2uw_P_chop(Word64 Rss) 470

convert_df2w
  Rd=convert_df2w(Rss)
    Word32 Q6_R_convert_df2w_P(Word64 Rss) 470
  Rd=convert_df2w(Rss):chop
    Word32 Q6_R_convert_df2w_P_chop(Word64 Rss) 470

convert_sf2d
  Rdd=convert_sf2d(Rs)
    Word64 Q6_P_convert_sf2d_R(Word32 Rs) 470
  Rdd=convert_sf2d(Rs):chop
    Word64 Q6_P_convert_sf2d_R_chop(Word32 Rs) 470

convert_sf2df
  Rdd=convert_sf2df(Rs)
    Word64 Q6_P_convert_sf2df_R(Word32 Rs) 466

convert_sf2ud
  Rdd=convert_sf2ud(Rs)
    Word64 Q6_P_convert_sf2ud_R(Word32 Rs) 470
  Rdd=convert_sf2ud(Rs):chop
    Word64 Q6_P_convert_sf2ud_R_chop(Word32 Rs) 470

convert_sf2uw
  Rd=convert_sf2uw(Rs)
    Word32 Q6_R_convert_sf2uw_R(Word32 Rs) 470
  Rd=convert_sf2uw(Rs):chop
    Word32 Q6_R_convert_sf2uw_R_chop(Word32 Rs) 470

convert_sf2w
  Rd=convert_sf2w(Rs)
    Word32 Q6_R_convert_sf2w_R(Word32 Rs) 470
  Rd=convert_sf2w(Rs):chop
    Word32 Q6_R_convert_sf2w_R_chop(Word32 Rs) 470

convert_ud2df
  Rdd=convert_ud2df(Rss)
    Word64 Q6_P_convert_ud2df_P(Word64 Rss) 467
```

```

convert_ud2sf
  Rd=convert_ud2sf(Rss)
  Word32 Q6_R_convert_ud2sf_P(Word64 Rss) 467

convert_uw2df
  Rdd=convert_uw2df(Rs)
  Word64 Q6_P_convert_uw2df_R(Word32 Rs) 467

convert_uw2sf
  Rd=convert_uw2sf(Rs)
  Word32 Q6_R_convert_uw2sf_R(Word32 Rs) 467

convert_w2df
  Rdd=convert_w2df(Rs)
  Word64 Q6_P_convert_w2df_R(Word32 Rs) 467

convert_w2sf
  Rd=convert_w2sf(Rs)
  Word32 Q6_R_convert_w2sf_R(Word32 Rs) 467

cround
  Rd=cround(Rs, #u5)
  Word32 Q6_R_cround_RI(Word32 Rs, Word32 Iu5) 357
  Rd=cround(Rs, Rt)
  Word32 Q6_R_cround_RR(Word32 Rs, Word32 Rt) 357
  Rdd=cround(Rss, #u6)
  Word64 Q6_P_cround_PI(Word64 Rss, Word32 Iu6) 357
  Rdd=cround(Rss, Rt)
  Word64 Q6_P_cround_PR(Word64 Rss, Word32 Rt) 357

ct0
  Rd=ct0(Rs)
  Word32 Q6_R_ct0_R(Word32 Rs) 412
  Rd=ct0(Rss)
  Word32 Q6_R_ct0_P(Word64 Rss) 412

ct1
  Rd=ct1(Rs)
  Word32 Q6_R_ct1_R(Word32 Rs) 412
  Rd=ct1(Rss)
  Word32 Q6_R_ct1_P(Word64 Rss) 412

```

**D**

```

dccleana
  dccleana(Rs)
  void Q6_dccleana_A(Address a) 320

dccleaninva
  dccleaninva(Rs)
  void Q6_dccleaninva_A(Address a) 320

dcfetch
  dcfetch(Rs)
  void Q6_dcfetch_A(Address a) 319

dcinva
  dcinva(Rs)
  void Q6_dcinva_A(Address a) 320

```

```
dczeroa
  dczeroa (Rs)
  void Q6_dczeroa_A(Address a) 316

deinterleave
  Rdd=deinterleave (Rss)
  Word64 Q6_P_deinterleave_P(Word64 Rss) 418

dfadd
  Rdd=dfadd (Rss,Rtt)
  Word64 Q6_P_dfadd_PP(Word64 Rss, Word64 Rtt) 461

dfclass
  Pd=dfclass (Rss,#u5)
  Byte Q6_p_dfclass_PI(Word64 Rss, Word32 Iu5) 462

dfcmp.eq
  Pd=dfcmp.eq (Rss,Rtt)
  Byte Q6_p_dfcmp_eq_PP(Word64 Rss, Word64 Rtt) 464

dfcmp.ge
  Pd=dfcmp.ge (Rss,Rtt)
  Byte Q6_p_dfcmp_ge_PP(Word64 Rss, Word64 Rtt) 464

dfcmp.gt
  Pd=dfcmp.gt (Rss,Rtt)
  Byte Q6_p_dfcmp_gt_PP(Word64 Rss, Word64 Rtt) 464

dfcmp.uo
  Pd=dfcmp.uo (Rss,Rtt)
  Byte Q6_p_dfcmp_uo_PP(Word64 Rss, Word64 Rtt) 464

dfmake
  Rdd=dfmake(#u10):neg
  Word64 Q6_P_dfmake_I_neg(Word32 Iu10) 478
  Rdd=dfmake(#u10):pos
  Word64 Q6_P_dfmake_I_pos(Word32 Iu10) 478

dfmax
  Rdd=dfmax (Rss,Rtt)
  Word64 Q6_P_dfmax_PP(Word64 Rss, Word64 Rtt) 479

dfmin
  Rdd=dfmin (Rss,Rtt)
  Word64 Q6_P_dfmin_PP(Word64 Rss, Word64 Rtt) 480

dfmpyfix
  Rdd=dfmpyfix (Rss,Rtt)
  Word64 Q6_P_dfmpyfix_PP(Word64 Rss, Word64 Rtt) 481

dfmpyhh
  Rxx+=dfmpyhh (Rss,Rtt)
  Word64 Q6_P_dfmpyhhacc_PP(Word64 Rxx, Word64 Rss, Word64 Rtt) 473

dfmpylh
  Rxx+=dfmpylh (Rss,Rtt)
  Word64 Q6_P_dfmpylhacc_PP(Word64 Rxx, Word64 Rss, Word64 Rtt) 473

dfmpyll
  Rdd=dfmpyll (Rss,Rtt)
```

Word64 Q6\_P\_dfmpyl1\_PP(Word64 Rss, Word64 Rtt) **481**

dfsub

Rdd=dfsub(Rss,Rtt)

Word64 Q6\_P\_dfsub\_PP(Word64 Rss, Word64 Rtt) **483**

dmsyncht

Rd=dmsyncht

Word32 Q6\_R\_dmsyncht() **329**

## E

extract

Rd=extract(Rs,#u5,#U5)

Word32 Q6\_R\_extract\_RII(Word32 Rs, Word32 Iu5, Word32 IU5) **414**

Rd=extract(Rs,Rtt)

Word32 Q6\_R\_extract\_RP(Word32 Rs, Word64 Rtt) **414**

Rdd=extract(Rss,#u6,#U6)

Word64 Q6\_P\_extract\_PII(Word64 Rss, Word32 Iu6, Word32 IU6) **414**

Rdd=extract(Rss,Rtt)

Word64 Q6\_P\_extract\_PP(Word64 Rss, Word64 Rtt) **414**

extractu

Rd=extractu(Rs,#u5,#U5)

Word32 Q6\_R\_extractu\_RII(Word32 Rs, Word32 Iu5, Word32 IU5) **414**

Rd=extractu(Rs,Rtt)

Word32 Q6\_R\_extractu\_RP(Word32 Rs, Word64 Rtt) **414**

Rdd=extractu(Rss,#u6,#U6)

Word64 Q6\_P\_extractu\_PII(Word64 Rss, Word32 Iu6, Word32 IU6) **414**

Rdd=extractu(Rss,Rtt)

Word64 Q6\_P\_extractu\_PP(Word64 Rss, Word64 Rtt) **414**

## F

fastcorner9

Pd=!fastcorner9(Ps,Pt)

Byte Q6\_p\_not\_fastcorner9\_pp(Byte Ps, Byte Pt) **196**

Pd=fastcorner9(Ps,Pt)

Byte Q6\_p\_fastcorner9\_pp(Byte Ps, Byte Pt) **196**

## I

insert

Rx=insert(Rs,#u5,#U5)

Word32 Q6\_R\_insert\_RII(Word32 Rx, Word32 Rs, Word32 Iu5, Word32 IU5) **417**

Rx=insert(Rs,Rtt)

Word32 Q6\_R\_insert\_RP(Word32 Rx, Word32 Rs, Word64 Rtt) **417**

Rxx=insert(Rss,#u6,#U6)

Word64 Q6\_P\_insert\_PII(Word64 Rxx, Word64 Rss, Word32 Iu6, Word32 IU6) **417**

Rxx=insert(Rss,Rtt)

Word64 Q6\_P\_insert\_PP(Word64 Rxx, Word64 Rss, Word64 Rtt) **417**

interleave

Rdd=interleave(Rss)

Word64 Q6\_P\_interleave\_P(Word64 Rss) **418**

## L

l2fetch

l2fetch(Rs,Rt)

```

    void Q6_l2fetch_AR(Address a, Word32 Rt) 326
    l2fetch(Rs,Rtt)
    void Q6_l2fetch_AP(Address a, Word64 Rtt) 326

lfs
    Rdd=lfs(Rss,Rtt)
    Word64 Q6_P_lfs_PP(Word64 Rss, Word64 Rtt) 419

lsl
    Rd=lsl(#s6,Rt)
    Word32 Q6_R_lsl_IR(Word32 Is6, Word32 Rt) 602
    Rd=lsl(Rs,Rt)
    Word32 Q6_R_lsl_RR(Word32 Rs, Word32 Rt) 602
    Rdd=lsl(Rss,Rt)
    Word64 Q6_P_lsl_PR(Word64 Rss, Word32 Rt) 602
    Rx&=lsl(Rs,Rt)
    Word32 Q6_R_lsland_RR(Word32 Rx, Word32 Rs, Word32 Rt) 608
    Rx+=lsl(Rs,Rt)
    Word32 Q6_R_lslacc_RR(Word32 Rx, Word32 Rs, Word32 Rt) 605
    Rx-=lsl(Rs,Rt)
    Word32 Q6_R_lslnac_RR(Word32 Rx, Word32 Rs, Word32 Rt) 605
    Rx|=lsl(Rs,Rt)
    Word32 Q6_R_lslor_RR(Word32 Rx, Word32 Rs, Word32 Rt) 608
    Rxx&=lsl(Rss,Rt)
    Word64 Q6_P_lsland_PR(Word64 Rxx, Word64 Rss, Word32 Rt) 608
    Rxx^=lsl(Rss,Rt)
    Word64 Q6_P_lslxacc_PR(Word64 Rxx, Word64 Rss, Word32 Rt) 608
    Rxx+=lsl(Rss,Rt)
    Word64 Q6_P_lslacc_PR(Word64 Rxx, Word64 Rss, Word32 Rt) 605
    Rxx-=lsl(Rss,Rt)
    Word64 Q6_P_lslnac_PR(Word64 Rxx, Word64 Rss, Word32 Rt) 605
    Rxx|=lsl(Rss,Rt)
    Word64 Q6_P_lslor_PR(Word64 Rxx, Word64 Rss, Word32 Rt) 609

lsr
    Rd=lsr(Rs,#u5)
    Word32 Q6_R_lsr_RI(Word32 Rs, Word32 Iu5) 589
    Rd=lsr(Rs,Rt)
    Word32 Q6_R_lsr_RR(Word32 Rs, Word32 Rt) 602
    Rdd=lsr(Rss,#u6)
    Word64 Q6_P_lsr_PI(Word64 Rss, Word32 Iu6) 590
    Rdd=lsr(Rss,Rt)
    Word64 Q6_P_lsr_PR(Word64 Rss, Word32 Rt) 602
    Rx&=lsr(Rs,#u5)
    Word32 Q6_R_lsrand_RI(Word32 Rx, Word32 Rs, Word32 Iu5) 596
    Rx&=lsr(Rs,Rt)
    Word32 Q6_R_lsrand_RR(Word32 Rx, Word32 Rs, Word32 Rt) 608
    Rx^=lsr(Rs,#u5)
    Word32 Q6_R_lsrxacc_RI(Word32 Rx, Word32 Rs, Word32 Iu5) 596
    Rx+=lsr(Rs,#u5)
    Word32 Q6_R_lsracc_RI(Word32 Rx, Word32 Rs, Word32 Iu5) 592
    Rx+=lsr(Rs,Rt)
    Word32 Q6_R_lsracc_RR(Word32 Rx, Word32 Rs, Word32 Rt) 605
    Rx=add(#u8,lsr(Rx,#U5))
    Word32 Q6_R_add_lsr_IRI(Word32 Iu8, Word32 Rx, Word32 IU5) 592
    Rx=and(#u8,lsr(Rx,#U5))
    Word32 Q6_R_and_lsr_IRI(Word32 Iu8, Word32 Rx, Word32 IU5) 596
    Rx-=lsr(Rs,#u5)

```

```

    Word32 Q6_R_lsrnac_RI(Word32 Rx, Word32 Rs, Word32 Iu5) 592
Rx-=lsr(Rs,Rt)
    Word32 Q6_R_lsrnac_RR(Word32 Rx, Word32 Rs, Word32 Rt) 605
Rx=or(#u8,lsr(Rx,#U5))
    Word32 Q6_R_or_lsr_IRI(Word32 Iu8, Word32 Rx, Word32 IU5) 596
Rx=sub(#u8,lsr(Rx,#U5))
    Word32 Q6_R_sub_lsr_IRI(Word32 Iu8, Word32 Rx, Word32 IU5) 592
Rx|=lsr(Rs,#u5)
    Word32 Q6_R_lsr_ror_RI(Word32 Rx, Word32 Rs, Word32 Iu5) 596
Rx|=lsr(Rs,Rt)
    Word32 Q6_R_lsr_ror_RR(Word32 Rx, Word32 Rs, Word32 Rt) 608
Rxx&=lsr(Rss,#u6)
    Word64 Q6_P_lsr_rand_PI(Word64 Rxx, Word64 Rss, Word32 Iu6) 596
Rxx&=lsr(Rss,Rt)
    Word64 Q6_P_lsr_rand_PR(Word64 Rxx, Word64 Rss, Word32 Rt) 608
Rxx^=lsr(Rss,#u6)
    Word64 Q6_P_lsr_xacc_PI(Word64 Rxx, Word64 Rss, Word32 Iu6) 596
Rxx^=lsr(Rss,Rt)
    Word64 Q6_P_lsr_xacc_PR(Word64 Rxx, Word64 Rss, Word32 Rt) 609
Rxx+=lsr(Rss,#u6)
    Word64 Q6_P_lsr_acc_PI(Word64 Rxx, Word64 Rss, Word32 Iu6) 592
Rxx+=lsr(Rss,Rt)
    Word64 Q6_P_lsr_acc_PR(Word64 Rxx, Word64 Rss, Word32 Rt) 605
Rxx-=lsr(Rss,#u6)
    Word64 Q6_P_lsr_nac_PI(Word64 Rxx, Word64 Rss, Word32 Iu6) 592
Rxx-=lsr(Rss,Rt)
    Word64 Q6_P_lsr_nac_PR(Word64 Rxx, Word64 Rss, Word32 Rt) 605
Rxx|=lsr(Rss,#u6)
    Word64 Q6_P_lsr_ror_PI(Word64 Rxx, Word64 Rss, Word32 Iu6) 597
Rxx|=lsr(Rss,Rt)
    Word64 Q6_P_lsr_ror_PR(Word64 Rxx, Word64 Rss, Word32 Rt) 609

```

**M**

## mask

```

Rd=mask(#u5,#U5)
    Word32 Q6_R_mask_II(Word32 Iu5, Word32 IU5) 588
Rdd=mask(Pt)
    Word64 Q6_P_mask_p(Byte Pt) 575

```

## max

```

Rd=max(Rs,Rt)
    Word32 Q6_R_max_RR(Word32 Rs, Word32 Rt) 349
Rdd=max(Rss,Rtt)
    Word64 Q6_P_max_PP(Word64 Rss, Word64 Rtt) 350

```

## maxu

```

Rd=maxu(Rs,Rt)
    UWord32 Q6_R_maxu_RR(Word32 Rs, Word32 Rt) 349
Rdd=maxu(Rss,Rtt)
    UWord64 Q6_P_maxu_PP(Word64 Rss, Word64 Rtt) 350

```

## memb

```

memb(Rx++#s4:0:circ(Mu))=Rt
    void Q6_memb_IMR_circ(void** StartAddress, Word32 Is4_0, Word32 Mu, Word32
    Rt, void* BaseAddress) 293
memb(Rx++I:circ(Mu))=Rt
    void Q6_memb_MR_circ(void** StartAddress, Word32 Mu, Word32 Rt, void*

```



```

        BaseAddress) 293
Rd=memb(Rx++#s4:0:circ(Mu))
    Word32 Q6_R_memb_IM_circ(void** StartAddress, Word32 Is4_0, Word32 Mu,
        void* BaseAddress) 227
Rd=memb(Rx++I:circ(Mu))
    Word32 Q6_R_memb_M_circ(void** StartAddress, Word32 Mu, void* BaseAddress)
    227

memd
memd(Rx++#s4:3:circ(Mu))=Rtt
    void Q6_memd_IMP_circ(void** StartAddress, Word32 Is4_3, Word32 Mu, Word64
        Rtt, void* BaseAddress) 289
memd(Rx++I:circ(Mu))=Rtt
    void Q6_memd_MP_circ(void** StartAddress, Word32 Mu, Word64 Rtt, void*
        BaseAddress) 289
Rdd=memd(Rx++#s4:3:circ(Mu))
    Word32 Q6_R_memd_IM_circ(void** StartAddress, Word32 Is4_3, Word32 Mu,
        void* BaseAddress) 223
Rdd=memd(Rx++I:circ(Mu))
    Word32 Q6_R_memd_M_circ(void** StartAddress, Word32 Mu, void* BaseAddress)
    223

memh
memh(Rx++#s4:1:circ(Mu))=Rt
    void Q6_memh_IMR_circ(void** StartAddress, Word32 Is4_1, Word32 Mu, Word32
        Rt, void* BaseAddress) 299
memh(Rx++#s4:1:circ(Mu))=Rt.H
    void Q6_memh_IMRh_circ(void** StartAddress, Word32 Is4_1, Word32 Mu, Word32
        Rt, void* BaseAddress) 299
memh(Rx++I:circ(Mu))=Rt
    void Q6_memh_MR_circ(void** StartAddress, Word32 Mu, Word32 Rt, void*
        BaseAddress) 299
memh(Rx++I:circ(Mu))=Rt.H
    void Q6_memh_MRh_circ(void** StartAddress, Word32 Mu, Word32 Rt, void*
        BaseAddress) 299
Rd=memh(Rx++#s4:1:circ(Mu))
    Word32 Q6_R_memh_IM_circ(void** StartAddress, Word32 Is4_1, Word32 Mu,
        void* BaseAddress) 237
Rd=memh(Rx++I:circ(Mu))
    Word32 Q6_R_memh_M_circ(void** StartAddress, Word32 Mu, void* BaseAddress)
    237

memub
Rd=memub(Rx++#s4:0:circ(Mu))
    Word32 Q6_R_memub_IM_circ(void** StartAddress, Word32 Is4_0, Word32 Mu,
        void* BaseAddress) 243
Rd=memub(Rx++I:circ(Mu))
    Word32 Q6_R_memub_M_circ(void** StartAddress, Word32 Mu, void*
        BaseAddress) 243

memuh
Rd=memuh(Rx++#s4:1:circ(Mu))
    Word32 Q6_R_memuh_IM_circ(void** StartAddress, Word32 Is4_1, Word32 Mu,
        void* BaseAddress) 247
Rd=memuh(Rx++I:circ(Mu))
    Word32 Q6_R_memuh_M_circ(void** StartAddress, Word32 Mu, void*
        BaseAddress) 247

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```

memw
    memw(Rx++#s4:2:circ(Mu))=Rt
        void Q6_memw_IMR_circ(void** StartAddress, Word32 Is4_2, Word32 Mu, Word32
            Rt, void* BaseAddress) 306
    memw(Rx++I:circ(Mu))=Rt
        void Q6_memw_MR_circ(void** StartAddress, Word32 Mu, Word32 Rt, void*
            BaseAddress) 306
    Rd=memw(Rx++#s4:2:circ(Mu))
        Word32 Q6_R_memw_IM_circ(void** StartAddress, Word32 Is4_2, Word32 Mu,
            void* BaseAddress) 251
    Rd=memw(Rx++I:circ(Mu))
        Word32 Q6_R_memw_M_circ(void** StartAddress, Word32 Mu, void* BaseAddress)
            251

min
    Rd=min(Rt,Rs)
        Word32 Q6_R_min_RR(Word32 Rt, Word32 Rs) 351
    Rdd=min(Rtt,Rss)
        Word64 Q6_P_min_PP(Word64 Rtt, Word64 Rss) 352

minu
    Rd=minu(Rt,Rs)
        UWord32 Q6_R_minu_RR(Word32 Rt, Word32 Rs) 351
    Rdd=minu(Rtt,Rss)
        UWord64 Q6_P_minu_PP(Word64 Rtt, Word64 Rss) 352

modwrap
    Rd=modwrap(Rs,Rt)
        Word32 Q6_R_modwrap_RR(Word32 Rs, Word32 Rt) 353

mpy
    Rd=mpy(Rs,Rt.H):<<1:rnd:sat
        Word32 Q6_R_mpy_RRh_s1_rnd_sat(Word32 Rs, Word32 Rt) 513
    Rd=mpy(Rs,Rt.H):<<1:sat
        Word32 Q6_R_mpy_RRh_s1_sat(Word32 Rs, Word32 Rt) 513
    Rd=mpy(Rs,Rt.L):<<1:rnd:sat
        Word32 Q6_R_mpy_RRl_s1_rnd_sat(Word32 Rs, Word32 Rt) 513
    Rd=mpy(Rs,Rt.L):<<1:sat
        Word32 Q6_R_mpy_RRl_s1_sat(Word32 Rs, Word32 Rt) 513
    Rd=mpy(Rs,Rt)
        Word32 Q6_R_mpy_RR(Word32 Rs, Word32 Rt) 513
    Rd=mpy(Rs,Rt):<<1
        Word32 Q6_R_mpy_RR_s1(Word32 Rs, Word32 Rt) 513
    Rd=mpy(Rs,Rt):<<1:sat
        Word32 Q6_R_mpy_RR_s1_sat(Word32 Rs, Word32 Rt) 513
    Rd=mpy(Rs,Rt):rnd
        Word32 Q6_R_mpy_RR_rnd(Word32 Rs, Word32 Rt) 513
    Rd=mpy(Rs.H,Rt.H)
        Word32 Q6_R_mpy_RhRh(Word32 Rs, Word32 Rt) 497
    Rd=mpy(Rs.H,Rt.H):<<1
        Word32 Q6_R_mpy_RhRh_s1(Word32 Rs, Word32 Rt) 497
    Rd=mpy(Rs.H,Rt.H):<<1:rnd
        Word32 Q6_R_mpy_RhRh_s1_rnd(Word32 Rs, Word32 Rt) 497
    Rd=mpy(Rs.H,Rt.H):<<1:rnd:sat
        Word32 Q6_R_mpy_RhRh_s1_rnd_sat(Word32 Rs, Word32 Rt) 497
    Rd=mpy(Rs.H,Rt.H):<<1:sat
        Word32 Q6_R_mpy_RhRh_s1_sat(Word32 Rs, Word32 Rt) 497
    Rd=mpy(Rs.H,Rt.H):rnd

```

Word32 Q6\_R\_mpy\_RhRh\_rnd(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.H,Rt.H):rnd:sat

Word32 Q6\_R\_mpy\_RhRh\_rnd\_sat(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.H,Rt.H):sat

Word32 Q6\_R\_mpy\_RhRh\_sat(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.H,Rt.L)

Word32 Q6\_R\_mpy\_RhRl(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.H,Rt.L):<<1

Word32 Q6\_R\_mpy\_RhRl\_sl(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.H,Rt.L):<<1:rnd

Word32 Q6\_R\_mpy\_RhRl\_sl\_rnd(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.H,Rt.L):<<1:rnd:sat

Word32 Q6\_R\_mpy\_RhRl\_sl\_rnd\_sat(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.H,Rt.L):<<1:sat

Word32 Q6\_R\_mpy\_RhRl\_sl\_sat(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.H,Rt.L):rnd

Word32 Q6\_R\_mpy\_RhRl\_rnd(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.H,Rt.L):rnd:sat

Word32 Q6\_R\_mpy\_RhRl\_rnd\_sat(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.H,Rt.L):sat

Word32 Q6\_R\_mpy\_RhRl\_sat(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.L,Rt.H)

Word32 Q6\_R\_mpy\_RlRh(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.L,Rt.H):<<1

Word32 Q6\_R\_mpy\_RlRh\_sl(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.L,Rt.H):<<1:rnd

Word32 Q6\_R\_mpy\_RlRh\_sl\_rnd(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.L,Rt.H):<<1:rnd:sat

Word32 Q6\_R\_mpy\_RlRh\_sl\_rnd\_sat(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.L,Rt.H):<<1:sat

Word32 Q6\_R\_mpy\_RlRh\_sl\_sat(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.L,Rt.H):rnd

Word32 Q6\_R\_mpy\_RlRh\_rnd(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.L,Rt.H):rnd:sat

Word32 Q6\_R\_mpy\_RlRh\_rnd\_sat(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.L,Rt.H):sat

Word32 Q6\_R\_mpy\_RlRh\_sat(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.L,Rt.L)

Word32 Q6\_R\_mpy\_RlRl(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.L,Rt.L):<<1

Word32 Q6\_R\_mpy\_RlRl\_sl(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.L,Rt.L):<<1:rnd

Word32 Q6\_R\_mpy\_RlRl\_sl\_rnd(Word32 Rs, Word32 Rt) 497  
Rd=mpy(Rs.L,Rt.L):<<1:rnd:sat

Word32 Q6\_R\_mpy\_RlRl\_sl\_rnd\_sat(Word32 Rs, Word32 Rt) 498  
Rd=mpy(Rs.L,Rt.L):<<1:sat

Word32 Q6\_R\_mpy\_RlRl\_sl\_sat(Word32 Rs, Word32 Rt) 498  
Rd=mpy(Rs.L,Rt.L):rnd

Word32 Q6\_R\_mpy\_RlRl\_rnd(Word32 Rs, Word32 Rt) 498  
Rd=mpy(Rs.L,Rt.L):rnd:sat

Word32 Q6\_R\_mpy\_RlRl\_rnd\_sat(Word32 Rs, Word32 Rt) 498  
Rd=mpy(Rs.L,Rt.L):sat

Word32 Q6\_R\_mpy\_RlRl\_sat(Word32 Rs, Word32 Rt) 498  
Rdd=mpy(Rs,Rt)

Word64 Q6\_P\_mpy\_RR(Word32 Rs, Word32 Rt) 515  
Rdd=mpy(Rs.H,Rt.H)

Word64 Q6\_P\_mpy\_RhRh(Word32 Rs, Word32 Rt) 498  
Rdd=mpy(Rs.H,Rt.H):<<1

Word64 Q6\_P\_mpy\_RhRh\_s1(Word32 Rs, Word32 Rt) **498**  
 Rdd=mpy(Rs.H,Rt.H):<<1:rnd  
 Word64 Q6\_P\_mpy\_RhRh\_s1\_rnd(Word32 Rs, Word32 Rt) **498**  
 Rdd=mpy(Rs.H,Rt.H):rnd  
 Word64 Q6\_P\_mpy\_RhRh\_rnd(Word32 Rs, Word32 Rt) **498**  
 Rdd=mpy(Rs.H,Rt.L)  
 Word64 Q6\_P\_mpy\_RhRl(Word32 Rs, Word32 Rt) **498**  
 Rdd=mpy(Rs.H,Rt.L):<<1  
 Word64 Q6\_P\_mpy\_RhRl\_s1(Word32 Rs, Word32 Rt) **498**  
 Rdd=mpy(Rs.H,Rt.L):<<1:rnd  
 Word64 Q6\_P\_mpy\_RhRl\_s1\_rnd(Word32 Rs, Word32 Rt) **498**  
 Rdd=mpy(Rs.H,Rt.L):rnd  
 Word64 Q6\_P\_mpy\_RhRl\_rnd(Word32 Rs, Word32 Rt) **498**  
 Rdd=mpy(Rs.L,Rt.H)  
 Word64 Q6\_P\_mpy\_RlRh(Word32 Rs, Word32 Rt) **498**  
 Rdd=mpy(Rs.L,Rt.H):<<1  
 Word64 Q6\_P\_mpy\_RlRh\_s1(Word32 Rs, Word32 Rt) **498**  
 Rdd=mpy(Rs.L,Rt.H):<<1:rnd  
 Word64 Q6\_P\_mpy\_RlRh\_s1\_rnd(Word32 Rs, Word32 Rt) **498**  
 Rdd=mpy(Rs.L,Rt.H):rnd  
 Word64 Q6\_P\_mpy\_RlRh\_rnd(Word32 Rs, Word32 Rt) **498**  
 Rdd=mpy(Rs.L,Rt.L)  
 Word64 Q6\_P\_mpy\_RlRl(Word32 Rs, Word32 Rt) **498**  
 Rdd=mpy(Rs.L,Rt.L):<<1  
 Word64 Q6\_P\_mpy\_RlRl\_s1(Word32 Rs, Word32 Rt) **498**  
 Rdd=mpy(Rs.L,Rt.L):<<1:rnd  
 Word64 Q6\_P\_mpy\_RlRl\_s1\_rnd(Word32 Rs, Word32 Rt) **498**  
 Rdd=mpy(Rs.L,Rt.L):rnd  
 Word64 Q6\_P\_mpy\_RlRl\_rnd(Word32 Rs, Word32 Rt) **498**  
 Rx+=mpy(Rs,Rt):<<1:sat  
 Word32 Q6\_R\_mpyacc\_RR\_s1\_sat(Word32 Rx, Word32 Rs, Word32 Rt) **513**  
 Rx+=mpy(Rs.H,Rt.H)  
 Word32 Q6\_R\_mpyacc\_RhRh(Word32 Rx, Word32 Rs, Word32 Rt) **498**  
 Rx+=mpy(Rs.H,Rt.H):<<1  
 Word32 Q6\_R\_mpyacc\_RhRh\_s1(Word32 Rx, Word32 Rs, Word32 Rt) **498**  
 Rx+=mpy(Rs.H,Rt.H):<<1:sat  
 Word32 Q6\_R\_mpyacc\_RhRh\_s1\_sat(Word32 Rx, Word32 Rs, Word32 Rt) **498**  
 Rx+=mpy(Rs.H,Rt.H):sat  
 Word32 Q6\_R\_mpyacc\_RhRh\_sat(Word32 Rx, Word32 Rs, Word32 Rt) **498**  
 Rx+=mpy(Rs.H,Rt.L)  
 Word32 Q6\_R\_mpyacc\_RhRl(Word32 Rx, Word32 Rs, Word32 Rt) **498**  
 Rx+=mpy(Rs.H,Rt.L):<<1  
 Word32 Q6\_R\_mpyacc\_RhRl\_s1(Word32 Rx, Word32 Rs, Word32 Rt) **498**  
 Rx+=mpy(Rs.H,Rt.L):<<1:sat  
 Word32 Q6\_R\_mpyacc\_RhRl\_s1\_sat(Word32 Rx, Word32 Rs, Word32 Rt) **498**  
 Rx+=mpy(Rs.H,Rt.L):sat  
 Word32 Q6\_R\_mpyacc\_RhRl\_sat(Word32 Rx, Word32 Rs, Word32 Rt) **498**  
 Rx+=mpy(Rs.L,Rt.H)  
 Word32 Q6\_R\_mpyacc\_RlRh(Word32 Rx, Word32 Rs, Word32 Rt) **498**  
 Rx+=mpy(Rs.L,Rt.H):<<1  
 Word32 Q6\_R\_mpyacc\_RlRh\_s1(Word32 Rx, Word32 Rs, Word32 Rt) **498**  
 Rx+=mpy(Rs.L,Rt.H):<<1:sat  
 Word32 Q6\_R\_mpyacc\_RlRh\_s1\_sat(Word32 Rx, Word32 Rs, Word32 Rt) **499**  
 Rx+=mpy(Rs.L,Rt.H):sat  
 Word32 Q6\_R\_mpyacc\_RlRh\_sat(Word32 Rx, Word32 Rs, Word32 Rt) **499**  
 Rx+=mpy(Rs.L,Rt.L)  
 Word32 Q6\_R\_mpyacc\_RlRl(Word32 Rx, Word32 Rs, Word32 Rt) **499**  
 Rx+=mpy(Rs.L,Rt.L):<<1

Word32 Q6\_R\_mpyacc\_RlRl\_s1(Word32 Rx, Word32 Rs, Word32 Rt) 499  
Rx+=mpy(Rs.L,Rt.L):<<1:sat

Word32 Q6\_R\_mpyacc\_RlRl\_s1\_sat(Word32 Rx, Word32 Rs, Word32 Rt) 499  
Rx+=mpy(Rs.L,Rt.L):sat

Word32 Q6\_R\_mpyacc\_RlRl\_sat(Word32 Rx, Word32 Rs, Word32 Rt) 499  
Rx-=mpy(Rs,Rt):<<1:sat

Word32 Q6\_R\_mpyacc\_RR\_s1\_sat(Word32 Rx, Word32 Rs, Word32 Rt) 513  
Rx-=mpy(Rs.H,Rt.H)

Word32 Q6\_R\_mpyacc\_RhRh(Word32 Rx, Word32 Rs, Word32 Rt) 499  
Rx-=mpy(Rs.H,Rt.H):<<1

Word32 Q6\_R\_mpyacc\_RhRh\_s1(Word32 Rx, Word32 Rs, Word32 Rt) 499  
Rx-=mpy(Rs.H,Rt.H):<<1:sat

Word32 Q6\_R\_mpyacc\_RhRh\_s1\_sat(Word32 Rx, Word32 Rs, Word32 Rt) 499  
Rx-=mpy(Rs.H,Rt.H):sat

Word32 Q6\_R\_mpyacc\_RhRh\_sat(Word32 Rx, Word32 Rs, Word32 Rt) 499  
Rx-=mpy(Rs.H,Rt.L)

Word32 Q6\_R\_mpyacc\_RhRl(Word32 Rx, Word32 Rs, Word32 Rt) 499  
Rx-=mpy(Rs.H,Rt.L):<<1

Word32 Q6\_R\_mpyacc\_RhRl\_s1(Word32 Rx, Word32 Rs, Word32 Rt) 499  
Rx-=mpy(Rs.H,Rt.L):<<1:sat

Word32 Q6\_R\_mpyacc\_RhRl\_s1\_sat(Word32 Rx, Word32 Rs, Word32 Rt) 499  
Rx-=mpy(Rs.H,Rt.L):sat

Word32 Q6\_R\_mpyacc\_RhRl\_sat(Word32 Rx, Word32 Rs, Word32 Rt) 499  
Rx-=mpy(Rs.L,Rt.H)

Word32 Q6\_R\_mpyacc\_RlRh(Word32 Rx, Word32 Rs, Word32 Rt) 499  
Rx-=mpy(Rs.L,Rt.H):<<1

Word32 Q6\_R\_mpyacc\_RlRh\_s1(Word32 Rx, Word32 Rs, Word32 Rt) 499  
Rx-=mpy(Rs.L,Rt.H):<<1:sat

Word32 Q6\_R\_mpyacc\_RlRh\_s1\_sat(Word32 Rx, Word32 Rs, Word32 Rt) 499  
Rx-=mpy(Rs.L,Rt.H):sat

Word32 Q6\_R\_mpyacc\_RlRh\_sat(Word32 Rx, Word32 Rs, Word32 Rt) 499  
Rx-=mpy(Rs.L,Rt.L)

Word32 Q6\_R\_mpyacc\_RlRl(Word32 Rx, Word32 Rs, Word32 Rt) 499  
Rx-=mpy(Rs.L,Rt.L):<<1

Word32 Q6\_R\_mpyacc\_RlRl\_s1(Word32 Rx, Word32 Rs, Word32 Rt) 499  
Rx-=mpy(Rs.L,Rt.L):<<1:sat

Word32 Q6\_R\_mpyacc\_RlRl\_s1\_sat(Word32 Rx, Word32 Rs, Word32 Rt) 499  
Rx-=mpy(Rs.L,Rt.L):sat

Word32 Q6\_R\_mpyacc\_RlRl\_sat(Word32 Rx, Word32 Rs, Word32 Rt) 499  
Rxx+=mpy(Rs,Rt)

Word64 Q6\_P\_mpyacc\_RR(Word64 Rxx, Word32 Rs, Word32 Rt) 515  
Rxx+=mpy(Rs.H,Rt.H)

Word64 Q6\_P\_mpyacc\_RhRh(Word64 Rxx, Word32 Rs, Word32 Rt) 499  
Rxx+=mpy(Rs.H,Rt.H):<<1

Word64 Q6\_P\_mpyacc\_RhRh\_s1(Word64 Rxx, Word32 Rs, Word32 Rt) 500  
Rxx+=mpy(Rs.H,Rt.L)

Word64 Q6\_P\_mpyacc\_RhRl(Word64 Rxx, Word32 Rs, Word32 Rt) 500  
Rxx+=mpy(Rs.H,Rt.L):<<1

Word64 Q6\_P\_mpyacc\_RhRl\_s1(Word64 Rxx, Word32 Rs, Word32 Rt) 500  
Rxx+=mpy(Rs.L,Rt.H)

Word64 Q6\_P\_mpyacc\_RlRh(Word64 Rxx, Word32 Rs, Word32 Rt) 500  
Rxx+=mpy(Rs.L,Rt.H):<<1

Word64 Q6\_P\_mpyacc\_RlRh\_s1(Word64 Rxx, Word32 Rs, Word32 Rt) 500  
Rxx+=mpy(Rs.L,Rt.L)

Word64 Q6\_P\_mpyacc\_RlRl(Word64 Rxx, Word32 Rs, Word32 Rt) 500  
Rxx+=mpy(Rs.L,Rt.L):<<1

Word64 Q6\_P\_mpyacc\_RlRl\_s1(Word64 Rxx, Word32 Rs, Word32 Rt) 500  
Rxx-=mpy(Rs,Rt)

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    Word64 Q6_P_mpynac_RR(Word64 Rxx, Word32 Rs, Word32 Rt) 515
Rxx-=mpy (Rs.H,Rt.H)
    Word64 Q6_P_mpynac_RhRh(Word64 Rxx, Word32 Rs, Word32 Rt) 500
Rxx-=mpy (Rs.H,Rt.H) :<<1
    Word64 Q6_P_mpynac_RhRh_s1(Word64 Rxx, Word32 Rs, Word32 Rt) 500
Rxx-=mpy (Rs.H,Rt.L)
    Word64 Q6_P_mpynac_RhRl(Word64 Rxx, Word32 Rs, Word32 Rt) 500
Rxx-=mpy (Rs.H,Rt.L) :<<1
    Word64 Q6_P_mpynac_RhRl_s1(Word64 Rxx, Word32 Rs, Word32 Rt) 500
Rxx-=mpy (Rs.L,Rt.H)
    Word64 Q6_P_mpynac_RlRh(Word64 Rxx, Word32 Rs, Word32 Rt) 500
Rxx-=mpy (Rs.L,Rt.H) :<<1
    Word64 Q6_P_mpynac_RlRh_s1(Word64 Rxx, Word32 Rs, Word32 Rt) 500
Rxx-=mpy (Rs.L,Rt.L)
    Word64 Q6_P_mpynac_RlRl(Word64 Rxx, Word32 Rs, Word32 Rt) 500
Rxx-=mpy (Rs.L,Rt.L) :<<1
    Word64 Q6_P_mpynac_RlRl_s1(Word64 Rxx, Word32 Rs, Word32 Rt) 500

mpyi
Rd=mpyi (Rs,#m9)
    Word32 Q6_R_mpyi_RI(Word32 Rs, Word32 Im9) 486
Rd=mpyi (Rs,Rt)
    Word32 Q6_R_mpyi_RR(Word32 Rs, Word32 Rt) 486
Rx+=mpyi (Rs,#u8)
    Word32 Q6_R_mpyiacc_RI(Word32 Rx, Word32 Rs, Word32 Iu8) 486
Rx+=mpyi (Rs,Rt)
    Word32 Q6_R_mpyiacc_RR(Word32 Rx, Word32 Rs, Word32 Rt) 486
Rx-=mpyi (Rs,#u8)
    Word32 Q6_R_mpyinac_RI(Word32 Rx, Word32 Rs, Word32 Iu8) 486
Rx-=mpyi (Rs,Rt)
    Word32 Q6_R_mpyinac_RR(Word32 Rx, Word32 Rs, Word32 Rt) 486

mpysu
Rd=mpysu (Rs,Rt)
    Word32 Q6_R_mpyisu_RR(Word32 Rs, Word32 Rt) 513

mpyu
Rd=mpyu (Rs,Rt)
    UWord32 Q6_R_mpyu_RR(Word32 Rs, Word32 Rt) 513
Rd=mpyu (Rs.H,Rt.H)
    UWord32 Q6_R_mpyu_RhRh(Word32 Rs, Word32 Rt) 504
Rd=mpyu (Rs.H,Rt.H) :<<1
    UWord32 Q6_R_mpyu_RhRh_s1(Word32 Rs, Word32 Rt) 504
Rd=mpyu (Rs.H,Rt.L)
    UWord32 Q6_R_mpyu_RhRl(Word32 Rs, Word32 Rt) 504
Rd=mpyu (Rs.H,Rt.L) :<<1
    UWord32 Q6_R_mpyu_RhRl_s1(Word32 Rs, Word32 Rt) 504
Rd=mpyu (Rs.L,Rt.H)
    UWord32 Q6_R_mpyu_RlRh(Word32 Rs, Word32 Rt) 504
Rd=mpyu (Rs.L,Rt.H) :<<1
    UWord32 Q6_R_mpyu_RlRh_s1(Word32 Rs, Word32 Rt) 504
Rd=mpyu (Rs.L,Rt.L)
    UWord32 Q6_R_mpyu_RlRl(Word32 Rs, Word32 Rt) 504
Rd=mpyu (Rs.L,Rt.L) :<<1
    UWord32 Q6_R_mpyu_RlRl_s1(Word32 Rs, Word32 Rt) 504
Rdd=mpyu (Rs,Rt)
    UWord64 Q6_P_mpyu_RR(Word32 Rs, Word32 Rt) 515
Rdd=mpyu (Rs.H,Rt.H)

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UWord64 Q6\_P\_mpyu\_RhRh(Word32 Rs, Word32 Rt) 504  
 Rdd=mpyu(Rs.H,Rt.H):<<1  
 UWord64 Q6\_P\_mpyu\_RhRh\_s1(Word32 Rs, Word32 Rt) 504  
 Rdd=mpyu(Rs.H,Rt.L)  
 UWord64 Q6\_P\_mpyu\_RhRl(Word32 Rs, Word32 Rt) 504  
 Rdd=mpyu(Rs.H,Rt.L):<<1  
 UWord64 Q6\_P\_mpyu\_RhRl\_s1(Word32 Rs, Word32 Rt) 504  
 Rdd=mpyu(Rs.L,Rt.H)  
 UWord64 Q6\_P\_mpyu\_RlRh(Word32 Rs, Word32 Rt) 504  
 Rdd=mpyu(Rs.L,Rt.H):<<1  
 UWord64 Q6\_P\_mpyu\_RlRh\_s1(Word32 Rs, Word32 Rt) 504  
 Rdd=mpyu(Rs.L,Rt.L)  
 UWord64 Q6\_P\_mpyu\_RlRl(Word32 Rs, Word32 Rt) 504  
 Rdd=mpyu(Rs.L,Rt.L):<<1  
 UWord64 Q6\_P\_mpyu\_RlRl\_s1(Word32 Rs, Word32 Rt) 504  
 Rx+=mpyu(Rs.H,Rt.H)  
 Word32 Q6\_R\_mpyuacc\_RhRh(Word32 Rx, Word32 Rs, Word32 Rt) 504  
 Rx+=mpyu(Rs.H,Rt.H):<<1  
 Word32 Q6\_R\_mpyuacc\_RhRh\_s1(Word32 Rx, Word32 Rs, Word32 Rt) 504  
 Rx+=mpyu(Rs.H,Rt.L)  
 Word32 Q6\_R\_mpyuacc\_RhRl(Word32 Rx, Word32 Rs, Word32 Rt) 504  
 Rx+=mpyu(Rs.H,Rt.L):<<1  
 Word32 Q6\_R\_mpyuacc\_RhRl\_s1(Word32 Rx, Word32 Rs, Word32 Rt) 504  
 Rx+=mpyu(Rs.L,Rt.H)  
 Word32 Q6\_R\_mpyuacc\_RlRh(Word32 Rx, Word32 Rs, Word32 Rt) 504  
 Rx+=mpyu(Rs.L,Rt.H):<<1  
 Word32 Q6\_R\_mpyuacc\_RlRh\_s1(Word32 Rx, Word32 Rs, Word32 Rt) 504  
 Rx+=mpyu(Rs.L,Rt.L)  
 Word32 Q6\_R\_mpyuacc\_RlRl(Word32 Rx, Word32 Rs, Word32 Rt) 504  
 Rx+=mpyu(Rs.L,Rt.L):<<1  
 Word32 Q6\_R\_mpyuacc\_RlRl\_s1(Word32 Rx, Word32 Rs, Word32 Rt) 504  
 Rx-=mpyu(Rs.H,Rt.H)  
 Word32 Q6\_R\_mpyunac\_RhRh(Word32 Rx, Word32 Rs, Word32 Rt) 504  
 Rx-=mpyu(Rs.H,Rt.H):<<1  
 Word32 Q6\_R\_mpyunac\_RhRh\_s1(Word32 Rx, Word32 Rs, Word32 Rt) 504  
 Rx-=mpyu(Rs.H,Rt.L)  
 Word32 Q6\_R\_mpyunac\_RhRl(Word32 Rx, Word32 Rs, Word32 Rt) 504  
 Rx-=mpyu(Rs.H,Rt.L):<<1  
 Word32 Q6\_R\_mpyunac\_RhRl\_s1(Word32 Rx, Word32 Rs, Word32 Rt) 505  
 Rx-=mpyu(Rs.L,Rt.H)  
 Word32 Q6\_R\_mpyunac\_RlRh(Word32 Rx, Word32 Rs, Word32 Rt) 505  
 Rx-=mpyu(Rs.L,Rt.H):<<1  
 Word32 Q6\_R\_mpyunac\_RlRh\_s1(Word32 Rx, Word32 Rs, Word32 Rt) 505  
 Rx-=mpyu(Rs.L,Rt.L)  
 Word32 Q6\_R\_mpyunac\_RlRl(Word32 Rx, Word32 Rs, Word32 Rt) 505  
 Rx-=mpyu(Rs.L,Rt.L):<<1  
 Word32 Q6\_R\_mpyunac\_RlRl\_s1(Word32 Rx, Word32 Rs, Word32 Rt) 505  
 Rxx+=mpyu(Rs,Rt)  
 Word64 Q6\_P\_mpyuacc\_RR(Word64 Rxx, Word32 Rs, Word32 Rt) 515  
 Rxx+=mpyu(Rs.H,Rt.H)  
 Word64 Q6\_P\_mpyuacc\_RhRh(Word64 Rxx, Word32 Rs, Word32 Rt) 505  
 Rxx+=mpyu(Rs.H,Rt.H):<<1  
 Word64 Q6\_P\_mpyuacc\_RhRh\_s1(Word64 Rxx, Word32 Rs, Word32 Rt) 505  
 Rxx+=mpyu(Rs.H,Rt.L)  
 Word64 Q6\_P\_mpyuacc\_RhRl(Word64 Rxx, Word32 Rs, Word32 Rt) 505  
 Rxx+=mpyu(Rs.H,Rt.L):<<1  
 Word64 Q6\_P\_mpyuacc\_RhRl\_s1(Word64 Rxx, Word32 Rs, Word32 Rt) 505  
 Rxx+=mpyu(Rs.L,Rt.H)

Word64 Q6\_P\_mpyuacc\_RlRh(Word64 Rxx, Word32 Rs, Word32 Rt) **505**  
 Rxx+=mpyu(Rs.L,Rt.H):<<1  
 Word64 Q6\_P\_mpyuacc\_RlRh\_s1(Word64 Rxx, Word32 Rs, Word32 Rt) **505**  
 Rxx+=mpyu(Rs.L,Rt.L)  
 Word64 Q6\_P\_mpyuacc\_RlRl(Word64 Rxx, Word32 Rs, Word32 Rt) **505**  
 Rxx+=mpyu(Rs.L,Rt.L):<<1  
 Word64 Q6\_P\_mpyuacc\_RlRl\_s1(Word64 Rxx, Word32 Rs, Word32 Rt) **505**  
 Rxx-=mpyu(Rs,Rt)  
 Word64 Q6\_P\_mpyunac\_RR(Word64 Rxx, Word32 Rs, Word32 Rt) **515**  
 Rxx-=mpyu(Rs.H,Rt.H)  
 Word64 Q6\_P\_mpyunac\_RhRh(Word64 Rxx, Word32 Rs, Word32 Rt) **505**  
 Rxx-=mpyu(Rs.H,Rt.H):<<1  
 Word64 Q6\_P\_mpyunac\_RhRh\_s1(Word64 Rxx, Word32 Rs, Word32 Rt) **505**  
 Rxx-=mpyu(Rs.H,Rt.L)  
 Word64 Q6\_P\_mpyunac\_RhRl(Word64 Rxx, Word32 Rs, Word32 Rt) **505**  
 Rxx-=mpyu(Rs.H,Rt.L):<<1  
 Word64 Q6\_P\_mpyunac\_RhRl\_s1(Word64 Rxx, Word32 Rs, Word32 Rt) **505**  
 Rxx-=mpyu(Rs.L,Rt.H)  
 Word64 Q6\_P\_mpyunac\_RlRh(Word64 Rxx, Word32 Rs, Word32 Rt) **505**  
 Rxx-=mpyu(Rs.L,Rt.H):<<1  
 Word64 Q6\_P\_mpyunac\_RlRh\_s1(Word64 Rxx, Word32 Rs, Word32 Rt) **505**  
 Rxx-=mpyu(Rs.L,Rt.L)  
 Word64 Q6\_P\_mpyunac\_RlRl(Word64 Rxx, Word32 Rs, Word32 Rt) **505**  
 Rxx-=mpyu(Rs.L,Rt.L):<<1  
 Word64 Q6\_P\_mpyunac\_RlRl\_s1(Word64 Rxx, Word32 Rs, Word32 Rt) **505**

**mpyui**  
 Rd=mpyui(Rs,Rt)  
 Word32 Q6\_R\_mpyui\_RR(Word32 Rs, Word32 Rt) **486**

**mux**  
 Rd=mux(Pu,#s8,#S8)  
 Word32 Q6\_R\_mux\_pII(Byte Pu, Word32 Is8, Word32 IS8) **174**  
 Rd=mux(Pu,#s8,Rs)  
 Word32 Q6\_R\_mux\_pIR(Byte Pu, Word32 Is8, Word32 Rs) **174**  
 Rd=mux(Pu,Rs,#s8)  
 Word32 Q6\_R\_mux\_pRI(Byte Pu, Word32 Rs, Word32 Is8) **174**  
 Rd=mux(Pu,Rs,Rt)  
 Word32 Q6\_R\_mux\_pRR(Byte Pu, Word32 Rs, Word32 Rt) **174**

**N**

**neg**  
 Rd=neg(Rs)  
 Word32 Q6\_R\_neg\_R(Word32 Rs) **160**  
 Rd=neg(Rs):sat  
 Word32 Q6\_R\_neg\_R\_sat(Word32 Rs) **354**  
 Rdd=neg(Rss)  
 Word64 Q6\_P\_neg\_P(Word64 Rss) **354**

**no mnemonic**  
 Pd=Ps  
 Byte Q6\_p\_equals\_p(Byte Ps) **203**  
 Pd=Rs  
 Byte Q6\_p\_equals\_R(Word32 Rs) **577**  
 Rd=#s16  
 Word32 Q6\_R\_equals\_I(Word32 Is16) **165**  
 Rd=Ps



```

    Word32 Q6_R_equals_p(Byte Ps) 577
Rd=Rs
    Word32 Q6_R_equals_R(Word32 Rs) 167
Rdd=#s8
    Word64 Q6_P_equals_I(Word32 Is8) 165
Rdd=Rss
    Word64 Q6_P_equals_P(Word64 Rss) 167
Rx.H=#u16
    Word32 Q6_Rh_equals_I(Word32 Rx, Word32 Iu16) 165
Rx.L=#u16
    Word32 Q6_Rl_equals_I(Word32 Rx, Word32 Iu16) 165

normamt
    Rd=normamt (Rs)
    Word32 Q6_R_normamt_R(Word32 Rs) 410
    Rd=normamt (Rss)
    Word32 Q6_R_normamt_P(Word64 Rss) 410

not
    Pd=not (Ps)
    Byte Q6_p_not_p(Byte Ps) 203
    Rd=not (Rs)
    Word32 Q6_R_not_R(Word32 Rs) 158
    Rdd=not (Rss)
    Word64 Q6_P_not_P(Word64 Rss) 343

O

or
    Pd=and(Ps,or(Pt,!Pu))
    Byte Q6_p_and_or_ppnp(Byte Ps, Byte Pt, Byte Pu) 203
    Pd=and(Ps,or(Pt,Pu))
    Byte Q6_p_and_or_ppp(Byte Ps, Byte Pt, Byte Pu) 203
    Pd=or(Ps,and(Pt,!Pu))
    Byte Q6_p_or_and_ppnp(Byte Ps, Byte Pt, Byte Pu) 203
    Pd=or(Ps,and(Pt,Pu))
    Byte Q6_p_or_and_ppp(Byte Ps, Byte Pt, Byte Pu) 203
    Pd=or(Ps,or(Pt,!Pu))
    Byte Q6_p_or_or_ppnp(Byte Ps, Byte Pt, Byte Pu) 204
    Pd=or(Ps,or(Pt,Pu))
    Byte Q6_p_or_or_ppp(Byte Ps, Byte Pt, Byte Pu) 204
    Pd=or(Pt,!Ps)
    Byte Q6_p_or_pnp(Byte Pt, Byte Ps) 204
    Pd=or(Pt,Ps)
    Byte Q6_p_or_pp(Byte Pt, Byte Ps) 204
    Rd=or(Rs,#s10)
    Word32 Q6_R_or_RI(Word32 Rs, Word32 Is10) 158
    Rd=or(Rs,Rt)
    Word32 Q6_R_or_RR(Word32 Rs, Word32 Rt) 158
    Rd=or(Rt,~Rs)
    Word32 Q6_R_or_RnR(Word32 Rt, Word32 Rs) 158
    Rdd=or(Rss,Rtt)
    Word64 Q6_P_or_PP(Word64 Rss, Word64 Rtt) 343
    Rdd=or(Rtt,~Rss)
    Word64 Q6_P_or_PnP(Word64 Rtt, Word64 Rss) 343
    Rx&=or(Rs,Rt)
    Word32 Q6_R_orand_RR(Word32 Rx, Word32 Rs, Word32 Rt) 347
    Rx^=or(Rs,Rt)

```

Word32 Q6\_R\_orxacc\_RR(Word32 Rx, Word32 Rs, Word32 Rt) **347**  
 Rx=or(Ru, and(Rx, #s10))  
 Word32 Q6\_R\_or\_and\_RRI(Word32 Ru, Word32 Rx, Word32 Is10) **347**  
 Rx|=or(Rs, #s10)  
 Word32 Q6\_R\_oror\_RI(Word32 Rx, Word32 Rs, Word32 Is10) **347**  
 Rx|=or(Rs, Rt)  
 Word32 Q6\_R\_oror\_RR(Word32 Rx, Word32 Rs, Word32 Rt) **347**

**P**

## packhl

Rdd=packhl(Rs, Rt)  
 Word64 Q6\_P\_packhl\_RR(Word32 Rs, Word32 Rt) **177**

## parity

Rd=parity(Rs, Rt)  
 Word32 Q6\_R\_parity\_RR(Word32 Rs, Word32 Rt) **420**  
 Rd=parity(Rss, Rtt)  
 Word32 Q6\_R\_parity\_PP(Word64 Rss, Word64 Rtt) **420**

## pmpyw

Rdd=pmpyw(Rs, Rt)  
 Word64 Q6\_P\_pmpyw\_RR(Word32 Rs, Word32 Rt) **509**  
 Rxx^=pmpyw(Rs, Rt)  
 Word64 Q6\_P\_pmpywxacc\_RR(Word64 Rxx, Word32 Rs, Word32 Rt) **509**

## popcount

Rd=popcount(Rss)  
 Word32 Q6\_R\_popcount\_P(Word64 Rss) **411**

**R**

## rol

Rd=rol(Rs, #u5)  
 Word32 Q6\_R\_rol\_RI(Word32 Rs, Word32 Iu5) **589**  
 Rdd=rol(Rss, #u6)  
 Word64 Q6\_P\_rol\_PI(Word64 Rss, Word32 Iu6) **590**  
 Rx&=rol(Rs, #u5)  
 Word32 Q6\_R\_roland\_RI(Word32 Rx, Word32 Rs, Word32 Iu5) **596**  
 Rx^=rol(Rs, #u5)  
 Word32 Q6\_R\_rolxacc\_RI(Word32 Rx, Word32 Rs, Word32 Iu5) **596**  
 Rx+=rol(Rs, #u5)  
 Word32 Q6\_R\_rolacc\_RI(Word32 Rx, Word32 Rs, Word32 Iu5) **592**  
 Rx-=rol(Rs, #u5)  
 Word32 Q6\_R\_rolnac\_RI(Word32 Rx, Word32 Rs, Word32 Iu5) **592**  
 Rx|=rol(Rs, #u5)  
 Word32 Q6\_R\_rolor\_RI(Word32 Rx, Word32 Rs, Word32 Iu5) **596**  
 Rxx&=rol(Rss, #u6)  
 Word64 Q6\_P\_roland\_PI(Word64 Rxx, Word64 Rss, Word32 Iu6) **596**  
 Rxx^=rol(Rss, #u6)  
 Word64 Q6\_P\_rolxacc\_PI(Word64 Rxx, Word64 Rss, Word32 Iu6) **597**  
 Rxx+=rol(Rss, #u6)  
 Word64 Q6\_P\_rolacc\_PI(Word64 Rxx, Word64 Rss, Word32 Iu6) **592**  
 Rxx-=rol(Rss, #u6)  
 Word64 Q6\_P\_rolnac\_PI(Word64 Rxx, Word64 Rss, Word32 Iu6) **592**  
 Rxx|=rol(Rss, #u6)  
 Word64 Q6\_P\_rolor\_PI(Word64 Rxx, Word64 Rss, Word32 Iu6) **597**

## round

Rd=round(Rs, #u5)

```

    Word32 Q6_R_round_RI(Word32 Rs, Word32 Iu5) 357
    Rd=round(Rs,#u5):sat
    Word32 Q6_R_round_RI_sat(Word32 Rs, Word32 Iu5) 357
    Rd=round(Rs,Rt)
    Word32 Q6_R_round_RR(Word32 Rs, Word32 Rt) 357
    Rd=round(Rs,Rt):sat
    Word32 Q6_R_round_RR_sat(Word32 Rs, Word32 Rt) 357
    Rd=round(Rss):sat
    Word32 Q6_R_round_P_sat(Word64 Rss) 357

```

**S**

```

sat
    Rd=sat(Rss)
    Word32 Q6_R_sat_P(Word64 Rss) 542

satb
    Rd=satb(Rs)
    Word32 Q6_R_satb_R(Word32 Rs) 542

sath
    Rd=sath(Rs)
    Word32 Q6_R_sath_R(Word32 Rs) 542

satub
    Rd=satub(Rs)
    Word32 Q6_R_satub_R(Word32 Rs) 542

satuh
    Rd=satuh(Rs)
    Word32 Q6_R_satuh_R(Word32 Rs) 542

setbit
    Rd=setbit(Rs,#u5)
    Word32 Q6_R_setbit_RI(Word32 Rs, Word32 Iu5) 422
    Rd=setbit(Rs,Rt)
    Word32 Q6_R_setbit_RR(Word32 Rs, Word32 Rt) 422

sfadd
    Rd=sfadd(Rs,Rt)
    Word32 Q6_R_sfadd_RR(Word32 Rs, Word32 Rt) 461

sfclass
    Pd=sfclass(Rs,#u5)
    Byte Q6_p_sfclass_RI(Word32 Rs, Word32 Iu5) 462

sfcmp.eq
    Pd=sfcmp.eq(Rs,Rt)
    Byte Q6_p_sfcmp_eq_RR(Word32 Rs, Word32 Rt) 464

sfcmp.ge
    Pd=sfcmp.ge(Rs,Rt)
    Byte Q6_p_sfcmp_ge_RR(Word32 Rs, Word32 Rt) 464

sfcmp.gt
    Pd=sfcmp.gt(Rs,Rt)
    Byte Q6_p_sfcmp_gt_RR(Word32 Rs, Word32 Rt) 464

sfcmp.uo
    Pd=sfcmp.uo(Rs,Rt)

```

Byte Q6\_p\_sfcmp\_uo\_RR(Word32 Rs, Word32 Rt) **464**

sffixupd  
Rd=sffixupd(Rs,Rt)  
Word32 Q6\_R\_sffixupd\_RR(Word32 Rs, Word32 Rt) **472**

sffixupn  
Rd=sffixupn(Rs,Rt)  
Word32 Q6\_R\_sffixupn\_RR(Word32 Rs, Word32 Rt) **472**

sffixupr  
Rd=sffixupr(Rs)  
Word32 Q6\_R\_sffixupr\_R(Word32 Rs) **472**

sfmake  
Rd=sfmake(#u10):neg  
Word32 Q6\_R\_sfmake\_I\_neg(Word32 Iu10) **478**  
Rd=sfmake(#u10):pos  
Word32 Q6\_R\_sfmake\_I\_pos(Word32 Iu10) **478**

sfmax  
Rd=sfmax(Rs,Rt)  
Word32 Q6\_R\_sfmax\_RR(Word32 Rs, Word32 Rt) **479**

sfmin  
Rd=sfmin(Rs,Rt)  
Word32 Q6\_R\_sfmin\_RR(Word32 Rs, Word32 Rt) **480**

sfmpy  
Rd=sfmpy(Rs,Rt)  
Word32 Q6\_R\_sfmpy\_RR(Word32 Rs, Word32 Rt) **481**  
Rx+=sfmpy(Rs,Rt,Pu):scale  
Word32 Q6\_R\_sfmpyacc\_RRp\_scale(Word32 Rx, Word32 Rs, Word32 Rt, Byte Pu) **474**  
Rx+=sfmpy(Rs,Rt)  
Word32 Q6\_R\_sfmpyacc\_RR(Word32 Rx, Word32 Rs, Word32 Rt) **473**  
Rx+=sfmpy(Rs,Rt):lib  
Word32 Q6\_R\_sfmpyacc\_RR\_lib(Word32 Rx, Word32 Rs, Word32 Rt) **476**  
Rx-=sfmpy(Rs,Rt)  
Word32 Q6\_R\_sfmpynac\_RR(Word32 Rx, Word32 Rs, Word32 Rt) **473**  
Rx-=sfmpy(Rs,Rt):lib  
Word32 Q6\_R\_sfmpynac\_RR\_lib(Word32 Rx, Word32 Rs, Word32 Rt) **476**

sfsb  
Rd=sfsb(Rs,Rt)  
Word32 Q6\_R\_sfsb\_RR(Word32 Rs, Word32 Rt) **483**

shuffeb  
Rdd=shuffeb(Rss,Rtt)  
Word64 Q6\_P\_shuffeb\_PP(Word64 Rss, Word64 Rtt) **556**

shuffeh  
Rdd=shuffeh(Rss,Rtt)  
Word64 Q6\_P\_shuffeh\_PP(Word64 Rss, Word64 Rtt) **556**

shuffob  
Rdd=shuffob(Rtt,Rss)  
Word64 Q6\_P\_shuffob\_PP(Word64 Rtt, Word64 Rss) **556**

shuffoh  
Rdd=shuffoh(Rtt,Rss)

Word64 Q6\_P\_shuffoh\_PP(Word64 Rtt, Word64 Rss) **556**

sub

Rd=add(Rs, sub(#s6, Ru))  
 Word32 Q6\_R\_add\_sub\_RIR(Word32 Rs, Word32 Is6, Word32 Ru) **335**

Rd=sub(#s10, Rs)  
 Word32 Q6\_R\_sub\_IR(Word32 Is10, Word32 Rs) **162**

Rd=sub(Rt, Rs)  
 Word32 Q6\_R\_sub\_RR(Word32 Rt, Word32 Rs) **162**

Rd=sub(Rt, Rs):sat  
 Word32 Q6\_R\_sub\_RR\_sat(Word32 Rt, Word32 Rs) **162**

Rd=sub(Rt.H, Rs.H):<<16  
 Word32 Q6\_R\_sub\_RhRh\_s16(Word32 Rt, Word32 Rs) **361**

Rd=sub(Rt.H, Rs.H):sat:<<16  
 Word32 Q6\_R\_sub\_RhRh\_sat\_s16(Word32 Rt, Word32 Rs) **361**

Rd=sub(Rt.H, Rs.L):<<16  
 Word32 Q6\_R\_sub\_RhRl\_s16(Word32 Rt, Word32 Rs) **361**

Rd=sub(Rt.H, Rs.L):sat:<<16  
 Word32 Q6\_R\_sub\_RhRl\_sat\_s16(Word32 Rt, Word32 Rs) **361**

Rd=sub(Rt.L, Rs.H)  
 Word32 Q6\_R\_sub\_RlRh(Word32 Rt, Word32 Rs) **361**

Rd=sub(Rt.L, Rs.H):<<16  
 Word32 Q6\_R\_sub\_RlRh\_s16(Word32 Rt, Word32 Rs) **361**

Rd=sub(Rt.L, Rs.H):sat  
 Word32 Q6\_R\_sub\_RlRh\_sat(Word32 Rt, Word32 Rs) **361**

Rd=sub(Rt.L, Rs.H):sat:<<16  
 Word32 Q6\_R\_sub\_RlRh\_sat\_s16(Word32 Rt, Word32 Rs) **361**

Rd=sub(Rt.L, Rs.L)  
 Word32 Q6\_R\_sub\_RlRl(Word32 Rt, Word32 Rs) **361**

Rd=sub(Rt.L, Rs.L):<<16  
 Word32 Q6\_R\_sub\_RlRl\_s16(Word32 Rt, Word32 Rs) **361**

Rd=sub(Rt.L, Rs.L):sat  
 Word32 Q6\_R\_sub\_RlRl\_sat(Word32 Rt, Word32 Rs) **361**

Rd=sub(Rt.L, Rs.L):sat:<<16  
 Word32 Q6\_R\_sub\_RlRl\_sat\_s16(Word32 Rt, Word32 Rs) **361**

Rdd=sub(Rtt, Rss)  
 Word64 Q6\_P\_sub\_PP(Word64 Rtt, Word64 Rss) **358**

Rx+=sub(Rt, Rs)  
 Word32 Q6\_R\_subacc\_RR(Word32 Rx, Word32 Rt, Word32 Rs) **359**

swiz

Rd=swiz(Rs)  
 Word32 Q6\_R\_swiz\_R(Word32 Rs) **544**

sxtb

Rd=sxtb(Rs)  
 Word32 Q6\_R\_sxtb\_R(Word32 Rs) **164**

sxth

Rd=sxth(Rs)  
 Word32 Q6\_R\_sxth\_R(Word32 Rs) **164**

sxtw

Rdd=sxtw(Rs)  
 Word64 Q6\_P\_sxtw\_R(Word32 Rs) **362**

**T**

tableidxb

Rx=tableidxb(Rs,#u4,#U5)

Word32 Q6\_R\_tableidxb\_RII(Word32 Rx, Word32 Rs, Word32 Iu4, Word32 IU5) **427**

tableidxd

Rx=tableidxd(Rs,#u4,#U5)

Word32 Q6\_R\_tableidxd\_RII(Word32 Rx, Word32 Rs, Word32 Iu4, Word32 IU5) **427**

tableidxh

Rx=tableidxh(Rs,#u4,#U5)

Word32 Q6\_R\_tableidxh\_RII(Word32 Rx, Word32 Rs, Word32 Iu4, Word32 IU5) **427**

tableidxw

Rx=tableidxw(Rs,#u4,#U5)

Word32 Q6\_R\_tableidxw\_RII(Word32 Rx, Word32 Rs, Word32 Iu4, Word32 IU5) **427**

tlbmatch

Pd=tlbmatch(Rss,Rt)

Byte Q6\_p\_tlbmatch\_PR(Word64 Rss, Word32 Rt) **576**

togglebit

Rd=togglebit(Rs,#u5)

Word32 Q6\_R\_togglebit\_RI(Word32 Rs, Word32 Iu5) **422**

Rd=togglebit(Rs,Rt)

Word32 Q6\_R\_togglebit\_RR(Word32 Rs, Word32 Rt) **422**

tstbit

Pd=!tstbit(Rs,#u5)

Byte Q6\_p\_not\_tstbit\_RI(Word32 Rs, Word32 Iu5) **578**

Pd=!tstbit(Rs,Rt)

Byte Q6\_p\_not\_tstbit\_RR(Word32 Rs, Word32 Rt) **578**

Pd=tstbit(Rs,#u5)

Byte Q6\_p\_tstbit\_RI(Word32 Rs, Word32 Iu5) **578**

Pd=tstbit(Rs,Rt)

Byte Q6\_p\_tstbit\_RR(Word32 Rs, Word32 Rt) **578****V**

vabsdiffb

Rdd=vabsdiffb(Rtt,Rss)

Word64 Q6\_P\_vabsdiffb\_PP(Word64 Rtt, Word64 Rss) **365**

vabsdiffh

Rdd=vabsdiffh(Rtt,Rss)

Word64 Q6\_P\_vabsdiffh\_PP(Word64 Rtt, Word64 Rss) **366**

vabsdiffub

Rdd=vabsdiffub(Rtt,Rss)

Word64 Q6\_P\_vabsdiffub\_PP(Word64 Rtt, Word64 Rss) **365**

vabsdiffw

Rdd=vabsdiffw(Rtt,Rss)

Word64 Q6\_P\_vabsdiffw\_PP(Word64 Rtt, Word64 Rss) **367**

vabsh

Rdd=vabsh(Rss)

Word64 Q6\_P\_vabsh\_P(Word64 Rss) **363**

Rdd=vabsh(Rss):sat

Word64 Q6\_P\_vabsh\_P\_sat(Word64 Rss) **363**

vabsw  
 Rdd=vabsw(Rss)  
 Word64 Q6\_P\_vabsw\_P(Word64 Rss) **364**  
 Rdd=vabsw(Rss):sat  
 Word64 Q6\_P\_vabsw\_P\_sat(Word64 Rss) **364**

vaddb  
 Rdd=vaddb(Rss,Rtt)  
 Word64 Q6\_P\_vaddb\_PP(Word64 Rss, Word64 Rtt) **378**

vaddh  
 Rd=vaddh(Rs,Rt)  
 Word32 Q6\_R\_vaddh\_RR(Word32 Rs, Word32 Rt) **168**  
 Rd=vaddh(Rs,Rt):sat  
 Word32 Q6\_R\_vaddh\_RR\_sat(Word32 Rs, Word32 Rt) **168**  
 Rdd=vaddh(Rss,Rtt)  
 Word64 Q6\_P\_vaddh\_PP(Word64 Rss, Word64 Rtt) **371**  
 Rdd=vaddh(Rss,Rtt):sat  
 Word64 Q6\_P\_vaddh\_PP\_sat(Word64 Rss, Word64 Rtt) **371**

vaddhub  
 Rd=vaddhub(Rss,Rtt):sat  
 Word32 Q6\_R\_vaddhub\_PP\_sat(Word64 Rss, Word64 Rtt) **373**

vaddub  
 Rdd=vaddub(Rss,Rtt)  
 Word64 Q6\_P\_vaddub\_PP(Word64 Rss, Word64 Rtt) **378**  
 Rdd=vaddub(Rss,Rtt):sat  
 Word64 Q6\_P\_vaddub\_PP\_sat(Word64 Rss, Word64 Rtt) **378**

vadduh  
 Rd=vadduh(Rs,Rt):sat  
 Word32 Q6\_R\_vadduh\_RR\_sat(Word32 Rs, Word32 Rt) **168**  
 Rdd=vadduh(Rss,Rtt):sat  
 Word64 Q6\_P\_vadduh\_PP\_sat(Word64 Rss, Word64 Rtt) **371**

vaddw  
 Rdd=vaddw(Rss,Rtt)  
 Word64 Q6\_P\_vaddw\_PP(Word64 Rss, Word64 Rtt) **379**  
 Rdd=vaddw(Rss,Rtt):sat  
 Word64 Q6\_P\_vaddw\_PP\_sat(Word64 Rss, Word64 Rtt) **379**

valignb  
 Rdd=valignb(Rtt,Rss,#u3)  
 Word64 Q6\_P\_valignb\_PPI(Word64 Rtt, Word64 Rss, Word32 Iu3) **545**  
 Rdd=valignb(Rtt,Rss,Pu)  
 Word64 Q6\_P\_valignb\_PPp(Word64 Rtt, Word64 Rss, Byte Pu) **545**

vaslh  
 Rdd=vaslh(Rss,#u4)  
 Word64 Q6\_P\_vaslh\_PI(Word64 Rss, Word32 Iu4) **611**  
 Rdd=vaslh(Rss,Rt)  
 Word64 Q6\_P\_vaslh\_PR(Word64 Rss, Word32 Rt) **616**

vaslw  
 Rdd=vaslw(Rss,#u5)  
 Word64 Q6\_P\_vaslw\_PI(Word64 Rss, Word32 Iu5) **617**  
 Rdd=vaslw(Rss,Rt)

Word64 Q6\_P\_vaslw\_PR(Word64 Rss, Word32 Rt) **618**

#### vasrh

Rdd=vasrh(Rss,#u4)  
 Word64 Q6\_P\_vasrh\_PI(Word64 Rss, Word32 Iu4) **611**  
 Rdd=vasrh(Rss,#u4):rnd  
 Word64 Q6\_P\_vasrh\_PI\_rnd(Word64 Rss, Word32 Iu4) **612**  
 Rdd=vasrh(Rss,Rt)  
 Word64 Q6\_P\_vasrh\_PR(Word64 Rss, Word32 Rt) **616**

#### vasrhub

Rd=vasrhub(Rss,#u4):rnd:sat  
 Word32 Q6\_R\_vasrhub\_PI\_rnd\_sat(Word64 Rss, Word32 Iu4) **614**  
 Rd=vasrhub(Rss,#u4):sat  
 Word32 Q6\_R\_vasrhub\_PI\_sat(Word64 Rss, Word32 Iu4) **614**

#### vasrw

Rd=vasrw(Rss,#u5)  
 Word32 Q6\_R\_vasrw\_PI(Word64 Rss, Word32 Iu5) **620**  
 Rd=vasrw(Rss,Rt)  
 Word32 Q6\_R\_vasrw\_PR(Word64 Rss, Word32 Rt) **620**  
 Rdd=vasrw(Rss,#u5)  
 Word64 Q6\_P\_vasrw\_PI(Word64 Rss, Word32 Iu5) **617**  
 Rdd=vasrw(Rss,Rt)  
 Word64 Q6\_P\_vasrw\_PR(Word64 Rss, Word32 Rt) **618**

#### vavgh

Rd=vavgh(Rs,Rt)  
 Word32 Q6\_R\_vavgh\_RR(Word32 Rs, Word32 Rt) **169**  
 Rd=vavgh(Rs,Rt):rnd  
 Word32 Q6\_R\_vavgh\_RR\_rnd(Word32 Rs, Word32 Rt) **169**  
 Rdd=vavgh(Rss,Rtt)  
 Word64 Q6\_P\_vavgh\_PP(Word64 Rss, Word64 Rtt) **381**  
 Rdd=vavgh(Rss,Rtt):crnd  
 Word64 Q6\_P\_vavgh\_PP\_crnd(Word64 Rss, Word64 Rtt) **381**  
 Rdd=vavgh(Rss,Rtt):rnd  
 Word64 Q6\_P\_vavgh\_PP\_rnd(Word64 Rss, Word64 Rtt) **381**

#### vavgub

Rdd=vavgub(Rss,Rtt)  
 Word64 Q6\_P\_vavgub\_PP(Word64 Rss, Word64 Rtt) **382**  
 Rdd=vavgub(Rss,Rtt):rnd  
 Word64 Q6\_P\_vavgub\_PP\_rnd(Word64 Rss, Word64 Rtt) **382**

#### vavguh

Rdd=vavguh(Rss,Rtt)  
 Word64 Q6\_P\_vavguh\_PP(Word64 Rss, Word64 Rtt) **381**  
 Rdd=vavguh(Rss,Rtt):rnd  
 Word64 Q6\_P\_vavguh\_PP\_rnd(Word64 Rss, Word64 Rtt) **381**

#### vavguw

Rdd=vavguw(Rss,Rtt)  
 Word64 Q6\_P\_vavguw\_PP(Word64 Rss, Word64 Rtt) **384**  
 Rdd=vavguw(Rss,Rtt):rnd  
 Word64 Q6\_P\_vavguw\_PP\_rnd(Word64 Rss, Word64 Rtt) **384**

#### vavgw

Rdd=vavgw(Rss,Rtt)



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    Word64 Q6_P_vavgw_PP(Word64 Rss, Word64 Rtt) 384
    Rdd=vavgw(Rss,Rtt):crnd
    Word64 Q6_P_vavgw_PP_crnd(Word64 Rss, Word64 Rtt) 384
    Rdd=vavgw(Rss,Rtt):rnd
    Word64 Q6_P_vavgw_PP_rnd(Word64 Rss, Word64 Rtt) 384

vclip
    Rdd=vclip(Rss,#u5)
    Word64 Q6_P_vclip_PI(Word64 Rss, Word32 Iu5) 385

vcmpb.eq
    Pd=!any8(vcmpb.eq(Rss,Rtt))
    Byte Q6_p_not_any8_vcmpb_eq_PP(Word64 Rss, Word64 Rtt) 581
    Pd=any8(vcmpb.eq(Rss,Rtt))
    Byte Q6_p_any8_vcmpb_eq_PP(Word64 Rss, Word64 Rtt) 581
    Pd=vcmpb.eq(Rss,#u8)
    Byte Q6_p_vcmpb_eq_PI(Word64 Rss, Word32 Iu8) 583
    Pd=vcmpb.eq(Rss,Rtt)
    Byte Q6_p_vcmpb_eq_PP(Word64 Rss, Word64 Rtt) 583

vcmpb.gt
    Pd=vcmpb.gt(Rss,#s8)
    Byte Q6_p_vcmpb_gt_PI(Word64 Rss, Word32 Is8) 583
    Pd=vcmpb.gt(Rss,Rtt)
    Byte Q6_p_vcmpb_gt_PP(Word64 Rss, Word64 Rtt) 583

vcmpb.gtu
    Pd=vcmpb.gtu(Rss,#u7)
    Byte Q6_p_vcmpb_gtu_PI(Word64 Rss, Word32 Iu7) 583
    Pd=vcmpb.gtu(Rss,Rtt)
    Byte Q6_p_vcmpb_gtu_PP(Word64 Rss, Word64 Rtt) 583

vcmph.eq
    Pd=vcmph.eq(Rss,#s8)
    Byte Q6_p_vcmph_eq_PI(Word64 Rss, Word32 Is8) 580
    Pd=vcmph.eq(Rss,Rtt)
    Byte Q6_p_vcmph_eq_PP(Word64 Rss, Word64 Rtt) 580

vcmph.gt
    Pd=vcmph.gt(Rss,#s8)
    Byte Q6_p_vcmph_gt_PI(Word64 Rss, Word32 Is8) 580
    Pd=vcmph.gt(Rss,Rtt)
    Byte Q6_p_vcmph_gt_PP(Word64 Rss, Word64 Rtt) 580

vcmph.gtu
    Pd=vcmph.gtu(Rss,#u7)
    Byte Q6_p_vcmph_gtu_PI(Word64 Rss, Word32 Iu7) 580
    Pd=vcmph.gtu(Rss,Rtt)
    Byte Q6_p_vcmph_gtu_PP(Word64 Rss, Word64 Rtt) 580

vcmpw.eq
    Pd=vcmpw.eq(Rss,#s8)
    Byte Q6_p_vcmpw_eq_PI(Word64 Rss, Word32 Is8) 585
    Pd=vcmpw.eq(Rss,Rtt)
    Byte Q6_p_vcmpw_eq_PP(Word64 Rss, Word64 Rtt) 585

vcmpw.gt
    Pd=vcmpw.gt(Rss,#s8)

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    Byte Q6_p_vcmpw_gt_PI(Word64 Rss, Word32 Is8) 585
    Pd=vcmpw.gt(Rss,Rtt)
    Byte Q6_p_vcmpw_gt_PP(Word64 Rss, Word64 Rtt) 585

vcmpw.gtu
    Pd=vcmpw.gtu(Rss,#u7)
    Byte Q6_p_vcmpw_gtu_PI(Word64 Rss, Word32 Iu7) 585
    Pd=vcmpw.gtu(Rss,Rtt)
    Byte Q6_p_vcmpw_gtu_PP(Word64 Rss, Word64 Rtt) 585

vcmpyi
    Rdd=vcmpyi(Rss,Rtt):<<1:sat
    Word64 Q6_P_vcmpyi_PP_sl_sat(Word64 Rss, Word64 Rtt) 448
    Rdd=vcmpyi(Rss,Rtt):sat
    Word64 Q6_P_vcmpyi_PP_sat(Word64 Rss, Word64 Rtt) 448
    Rxx+=vcmpyi(Rss,Rtt):sat
    Word64 Q6_P_vcmpyiacc_PP_sat(Word64 Rxx, Word64 Rss, Word64 Rtt) 448

vcmpyr
    Rdd=vcmpyr(Rss,Rtt):<<1:sat
    Word64 Q6_P_vcmpyr_PP_sl_sat(Word64 Rss, Word64 Rtt) 448
    Rdd=vcmpyr(Rss,Rtt):sat
    Word64 Q6_P_vcmpyr_PP_sat(Word64 Rss, Word64 Rtt) 448
    Rxx+=vcmpyr(Rss,Rtt):sat
    Word64 Q6_P_vcmpyracc_PP_sat(Word64 Rxx, Word64 Rss, Word64 Rtt) 448

vcnegh
    Rdd=vcnegh(Rss,Rt)
    Word64 Q6_P_vcnegh_PR(Word64 Rss, Word32 Rt) 386

vconj
    Rdd=vconj(Rss):sat
    Word64 Q6_P_vconj_P_sat(Word64 Rss) 450

vcrotate
    Rdd=vcrotate(Rss,Rt)
    Word64 Q6_P_vcrotate_PR(Word64 Rss, Word32 Rt) 452

vdmpy
    Rd=vdmpy(Rss,Rtt):<<1:rnd:sat
    Word32 Q6_R_vdmpy_PP_sl_rnd_sat(Word64 Rss, Word64 Rtt) 521
    Rd=vdmpy(Rss,Rtt):rnd:sat
    Word32 Q6_R_vdmpy_PP_rnd_sat(Word64 Rss, Word64 Rtt) 521
    Rdd=vdmpy(Rss,Rtt):<<1:sat
    Word64 Q6_P_vdmpy_PP_sl_sat(Word64 Rss, Word64 Rtt) 518
    Rdd=vdmpy(Rss,Rtt):sat
    Word64 Q6_P_vdmpy_PP_sat(Word64 Rss, Word64 Rtt) 518
    Rxx+=vdmpy(Rss,Rtt):<<1:sat
    Word64 Q6_P_vdmpyacc_PP_sl_sat(Word64 Rxx, Word64 Rss, Word64 Rtt) 518
    Rxx+=vdmpy(Rss,Rtt):sat
    Word64 Q6_P_vdmpyacc_PP_sat(Word64 Rxx, Word64 Rss, Word64 Rtt) 518

vdmpybsu
    Rdd=vdmpybsu(Rss,Rtt):sat
    Word64 Q6_P_vdmpybsu_PP_sat(Word64 Rss, Word64 Rtt) 525
    Rxx+=vdmpybsu(Rss,Rtt):sat
    Word64 Q6_P_vdmpybsuacc_PP_sat(Word64 Rxx, Word64 Rss, Word64 Rtt) 525

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vitpack
  Rdd=vitpack(Ps,Pt)
  Word32 Q6_R_vitpack_PP(Byte Ps, Byte Pt) 586

vlslh
  Rdd=vlslh(Rss,Rt)
  Word64 Q6_P_vlslh_PR(Word64 Rss, Word32 Rt) 616

vlslw
  Rdd=vlslw(Rss,Rt)
  Word64 Q6_P_vlslw_PR(Word64 Rss, Word32 Rt) 618

vlsrh
  Rdd=vlsrh(Rss,#u4)
  Word64 Q6_P_vlsrh_PI(Word64 Rss, Word32 Iu4) 611
  Rdd=vlsrh(Rss,Rt)
  Word64 Q6_P_vlsrh_PR(Word64 Rss, Word32 Rt) 616

vlsrw
  Rdd=vlsrw(Rss,#u5)
  Word64 Q6_P_vlsrw_PI(Word64 Rss, Word32 Iu5) 617
  Rdd=vlsrw(Rss,Rt)
  Word64 Q6_P_vlsrw_PR(Word64 Rss, Word32 Rt) 618

vmaxb
  Rdd=vmaxb(Rtt,Rss)
  Word64 Q6_P_vmaxb_PP(Word64 Rtt, Word64 Rss) 388

vmaxh
  Rdd=vmaxh(Rtt,Rss)
  Word64 Q6_P_vmaxh_PP(Word64 Rtt, Word64 Rss) 389

vmaxub
  Rdd=vmaxub(Rtt,Rss)
  Word64 Q6_P_vmaxub_PP(Word64 Rtt, Word64 Rss) 388

vmaxuh
  Rdd=vmaxuh(Rtt,Rss)
  Word64 Q6_P_vmaxuh_PP(Word64 Rtt, Word64 Rss) 389

vmaxuw
  Rdd=vmaxuw(Rtt,Rss)
  Word64 Q6_P_vmaxuw_PP(Word64 Rtt, Word64 Rss) 394

vmaxw
  Rdd=vmaxw(Rtt,Rss)
  Word64 Q6_P_vmaxw_PP(Word64 Rtt, Word64 Rss) 394

vminb
  Rdd=vminb(Rtt,Rss)
  Word64 Q6_P_vminb_PP(Word64 Rtt, Word64 Rss) 395

vminh
  Rdd=vminh(Rtt,Rss)
  Word64 Q6_P_vminh_PP(Word64 Rtt, Word64 Rss) 397

vminub
  Rdd=vminub(Rtt,Rss)
  Word64 Q6_P_vminub_PP(Word64 Rtt, Word64 Rss) 395
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vminuh
  Rdd=vminuh(Rtt,Rss)
  Word64 Q6_P_vminuh_PP(Word64 Rtt, Word64 Rss) 397

vminuw
  Rdd=vminuw(Rtt,Rss)
  Word64 Q6_P_vminuw_PP(Word64 Rtt, Word64 Rss) 402

vminw
  Rdd=vminw(Rtt,Rss)
  Word64 Q6_P_vminw_PP(Word64 Rtt, Word64 Rss) 402

vmpybsu
  Rdd=vmpybsu(Rs,Rt)
  Word64 Q6_P_vmpybsu_RR(Word32 Rs, Word32 Rt) 537
  Rxx+=vmpybsu(Rs,Rt)
  Word64 Q6_P_vmpybsuacc_RR(Word64 Rxx, Word32 Rs, Word32 Rt) 537

vmpybu
  Rdd=vmpybu(Rs,Rt)
  Word64 Q6_P_vmpybu_RR(Word32 Rs, Word32 Rt) 537
  Rxx+=vmpybu(Rs,Rt)
  Word64 Q6_P_vmpybuacc_RR(Word64 Rxx, Word32 Rs, Word32 Rt) 537

vmpyeh
  Rdd=vmpyeh(Rss,Rtt):<<1:sat
  Word64 Q6_P_vmpyeh_PP_s1_sat(Word64 Rss, Word64 Rtt) 527
  Rdd=vmpyeh(Rss,Rtt):sat
  Word64 Q6_P_vmpyeh_PP_sat(Word64 Rss, Word64 Rtt) 527
  Rxx+=vmpyeh(Rss,Rtt)
  Word64 Q6_P_vmpyehacc_PP(Word64 Rxx, Word64 Rss, Word64 Rtt) 527
  Rxx+=vmpyeh(Rss,Rtt):<<1:sat
  Word64 Q6_P_vmpyehacc_PP_s1_sat(Word64 Rxx, Word64 Rss, Word64 Rtt) 527
  Rxx+=vmpyeh(Rss,Rtt):sat
  Word64 Q6_P_vmpyehacc_PP_sat(Word64 Rxx, Word64 Rss, Word64 Rtt) 527

vmpyh
  Rd=vmpyh(Rs,Rt):<<1:rnd:sat
  Word32 Q6_R_vmpyh_RR_s1_rnd_sat(Word32 Rs, Word32 Rt) 531
  Rd=vmpyh(Rs,Rt):rnd:sat
  Word32 Q6_R_vmpyh_RR_rnd_sat(Word32 Rs, Word32 Rt) 531
  Rdd=vmpyh(Rs,Rt):<<1:sat
  Word64 Q6_P_vmpyh_RR_s1_sat(Word32 Rs, Word32 Rt) 529
  Rdd=vmpyh(Rs,Rt):sat
  Word64 Q6_P_vmpyh_RR_sat(Word32 Rs, Word32 Rt) 529
  Rxx+=vmpyh(Rs,Rt)
  Word64 Q6_P_vmpyhacc_RR(Word64 Rxx, Word32 Rs, Word32 Rt) 529
  Rxx+=vmpyh(Rs,Rt):<<1:sat
  Word64 Q6_P_vmpyhacc_RR_s1_sat(Word64 Rxx, Word32 Rs, Word32 Rt) 529
  Rxx+=vmpyh(Rs,Rt):sat
  Word64 Q6_P_vmpyhacc_RR_sat(Word64 Rxx, Word32 Rs, Word32 Rt) 529

vmpyhsu
  Rdd=vmpyhsu(Rs,Rt):<<1:sat
  Word64 Q6_P_vmpyhsu_RR_s1_sat(Word32 Rs, Word32 Rt) 532
  Rdd=vmpyhsu(Rs,Rt):sat
  Word64 Q6_P_vmpyhsu_RR_sat(Word32 Rs, Word32 Rt) 532
  Rxx+=vmpyhsu(Rs,Rt):<<1:sat

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Word64 Q6_P_vmpyhsuacc_RR_sl_sat(Word64 Rxx, Word32 Rs, Word32 Rt) 532
Rxx+=vmpyhsu(Rs,Rt):sat
Word64 Q6_P_vmpyhsuacc_RR_sat(Word64 Rxx, Word32 Rs, Word32 Rt) 532

vmpyweh
Rdd=vmpyweh(Rss,Rtt):<<1:rnd:sat
Word64 Q6_P_vmpyweh_PP_sl_rnd_sat(Word64 Rss, Word64 Rtt) 489
Rdd=vmpyweh(Rss,Rtt):<<1:sat
Word64 Q6_P_vmpyweh_PP_sl_sat(Word64 Rss, Word64 Rtt) 489
Rdd=vmpyweh(Rss,Rtt):rnd:sat
Word64 Q6_P_vmpyweh_PP_rnd_sat(Word64 Rss, Word64 Rtt) 489
Rdd=vmpyweh(Rss,Rtt):sat
Word64 Q6_P_vmpyweh_PP_sat(Word64 Rss, Word64 Rtt) 489
Rxx+=vmpyweh(Rss,Rtt):<<1:rnd:sat
Word64 Q6_P_vmpywehacc_PP_sl_rnd_sat(Word64 Rxx, Word64 Rss, Word64 Rtt)
490
Rxx+=vmpyweh(Rss,Rtt):<<1:sat
Word64 Q6_P_vmpywehacc_PP_sl_sat(Word64 Rxx, Word64 Rss, Word64 Rtt) 490
Rxx+=vmpyweh(Rss,Rtt):rnd:sat
Word64 Q6_P_vmpywehacc_PP_rnd_sat(Word64 Rxx, Word64 Rss, Word64 Rtt) 490
Rxx+=vmpyweh(Rss,Rtt):sat
Word64 Q6_P_vmpywehacc_PP_sat(Word64 Rxx, Word64 Rss, Word64 Rtt) 490

vmpyweuh
Rdd=vmpyweuh(Rss,Rtt):<<1:rnd:sat
Word64 Q6_P_vmpyweuh_PP_sl_rnd_sat(Word64 Rss, Word64 Rtt) 493
Rdd=vmpyweuh(Rss,Rtt):<<1:sat
Word64 Q6_P_vmpyweuh_PP_sl_sat(Word64 Rss, Word64 Rtt) 493
Rdd=vmpyweuh(Rss,Rtt):rnd:sat
Word64 Q6_P_vmpyweuh_PP_rnd_sat(Word64 Rss, Word64 Rtt) 493
Rdd=vmpyweuh(Rss,Rtt):sat
Word64 Q6_P_vmpyweuh_PP_sat(Word64 Rss, Word64 Rtt) 494
Rxx+=vmpyweuh(Rss,Rtt):<<1:rnd:sat
Word64 Q6_P_vmpyweuhacc_PP_sl_rnd_sat(Word64 Rxx, Word64 Rss, Word64 Rtt)
494
Rxx+=vmpyweuh(Rss,Rtt):<<1:sat
Word64 Q6_P_vmpyweuhacc_PP_sl_sat(Word64 Rxx, Word64 Rss, Word64 Rtt) 494
Rxx+=vmpyweuh(Rss,Rtt):rnd:sat
Word64 Q6_P_vmpyweuhacc_PP_rnd_sat(Word64 Rxx, Word64 Rss, Word64 Rtt) 494
Rxx+=vmpyweuh(Rss,Rtt):sat
Word64 Q6_P_vmpyweuhacc_PP_sat(Word64 Rxx, Word64 Rss, Word64 Rtt) 494

vmpywoh
Rdd=vmpywoh(Rss,Rtt):<<1:rnd:sat
Word64 Q6_P_vmpywoh_PP_sl_rnd_sat(Word64 Rss, Word64 Rtt) 490
Rdd=vmpywoh(Rss,Rtt):<<1:sat
Word64 Q6_P_vmpywoh_PP_sl_sat(Word64 Rss, Word64 Rtt) 490
Rdd=vmpywoh(Rss,Rtt):rnd:sat
Word64 Q6_P_vmpywoh_PP_rnd_sat(Word64 Rss, Word64 Rtt) 490
Rdd=vmpywoh(Rss,Rtt):sat
Word64 Q6_P_vmpywoh_PP_sat(Word64 Rss, Word64 Rtt) 490
Rxx+=vmpywoh(Rss,Rtt):<<1:rnd:sat
Word64 Q6_P_vmpywohacc_PP_sl_rnd_sat(Word64 Rxx, Word64 Rss, Word64 Rtt)
490
Rxx+=vmpywoh(Rss,Rtt):<<1:sat
Word64 Q6_P_vmpywohacc_PP_sl_sat(Word64 Rxx, Word64 Rss, Word64 Rtt) 490
Rxx+=vmpywoh(Rss,Rtt):rnd:sat
Word64 Q6_P_vmpywohacc_PP_rnd_sat(Word64 Rxx, Word64 Rss, Word64 Rtt) 490
Rxx+=vmpywoh(Rss,Rtt):sat

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Word64 Q6\_P\_vmpywouhacc\_PP\_sat(Word64 Rxx, Word64 Rss, Word64 Rtt) **490**

#### vmpywouh

Rdd=vmpywouh(Rss,Rtt):<<1:rnd:sat

Word64 Q6\_P\_vmpywouh\_PP\_sl\_rnd\_sat(Word64 Rss, Word64 Rtt) **494**

Rdd=vmpywouh(Rss,Rtt):<<1:sat

Word64 Q6\_P\_vmpywouh\_PP\_sl\_sat(Word64 Rss, Word64 Rtt) **494**

Rdd=vmpywouh(Rss,Rtt):rnd:sat

Word64 Q6\_P\_vmpywouh\_PP\_rnd\_sat(Word64 Rss, Word64 Rtt) **494**

Rdd=vmpywouh(Rss,Rtt):sat

Word64 Q6\_P\_vmpywouh\_PP\_sat(Word64 Rss, Word64 Rtt) **494**

Rxx+=vmpywouh(Rss,Rtt):<<1:rnd:sat

Word64 Q6\_P\_vmpywouhacc\_PP\_sl\_rnd\_sat(Word64 Rxx, Word64 Rss, Word64 Rtt) **494**

Rxx+=vmpywouh(Rss,Rtt):<<1:sat

Word64 Q6\_P\_vmpywouhacc\_PP\_sl\_sat(Word64 Rxx, Word64 Rss, Word64 Rtt) **494**

Rxx+=vmpywouh(Rss,Rtt):rnd:sat

Word64 Q6\_P\_vmpywouhacc\_PP\_rnd\_sat(Word64 Rxx, Word64 Rss, Word64 Rtt) **494**

Rxx+=vmpywouh(Rss,Rtt):sat

Word64 Q6\_P\_vmpywouhacc\_PP\_sat(Word64 Rxx, Word64 Rss, Word64 Rtt) **494**

#### vmux

Rdd=vmux(Pu,Rss,Rtt)

Word64 Q6\_P\_vmux\_pPP(Byte Pu, Word64 Rss, Word64 Rtt) **587**

#### vnavgh

Rd=vnavgh(Rt,Rs)

Word32 Q6\_R\_vnavgh\_RR(Word32 Rt, Word32 Rs) **169**

Rdd=vnavgh(Rtt,Rss)

Word64 Q6\_P\_vnavgh\_PP(Word64 Rtt, Word64 Rss) **381**

Rdd=vnavgh(Rtt,Rss):crnd:sat

Word64 Q6\_P\_vnavgh\_PP\_crnd\_sat(Word64 Rtt, Word64 Rss) **381**

Rdd=vnavgh(Rtt,Rss):rnd:sat

Word64 Q6\_P\_vnavgh\_PP\_rnd\_sat(Word64 Rtt, Word64 Rss) **381**

#### vnavgw

Rdd=vnavgw(Rtt,Rss)

Word64 Q6\_P\_vnavgw\_PP(Word64 Rtt, Word64 Rss) **384**

Rdd=vnavgw(Rtt,Rss):crnd:sat

Word64 Q6\_P\_vnavgw\_PP\_crnd\_sat(Word64 Rtt, Word64 Rss) **384**

Rdd=vnavgw(Rtt,Rss):rnd:sat

Word64 Q6\_P\_vnavgw\_PP\_rnd\_sat(Word64 Rtt, Word64 Rss) **384**

#### vpmpyh

Rdd=vpmpyh(Rs,Rt)

Word64 Q6\_P\_vpmpyh\_RR(Word32 Rs, Word32 Rt) **539**

Rxx^=vpmpyh(Rs,Rt)

Word64 Q6\_P\_vpmpyhxacc\_RR(Word64 Rxx, Word32 Rs, Word32 Rt) **539**

#### vradddh

Rd=vradddh(Rss,Rtt)

Word32 Q6\_R\_vradddh\_PP(Word64 Rss, Word64 Rtt) **376**

#### vradddub

Rdd=vradddub(Rss,Rtt)

Word64 Q6\_P\_vradddub\_PP(Word64 Rss, Word64 Rtt) **374**

Rxx+=vradddub(Rss,Rtt)

Word64 Q6\_P\_vradddubacc\_PP(Word64 Rxx, Word64 Rss, Word64 Rtt) **374**

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vradduh
  Rd=vradduh(Rss,Rtt)
    Word32 Q6_R_vradduh_PP(Word64 Rss, Word64 Rtt) 376

vrcmpys
  Rd=vrcmpys(Rss,Rt):<<1:rnd:sat
    Word32 Q6_R_vrcmpys_PR_s1_rnd_sat(Word64 Rss, Word32 Rt) 457
  Rdd=vrcmpys(Rss,Rt):<<1:sat
    Word64 Q6_P_vrcmpys_PR_s1_sat(Word64 Rss, Word32 Rt) 454
  Rxx+=vrcmpys(Rss,Rt):<<1:sat
    Word64 Q6_P_vrcmpysacc_PR_s1_sat(Word64 Rxx, Word64 Rss, Word32 Rt) 454

vrcnegh
  Rxx+=vrcnegh(Rss,Rt)
    Word64 Q6_P_vrcneghacc_PR(Word64 Rxx, Word64 Rss, Word32 Rt) 386

vrcrotate
  Rdd=vrcrotate(Rss,Rt,#u2)
    Word64 Q6_P_vrcrotate_PRI(Word64 Rss, Word32 Rt, Word32 Iu2) 460
  Rxx+=vrcrotate(Rss,Rt,#u2)
    Word64 Q6_P_vrcrotateacc_PRI(Word64 Rxx, Word64 Rss, Word32 Rt, Word32 Iu2)
    460

vrmaxh
  Rxx=vrmaxh(Rss,Ru)
    Word64 Q6_P_vrmaxh_PR(Word64 Rxx, Word64 Rss, Word32 Ru) 390

vrmaxuh
  Rxx=vrmaxuh(Rss,Ru)
    Word64 Q6_P_vrmaxuh_PR(Word64 Rxx, Word64 Rss, Word32 Ru) 390

vrmaxuw
  Rxx=vrmaxuw(Rss,Ru)
    Word64 Q6_P_vrmaxuw_PR(Word64 Rxx, Word64 Rss, Word32 Ru) 392

vrmaxw
  Rxx=vrmaxw(Rss,Ru)
    Word64 Q6_P_vrmaxw_PR(Word64 Rxx, Word64 Rss, Word32 Ru) 392

vrminh
  Rxx=vrminh(Rss,Ru)
    Word64 Q6_P_vrminh_PR(Word64 Rxx, Word64 Rss, Word32 Ru) 398

vrminuh
  Rxx=vrminuh(Rss,Ru)
    Word64 Q6_P_vrminuh_PR(Word64 Rxx, Word64 Rss, Word32 Ru) 398

vrminuw
  Rxx=vrminuw(Rss,Ru)
    Word64 Q6_P_vrminuw_PR(Word64 Rxx, Word64 Rss, Word32 Ru) 400

vrminw
  Rxx=vrminw(Rss,Ru)
    Word64 Q6_P_vrminw_PR(Word64 Rxx, Word64 Rss, Word32 Ru) 400

vrmpybsu
  Rdd=vrmpybsu(Rss,Rtt)
    Word64 Q6_P_vrmpybsu_PP(Word64 Rss, Word64 Rtt) 523
  Rxx+=vrmpybsu(Rss,Rtt)
    Word64 Q6_P_vrmpybsuacc_PP(Word64 Rxx, Word64 Rss, Word64 Rtt) 523

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vrmpybu
  Rdd=vrmpybu (Rss,Rtt)
    Word64 Q6_P_vrmpybu_PP(Word64 Rss, Word64 Rtt) 523
  Rxx+=vrmpybu (Rss,Rtt)
    Word64 Q6_P_vrmpybuacc_PP(Word64 Rxx, Word64 Rss, Word64 Rtt) 523

vrmpyh
  Rdd=vrmpyh (Rss,Rtt)
    Word64 Q6_P_vrmpyh_PP(Word64 Rss, Word64 Rtt) 534
  Rxx+=vrmpyh (Rss,Rtt)
    Word64 Q6_P_vrmpyhacc_PP(Word64 Rxx, Word64 Rss, Word64 Rtt) 534

vrmpyweh
  Rdd=vrmpyweh (Rss,Rtt)
    Word64 Q6_P_vrmpyweh_PP(Word64 Rss, Word64 Rtt) 511
  Rdd=vrmpyweh (Rss,Rtt) :<<1
    Word64 Q6_P_vrmpyweh_PP_s1(Word64 Rss, Word64 Rtt) 511
  Rxx+=vrmpyweh (Rss,Rtt)
    Word64 Q6_P_vrmpywehacc_PP(Word64 Rxx, Word64 Rss, Word64 Rtt) 511
  Rxx+=vrmpyweh (Rss,Rtt) :<<1
    Word64 Q6_P_vrmpywehacc_PP_s1(Word64 Rxx, Word64 Rss, Word64 Rtt) 511

vrmpywoh
  Rdd=vrmpywoh (Rss,Rtt)
    Word64 Q6_P_vrmpywoh_PP(Word64 Rss, Word64 Rtt) 511
  Rdd=vrmpywoh (Rss,Rtt) :<<1
    Word64 Q6_P_vrmpywoh_PP_s1(Word64 Rss, Word64 Rtt) 511
  Rxx+=vrmpywoh (Rss,Rtt)
    Word64 Q6_P_vrmpywohacc_PP(Word64 Rxx, Word64 Rss, Word64 Rtt) 511
  Rxx+=vrmpywoh (Rss,Rtt) :<<1
    Word64 Q6_P_vrmpywohacc_PP_s1(Word64 Rxx, Word64 Rss, Word64 Rtt) 511

vrndwh
  Rd=vrndwh (Rss)
    Word32 Q6_R_vrndwh_P(Word64 Rss) 547
  Rd=vrndwh (Rss) :sat
    Word32 Q6_R_vrndwh_P_sat(Word64 Rss) 547

vrsadub
  Rdd=vrsadub (Rss,Rtt)
    Word64 Q6_P_vrsadub_PP(Word64 Rss, Word64 Rtt) 404
  Rxx+=vrsadub (Rss,Rtt)
    Word64 Q6_P_vrsadubacc_PP(Word64 Rxx, Word64 Rss, Word64 Rtt) 404

vsathb
  Rd=vsathb (Rs)
    Word32 Q6_R_vsathb_R(Word32 Rs) 550
  Rd=vsathb (Rss)
    Word32 Q6_R_vsathb_P(Word64 Rss) 550
  Rdd=vsathb (Rss)
    Word64 Q6_P_vsathb_P(Word64 Rss) 553

vsathub
  Rd=vsathub (Rs)
    Word32 Q6_R_vsathub_R(Word32 Rs) 550
  Rd=vsathub (Rss)
    Word32 Q6_R_vsathub_P(Word64 Rss) 550
  Rdd=vsathub (Rss)
    Word64 Q6_P_vsathub_P(Word64 Rss) 553

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vsatwh
  Rd=vsatwh(Rss)
    Word32 Q6_R_vsathw_P(Word64 Rss) 550
  Rdd=vsatwh(Rss)
    Word64 Q6_P_vsathw_P(Word64 Rss) 553

vsatwuh
  Rd=vsatwuh(Rss)
    Word32 Q6_R_vsathwuh_P(Word64 Rss) 550
  Rdd=vsatwuh(Rss)
    Word64 Q6_P_vsathwuh_P(Word64 Rss) 553

vsplatb
  Rd=vsplatb(Rs)
    Word32 Q6_R_vsplatb_R(Word32 Rs) 557
  Rdd=vsplatb(Rs)
    Word64 Q6_P_vsplatb_R(Word32 Rs) 557

vsplath
  Rdd=vsplath(Rs)
    Word64 Q6_P_vsplath_R(Word32 Rs) 558

vspliceb
  Rdd=vspliceb(Rss,Rtt,#u3)
    Word64 Q6_P_vspliceb_PPI(Word64 Rss, Word64 Rtt, Word32 Iu3) 559
  Rdd=vspliceb(Rss,Rtt,Pu)
    Word64 Q6_P_vspliceb_PPp(Word64 Rss, Word64 Rtt, Byte Pu) 559

vsubb
  Rdd=vsubb(Rss,Rtt)
    Word64 Q6_P_vsubb_PP(Word64 Rss, Word64 Rtt) 407

vsubh
  Rd=vsubh(Rt,Rs)
    Word32 Q6_R_vsubh_RR(Word32 Rt, Word32 Rs) 170
  Rd=vsubh(Rt,Rs):sat
    Word32 Q6_R_vsubh_RR_sat(Word32 Rt, Word32 Rs) 170
  Rdd=vsubh(Rtt,Rss)
    Word64 Q6_P_vsubh_PP(Word64 Rtt, Word64 Rss) 405
  Rdd=vsubh(Rtt,Rss):sat
    Word64 Q6_P_vsubh_PP_sat(Word64 Rtt, Word64 Rss) 405

vsubub
  Rdd=vsubub(Rtt,Rss)
    Word64 Q6_P_vsubub_PP(Word64 Rtt, Word64 Rss) 407
  Rdd=vsubub(Rtt,Rss):sat
    Word64 Q6_P_vsubub_PP_sat(Word64 Rtt, Word64 Rss) 407

vsubuh
  Rd=vsubuh(Rt,Rs):sat
    Word32 Q6_R_vsubuh_RR_sat(Word32 Rt, Word32 Rs) 170
  Rdd=vsubuh(Rtt,Rss):sat
    Word64 Q6_P_vsubuh_PP_sat(Word64 Rtt, Word64 Rss) 405

vsubw
  Rdd=vsubw(Rtt,Rss)
    Word64 Q6_P_vsubw_PP(Word64 Rtt, Word64 Rss) 408
  Rdd=vsubw(Rtt,Rss):sat
    Word64 Q6_P_vsubw_PP_sat(Word64 Rtt, Word64 Rss) 408

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vsxtbh
  Rdd=vsxtbh(Rs)
  Word64 Q6_P_vsxtbh_R(Word32 Rs) 561

vsxthw
  Rdd=vsxthw(Rs)
  Word64 Q6_P_vsxthw_R(Word32 Rs) 561

vtrunehb
  Rd=vtrunehb(Rss)
  Word32 Q6_R_vtrunehb_P(Word64 Rss) 564
  Rdd=vtrunehb(Rss,Rtt)
  Word64 Q6_P_vtrunehb_PP(Word64 Rss, Word64 Rtt) 564

vtrunewh
  Rdd=vtrunewh(Rss,Rtt)
  Word64 Q6_P_vtrunewh_PP(Word64 Rss, Word64 Rtt) 564

vtrunohb
  Rd=vtrunohb(Rss)
  Word32 Q6_R_vtrunohb_P(Word64 Rss) 564
  Rdd=vtrunohb(Rss,Rtt)
  Word64 Q6_P_vtrunohb_PP(Word64 Rss, Word64 Rtt) 564

vtrunowh
  Rdd=vtrunowh(Rss,Rtt)
  Word64 Q6_P_vtrunowh_PP(Word64 Rss, Word64 Rtt) 564

vxaddsubh
  Rdd=vxaddsubh(Rss,Rtt):rnd:>>1:sat
  Word64 Q6_P_vxaddsubh_PP_rnd_rsl_sat(Word64 Rss, Word64 Rtt) 430
  Rdd=vxaddsubh(Rss,Rtt):sat
  Word64 Q6_P_vxaddsubh_PP_sat(Word64 Rss, Word64 Rtt) 430

vxaddsubw
  Rdd=vxaddsubw(Rss,Rtt):sat
  Word64 Q6_P_vxaddsubw_PP_sat(Word64 Rss, Word64 Rtt) 431

vxsubaddh
  Rdd=vxsubaddh(Rss,Rtt):rnd:>>1:sat
  Word64 Q6_P_vxsubaddh_PP_rnd_rsl_sat(Word64 Rss, Word64 Rtt) 430
  Rdd=vxsubaddh(Rss,Rtt):sat
  Word64 Q6_P_vxsubaddh_PP_sat(Word64 Rss, Word64 Rtt) 430

vxsubaddw
  Rdd=vxsubaddw(Rss,Rtt):sat
  Word64 Q6_P_vxsubaddw_PP_sat(Word64 Rss, Word64 Rtt) 431

vzxtbh
  Rdd=vzxtbh(Rs)
  Word64 Q6_P_vzxtbh_R(Word32 Rs) 565

vzxthw
  Rdd=vzxthw(Rs)
  Word64 Q6_P_vzxthw_R(Word32 Rs) 565

X

xor
  Pd=xor(Ps,Pt)

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Byte Q6\_p\_xor\_pp(Byte Ps, Byte Pt) **204**  
Rd=xor(Rs,Rt)  
Word32 Q6\_R\_xor\_RR(Word32 Rs, Word32 Rt) **158**  
Rdd=xor(Rss,Rtt)  
Word64 Q6\_P\_xor\_PP(Word64 Rss, Word64 Rtt) **343**  
Rx&=xor(Rs,Rt)  
Word32 Q6\_R\_xorand\_RR(Word32 Rx, Word32 Rs, Word32 Rt) **347**  
Rx^=xor(Rs,Rt)  
Word32 Q6\_R\_xorxacc\_RR(Word32 Rx, Word32 Rs, Word32 Rt) **347**  
Rx|=xor(Rs,Rt)  
Word32 Q6\_R\_xoror\_RR(Word32 Rx, Word32 Rs, Word32 Rt) **347**  
Rxx^=xor(Rss,Rtt)  
Word64 Q6\_P\_xorxacc\_PP(Word64 Rxx, Word64 Rss, Word64 Rtt) **345**

## Z

zxtb  
Rd=zxtb(Rs)  
Word32 Q6\_R\_zxtb\_R(Word32 Rs) **171**

zxth  
Rd=zxth(Rs)  
Word32 Q6\_R\_zxth\_R(Word32 Rs) **171**